Computer Structure and Language

8086/8088 Hardware Design

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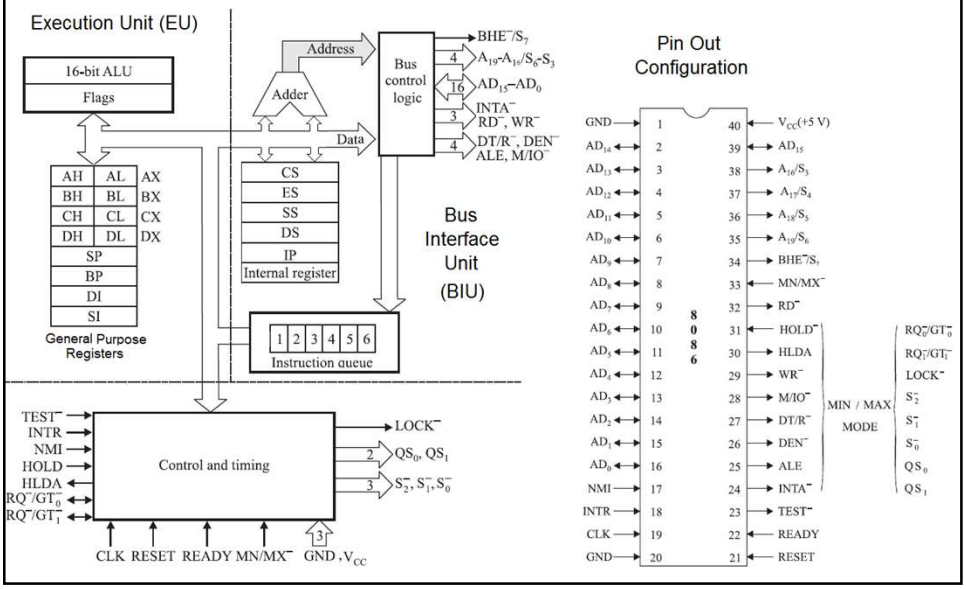
Intel's 8086 Pins/Signals



8086 Status Signals Encoding

Encoding of $\overline{BHE}$ and A_0

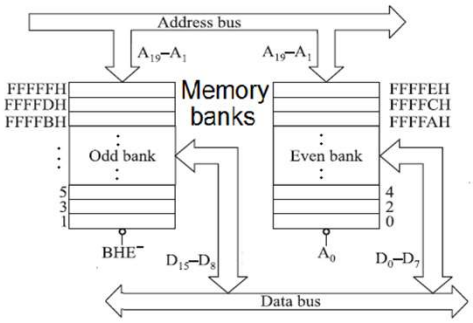
$\overline{BHE}$	A_0	Operation
0	0	Word (16-bit) will be access
0	1	Upper or odd byte will be access
1	0	Lower or even byte will be access
1	1	None

Encoding of S_4 and S_3

S_4	S_3	Segment in use
0	0	Alternate data (ES)
0	1	Stack (SS)
1	0	Code (CS) or none
1	1	Data (DS)

Encoding of S_2, S_1 and S_0

$\overline{S_2}$	$\overline{S_1}$	$\overline{S_0}$	Operation
0	0	0	Interrupt acknowledge
0	0	1	Read I/Q port
0	1	0	Write I/Q port
0	1	1	Halt
1	0	0	Code access
1	0	1	Read memory
1	1	0	Write memory
1	1	1	Passive



Encoding of QS_1 and QS_0

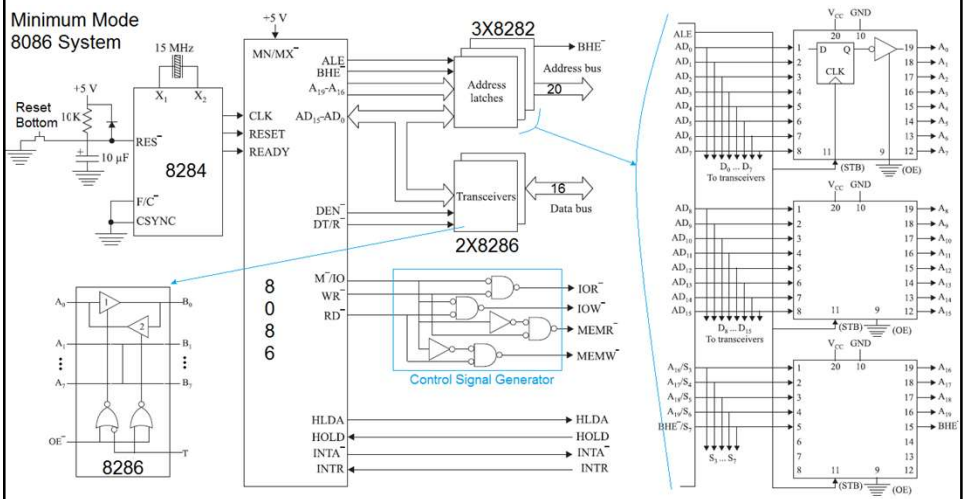
QS_1	QS_0	Characteristics
0	0	No operation
0	1	First byte of op-code from queue
1	0	Empty the queue
1	1	Subsequent byte from queue

S_5: Reports IF content to outside.

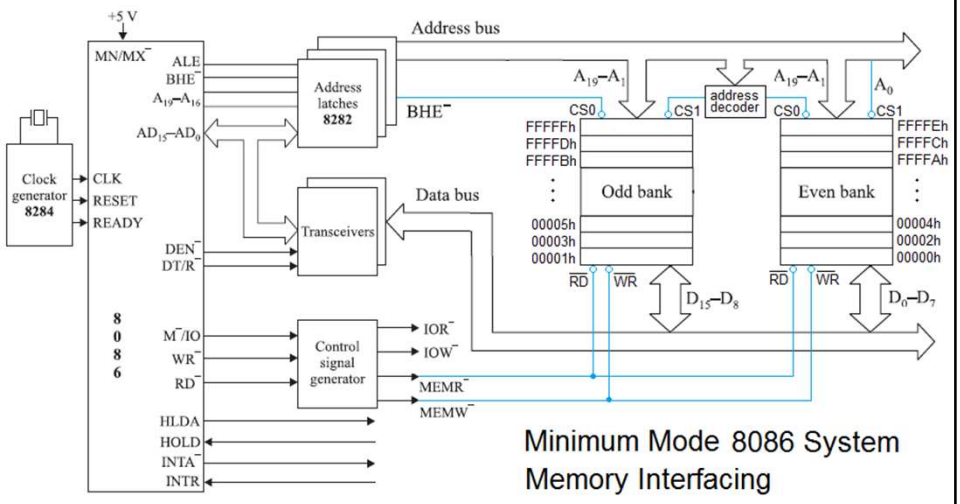
S_6: Indicates if 8086 is bus master. It is always low indicating that 8086 is bus master. If tristated, another bus master has taken control of the bus.

S_7: Used by 8087 to know if CPU is 8086 or 8088.

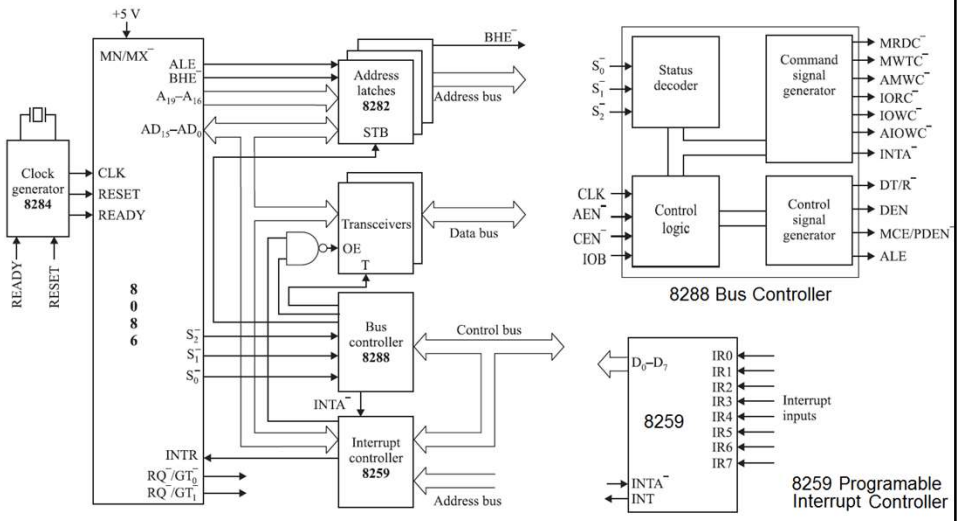
8086 System in Minimum Mode



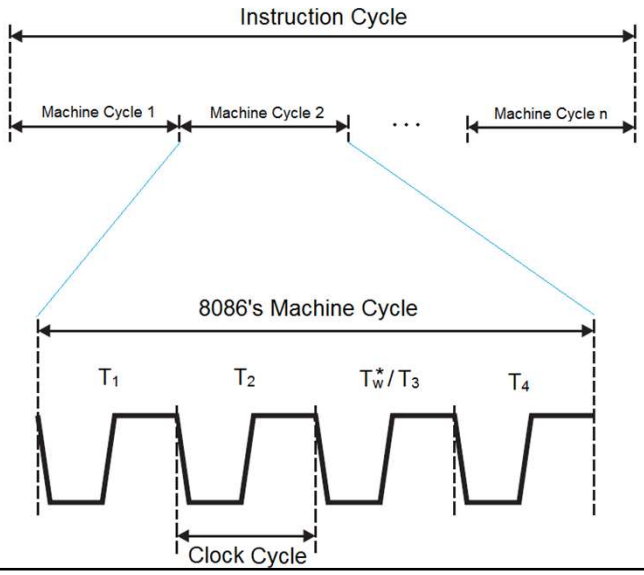
Memory Interface in Minimum Mode



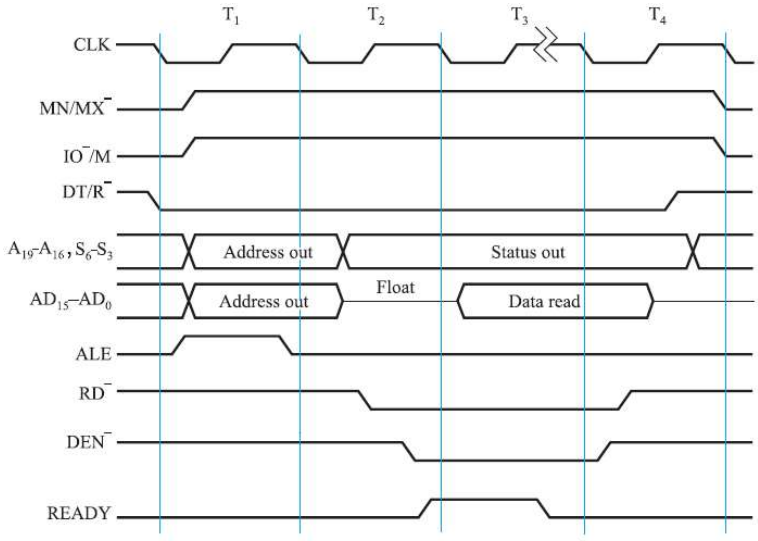
8086 System in Maximum Mode



Instruction, Machine (Bus), and Clock Cycles

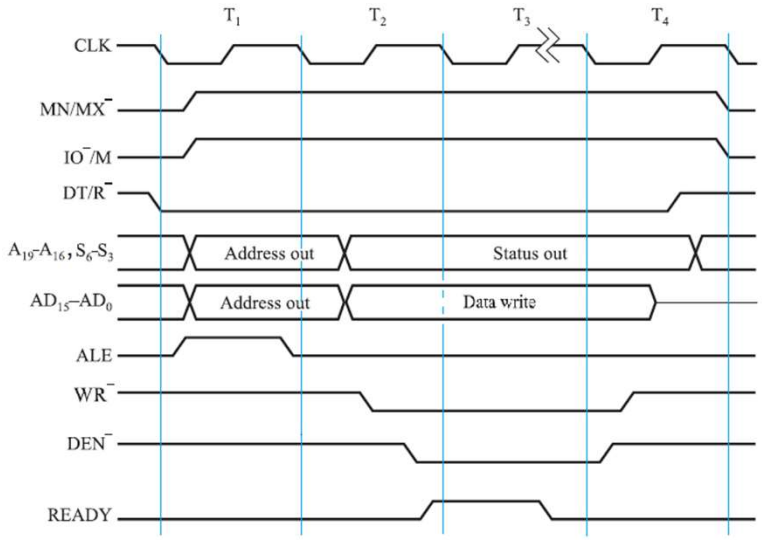


8086's Memory Read Cycle (Minimum Mode)



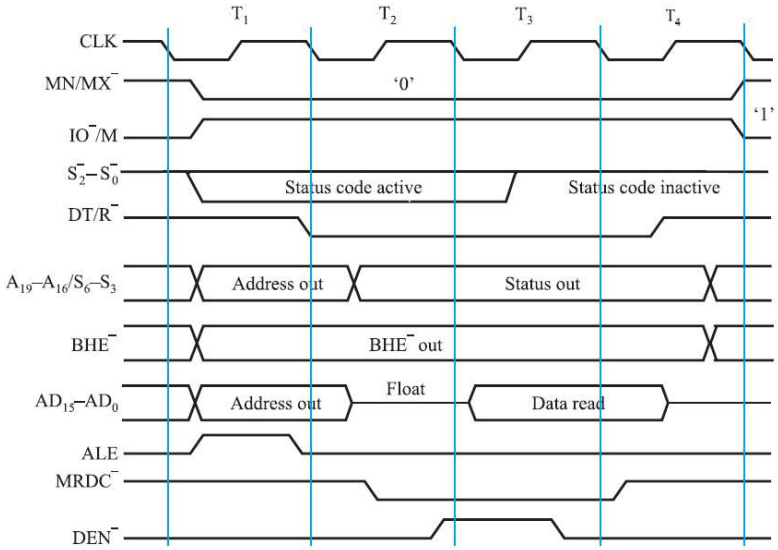
Timing diagram for 8086 minimum mode memory read

8086's Memory Write Cycle (Minimum Mode)



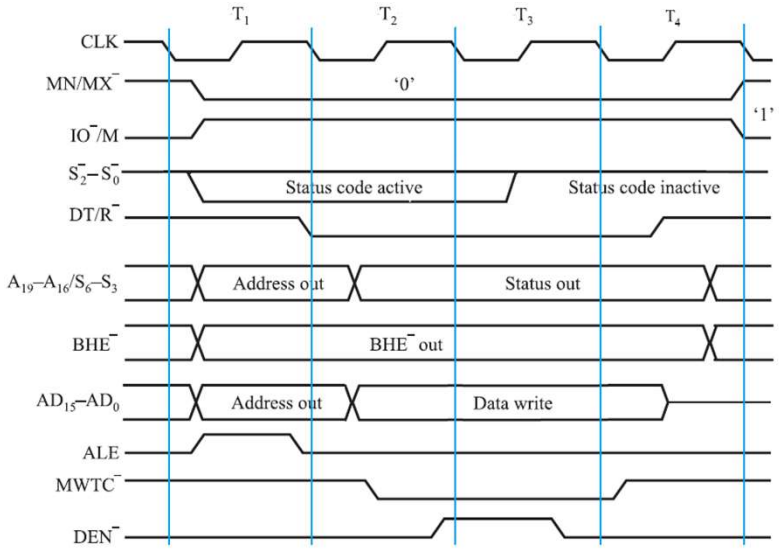
The timing diagram for 8086 minimum mode memory write

8086's Memory Read Cycle (Maximum Mode)



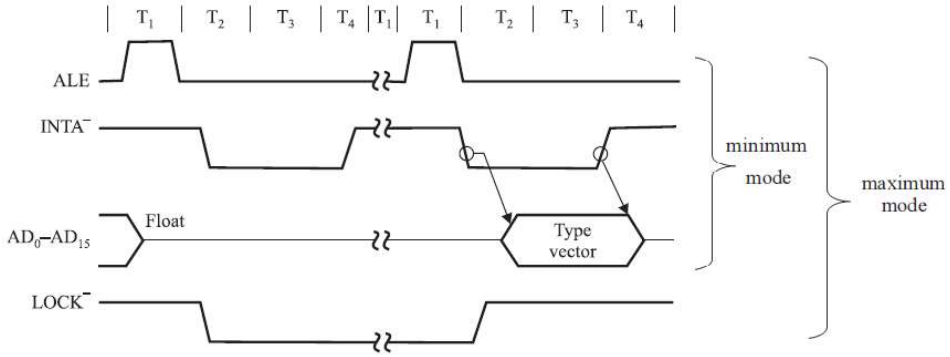
Maximum mode memory read bus cycle of 8086 system

8086's Memory Write Cycle (Maximum Mode)



Maximum mode memory write bus cycle of 8086 system

Interrupt Acknowledge Cycle



Interrupt acknowledge machine cycle

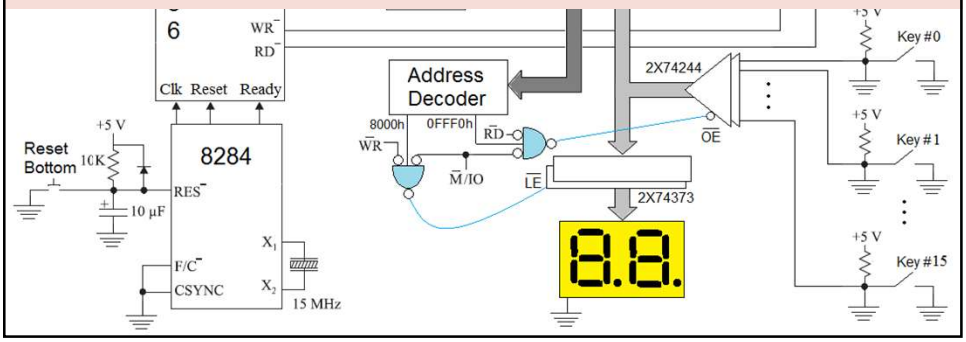
Input/Output Device Interfacing:

- **Memory Mapped:** I/O device is treated as a location of memory address space.
- **I/O Mapped:** I/O device is accessed in I/O address space.

Memory Mapped I/O device interfacing

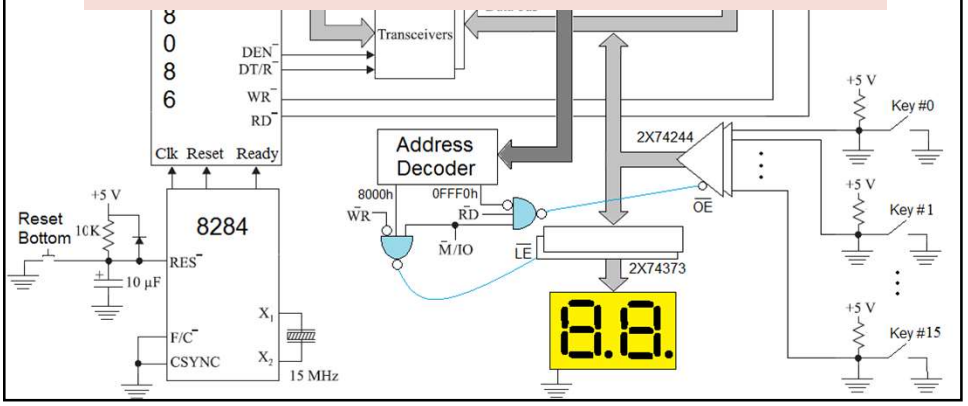
Assume (DS)=0:

```
mov  cx, word ptr [0fff0h];      == read from input device into cx
mov  word ptr [8000h], bx;    == write (bx) to the output device
add   dx, word ptr [0fff0h]      == read from input device & add to dx
```

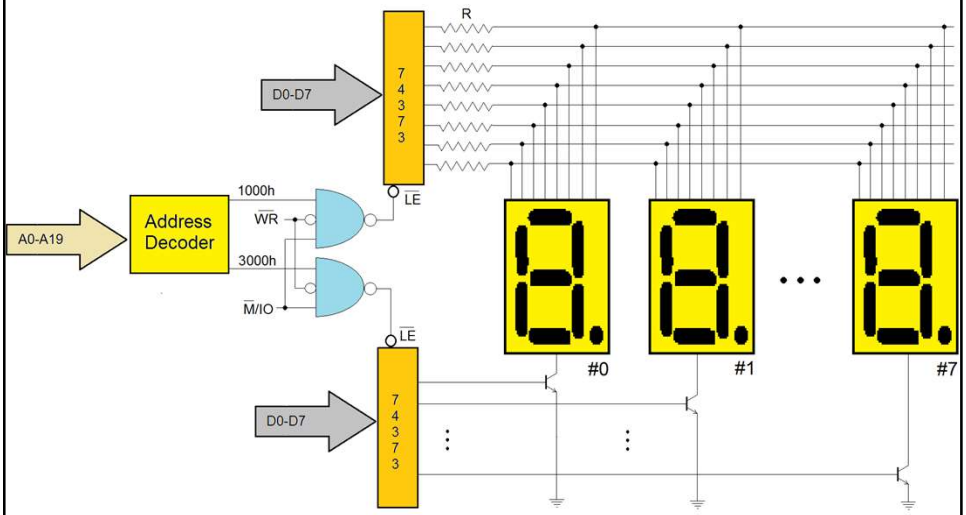


I/O Mapped I/O device interfacing

```
mov   dx,8000h
out   dx,ax ;     write (ax) to output device
mov   dx,0fff0h
in    ax,dx ;     read from input device into ax
```

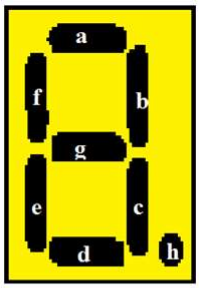


Design Example 1: Build a power-efficient 8-digit 7-segment display.



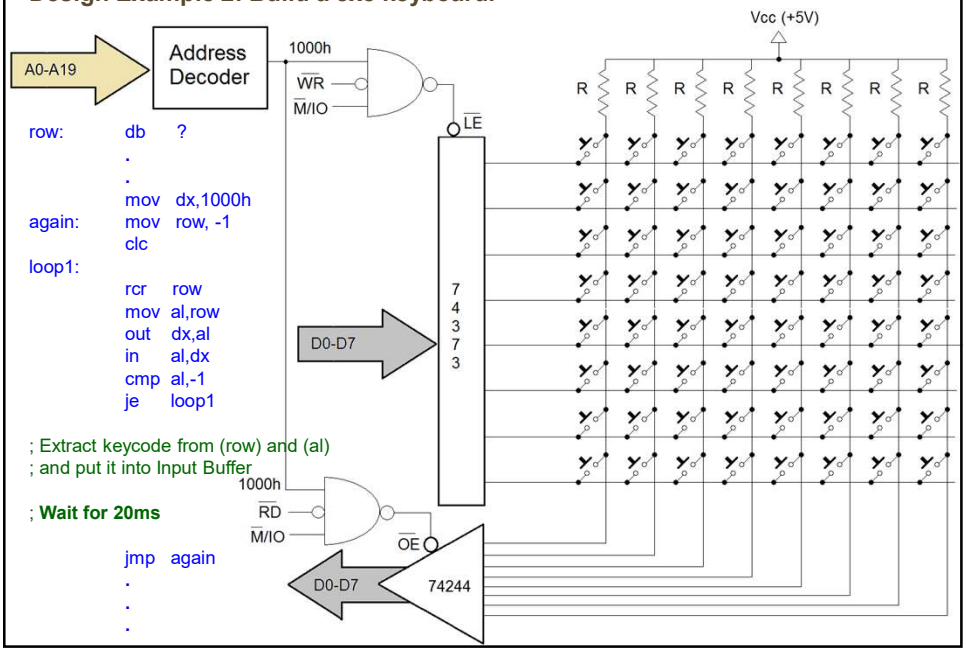
Design Example 1: Build a power-efficient 8-digit 7-segment display. (cont.)

```
display     proc        far

            %pushr      (bp,bx,ax,si,dx)

            mov         bp,sp
            lea         bx,seven_seg
            mov         si,[bp+14]
            mov         cx,8
next:       mov         al, byte ptr [si+7]
            xlat
            mov         dx,1000h
            out         dx,al
            mov         dx,3000h
            mov         al,order
            out         dx,al
            rol         order,1
            dec         si
            call        delay20ms
            loop        next

            %popr       (dx,si,ax,bx,bp)
            ret         2

seven_seg:  db      0,60h,0DAh,0F2h,66h,0B6h,0BEh,0E0h,0FEh,0F6h,1
    order:  db      0FEh

display     endp
```



digit	abcdefgh
0	11111100
1	01100000
2	11011010
3	11110010
4	01100110
5	10110110
6	10111110
7	11100000
8	11111110
9	11110110
.	00000001

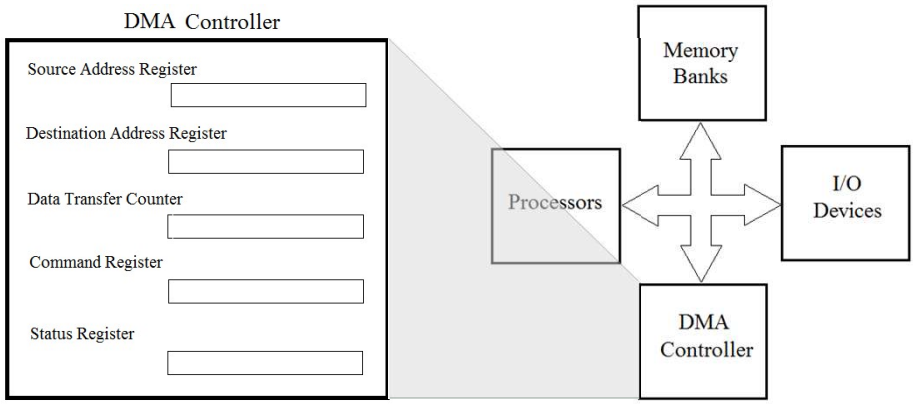
Design Example 2: Build a 8x8 keyboard.



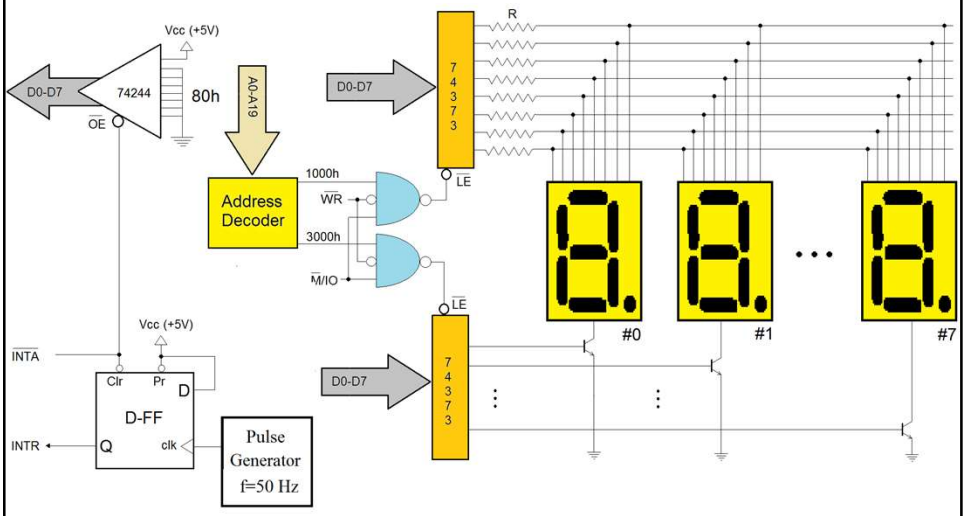

```
row:        db      ?
            .
            .
            mov     dx,1000h
again:      mov     row, -1
            clc
loop1:
            rcr     row
            mov     al,row
            out     dx,al
            in      al,dx
            cmp     al,-1
            je      loop1

; Extract keycode from (row) and (al)
; and put it into Input Buffer

; Wait for 20ms

            jmp     again
            .
            .
            .
```

I/O Schemes:

1. **Programmed I/O: CPU is involved in checking for data availability and for data movement.**
2. **Interrupted I/O: CPU is involved for data movement.**
3. **Direct memory Access (DMA): CPU is not involved in I/O operation.**



Design Example 1: Build a power-efficient 8-digit 7-segment display (Interrupted).

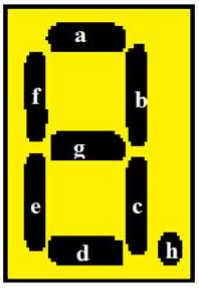


Design Example 1: Build a power-efficient 8-digit 7-segment display (cont).

```
display_buf:  db   8 dup (0)   ; a global buffer accessed by producer & displayer

displayer   proc      far     ; this is executed as timer  interrupt service routine

            %push     (flags, dx, ax, bx)  ; push used registers
            mov       si,index
            mov       al, byte ptr display_buf [si]
            lea       bx,seven_seg
            xlat
            mov       dx,1000h
            out       dx,al
            mov       dx,3000h
            mov       al,order
            out       dx,al
            rol       order,1
            dec       index
            and       index,7

            %pop      (bx,ax,dx,flags)    ; pop used registers
            iret

      index:  dw      7
seven_seg:    db      0,60h,0DAh,0F2h,66h,0B6h,0BEh,0E0h,0FEh,0F6h,1
      order:  db      0FEh
displayer   endp
```

In main program, at the beginning:

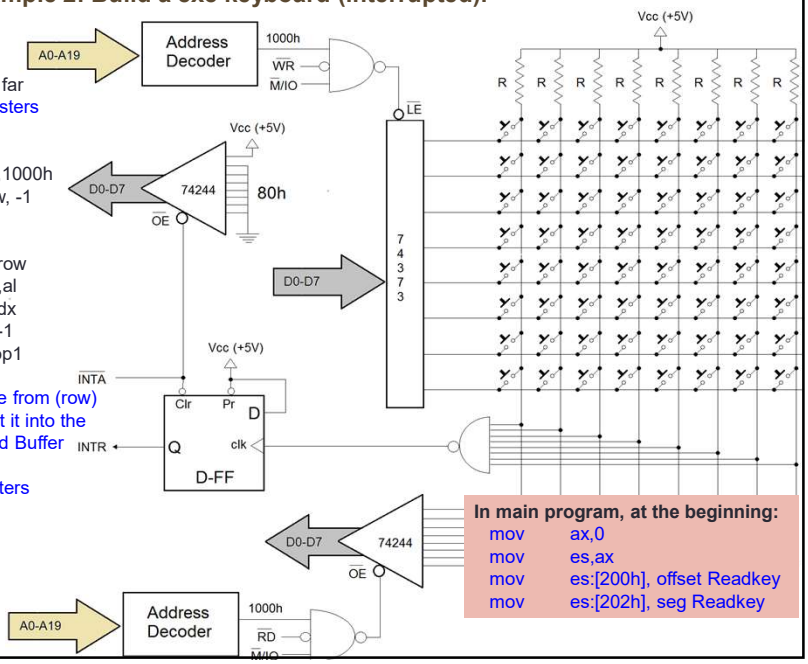
```
mov     ax,0
mov     es,ax
mov     es:[200h], offset displayer
mov     es:[202h], seg displayer
```



digit	abcdefgh
0	11111100
1	01100000
2	11011010
3	11110010
4	01100110
5	10110110
6	10111110
7	11100000
8	11111110
9	11110110
.	00000001

Design Example 2: Build a 8x8 keyboard (interrupted).

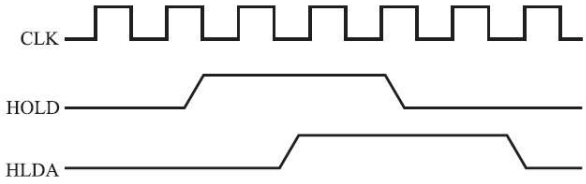
```
Readkey   proc  far
; Push used registers
; Wait for 20 ms

          mov   dx,1000h
again:    mov   row, -1
          clc
loop1:    rcr
          mov   al,row
          out   dx,al
          in    al,dx
          cmp   al,-1
          je    loop1

; Extract keycode from (row)
; and (al) and put it into the
; global Keyboard Buffer

; Pop used registers

          iret

Readkey   endp
```



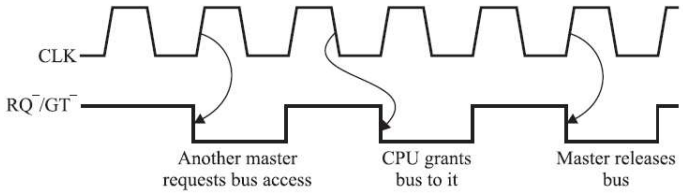
In main program, at the beginning:

```
mov     ax,0
mov     es,ax
mov     es:[200h], offset Readkey
mov     es:[202h], seg Readkey
```

Taking the control of the bus



Bus request and bus grant timings in minimum mode system



Bus request and bus grant timings in maximum mode system

8087 Coprocessor

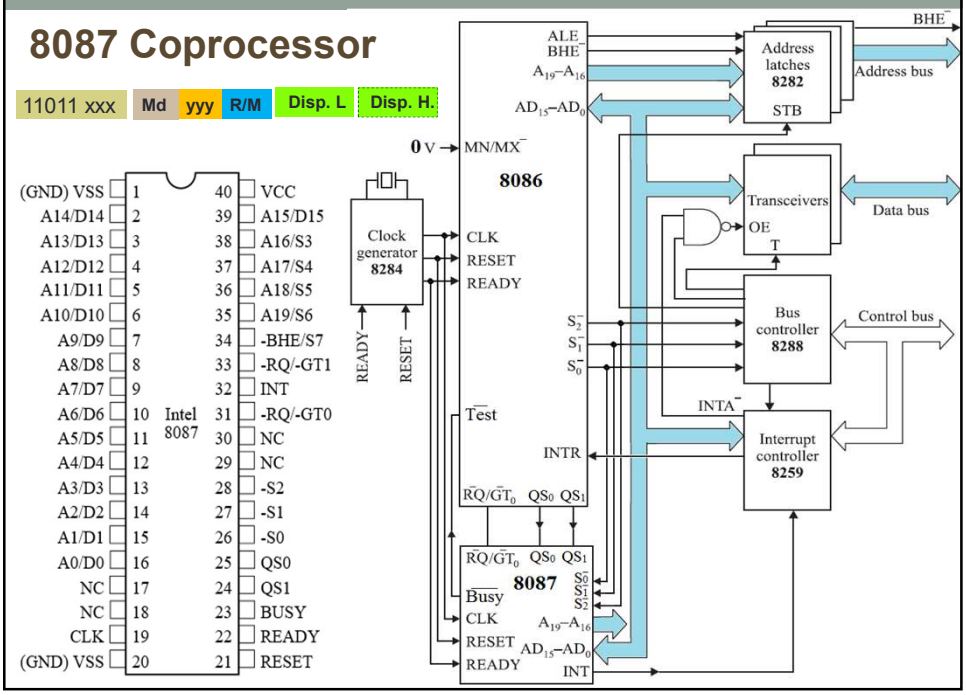
Introduced by Intel in 1980 to speed up computations for floating-point arithmetic, such as addition, subtraction, multiplication, division, and square root. It also realize exponential, logarithmic and trigonometric calculations, and besides floating-point numbers it could also operate on large binary and decimal integers.

Main Features:

- Floating point (32-bit short, 64-bit long, 80-bit extended) numbers
- 18 digit decimal numbers
- Eight 80-bit internal registers

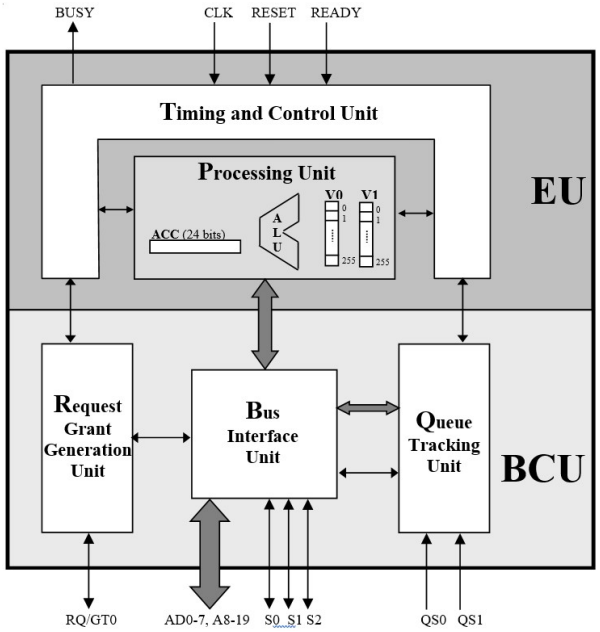
8087 Coprocessor

11011 xxx | Md | yyy | R/M | Disp. L | Disp. H.



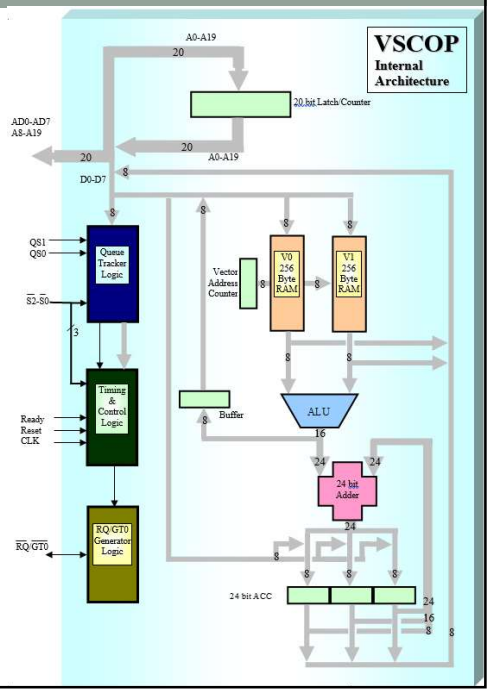
VSCOP:
Vector-Signal Coprocessor

Functional block diagram of VSCOP



VSCOP:
Vector-Signal Coprocessor

Internal Architecture

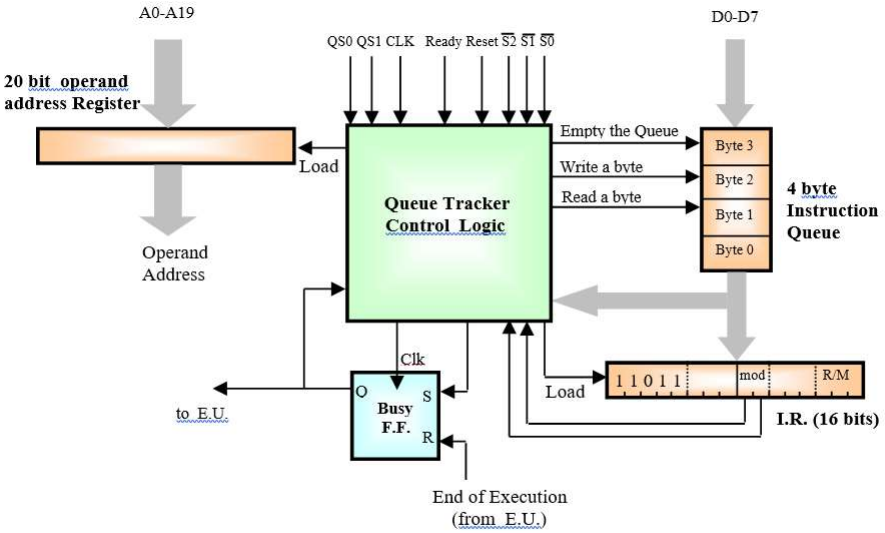


VSCOP: Vector-Signal Coprocessor

List of instructions

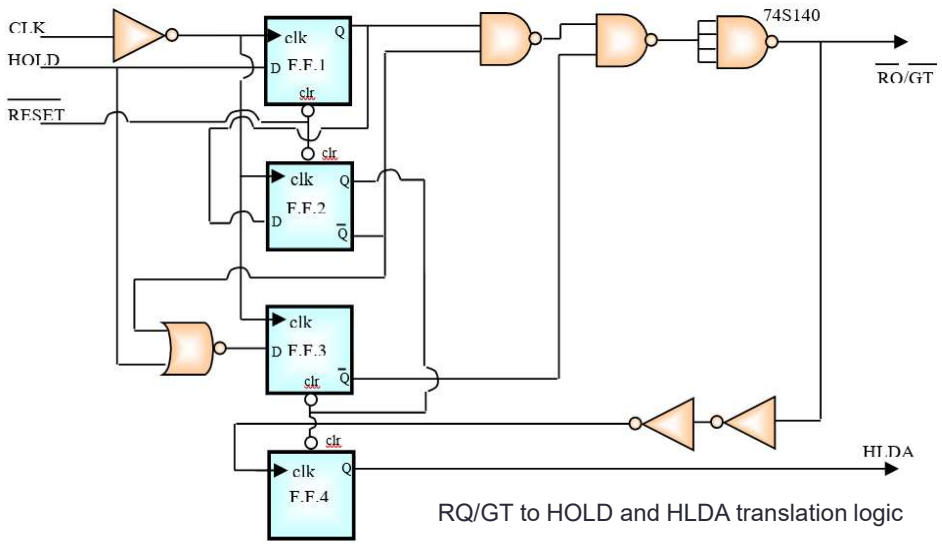
type	Instruction	Mnemonic	Code	Function	Execution time (clock cycle)
INTERNAL	Vector Clear	vclr V0 vclr V1	DBE8 DBE9	V0 ← 0 V1 ← 0	1024 1024
	Vector Addition	vadd V0 vadd V1	DBEA DBEB	V0 ← V0 + V1; ACC ← ACC + $\sum_{i=0}^{255} V0[i] + V1[i]$; V1 ← V0 + V1; ACC ← ACC + $\sum_{i=0}^{255} V0[i] + V1[i]$;	1024 1024
	Vector Subtraction	vsub V0 vsub V1	DBEC DBED	V0 ← V0 − V1; ACC ← ACC + $\sum_{i=0}^{255} V0[i] - V1[i]$; V1 ← V0 − V1; ACC ← ACC + $\sum_{i=0}^{255} V0[i] - V1[i]$;	1024 1024
	Vector Multiplication	vcon	DBEE/DBEF	ACC ← ACC + $\sum_{i=0}^{255} V0[i] * V1[i]$;	1024
MEMORY	Vector Load	vld V0,Addr vld V1,Addr	D9 mod* 001 r/m* DD mod 001 r/m	V0 ← Memory[Addr]; V1 ← Memory[Addr];	1024 1024
	Vector Store	vst V0,Addr vst V1,Addr	DB mod 001 r/m DF mod 001 r/m	Memory[Addr] ← V0; Memory[Addr] ← V1;	1024 1024
	Load Accumulator	ldacc Addr	DB mod 100 r/m	ACC ← Memory[Addr];	12
	Store Accumulator	stdacc Addr	DB mod 110 r/m	Memory[Addr] ← ACC;	12

VSCOP: Vector-Signal Coprocessor



A block diagram of the logical circuit for copying the four-byte instruction queue in 8088

VSCOP: Vector-Signal Coprocessor



RQ/GT to HOLD and HLDA translation logic

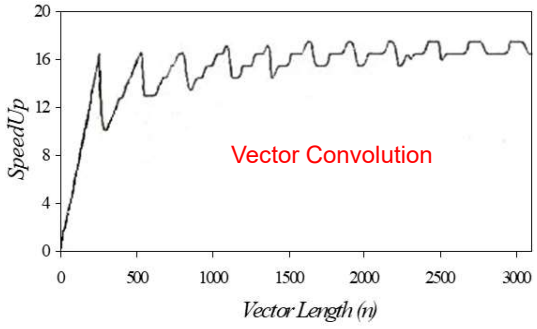
VSCOP:
Vector-Signal Coprocessor

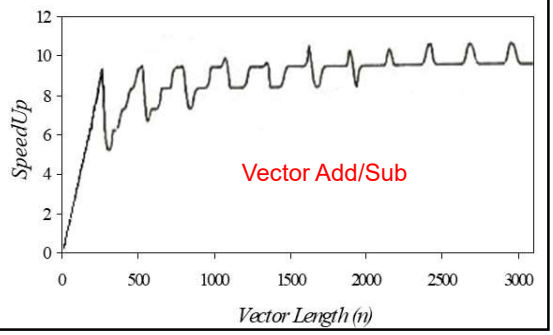
Discrete implementation using TTL MSI/LSI ICs

VSCOP:
Vector-Signal Coprocessor

Performance Evaluation

$$SpeedUp(F,n) = \frac{T_{8088}(F,n)}{T_{8088+VSCO}(F,n)}$$





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