# Report Lab Assignment Running x264 on Host Processor Extracting and Accelerating Kernel on Co-processor

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#### Samenvatting

During this lab, we are speeding up the x264 application by extracting the pixel\_satd\_8x4 and the pixel\_satd\_4x4 kernel for parallel processing on a  $\rho$ -VEX co-processor. This involves running the x264 application on a FPGA which runs a MicroBlaze operating system. The required instructions and data for the extracted kernels is being written to the instruction- and data memory of the  $\rho$ -VEX. This lab report describes in detail our chosen approach, problems we encountered along the way, test results measuring the execution time before and after extracting the kernels and an overal evaluation.

# 1 Wat moet er nog gebeuren

- debuggen van de rovex uitleggen met nep pixels en bewijs van werking
- Stride en manier van het schrijven naar het geheugen uitleggen
- Gebruik van lseek uitleggen
- uitleg van de kernel functie met plaatjes
- communicatie van pvex en microblaze
- big or little endiannes, en aanpak in code
- speedup & theoretical calculation speedup
- results of additional assignment and if more time which we would have chosen

## 2 Introduction

For the lab a computational intensive kernel has to be extracted from the  $\times 264$  software application: a free software library for encoding video stream into H.236/MPEG-4 AVC format. It is able to use Periodic Intra Refresh instead of keyframes as used by h.264, refreshing the image of the video by moving a column of intra blocks from one side of the screen to the other. This hides the refreshing effect from the user while the frame loads. The  $\times 264$  application is executed on a FPGA running a MicroBlaze host processor. By running the particular kernel on the  $\rho$ -VEX co-processor, the execution time can be decreased by making use of parallel processing.

# 3 Exploring the x264 application

Duing the first lab we got familiar with the x264 application, a powerfull coding tool to remove temporal redundancy by using motion estimation. In Section 3.1 we will discuss the various computational intenstive kernels within the x264 algorithm. The choosen kernel will be explained in Section 3.2, the provided memory layout of the FPGA architecture in Section 3.3 and the actual implementation will be discussed in Section 3.4.

# 3.1 Detecting the Computationally Most Intensive Kernel

A few input files are provided in the first lab to demonstrate the compile and run commands of the x264 application. The .y4m input files are decoded to create a short .mkv movie using this command in the terminal: ./x264 eledream\_64x32\_3.y4m -o testframe.mkv. By adding the gprof flag to the compile command, a list is created of all functions ordered by their share of the total execution time (in percentage). The number of function calls is also shown, as well as the total execution time for each input file. These five input .y4m files, vary in resolution (either 64x32 or 640x320 pixels) and amount of frames (either 1, 3, 8, 32 or 128). Profiling the x264 execution for the .y4m files that are provided by the lab leads to the ranking show in table 1.

Input file	Time (sec)	Share $(\%)$	$\mathbf{Calls}$	Kernel name
eledream_32x18_1.y4m	0.02	100.00	71	x264_analyse_init_costs
		0.00	1646	$x264$ _free
		0.00	784	$x264\_cabac\_encode\_desicion\_c$
eledream_64x32_3.y4m	0.03	66.67	71	x264_analyse_init_costs
		33.33	1971	$x264$ _pixel_satd_ $4x4$
		0.00	4830	$x264$ _pixel_satd_8x4
eledream_640x320_8.y4m	1.61	14.29	1599044	x264_pixel_satd_8x4
		11.80	570708	$x264\_get\_ref$
		4.97	38770	$x264$ _pixel_satd_ $x4$ _ $16x16$
eledream_640x320_32.y4m	8.54	20.61	7228633	get_ref
		13.23	15386501	$x264$ _pixel_satd_8x4
		4.57	1009690	$x264$ _pixel_satd_ $x4$ _8 $x8$
eledream_640x320_128.y4m	29.86	17.48	21956292	get_ref
		14.17	49862831	$x264$ _pixel_satd_8x4
		6.56	1023315	$x264$ _pixel_satd_ $x4$ _16 $x16$

Tabel 1: Chart with computationally most intensive kernels for each input stream.

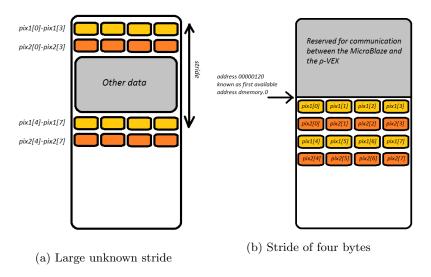
## 3.2 The pixel\_satd\_8x4 kernel

Given these statistics, we decide to extract the x264\_pixel\_satd\_8x4 kernel, since its share in execution time increases as the files become larger. This kernel evaluates the Sum of Transformed Differences (SATD) between a 8x4 pixel block from the input stream and reference blocks. The SATD is a metric used for video compression where the differences between the pixels are taken and put into a frequency transform, usually a Hadamard transform. These blocks can have various sizes depending on the inputs and constraint of the system varying from 4x4 pixel blocks to 8x4 blocks and even 16x16 blocks. The reference block with the lowest SATD value can be used to estimate motion in video coding to remove temporal redundancy.

## 3.3 Memory layout of the MicroBlaze

When writing pixels to the  $\rho$ -VEX for SATD calculation, these pixels need to be picked from the MicroBlaze memory. However, these pixels are not written as one block of contiguous bytes. Rows of four bytes

(MicroBlaze has 32-bit registers) are interrupted by a fixed amount of memory called a stride. This value is often given in bytes to be skipped in order to continue reading the particular data file. Figure 1a shows the concept of strides for two pixels from the MicroBlaze memory, which are splitted into two rows of 4 bytes since a pixels consists of eight bytes. When writing data to the  $\rho$ -VEX memory we use a stride of four bytes, resulting in a configuration as shown in 1b

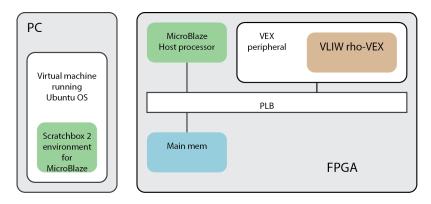


Figuur 1: MicroBlaze memory using strides to split up data files

# 3.4 Executing an Application on the Development Board and $\rho$ -VEX

When executing ./configure, a file is created for configuration of the application. However, the resulting config.mak file is made for applications running on the guest (Ubuntu). In order to configure for MicroBlaze, the config.mak file has to altered. All references to m32 have to be removed and the --DWORDS\_BIGENDIAN flag has to be added to the CFLAGS variable. This has to be done everytime when configurating the application for MicroBlaze

After doing this, the application now can be 'made' for MicroBlaze by first moving to the Scratchbox 2 environment for MicroBlaze and then execute the make command. Fig. 2 shows a block diagram of both the platforms.



Figuur 2: Block diagram of the Virtual Machine and the ERA platform

In order to run the application on the MicroBlaze host processor, the x264 executable must be compressed along with an input stream in a tar.gz file. This file can be put, via the FPGA host machine, on one of the three FPGAs using the scp, ftp and put command. For the programmer to connect to the development board, one must use the ssh and telnet command. See also fig 3.



Figuur 3: Ubuntu commands for putting files on and navigating to the FPGAs

For the first lab, the satd\_8x4 kernel has to be:

- extracted into a separate file
- supplemented to a proper .c file which can be compiled and run using the  $\rho$ -VEX
- included in the makefile that creates a bytecode and bytedata file of the kernel
- debugged

Looking back, the debugging part has been chasing us all the way through the lab. Not only did we suffer from bugs in our source code, the  $\rho$ -VEX itself did also have shortcomings that had to be circumvented by downloading several fixes, competing for FPGA availability and break downs of the entire host machine due to unsufficient capacity for the amount of students.

# 4 Approach

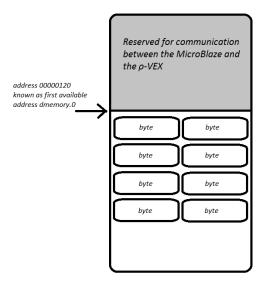
The kernel can run in the  $\rho$ -VEX by putting the bytecode generated by the makefile into the instruction memory of the  $\rho$ -VEX and placing all required data in its data memory dynamically. By executing make byte-command two files are created:

- bytecode, containing all the interructions to be executed by the  $\rho$ -VEX
- bytedata, containing the pixels of the input stream

In order to make the extracted kernel qualified for compilation and execution, a few things have to be altered. First, the new file has to be recognized by the makefile in the rovex-examples directory. Second, some type definitions have to be made. Originally, pixel\_satd\_8x4 resides in the pixel.c file of the x264 application. When extracting this kernel, all prior knowledge is lost and has to be defined again. Then, in order to make the kernel compile and run on the  $\rho$ -VEX, the development board has to be reset and started. Bytecode has to be written to the instruction memory (rex-imemory) and bytedata to the data memory (rvex-dmemory). Finally, the calculated result should be returned to the host. Figure 4 shows the  $\rho$ -VEX memory layout for our kernel.

#### 4.1 Adjusting the makefile

We added the kernel to the EXECUTABLES defined in the makefile. Other files defined here can be removed since we will not need them for our application. An important issue is the difference between logical memory and physical memory. While the first address of a logical memory is obviously zero, the physical memory can have the first part of the register being occupied by the operating system. For the  $\rho$ -VEX, the first address



Figuur 4: Memory layout of the  $\rho$ -VEX

which is allowed to be written to is 120. Thus, when writing the result of the kernel to the logical address '0', it is actually writing to the physical address '120' of the data memory.

Despite the fact that this is a rather simplistic operation, it took us some struggling to have this action confirmed as correctly. For example, we were told to remove the autoinline flag which resulted in a lot of wrong hexdumps. Also, half way the lab a fix had been made to eliminate the issue of the logical and physical addresses. Since the purpose of \_\_DATA\_START was now unclear we decided to remove all references to it, which felt like we had wasted lots of time. As it turned out, the fix had only affected the 'home' directory. In order to avoid conflicting files, we had created a separate folder named 'lab2' from where we executed the application. Due to this, the fix did not reach our code and thus did not eliminate the start address issue untill this problem was discovered very late into the project.

## 4.2 Constraints of the $\rho$ -VEX

Using the  $\rho$ -VEX as a co-processor comes with some limitations. First, the  $\rho$ -VEX supports a very old and basic C compiler, so using printf statements to find bugs in the extracted kernel file is not possible. Everytime the application is being run concurrently on the  $\rho$ -VEX, the status is being inspected by placing printf statements in the pixel.c file. Also, the  $\rho$ -VEX is very strict in the order of variable declarations and initialization. One of our bugs involved int i being declared and initialized, followed by a declaration of int result. This is not supported by the  $\rho$ -VEX, since it wants to have int i, result; both to be declared first, before initializing i = 4;.

When running an extracted kernel on the  $\rho$ -VEX, the kernel is unable to obtain information from previous code. Consequently, all parameters have to be re-defined. Table 2 shows all required type definitions.

## 4.3 Communication between the MicroBlaze and the $\rho$ -VEX

In order to delegate the satd\_8x4 kernel from the MicroBlaze to the  $\rho$ -VEX, both the environments need to communicate with each other. When executing the x264 application, the MicroBlaze has to load the instructions of the extracted kernel into the instruction memory of the  $\rho$ -VEX and the data (for which the SATD has to be evaluated) into the data memory. This is when the source code of x264 becomes involved. The pixel\_satd\_8x4 kernel, which is residing in the pixel\_c file of the application, needs to be adjusted. Instead of calculating the SATD, it should send the instructions and input to the  $\rho$ -VEX. The commands 5a

Parameter	Type definition	Motivation
intptr_t	unsigned int	This parameter represents the stride, which cannot be <0
pixel	unsigned char	Pixels are made out of bytes, which have 8 bits (like a char)
sum_t	short int	Same type definition as in the source code (16 bits)
sum2_t	long int	Same type definition as in the source code (32 bits)
BIT_DEPTH	defined as 8	Also in the source code, plus the kernel handles pixels (bytes)
BIT_PER_SUM	$(8 * sizeof(sum_t))$	Also predefined as in the source code, being 16

Tabel 2: Type definitions in the extracted kernel file.

and 5b are therefore added to the pixel.c file of the x264 application. These pixels are written at physical address 120, the first address on the  $\rho$ -VEX that is not set apart for communication between the driver and the  $\rho$ -VEX, defined as address 0. To be sure that the data is always written to the same location we use the lseek() function to specify the start location of data as 0x00 from the SEEK\_SET, or the beginning of the file

HOEVEEL TOEGEVOEGDE WAARDE HEEFT DIT? IS NAUWELIJKS LEESBAAR

```
char temp;
// Write the reset command to rvex
                                                                                                                          // Open the data register and check whether it has opened correc
// point data mgm to 120 on the more, address 0 on the microble
// write pixel data to memory written in a sequential manner to
// the data memory and start at 0
int data = open("dew/ruse_dmemory.0", O_REWR);
if (data == -1){
    printf("domain is open or error has occurred\n");
    return -1;
}
                                                                                                                                                                                                                                                        temp = '2';
write(ctrl, &temp, 1);
                                                                                                                                                                                                                                                         // Write the start command to rvex
 // Open the imem and fill it with bytecode
// Open the lamm and if if the ten systems
// Write the Dyfacode into the instruction memory
int instr = open("/dex/ryx-lmemory.o", O_WRONLT);
if (instr = -1){
    printf("\ninstr is already open or error has occured\n")
    return -1;
                                                                                                                                                                                                                                                        temp = '1';
write(ctrl, &temp, 1);
                                                                                                                                                                                                                                                         // Clean temp
                                                                                                                           if (lseek(data, 0, SEEK SET) == -1)
                                                                                                                                                                                                                                                         temp = 0;
printf("\nafter running code g the rovex");
                                                                                                                                 printf("\nset data location to start data memory failed");
                                                                                                                                                                                                                                                        // Wait for the rvex to finish by reading the status
 write(instr, codebuffer, 2000);
close(instr);
printf("\nafter writing bytecode");
                                                                                                                               or (i = 0; i < 4; i++, pix1 += i_pix1, pix2 += i_pix2)
                             (a) Bytecode
                                                                                                                                                                                                                                                     (c) Reset, set and get the status of
                                                                                                                                                           (b) Bytedata
                                                                                                                                                                                                                                                     the \rho-VEX
```

Figurr 5: Commands for reading from and writing to the  $\rho$ -VEX

To control the  $\rho$ -VEX from the x264 application, the control registers can be written to. By writing '2' to the control registers, the  $\rho$ -VEX is being reset. It can then be started by writing '1' to the register, telling the  $\rho$ -VEX to start running the instructions available in the instruction memory, calculating the satd of the two pixels. While calculating, the status of the  $\rho$ -VEX can be checked in a while-loop by reading the status variable of the status memory (rvex-smemory). When the satd kernel is finished, the result can be read from the data memory. See also figure 5c.

## 4.4 Result Hyphothesis

By having the computational intensive kernel being run concurrently on a co-processor, one would expect an overal speedup of the application. A problem is, however, that the bytecode and new data have to be sent to the  $\rho$ -VEX every time the kernel is being called. The data contains two new pixels of which the SATD has to be calculated, bytecode the kernel instructions. Because of this constant data traffic between the MicroBlaze and the  $\rho$ -VEX we don't think the intended speed up is achieved.

The reason that bytecode has to be sent every time the kernel is called, is that the three FPGAs are shared among a lot of students. When running their application concurrently, the instruction memory is constantly overwritten by another group. If there was a one-to-one setup, bytecode could have to been placed in main.c, written tot the imemory of the  $\rho$ -VEX once when starting the application.

# 5 Implementation

This section will cover some difficulties that came along during the lab. Section 5.1 will discuss our method for debugging and testing the working manner of the  $\rho$ -VEX, we will discuss some limitations we came accross in Section 5.2 and ??. Important features to be considered were the difference in Endianess as will be discussed in Section 5.3 and the reliability of the code, discussed in Section 5.4.

The approach was to first let the pixel.c write mock up pixel data to the data memory. The instruction memory would already contain the bytecode of our extracted kernel as it was put in manually. The kernel would only contain a single public declared variable pixel datamem[128];. This array of pixels points towards the first address of the data memory as no other variables are declared and thus there are no competing variables after this data memory location. The pixel.map file and the hexdump proved this to be correct. By declaring a pointer to this data memory location we could determine the result on the  $\rho$ -VEX and write the result into this array at a set location where the pixel.c can read and return it.

## 5.1 Mock up pixels

In order to get some insight in how to communicate with the  $\rho$ -VEX we made hexdumps of the data memory and created a mock up version of the satd function using 2 pixels we took from the original application. In order to check where the  $\rho$ -VEX is writing the data, a debugging file Debugging\_pixel.c is created. When using this debugging file, bytecode is manually put into the instruction memory using scp, ftp and put. The data that is being sent from the MicroBlaze are two pre-defined pixels. By knowing the expected result and the possibility to write easy recognizable hexadecimal numbers such as 0xdead, we can use the hexdump to examine the registers. The hexdump in figure 6 portrays all but the first two bytes or our mock up bytes. These first bytes display the correct result written to a convenient addres 0x00.

#### 5.2 Issues with the fix

Halfway the lab, a fix was introduced. The physical address was supposed to be set to 0x120 and did not have to be passed on using \_DATA\_START. In order to make the  $\rho$ -VEX read the pixels in the data memory correct a pointer to the first logical address 0 would be enough. However as we would later discover this was not the case. The fix we applied did not adjust our working directory as it was designed to alter the tools folder in the home directory and we were working in a seperate folder. When we did find out, we changed the script of the fix to match our own tools folder, but it did cause some days delay.

Figuur 6: Hexdump of the kernel run with two mock up pixels

#### 5.3 Endianness

The MicroBlaze and the  $\rho$ -VEX understand basic operations like open(), close(), read() and write(). However, they differ in Endiannes.  $\rho$ -VEX operates in Little Endian, meaning the bytes are read from right to left, as seen in figure 7a. MicroBlaze, on the other hand, operates in Big Endian and thus reads the bytes from left to right (see figure 7b). This becomes an issue when reading the result variable from the data

memory, as we will get the bytes in reversed order. This problem is solved by manually reversing the order of the result bytes, realized by the piece of code in figure 7c.

```
Byte[3] | Byte[2] | Byte[1] | Byte[0]

(a) Little Endian

Byte[0] | Byte[1] | Byte[2] | Byte[3]

(b) Big Endian

// write result in reverse endianess datamem[0x00] = (result >> 24) & 0xff; datamem[0x01] = (result >> 8) & 0xff; datamem[0x02] = (result >> 8) & 0xff; datamem[0x03] = result & 0xff;

(c) Reversing the order of bytes
```

. , \_ \_ \_

Figuur 7: Different reading direction due to Endianness

5.4 Reliability of the code

When writing pixels to the data memory of the  $\rho$ -VEX, the data pointer is automatically increased. To ensure that pixel.c reads the correct result from the memory location we can use lseek, a system call that is used to change the location of the read/write pointer of a file descriptor. Using lseek (data, 0, SEEK\_SET), the pointer in the  $\rho$ -VEX moves to the address that is the start address from his point of view (thus, physical address 0x120). Translating the logical address 0 to the physical address 0x120 is done by the linker and is not of our concern. We also added some if statements that ensure that if memory locations are not opened correctly, they return 0xff and print an error to the screen. These measures increase the reliability of the code.

# 6 Extracting a Second Kernel

For lab 3, a second kernel has to be extracted in order to achieve an even higher speed up. For this assignment, we wanted to choose our kernel more wisely. Given the constant competition for the use of the FPGAs, bytecode has to be written to the instruction memory everytime the kernel is being called. In combination with the data to be processed by the kernel, this leads to huge data traffic which is actually slowing down the application. This is why we wanted, at first, to pick a kernel that is not being called too often but requires a lot of heavy computation. A candidate for this would be the halfpel kernel, which has quite its share in execution time but not that many functions calls (halfpel zit helaas niet in de top 10 kernels, dus kan harde getallen noemen. Unfortunately, there is no time for writing an entire new file for the extraction of the halfpel kernel. Instead, we will pick a kernel similar to pixel\_satd\_8x4 which also have a large share in the execution time: pixel\_satd\_4x4.

#### 6.1 Extract pixel\_satd\_4x4

\*\*\* uitleg aanvulling rovex code \*\*\*

## 7 Evaluation

• How much speedup did you obtain?

- Is this what you had expected?
- Give a theoretical calculation for the speedup you should have expected and compare it to the practical result
- What were the results of the additional assignment? How did it affect the speed of the application?

# A Extracted kernels pixel\_satd\_8x4 and pixel\_satd\_4x4

```
\ensuremath{//} Defining types for the expected input arguments
// changed intptr_t to int instead of unsigned int and sum2 from us long to us int
typedef unsigned int intptr_t;
typedef unsigned char pixel;
typedef unsigned short sum_t;
typedef unsigned long sum2_t;
// Define an inputarray that points to the start of the data memory
pixel datamem[128];
// BIT_DEPTH is 8 always, so if else statement is removed
#define BIT_DEPTH 8
// #define BITS_PER_SUM (8 * sizeof(sum_t))
#define BITS_PER_SUM (8 * sizeof(sum_t))
// HADAMARD4 function exported from the original pixel.c file
#define HADAMARD4(d0, d1, d2, d3, s0, s1, s2, s3) {\
    sum2_t t0 = s0 + s1;
   sum2_t t1 = s0 - s1;
   sum2_t t2 = s2 + s3;
   sum2_t t3 = s2 - s3;
   d0 = t0 + t2;
   d2 = t0 - t2;
   d1 = t1 + t3; \setminus
    d3 = t1 - t3;
}
// abs function exported from the original pixel.c file
static sum2_t abs2(sum2_t a);
// x264-pixel_satd_8x4 exported from the original pixel.c file
int x264_pixel_satd_8x4();
// x264_pixel_satd_4x4 exported from the original pixel.c file
int x264_pixel_satd_4x4();
char main()
    // run the x264_pixel_satd_8x4 by passing on the pixel values.
    int result, i;
   result = 0;
    i = 0;
    // calculate the result
    if (datamem[0x40] == 0x44)
        result = x264-pixel_satd_4x4();
    else if (datamem[0x40] == 0x84)
        result = x264-pixel_satd_8x4();
```

```
else result = 0xdead;
    // clean result array
    for (i = 0; i < 4; i++)
        datamem[i] = 0x00;
    }
    // write result in reverse endianess
    datamem[0x00] = (result >> 24) & 0xFF;
    datamem[0x01] = (result >> 16) \& 0xFF;
    datamem[0x02] = (result >> 8) & 0xFF;
    datamem[0x03] = result & 0xFF;
    // DEBUG
    // datamem[0x40] = 0xfe;
    // datamem[0x41] = 0xed;
    return 0;
}
static sum2_t abs2(sum2_t a)
    sum2_t s = ((a)>(BITS_PER_SUM-1))&(((sum2_t)1<<BITS_PER_SUM)+1))*((sum_t)-1);
    return (a+s)^s;
}
int x264_pixel_satd_8x4()
{
        sum2_t tmp[4][4];
        sum2_t a0, a1, a2, a3;
    sum2_t sum = 0;
    int i = 0;
    // Strides are set values and do not have to be given as input arguments
    // Stride is 16 to compute 2 rows of 8 bytes of pixel array elements
    intptr_t stride = 16;
    // Pointer to data memory
    pixel* datapoint = datamem;
    // Adjust the for loop so it can fetch data from data memory
    for(i = 0; i < 4; i++, datapoint += stride)
    {
        a0 = (datapoint[0] - datapoint[8]) + ((sum2_t) (datapoint[4] - datapoint[12]) << BITS_PER_SUM);</pre>
        a1 = (datapoint[1] - datapoint[9]) + ((sum2_t)(datapoint[5] - datapoint[13]) << BITS_PER_SUM);</pre>
         \texttt{a2} = (\texttt{datapoint[2]} - \texttt{datapoint[10]}) + ((\texttt{sum2\_t})(\texttt{datapoint[6]} - \texttt{datapoint[14]}) << \texttt{BITS\_PER\_SUM}); 
        a3 = (datapoint[3] - datapoint[11]) + ((sum2_t)(datapoint[7] - datapoint[15]) << BITS_PER_SUM);
        HADAMARD4( tmp[i][0], tmp[i][1], tmp[i][2], tmp[i][3], a0,a1,a2,a3 );
    for(i = 0; i < 4; i++)
        HADAMARD4( a0, a1, a2, a3, tmp[0][i], tmp[1][i], tmp[2][i], tmp[3][i] );
        sum += abs2(a0) + abs2(a1) + abs2(a2) + abs2(a3);
    return (((sum_t)sum) + (sum>>BITS_PER_SUM)) >> 1;;
int x264_pixel_satd_4x4()
    sum2_t tmp[4][2];
    sum2_t a0, a1, a2, a3, b0, b1;
    sum2_t sum = 0;
```

```
int i = 0;
    \ensuremath{//} Strides are set values and do not have to be given as input arguments
// Stride is 8 to compute 2 rows of 4 bytes of pixel array elements
intptr_t stride = 8;
// Pointer to data memory
pixel* datapoint = datamem;
// instead of 8x4 method, it writes per 4 bytes
for (i = 0; i < 4; i++, datapoint +=stride)
    a0 = datapoint[0] - datapoint[4];
    a1 = datapoint[1] - datapoint[5];
    b0 = (a0+a1) + ((a0-a1) << BITS_PER_SUM);
    a2 = datapoint[2] - datapoint[6];
a3 = datapoint[3] - datapoint[7];
    b1 = (a2+a3) + ((a2-a3) << BITS_PER_SUM);
    tmp[i][0] = b0 + b1;
    tmp[i][1] = b0 - b1;
for( i = 0; i < 2; i++)
    {\tt HADAMARD4(a0,a1,a2,a3,tmp[0][i],tmp[1][i],tmp[2][i],tmp[3][i]);}\\
    a0 = abs2(a0) + abs2(a1) + abs2(a2) + abs2(a3);
    sum += ((sum_t)a0) + (a0>>BITS_PER_SUM);
return sum >> 1;
```