

Best Response, Nash Equilibrium

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Example

	X	Y	Z
A	2, 1	0, 1	1, 0
B	0, 1	2, 1	1, 0
C	1, 1	1, 0	0, 0
D	1, 0	0, 1	0, 0

- Weakly Dominated strategy for player1?
- Weakly Dominated strategy for player2?

Example

	X	Y	Z
A	2, 1	0, 1	1, 0
B	0, 1	2, 1	1, 0
C	1, 1	1, 0	0, 0
D	1, 0	0, 1	0, 0

- Weakly Dominated strategy for player1? D is Weakly Dominated by A.
- Weakly Dominated strategy for player2? Z is Weakly Dominated by X or Y.

Example

- After elimination D,Z:

	X	Y
A	2, 1	0, 1
B	0, 1	2, 1
C	1, 1	1, 0

- Now, Weakly Dominated strategy for player1?
- Weakly Dominated strategy for player2?

Example

- After elimination D,Z:

	X	Y
A	2, 1	0, 1
B	0, 1	2, 1
C	1, 1	1, 0

- **Now**, Weakly Dominated strategy for player1? **No**
- Weakly Dominated strategy for player2? Y is **W.D** by X

Example

- After elimination Y:

	X
A	2, 1
B	0, 1
C	1, 1

- **Now**, Weakly Dominated strategy for player1?
- Weakly Dominated strategy for player2?

Example

- After elimination Y:

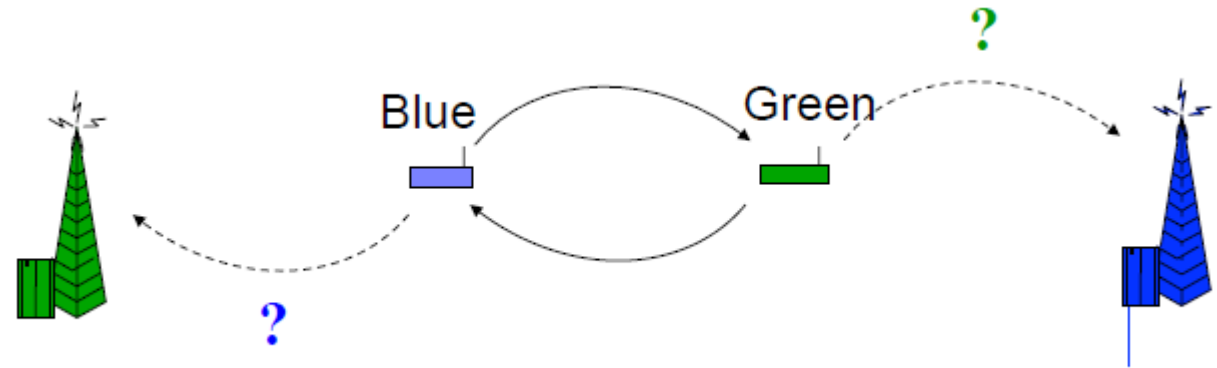
	X
A	2, 1
B	0, 1
C	1, 1

- Now, Weakly Dominated strategy for player1? **B is W.D by A, then C is W.D by A.**
- Weakly Dominated strategy for player2? **No**
- Finally, the reduced game is

	X
A	2, 1

Forwarder Game

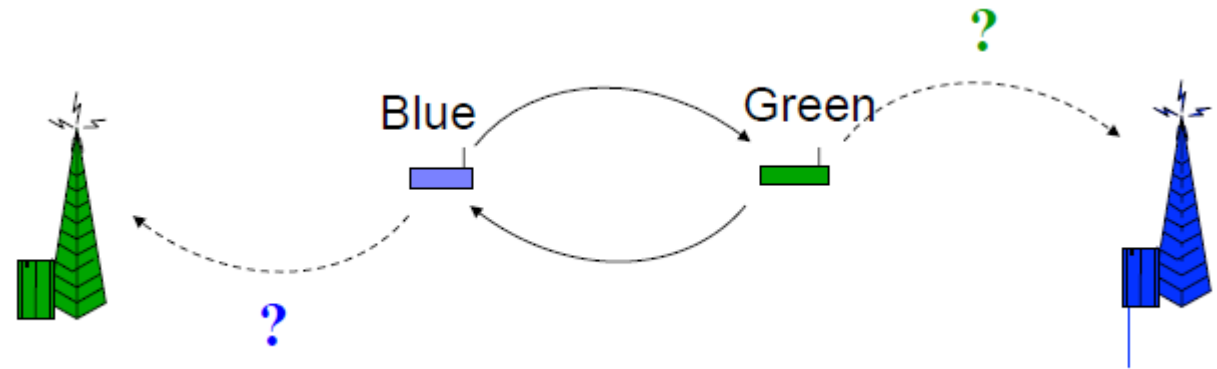
- Reward for packet delivering to destination: 1
- Cost of packet forwarding: c
- $0 < c \ll 1$



- Model this situation in normal form game.

Forwarder Game

- Reward for packet delivering to destination: 1
- Cost of packet forwarding: c
- $0 < c \ll 1$



- Model this situation in normal form game.

		Green	
		Forward	Drop
Blue	Forward	$(1-c, 1-c)$	$(-c, 1)$
	Drop	$(1, -c)$	$(0, 0)$

Forwarder Game

- Is it a symmetric game?

		Green	
		Forward	Drop
Blue	Forward	$(1-c, 1-c)$	$(-c, 1)$
	Drop	$(1, -c)$	$(0, 0)$

Forwarder Game

- Is it a symmetric game? **Yes**

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	Drop	$(1, -c)$	$(0, 0)$

Forwarder Game

- Is it a symmetric game? **Yes**
- What is best responses?
 - For player1:
 - For player2:

		Green	
		Forward	Drop
Blue	Forward	$(1-c, 1-c)$	$(-c, 1)$
	Drop	$(1, -c)$	$(0, 0)$

Forwarder Game

- Is it a symmetric game? **Yes**
- What is best responses?
 - For player1: $br_1(\text{Forward}) = \text{Drop}$ and $br_1(\text{Drop}) = \text{Drop}$
 - For player2: $br_2(\text{Forward}) = \text{Drop}$ and $br_2(\text{Drop}) = \text{Drop}$

		Green	
		Forward	Drop
Blue	Forward	$(1-c, 1-c)$	$(-c, 1)$
	Drop	$(1, -c)$	$(0, 0)$

Forwarder Game

- Is it a symmetric game? **Yes**
- What is best responses?
 - For player1: $br_1(\text{Forward}) = \text{Drop}$ and $br_1(\text{Drop}) = \text{Drop}$
 - For player2: $br_2(\text{Forward}) = \text{Drop}$ and $br_2(\text{Drop}) = \text{Drop}$
- Similar to witch previous games?

		Green	
		Forward	Drop
Blue	Forward	$(1-c, 1-c)$	$(-c, 1)$
	Drop	$(1, -c)$	$(0, 0)$

Forwarder Game

- Is it a symmetric game? **Yes**
- What is best responses?
 - For player1: $br_1(\text{Forward}) = \text{Drop}$ and $br_1(\text{Drop}) = \text{Drop}$
 - For player2: $br_2(\text{Forward}) = \text{Drop}$ and $br_2(\text{Drop}) = \text{Drop}$
- Similar to which previous games?
 - Prisoner's Dilemma

		Green	
Blue		Forward	Drop
Forward	$(1-c, 1-c)$	$(-c, 1)$	
Drop	$(1, -c)$	$(0, 0)$	

Best Response

Definition: Best Response

Player i 's strategy \hat{s}_i is a BR to strategy s_{-i} of other players if:

$$u_i(\hat{s}_i, s_{-i}) \geq u_i(s'_i, s_{-i}) \text{ for all } s'_i \text{ in } S_i$$

or

$$\hat{s}_i \text{ solves } \max u_i(s_i, s_{-i})$$

- Question: This definition is similar to which of the previous definitions?

Reminder

Definition: Best Response

Player i 's strategy \hat{s}_i is a BR to strategy s_{-i} of other players if:

$$u_i(\hat{s}_i, s_{-i}) \geq u_i(s'_i, s_{-i}) \text{ for all } s'_i \text{ in } S_i$$

or

$$\hat{s}_i \text{ solves } \max u_i(s_i, s_{-i})$$

Definition: Strict dominance

We say player i 's strategy s_i' is strictly dominated by player i 's strategy s_i if:

$$u_i(s_i, s_{-i}) > u_i(s_i', s_{-i}) \text{ for all } s_{-i}$$

Best response vs. Strict dominance

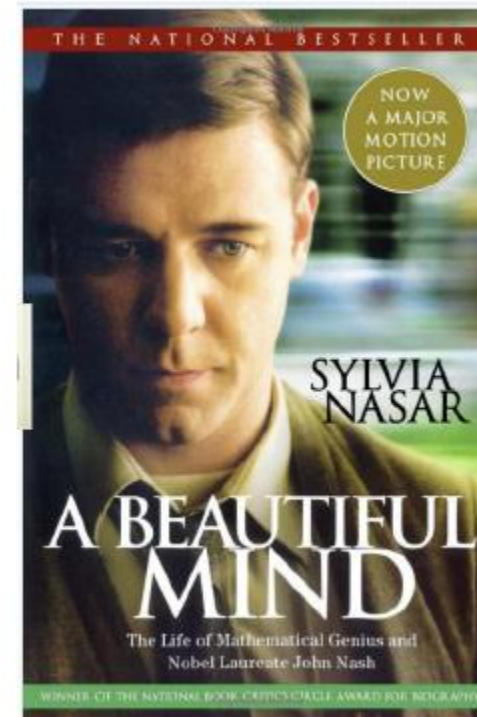
- $S_{-i} = [S_1, \dots, S_{i-1}, S_{i+1}, \dots, S_n]$
- $S = [S_{-i}, S_i]$
- S.D: for all other players strategies (S_{-i})
- BR: for all player i's strategies (S'_i in S_i)

Main Lesson

Rational players don't choose a strategy
that is **never** a **Best Response**!

About Nash

- John Nash (1928, 2015)
- Princeton Mathematics Department
- Economic Nobel prize at 1994
- Abel Prize at 2015



Nash Equilibrium

- All players simultaneously play best response to others

Definition (1): Nash Equilibrium

A strategy profile $(s_1^*, s_2^*, \dots, s_N^*)$ is a **Nash Equilibrium (NE)** if, for each i , her choice s_i^* is a best response to the other players' choices s_{-i}^*

Nash Equilibrium = Mutual best responses

Nash Equilibrium

Definition (2): Nash Equilibrium

At **Nash Equilibrium** no player can increase its payoff by deviating unilaterally.



No regret for every player!

Nash Equilibrium

Definition (3): Nash Equilibrium

Strategy profile s^* constitutes a **Nash Equilibrium** if, for each player i ,

Where: $u_i(s_i^*, s_{-i}^*) \geq u_i(s_i, s_{-i}^*), \forall s_i \in S_i$
 $u_i \in U$ utility function of player i
 $s_i \in S_i$ strategy of player i

Challenges

- Does any game have a Nash equilibrium?
- Is there a game with more than one Nash equilibrium?

Finding Nash Equilibrium

		Player 2		
		l	c	r
Player 1	U	0,4	4,0	5,3
	M	4,0	0,4	5,3
	D	3,5	3,5	6,6

- NE= ?

Finding Nash Equilibrium

		Player 2		
		l	c	r
Player 1	U	0,4	4,0	5,3
	M	4,0	0,4	5,3
	D	3,5	3,5	6,6

- Definition 1: What is best responses?

$$\begin{array}{ll}
 \diamond BR_1(l) = & BR_2(U) = \\
 \diamond BR_1(c) = & BR_2(M) = \\
 \diamond BR_1(r) = & BR_2(D) =
 \end{array}$$

Finding Nash Equilibrium

		Player 2		
		l	c	r
Player 1	U	0,4	4,0	5,3
	M	4,0	0,4	5,3
	D	3,5	3,5	6,6

- Definition 1: What is best responses?

$$\begin{array}{ll} \diamond BR_1(l) = M & BR_2(U) = l \\ \diamond BR_1(c) = U & BR_2(M) = c \\ \diamond BR_1(r) = D & BR_2(D) = r \end{array}$$

- NE = (D,r)

Finding Nash Equilibrium

		Player 2		
		l	c	r
Player 1	U	0,4	4,0	5,3
	M	4,0	0,4	5,3
	D	3,5	3,5	6,6

- **Definition2:** deviation unilaterally?
 - If $s^* = (U, l)$, then deviation is profitable?
 - If $s^* = (U, c)$, then deviation is profitable?
 - If $s^* = (U, r)$, then deviation is profitable?
 -
 - If $s^* = (D, c)$, then deviation is profitable?
 - If $s^* = (D, r)$, then deviation is profitable?

Finding Nash Equilibrium

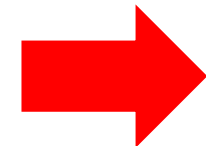
		Player 2		
		l	c	r
Player 1	U	0,4	4,0	5,3
	M	4,0	0,4	5,3
	D	3,5	3,5	6,6

- **Definition2:** deviation unilaterally?

- If $s^* = (U, l)$, then deviation is profitable? **Yes**
- If $s^* = (U, c)$, then deviation is profitable? **Yes**
- If $s^* = (U, r)$, then deviation is profitable? **Yes**
-

- If $s^* = (D, c)$, then deviation is profitable? **Yes**

- If $s^* = (D, r)$, then deviation is profitable? **No for each player**



NE

Finding Nash Equilibrium

- NE for Prisoner's Dilemma game?
- NE for Driving game (coordination game)?

	C	D
C	$-1, -1$	$-4, 0$
D	$0, -4$	$-3, -3$

	Left	Right
Left	$1, 1$	$0, 0$
Right	$0, 0$	$1, 1$

Examples

- NE for Prisoner's Dilemma game?
 - NE = (D,D)
- NE for Driving game (coordination game)?
 - NE = (Left, Left) and (Right, Right)

	<i>C</i>	<i>D</i>
<i>C</i>	-1, -1	-4, 0
<i>D</i>	0, -4	-3, -3

	Left	Right
Left	1, 1	0, 0
Right	0, 0	1, 1

Examples

- Pick a number game
- What is NE for n players?
 - $NE = (1, 1, \dots, 1)$

Examples

- Matching Penny game
 - One player wants to match, other player wants to mismatch
- What is NE?

	Heads	Tails
Heads	1, -1	-1, 1
Tails	-1, 1	1, -1

Examples

- Matching Penny game
- What is NE?
 - No pure NE

	Heads	Tails
Heads	1, -1	-1, 1
Tails	-1, 1	1, -1

Examples

- Battle of the sexes (BS) game
- What is NE?

	B	F
B	2, 1	0, 0
F	0, 0	1, 2

Examples

- Battle of the sexes (BS) game
- What is NE?
 - NE = (B, B) and (F, F)

	B	F
B	2, 1	0, 0
F	0, 0	1, 2

First Theorem

- Iterated Elimination of Strictly Dominated Strategies (IESDS)

If G is a finite game and solved by IESDS, then
the resulting outcome is unique NE of G .
(order independent)

Second Theorem

- Iterated Elimination of Weakly Dominated Strategies (IEWDS)

If G is a finite game and solved by IEWDS, then the resulting outcome is a Nash equilibria of G .

This outcome does **not** need to be **unique**.
(order dependent)

Third Theorem

- Iterated Elimination of Never Best Responses (IENBR)

If G is a finite game and solved by IENBR,
then the resulting outcome is **unique** NE.
(order independent)

Analyzing Games

- From the point of view of an outside observer, can some strategy profiles of a game be said to be **better** than others?
- Are there ways to prefer one profile to another?

Pareto Optimality

- Informal Definition

Definition

A strategy profile is pareto optimal if it is not possible to increase the payoff of any player without decreasing the payoff of another player.

Example

- Driving game
- Is there any pareto optimal?

	Left	Right
Left	1, 1	0, 0
Right	0, 0	1, 1

Example

- Driving game
- Is there any pareto optimal?
 - PO = (left,left) and (right,right)

	Left	Right
Left	1, 1	0, 0
Right	0, 0	1, 1

Example

- PD game
- Is there any pareto optimal?

	C	D
C	$-1, -1$	$-4, 0$
D	$0, -4$	$-3, -3$

Example

- PD game
- Is there any pareto optimal?
 - PO = (C,C) and (C,D) and (D,D)

	<i>C</i>	<i>D</i>
<i>C</i>	-1, -1	-4, 0
<i>D</i>	0, -4	-3, -3

Pareto Dominance

- Sometimes, one strategy profile s is at least as good for **every** agent as another profile s' , and there is **some agent** who **strictly** prefers s to s' .
- In this case, it seems reasonable to say that s is better than s' we say that s **Pareto-dominates** s' .

Pareto Optimality

- Formal Definition

Definition (Pareto Optimality)

A profile s^* is Pareto-optimal if there is no other profile that Pareto-dominates it.