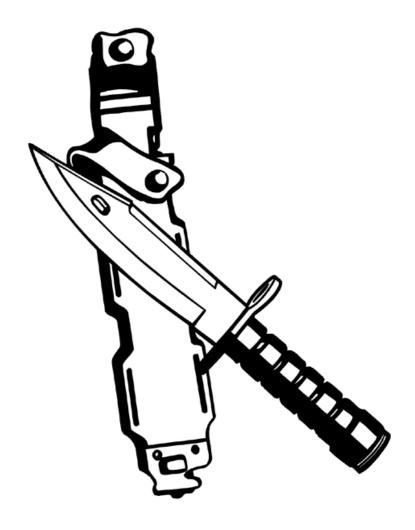
TiZi ACM-ICPC Notebook (2016)

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1 Data Structures

1.1 Union Find Disjoint Set - Kruskal

```
#include <iostream>
#include <vector>
#include <algorithm>
#include <sstream>
#include <string>
using namespace std;
typedef vector<int> vi;
struct DisJointSet {
         vi par, rnk, cnt; int numOfSets;
         DisJointSet(int n = 0) {
                  par.assign(n,-1); rnk.assign(n,0); cnt.assign(n,1); //par==parent
                   numOfSets=n; // if we wanna count number of disjoint sets
         int find(int a) {
                  int i=a, j=a, tmp;
                  while (par[i]!=-1) { i=par[i]; }
                  while (par[j]!=-1) { tmp=par[j]; par[j]=i; j=tmp; } //path compression
                  return i;
         int uni(int a, int b) {
   int A=find(a), B=find(b);
                  if(A!=B) {
                            if(rnk[A]<rnk[B]) swap(A,B); // union using rank</pre>
                            if (rnk [A] == rnk [B]) rnk [A] ++;
                            cnt[A] += cnt[B]; // if we wanna count each set size
                            numOfSets--; // if we wanna count number of disjoint sets
                  return cnt[A]; // if we wanna count each set size
};
struct Edge { int u, v, w;
    Edge(int u=0, int v=0, int w=0):u(u), v(v), w(w) {}
         bool operator<(const Edge& b) const { return w < b.w; }</pre>
         string toString(){
                  stringstream sstr;
                  sstr << u << "," << v << "," << w;
string str; sstr >> str;
                  return str;
};
typedef vector<Edge> ve;
struct Kruskal {
         ve edges; vi marked; DisJointSet st;
         Kruskal(int n, ve& edges):edges(edges) { st = DisJointSet(n); }
int run() { int result; sort(edges.begin(), edges.end());
                  for (int i=0 ; i < edges.size() ; i++) { Edge e = edges[i];
      if (st.find(e.u) != st.find(e.v)) {</pre>
                                     st.uni(e.u, e.v); result += e.w; marked.push_back(i);
                  return result;
         void printSelectedEdges() {
     cout << "MST edges:" << endl;</pre>
                  for(int i=0 ; i<marked.size() ; i++) {</pre>
                            Edge e = edges[marked[i]]:
                            cout << e.toString() << endl;</pre>
         } // remove
};
int main() {
         int n, m;
         ve edges;
         for (int i=0 ; i<m ; i++) {</pre>
                  int u, v, w;
                  cin >> u >> v >> w;
                  edges.push_back(Edge(u, v, w));
         for(int i=0 ; i<edges.size() ; i++) {</pre>
                  cout << edges[i].toString() << endl;</pre>
```

```
Kruskal kruskal.(n, edges);
cout << kruskal.run() << endl;
kruskal.printSelectedEdges();
}

/*
IN:
5 6
1 3 5
4 5 0
2 1 3
3 2 1
4 3 4
4 2 2

OUT:
4,5,0
3,2,1
4,2,2
2,1,3
*/</pre>
```

1.2 Segment Tree, RSQ, RMQ

/***************************

```
* Compilation: javac SegmentTree.java
   * Execution:
                                      java SegmentTree
   * A segment tree data structure.
   import java.util.Arrays:
class SegmentTree {
         private Node[] heap;
         private int[] array;
         private int size:
         public SegmentTree(int[] array) {
                  this.array = Arrays.copyOf(array, array.length);
                   //The max size of this array is about 2 * 2 \log_2(n) + 1
                  size = (int) (2 * Math.pow(2.0, Math.floor((Math.log((double) array.length) / Math.log(2.0)) +
                                  1)));
                  heap = new Node[size];
                  build(1, 0, array.length);
         public int size() {
                  return array.length;
          //Initialize the Nodes of the Segment tree
         private void build(int v, int from, int size) {
                  heap[v] = new Node();
                   heap[v].from = from;
                  heap[v].to = from + size - 1;
                  if (size == 1) {
                            heap[v].sum = array[from];
                           heap[v].min = array[from];
heap[v].minId = from;
                  } else {
                            //Build childs
                            build(2 * v, from, size / 2);
                            build(2 * v + 1, from + size / 2, size - size / 2);
                            heap[v].sum = heap[2 * v].sum + heap[2 * v + 1].sum;
                            heap[v].min = Math.min(heap[2 * v].min, heap[2 * v + 1].min); //min = min of the children
                            heap[v].minId = (heap[2 * v].min < heap[2 * v + 1].min ? heap[2 * v].minId : heap[2 * v + 1].min ? heap[2 * v].minId : heap[2 * v + 1].min ? heap[2 * v].minId : heap[2 * v + 1].min ? heap[2 * v].minId : heap[2 * v + 1].min ? heap[2 * v].minId : heap[2 * v + 1].min ? heap[2 * v].minId : heap[2 * v + 1].min ? heap[2 * v].minId : heap[2 * v + 1].min ? heap[2 * v].minId : heap[2 * v].minId : heap[2 * v + 1].min ? heap[2 * v].minId : heap[2 * v + 1].min ? heap[2 * v].minId : heap[2 * 
                                          1].minId);
         public int rsq(int from, int to) {
                   return rsq(1, from, to);
         private int rsq(int v, int from, int to) {
                  Node n = heap[v];
                   //If you did a range update that contained this node, you can infer the Sum without going down
                                   the tree
```

```
if (n.pendingVal != null && contains(n.from, n.to, from, to)) {
         return (to - from + 1) * n.pendingVal;
    if (contains(from, to, n.from, n.to)) {
        return heap[v].sum;
    if (intersects(from, to, n.from, n.to)) {
         propagate(v);
        int leftSum = rsq(2 * v, from, to);
int rightSum = rsq(2 * v + 1, from, to);
        return leftSum + rightSum:
public int rMinQ(int from, int to) {
    return rMinQ(1, from, to);
private int rMinQ(int v, int from, int to) {
    Node n = heap[v];
    //If you did a range update that contained this node, you can infer the Min value without
          going down the tree
    if (n.pendingVal != null && contains(n.from, n.to, from, to)) {
        return n.pendingVal;
    if (contains(from, to, n.from, n.to)) {
         return heap[v].min;
    if (intersects(from, to, n.from, n.to)) {
         propagate(v);
        int leftMin = rMinQ(2 * v, from, to);
int rightMin = rMinQ(2 * v + 1, from, to);
        return Math.min(leftMin, rightMin):
    return Integer.MAX_VALUE;
public int rMinIdQ(int from, int to) {
    return rMinIdQ(1, from, to);
private int rMinIdQ(int v, int from, int to) {
    Node n = heap[v]:
    //If you did a range update that contained this node, you can infer the Min value without
           going down the tree
    if (n.pendingVal != null && contains(n.from, n.to, from, to)) {
        return n.pendingVal;
    if (contains(from, to, n.from, n.to)) {
         return heap[v].minId;
    if (intersects(from, to, n.from, n.to)) {
         propagate(v);
         int leftMinId = rMinIdQ(2 * v, from, to);
        int rightMinId = rMinIdQ(2 * v + 1, from, to);
if (leftMinId == Integer.MAX_VALUE) return rightMinId;
if (rightMinId == Integer.MAX_VALUE) return leftMinId;
        return (array[leftMinId] < array[rightMinId] ? leftMinId : rightMinId);</pre>
    return Integer.MAX_VALUE;
public void update(int from, int to, int value) {
    update(1, from, to, value);
private void update(int v. int from, int to, int value) {
    //The Node of the heap tree represents a range of the array with bounds: [n.from, n.to]
    Node n = heap[v]:
    if (contains(from, to, n.from, n.to)) {
        change(n. value);
    if (n.size() == 1) return;
```

```
if (intersects(from, to, n.from, n.to)) {
         propagate(v);
         update(2 * v, from, to, value);
         update(2 * v + 1, from, to, value);
         n.sum = heap[2 * v].sum + heap[2 * v + 1].sum;
         n.min = Math.min(heap[2 * v].min, heap[2 * v + 1].min);
          \texttt{n.minId} = (\texttt{heap[2 * v].min} < \texttt{heap[2 * v + 1].min} ? \texttt{heap[2 * v].minId} : \texttt{heap[2 * v + 1].} 
//Propagate temporal values to children
private void propagate(int v) {
   Node n = heap[v];
    if (n.pendingVal != null) {
        change(heap[2 * v], n.pendingVal);
change(heap[2 * v + 1], n.pendingVal);
         n.pendingVal = null; //unset the pending propagation value
//Save the temporal values that will be propagated lazily
private void change(Node n, int value) {
    n.pendingVal = value:
    n.sum = n.size() * value;
    n.min = value;
    n.minId = n.from;
    array[n.from] = value;
//Test if the rangel contains range2
private boolean contains(int from1, int to1, int from2, int to2) {
    return from2 >= from1 && to2 <= to1;</pre>
//check inclusive intersection, test if rangel[from1, to1] intersects range2[from2, to2]
private boolean intersects(int from1, int to1, int from2, int to2) {
    return from1 <= from2 && to1 >= from2 // (.[..]..) or (.[..]..)
    || from1 >= from2 && from1 <= to2; // [.(..]..) or [.(.)..</pre>
//The Node class represents a partition range of the array.
static class Node {
    int sum;
    int min;
    int minId:
    //Here We store the value that will be propagated lazily
    Integer pendingVal = null;
    int from:
    int to:
    int size() {
        return to - from + 1:
public static void main(String[] args) {
    int[] a = new int[]{2, 3, 5, 1, 8, 4, 10};
    SegmentTree segmentTree = new SegmentTree(a);
    System.out.println(segmentTree.rsq(1, 4));
    System.out.println(segmentTree.rMinQ(1, 4));
    System.out.println(segmentTree.rMinIdQ(1, 4));
    segmentTree.update(3,3,4);
    System.out.println(segmentTree.rsq(1, 4));
    System.out.println(segmentTree.rMinQ(1, 4));
    System.out.println(segmentTree.rMinIdQ(1, 4));
```

1.3 Segment Tree, 2D-RMQ

```
import java.util.*;
public class SegmentTree2D {
  public static int max(int[][] t, int x1, int y1, int x2, int y2) {
    int n = t.length >> 1;
    x1 += n;
    x2 += n;
```

```
int m = t[0].length >> 1;
 y1 += m;
  y2 += m;
  int res = Integer.MIN_VALUE;
 for (int 1x = x1, rx = x2; 1x <= rx; 1x = (1x + 1) >> 1, rx = (rx - 1) >> 1)
   for (int 1y = y1, ry = y2; 1y <= ry; 1y = (1y + 1) >> 1, ry = (ry - 1) >> 1) {
     if ((lx & 1) != 0 && (ly & 1) != 0) res = Math.max(res, t[lx][ly]);
      if ((lx & 1) != 0 && (ry & 1) == 0) res = Math.max(res, t[lx][ry]);
     if ((rx & 1) == 0 && (ly & 1) != 0) res = Math.max(res, t[rx][ly]);
     if ((rx \& 1) == 0 \&\& (ry \& 1) == 0) res = Math.max(res, t[rx][ry]);
 return res;
public static void add(int[][] t, int x, int y, int value) {
 x += t.length >> 1;
 y += t[0].length >> 1;
  t[x][y] += value;
 for (int tx = x; tx > 0; tx >>= 1)
    for (int ty = y; ty > 0; ty >>= 1) {
     if (tx > 1) t[tx >> 1][ty] = Math.max(t[tx][ty], t[tx ^ 1][ty]);
     if (ty > 1) t[tx][ty >> 1] = Math.max(t[tx][ty], t[tx][ty ^ 1]);
public static void main(String[] args) {
 int[][] t = new int[sx * 2][sy * 2];
 add(t, x, y, v);//tree-x-y-value
int res1 = max(t, x1, y1, x2, y2);//t-[x1,y1]*[x2,y2]
```

1.4 Static RMQ, Lookup Table

```
/ keep code simple.
int lookup[MAX][LOGMAX];
struct Ouerv
  int L. R:
1:
void preprocess(int arr[], int n)
  // Initialize M for the intervals with length 1
  for (int i = 0; i < n; i++)</pre>
    lookup[i][0] = i;
  for (int j = 1; (1 << j) <= n; j++)
    for (int i = 0; (i + (1 << j) - 1) < n; i++)
      if (arr[lookup[i][j - 1]] < arr[lookup[i + (1 << (j - 1))][j - 1]])
lookup[i][j] = lookup[i][j - 1];</pre>
      else
         lookup[i][j] = lookup[i + (1 << (j - 1))][j - 1];
// Returns minimum of arr[L..R]
int query(int arr[], int L, int R)
  int j = (int) log2(R - L + 1);
   \textbf{if} \ (arr[lookup[L][j]] \mathrel{<=} arr[lookup[R - (\textbf{int})pow(2, j) + 1][j]]) \\
    return arr[lookup[L][j]];
  else return arr[lookup[R - (int)pow(2, j) + 1][j]];
void RMQ(int arr[], int n, Query q[], int m)
  // Fills table lookup[n][Log n]
  preprocess(arr, n);
  for (int i = 0; i < m; i++)
     // Left and right boundaries of current range
    int L = q[i].L, R = q[i].R;
// Print sum of current query range
    cout << "Minimum of [" << L << ",
      << R << "] is " << query(arr, L, R) << endl;
```

```
int main()
{
  int a[] = { 7, 2, 3, 0, 5, 10, 3, 12, 18 };
  int n = sizeof(a) / sizeof(a[0]);
  Query q[] = { { 0, 4 }, { 4, 7 }, { 7, 8 } };
  int m = sizeof(q) / sizeof(q[0]);
  RMQ(a, n, q, m);
  return 0;
```

1.5 Fenwick Tree, Inversions

```
#include <iostream>
#include <vector>
#include <algorithm>
#include <cstdio>
using namespace std;
typedef long long int64;
typedef vector<int64> vi;
// vector (vi), iostream, algo,
#define LSOne(i) (i & (-i))
struct FenwickTree {
        vi ft; FenwickTree() {}
        FenwickTree(int n) { ft.assign(n + 1, 0); } // init n + 1 zeroes
        int rsq(int b) { // returns RSQ(1, b), pass b >= 1
                 int sum = 0; for (; b; b -= LSOne(b)) sum += ft[b];
        int rsq(int a, int b) { // returns RSQ(a, b)
                 return rsq(b) - (a == 1 ? 0 : rsq(a - 1));
         // adjusts value of the k-th element by v (v can be +ve/inc or -ve/dec)
        void update(int k, int v) { // note: n = ft.size()
                 for (; k < (int)ft.size(); k += LSOne(k)) ft[k] += v;</pre>
1:
int main() {
        int f[] = { 2,4,5,5,6,6,6,7,7,8,9 }; // m = 11 scores
        FenwickTree ft(10); // declare a Fenwick Tree for range [1..10]
                                                  // insert these scores manually one by one into an
                                                          empty Fenwick Tree
        for (int i = 0; i < 11; i++) ft.update(f[i], 1); // this is O(k log n)</pre>
        printf("%d\n", ft.rsq(1, 1)); // 0 => ft[1] = (
printf("%d\n", ft.rsq(1, 2)); // 1 => ft[2] = 1
        printf("%d\n", ft.rsq(1, 6)); // 7 => ft[6] + ft[4] = 5 + 2 = 7
printf("%d\n", ft.rsq(1, 10)); // 11 => ft[10] + ft[8] = 1 + 10 = 11
        printf("%d\n", ft.rsq(3, 6)); // 6 => rsq(1, 6) - rsq(1, 2) = 7 - 1
        ft.update(5, 2); // update demo
        printf("%d\n", ft.rsq(1, 10)); // now 13
} // return 0;
/* extera
// get largest value with cumulative sum less than or equal to x;
// for smallest, pass x-1 and add 1 to result
int getind(int x) {// ***Change Needed***
  int idx = 0, mask = TREE_SIZE; //(must be a power of 2)
  while (mask && idx < TREE_SIZE)
    int t = idx + mask;
    if(x \ge tree[t]) \{idx = t; x = tree[t]; \}
    mask >>= 1:
  return idx:
// how to count inversions
int main(){
        while(cin >> n){ // count inversions, (Change Needed)
                 a.assign(n,0); b.assign(n,0); tree.assign(n,0);
                 for(int i=0; i<n; i++)
                         cin >> a[i]; b[i]=a[i];
                 sort(b.begin(),b.end());
                 for(int i=0 ; i<n ; i++) {
                         int rank=(int)(lower_bound(b.begin(),b.end(),a[i])-b.begin());
                         a[i]=rank+1:
                 int64 invs=0;//num of inversions
                 for(int i=n-1; i>=0; i--){
                         invs+=read(a[i]-1);
```

```
update(a[i],1);

}
cout << invs << endl;
}
}</pre>
```

1.6 Fenwick Tree, 2D

```
public class FenwickTree2D {
  public static void add(int[][] t, int r, int c, int value) {
    for (int i = r; i < t.length; i |= i + 1)
for (int j = c; j < t[0].length; j |= j + 1)</pre>
        t[i][j] += value;
  // sum[(0, 0), (r, c)]
  public static int sum(int[][] t, int r, int c) {
    for (int i = r; i >= 0; i = (i & (i + 1)) - 1)
      for (int j = c; j >= 0; j = (j & (j + 1)) - 1)
        res += t[i][j];
    return res;
  // sum[(r1, c1), (r2, c2)]
  public static int sum(int[][] t, int r1, int c1, int r2, int c2) {
    return sum(t, r2, c2) - sum(t, r1 - 1, c2) - sum(t, r2, c1 - 1) + sum(t, r1 - 1, c1 - 1);
  public static int get(int[][] t, int r, int c) {
    return sum(t, r, c, r, c);
  public static void set(int[][] t, int r, int c, int value) {
    add(t, r, c, -get(t, r, c) + value);
  // Usage example
  public static void main(String[] args) {
    int[][] t = new int[10][20];
    add(t, 0, 0, 1);
    add(t, 9, 19, -2);
    System.out.println(-1 == sum(t, 0, 0, 9, 19));
```

1.7 Fenwick Tree, Extended

```
public class FenwickTreeExtended {
  // T[i] += value
  public static void add(int[] t, int i, int value) {
    for (; i < t.length; i |= i + 1)</pre>
      t[i] += value;
  // sum[0..i]
  public static int sum(int[] t, int i) {
    int res = 0;
    for (; i >= 0; i = (i & (i + 1)) - 1)
      res += t[i]:
    return res:
  public static int[] createTreeFromArray(int[] a) {
    int[] res = new int[a.length];
    for (int i = 0; i < a.length; i++) {</pre>
      res[i] += a[i];
      int j = i | (i + 1);
if (j < a.length)</pre>
        res[j] += res[i];
    return res;
  // sum[a..b]
  public static int sum(int[] t, int a, int b) {
    return sum(t, b) - sum(t, a - 1);
```

```
public static int get(int[] t, int i) {
 int res = t[i];
 if (i > 0) {
   int lca = (i & (i + 1)) - 1;
   for (--i; i != lca; i = (i & (i + 1)) - 1)
     res -= t[i];
 return res;
public static void set(int[] t, int i, int value) {
 add(t, i, -get(t, i) + value);
// interval add
public static void add(int[] t, int a, int b, int value) {
 add(t, a, value);
 add(t, b + 1, -value);
// point query
public static int get1(int[] t, int i) {
 return sum(t, i);
// interval add
public static void add(int[] t1, int[] t2, int a, int b, int value) {
 add(t1, a, value);
 add(t1, b, -value);
 add(t2, a, -value * (a - 1));
 add(t2, b, value \star b);
// interval query
public static int sum(int[] t1, int[] t2, int i) {
 return sum(t1, i) * i + sum(t2, i);
// Returns min(p|sum[0,p]>=sum)
public static int lower_bound(int[] t, int sum) {
  --sum:
 int pos = -1;
 for (int blockSize = Integer highestOneBit(t.length); blockSize != 0; blockSize >>= 1) {
   int nextPos = pos + blockSize;
if (nextPos < t.length && sum >= t[nextPos]) {
      sum -= t[nextPos];
      pos = nextPos;
 return pos + 1;
// Usage example
public static void main(String[] args) {
 int[] t = new int[10];
 set(t, 0, 1);
 add(t, 9, -2);
 System.out.println(-1 == sum(t, 0, 9));
 t = createTreeFromArray(new int[] {1, 2, 3, 4, 5, 6});
 for (int i = 0; i < t.length; i++)</pre>
   System.out.print(get(t, i) + " ");
 System.out.println();
t = createTreeFromArray(new int[] {0, 0, 1, 0, 0, 1, 0, 0});
 System.out.println(5 == lower_bound(t, 2));
 int[] t1 = new int[10]:
 int[] t2 = new int[10];
 add(t1, t2, 0, 9, 1);
add(t1, t2, 0, 0, -2);
 System.out.println(sum(t1, t2, 9));
```

1.8 KDTree

```
import java.util.*;
public class KdTreePointQuery {
  public static class Point {
   int x, y;
   public Point(int x, int y) {
```

6

```
this.x = x;
   this.y = y;
int[] tx;
int[] ty;
public KdTreePointQuery(Point[] points) {
  int n = points.length;
  tx = new int[n];
  ty = new int[n];
 build(0, n, true, points);
void build(int low, int high, boolean divX, Point[] points) {
 if (low >= high)
   return;
  int mid = (low + high) >> > 1;
  nth_element(points, low, high, mid, divX);
  tx[mid] = points[mid].x;
  ty[mid] = points[mid].y;
 build(low, mid, !divX, points);
  build(mid + 1, high, !divX, points);
static void nth_element(Point[] a, int low, int high, int n, boolean divX) {
  while (true) {
    int k = randomizedPartition(a, low, high, divX);
   if (n < k)
    else if (n > k)
    else
      return;
static final Random rnd = new Random();
static int randomizedPartition(Point[] a, int low, int high, boolean divX) {
  swap(a, low + rnd.nextInt(high - low), high - 1);
  int v = divX ? a[high - 1].x : a[high - 1].y;
  int i = low - 1;
  for (int j = low; j < high; j++)
   if (divX ? a[j].x <= v : a[j].y <= v)</pre>
      swap(a, ++i, j);
  return i;
static void swap(Point[] a, int i, int j) {
 Point t = a[i];
  a[i] = a[j];
 a[j] = t;
long bestDist;
int bestNode;
public int findNearestNeighbour(int x, int y) {
  bestDist = Long.MAX_VALUE;
  findNearestNeighbour(0, tx.length, x, y, true);
  return bestNode;
void findNearestNeighbour(int low, int high, int x, int y, boolean divX) {
 if (low >= high)
   return:
  int mid = (low + high) >> > 1;
  long dx = x - tx[mid];
  long dy = y - ty[mid];
long dist = dx * dx + dy * dy;
  if (bestDist > dist) {
   bestDist = dist;
   bestNode = mid;
  long delta = divX ? dx : dy;
  long delta2 = delta * delta;
  if (delta <= 0) {
    findNearestNeighbour(low, mid, x, y, !divX);
    if (delta2 < bestDist)</pre>
      findNearestNeighbour(mid + 1, high, x, y, !divX);
    findNearestNeighbour(mid + 1, high, x, y, !divX);
    if (delta2 < bestDist)</pre>
      findNearestNeighbour(low, mid, x, y, !divX);
```

```
public static void main(String[] args) {
  Point[] points = new Point[n];
  //fill points
  //build tree
  KdTreePointQuery kdTree = new KdTreePointQuery(points);
  int index = kdTree.findNearestNeighbour(qx, qy);
  Point p = points[index];
}
```

2 Dynamic Programming

2.1 Classic DP problems

```
// ---- Maximum Subrectangle Sum
int main(){
        for(int i=1 ; i<n ; i++)//preprocess</pre>
                                  for(int j=0 ; j<n ; j++)</pre>
                                          a[i][j]+=a[i-1][j];
        int Max=0, ans=0;
for(int k=0; k<n; k++){//calc</pre>
                 for(int i=0 ; i<n-k ; i++) { Max=0;</pre>
                         for(int j=0; j<n; j++) {
                                  if(Max<0) Max=a[i+k][j]-a[i][j];</pre>
                                  else Max+=a[i+k][j]-a[i][j];
                                  if (Max>ans) ans=Max;
        //sub array, finsh and start point p=(val, startidx, finishidx)
         p ans=p(-1,0,0); int sum=0,id=1;
        for(int i=1 ; i<n ; i++) {</pre>
                                  if(sum<0) {sum=0; id=i;}</pre>
                                  sum+=a[i]:
                                  p tmp=p(sum,id,i+1); ans=Max(ans,tmp);
// ---- Optimal Array Multiplication Sequence (Print Path)
int n,a[10+5],p[10+5][10+5],dp[10+5][10+5];
int solve(int L, int R) {
        if(L==R) { return 0; }
        if (dp[L][R]!=-1) return dp[L][R];
        int Min=INF;
        for(int i=L ; i<R ; i++) {</pre>
                 int slv=solve(L,i)+solve(i+1,R)+a[(L-1)]*a[i]*a[R];
                 if(Min>slv)
                                 Min=slv; p[L][R]=i;
        return dp[L][R]=Min;
//prints like this => (A1 x (A2 x A3))
void print(int L, int R) {
        if(L==R) { cout << "A" <<L; return; }</pre>
        cout << "("; print(L,p[L][R]);</pre>
        cout << " x ";
        print(p[L][R]+1,R); cout << ")";
int main(){
               int t=1:
        while (cin >> n && n) {
                 for(int i=1; i<=n; i++)cin >> a[i-1] >> a[i];
                 memset (dp,-1,sizeof dp);
solve(1,n);//cout << solve(1,n) << endl;</pre>
                 printf("Case %d: ",t++); print(1,n); printf("\n");
        return 0:
// ---- LIS
int main(){
        vector<int> v;
        v.push_back(inf);
        for (int i = 0; i<n; i++) {</pre>
                 int x = dolls[i].w; // array element
                 int id = lower_bound(v.begin(), v.end(), x + 1) - v.begin();
                 if (id == v.size() - 1) v.push back(inf); v[id] = x;
        cout << v.size() - 1 << endl;
```

```
// ---- LCS
int main(){
        dp[MAX][MAX] = {0};
        for (int i=1 ; i<=n ; i++) {</pre>
                                  for(int j=1 ; j<=n ; j++) {</pre>
                                          if(a[i-1]==b[j-1]) dp[i][j]=dp[i-1][j-1]+1;
                                          else dp[i][j]=max(dp[i-1][j],dp[i][j-1]);
        cout << dp[n][n] << endl;</pre>
// ---- TSP
p a[15]; int n, dp[15][1<<15];</pre>
int solve(int pos, int bitset){
        int& dpp=dp[pos][bitset]; //dpp = dp poniter
        if (bitset==(1<<n)-1) return dist(a[pos],a[0]);</pre>
        if(dpp!=-1) return dpp;
        for (int i=0 ; i<n ; i++) {
                if(!(bitset&(1<<i))) dpp=min(dpp,solve(i,bitset|(1<<i))+dist(a[pos],a[i]));</pre>
        return dpp:
int main(){
        int to: cin >> to:
        while(tc--){
                cin >> a[0].X >> a[0].Y; cin >> n; n++;
                for(int i=1; i<n; i++) cin >> a[i].X >> a[i].Y;
                memset (dp, -1, sizeof dp);
                cout << solve(0,1) << endl;
        return 0;
```

2.2 Edit Distance

```
#include <iostream>
#include <cstdio>
#include <cstring>
#include <algorithm>
#include <string>
#define MAX (100+10)
using namespace std;
string str1,str2; int d[MAX][MAX]; char path[MAX][MAX];
int Step, DelCnt, InsCnt;
void PrintPath(int i, int j) {
        //cout << path[i][j] << endl;
if(path[i][j]==0) PrintPath(i-1, j-1);
        else if(path[i][j]=='d'){
                 PrintPath(i-1, j);
                 printf("%d Delete %d\n", Step++, i + InsCnt - DelCnt);
                 DelCnt++;
        else if(path[i][j]=='i'){
                 PrintPath(i,j-1);
                 printf("%d Insert %d,%c\n", Step++, j, str2[j]);
                 PrintPath(i-1, j-1);
                 printf("%d Replace %d,%c\n", Step++, i + InsCnt - DelCnt, str2[j]);
int solve(){
         memset(d, 0, sizeof d); memset(path, 0, sizeof path);
        for(int i=1 ; i<=str1.length() ; i++) {</pre>
                 d[i][0]=i; path[i][0]='d';
         for(int j=1; j<=str2.length(); j++){</pre>
                 d[0][j]=j; path[0][j]='i';
         int del,ins,Rplc;
         for(int i=1 ; i<=str1.length() ; i++) {</pre>
                 for(int j=1 ; j<=str2.length() ; j++) {</pre>
                          del=d[i-1][j]+1;
                          ins=d[i][j-1]+1;
                          Rplc=d[i-1][j-1]+(str1[i]!=str2[j]);
d[i][j]=min(min(del,ins),Rplc);
                          if(d[i][j]==del) path[i][j]='d';
                          else if(d[i][j]==ins) path[i][j]='i';
                          else if(str1[i]!=str2[j]) path[i][j]='r';
```

```
}
return d[str1.length()][str2.length()];

int main(){
    int t=1;
    while(getline(cin,str1) && getline(cin,str2)){
        if(t++!=1) cout << endl; str1=' '+str1; str2=' '+str2;
        cout << solve() << endl; Step=1; InsCnt=DelCnt=0;
        PrintPath(str1.length(), str2.length());
}
```

3 Graph Algorithms

3.1 Articulation Points, Bridges

```
int n, lev, dfsRoot, rootChilds;
int dfsLow[MAX], dfsNum[MAX], parent[MAX];
vvi adj; set<pii> bridges; set<int> artPoints;
         dfsLow[u] = dfsNum[u] = lev++;
         for (int i = 0; i < adj[u].size(); i++) {</pre>
                  int v = adj[u][i];
                  if (dfsNum[v] == 0) {
                           if (u == dfsRoot) rootChilds++;
                           parent[v] = u; dfs(v);
                           if (dfsLow[v] >= dfsNum[u] && u != dfsRoot) //u is articulation point
                                    artPoints.insert(u);
                           if (dfsLow[v] > dfsNum[u]) {
                                    bridges.insert(pii(v, u));
                                    bridges.insert(pii(u, v));
                           dfsLow[u] = min(dfsLow[u], dfsLow[v]);
                  else if (parent[u] != v)
                           dfsLow[u] = min(dfsLow[u], dfsNum[v]);
int main() {
         while (cin >> n) {
                 adj.assign(n, vi()); //initialization
memset(dfsLow, 0, sizeof dfsLow);
memset(dfsNum, 0, sizeof dfsNum);
memset(parent, 0, sizeof parent);
bridges.clear(); artPoints.clear();
                  lev = 1; int tmp, u, m;
                  for (int i = 0; i<n; i++) { // construct the graph
                           scanf("%d (%d", &u, &m); cin.ignore();
                           for (int i = 0; i<m; i++) {
                                   cin >> tmp; adj[u].push_back(tmp);
                  for (int i = 0; i < n; i++) {
                           if (dfsNum[i] == 0) {
                                    dfsRoot = i; rootChilds = 0; dfs(i);
                                    if (rootChilds >= 2) artPoints.insert(dfsRoot);
                  printf("%d critical links\n", bridges.size());
                  set<pii>>::iterator itr; // print answer
                  for (itr = bridges.begin(); itr != bridges.end(); itr++)
                           printf("%d - %d\n", itr->first, itr->second);
         return 0;}
```

3.2 Strongly Connected Component, Kosaraju

```
#include <iostream>
#include <string>
#include <sstream>
#include <vector>
```

```
#include <map>
using namespace std;
typedef vector<int> vi;
typedef vector<vi > vvi;
map<string, int> mpsi;
map<int, string> mpis;
int mapId;
int ID(string name){
        if(mpsi.count(name)) return mpsi[name];
        mpsi[name] = mapId;
mpis[mapId] = name;
        return mapId++;
vvi adjOrg, adjRev; vi vis, ord, col; int n, m;
void dfsOrg(int u) {
        if (vis[u]) return; vis[u] = true;
        for (int i = 0; i < adjOrg[u].size(); i++) {</pre>
                dfsOrg(adjOrg[u][i]);
        ord.push_back(u);
int dfsRev(int u, int color){
        if (col[u]) return 0; col[u] = color;
        for (int i = 0; i < adjRev[u].size(); i++) {</pre>
                ret += dfsRev(adjRev[u][i], color);
        return ret;
int main(){
        int cs = 1;
        while (cin >> n >> m && (n+m)) {
                mpsi.clear(); mpis.clear(); mapId = 0;
                adjOrg.assign(n, vi());
                adjRev.assign(n, vi());
                for (int i=0 ; i<m ; i++) {</pre>
                         string uname, vname;
                        cin >> uname >> vname;
                         int u = ID(uname);
                        int v = ID(vname);
                        adjOrg[u].push_back(v);
                         adjRev[v].push_back(u);
                ord.clear();
                vis.assign(n, 0);
                for (int u = 0; u < n; u++) {
                        if (!vis[u]) dfsOrg(u);
                int color = 1;
                col.assign(n, 0);
                while (!ord.empty()){
                         int u = ord.back();
                         if (!col[u]){
                                 int size = dfsRev(u, color); // SCC Size
                                 color++;
                         ord.pop_back();
                if (cs != 1) cout << endl;
                cout << "Calling circles for data set " << cs++ << ":" << endl;
                for (int c = 1; c<color ; c++) {</pre>
                         for (int u=0; u<n; u++) {
                                 if(col[u] == c) cout << ws << mpis[u], ws=", ";</pre>
                         cout << endl;
```

3.3 2-Satisfiability (2SAT)

#include <iostream>

```
#include <cstdio>
#include <vector>
#include <algorithm>
using namespace std;
#define FOR(i,x,y) for(int i = (x); i \le (y); ++i)
#define SZ(x) ((int)(x).size())
#define ALL(x) (x).begin(),(x).end()
#define UNIQUE(V) (V).erase(unique((V).begin(),(V).end()),(V).end())
typedef vector<int> vi;
typedef vector<vi > vvi;
struct Episode {
        Episode(int s=0, int e=0): s(s), e(e) {}
#define SCCNODE 1000
struct SCC{
    int num[SCCNODE], low[SCCNODE], col[SCCNODE], cycle[SCCNODE], st[SCCNODE];
    int tail, cnt, cc;
    vi adj[SCCNODE];
    SCC():tail(0),cnt(0),cc(0) {}
    void clear ( int n ) {
       cc += 3;
        FOR(i,0,n) adj[i].clear();
        tail = 0;
    void tarjan ( int s ) {
       num[s] = low[s] = cnt++;
col[s] = cc + 1;
        st[tail++] = s;
        FOR(i,0,SZ(adj[s])-1) {
            int t = adj[s][i];
            if ( col[t] <= cc ) {</pre>
                tarjan (t);
                low[s]=min(low[s],low[t]);
            /*Back edge*/
            else if (col[t] ==cc+1)
                low[s]=min(low[s],low[t]);
        if ( low[s] == num[s] ) {
            while ( 1 ) {
                int temp=st[tail-1];
                tail--;
                col[temp] = cc + 2;
                cycle[temp] = s;
                if ( s == temp ) break;
    void shrink( int n ) {
        FOR(i,0,n){
            FOR(j,0,SZ(adj[i])-1){
                adj[i][j] = cycle[adj[i][j]]; ///Careful. This will create self-loop
        FOR(i,0,n){
            if ( cycle[i] == i ) continue;
            int u = cycle[i];
            FOR(j,0,SZ(adj[i])-1){
                int v = adj[i][j];
                adj[u].push_back( v );
            adj[i].clear();
        FOR(i,0,n){ ///Not always necessary
            sort ( ALL(adj[i]) );
            UNIQUE (adj[i]);
    void findSCC( int n ) {
        FOR(i,0,n) {
            if ( col[i] <= cc ) {</pre>
                tarjan ( i );
};
#define SAT2NODE 2000
1. The nodes need to be split. So change convert() accordingly.
2. Using clauses, populate scc edges.
3. Call possible, to find if a valid solution is possible or not.
4. Dont forget to keep space for !A variables
struct SAT2 {
```

```
SCC scc;
    SAT2(): bfscc(1) {}
    void clear( int n ) { scc.clear( n ); }
    int convert ( int n ) { return n \star 2; } ///Change here. Depends on how input is provided
    void mustTrue ( int a ) { scc.adj[a^1].push_back( a ); } ///A is True
    void andClause ( int a, int b ) { scc.adj[a].push_back( b ); scc.adj[b].push_back( a ); }
    void orClause ( int a, int b ) { scc.adj[a^1].push_back( b ); scc.adj[b^1].push_back( a ); } /// A
           || B clause //!a->b !b->a
    void xorclause ( int a, int b ) { orclause(a, b ); orclause(a^1, b^1); }
void notAndClause ( int a, int b ) { scc.adj[a].push_back(b ^ 1 ); scc.adj[b].push_back(a ^ 1 );
    void atMostOneClause ( int a[], int n, int flag ) { /// Out of all possible option, only one is
          true
        if ( flag == 0 ) { /// At most one can be false
            FOR(i,0,n){
                a[i] = a[i] ^ 1;
        FOR(i,0,n) {
            FOR(j,i+1,n) {
                orClause(a[i] ^ 1, a[j] ^ 1 ); /// !a || !b both being true not allowed
    ///Send n, total number of nodes, after expansion
    bool possible(int n ) {
        scc.findSCC( n );
        FOR(i,0,n) {
            int a = i, b = i^1;
             ///Falls on same cycle a and !a.
            if ( scc.cycle[a] == scc.cycle[b] ) return false;
        ///Valid solution exists
        return true;
    ///To determine if A can be true. It cannot be true, if a path exists from A to !A.
    int vis[SAT2NODE], qqq[SAT2NODE], bfscc;
    void bfs( int s ) {
        bfscc++;
        int qs = 0, qt = 0;
        vis[s] = bfscc;
        qqq[qt++] = s;
        while ( qs < qt ) {
            s = qqq[qs++];
            FOR(i,0,SZ(scc.adj[s])-1) {
                int t = scc.adj[s][i];
                if ( vis[t] != bfscc ) {
                    vis[t] = bfscc;
                    qqq[qt++] = t;
} sat2;
vector< Episode > ep:
void build_graph(int 1, int r) {
        for(int i=1; i<=r; i+=2) {
                sat2.xorClause(sat2.convert(i), sat2.convert(i+1));
        for (int i=1 ; i<=r ; i++) {</pre>
                for(int j=i+1; j<=r; j++) {</pre>
                         if(ep[i].e < ep[j].s || ep[j].e < ep[i].s) continue;</pre>
                         sat2.notAndClause(sat2.convert(i), sat2.convert(j));
int main () {
        int tc;
        cin >> tc;
        while(tc--) {
                int n;
                cin >> n:
                n = n * 2:
                ep.assign(n, Episode());
                for(int i=0 ; i<n ; i++) cin >> ep[i].s >> ep[i].e;
                int ansL = -1, ansR = -1, maxDist = -1;
                int 1 = 0;
                for (int r = 1; r < n; r + = 2) {
```

```
sat2.clear(n * 2);
build_graph(l, r);

while(sat2.possible(n * 2) == false) {
    l+=2;
        sat2.clear(n * 2);
        build_graph(l, r);
}

if (r - 1 > maxDist) {
        ansR = r/2;
        ansL = 1/2;
        maxDist = r - 1;
}

cout << ansL + 1 << " " << ansR + 1 << endl;
}
</pre>
```

3.4 Shortest Path, Dijkstra

```
#include <iostream>
#include <vector>
#include <queue>
using namespace std;
const int INF = 1e8;
struct ToNode{
        int v, w;
        ToNode (int v. int w)
                :v(v), w(w){}
};
struct QEntry{
        int node, cost:
        QEntry(int node, int cost)
                :node(node), cost(cost){}
        bool operator<(const QEntry& op) const {</pre>
                return cost < op.cost;
};
typedef vector<int>
typedef vector<ToNode> vtn;
typedef vector<vtn > vvtn;
int n, m; vvtn adj;
int dijkstra(int s, int t, vi& dist){
        dist.assign(n, INF);
        priority_queue<QEntry> q;
        q.push(QEntry(s, 0)); dist[s] = 0;
        while (!q.empty()){
                QEntry u = q.top(); q.pop();
                if (u.node == t) return u.cost;
                if (u.cost > dist[u.node]) continue;
                for (int i = 0; i < adj[u.node].size(); i++) {</pre>
                        QEntry v(adj[u.node][i].v, u.cost + adj[u.node][i].w);
                        if (dist[v.node] > v.cost){
                                dist[v.node] = v.cost; q.push(v);
        return INF:
```

3.5 LCA Tree Distance

```
if (n == 0)
    return -1;
  int p = 0;
  if (n >= 1 << 16) { n >>= 16; p += 16; }
  if (n >= 1 << 8) { n >>= 8; p += 8; }
if (n >= 1 << 4) { n >>= 4; p += 4; }
  if (n >= 1 << 2) { n >>= 2; p += 2; }
  if (n >= 1 << 1) { p += 1; }
  return p;
void DFS(int i, int 1)
  L[i] = 1;
  for (int j = 0; j < children[i].size(); j++)</pre>
    DFS(children[i][j], 1 + 1);
int LCA(int p, int q)
  // ensure node p is at least as deep as node q
  if (L[p] < L[q])
    swap(p, q);
  // "binary search" for the ancestor of node p situated on the same level as \boldsymbol{q}
  for (int i = log_num_nodes; i >= 0; i--)
    if (L[p] - (1 << i) >= L[q])
     p = A[p][i];
  if (p == q)
    return p;
  // "binary search" for the LCA
  for (int i = log_num_nodes; i >= 0; i--)
    if (A[p][i] != -1 && A[p][i] != A[q][i])
      q = A[q][i];
  return A[p][0];
int main(int argc, char* argv[])
  // read num_nodes, the total number of nodes
  log_num_nodes = lb(num_nodes);
  for (int i = 0; i < num_nodes; i++)</pre>
    int p;
    // read p, the parent of node i or -1 if node i is the root
    A[i][0] = p;
    if (p != -1)
      children[p].push_back(i);
    else
      root = i:
  // precompute A using dynamic programming
  for (int j = 1; j <= log_num_nodes; j++)</pre>
    for (int i = 0; i < num_nodes; i++)
if (A[i][j - 1] != -1)</pre>
        A[i][j] = A[A[i][j-1]][j-1];
      else
        A[i][j] = -1;
  // precompute L
  DFS(root, 0);
  return 0;
```

3.6 Graphic Sequence

```
// given a sequence of integers see if it s a sequence of degrees of graph or not.
int a[10010]; long long sum,Min;;
int main(){
   int n;
   while(cin >> n && n){
      for(int i=0; i<n; i++) scanf("%d",&a[i]);
      sort(a,a+n, ::greater<int>() );
      bool possible=true; sum=0;
   for(int i=0; i<n; i++){</pre>
```

```
sum+=a[i]; Min=0;
for(int j=i+1; j<n; j++) Min+=min(a[j],i+1);
if(sum>i+(i+1)+Min){
    possible=false;
    break;
}
if(!possible || sum%2) cout << "Not possible" << endl;
else cout << "Possible" << endl;
}
return 0;}</pre>
```

3.7 Floyd Warshal, Print Path

```
#define MAX (100+10)
int adj[MAX] [MAX], path[MAX] [MAX]; int n;
void print(int i, int j) {
       if(i!=j){
               printf(" %d",i );
               print (path[i][j], j);
int main(){
       int tc; cin >> tc;
       while(tc--){
               cin >> n;
               path[i][j]=j;//initial parent
               for(int k=0 ; k<n ; k++) {</pre>
                       for (int i=0 ; i < n ; i++) {</pre>
                              for(int j=0; j<n; j++) {
                                      if(adj[i][j]>adj[i][k]+adj[k][j]){
                                              adj[i][j]=adj[i][k]+adj[k][j];
                                              path[i][j]=path[i][k];//set parent
               int s,d;
cin >> s >> d;
               printf("%d euros\n",adj[s][d]);
               //this prints the path even if source and distinaion are same
               printf("%d",s); print(path[s][d],d); printf(" %d\n",d);
       return 0;
```

3.8 Max Flow, Edmonds Karp

```
//UVa 820 - Internet Bandwidth
#define INF (int) 1e9
#define MAX 100+10
using namespace std:
int res[MAX][MAX], mf, f, s, t, n, m, par[MAX]; vector<int> dist, adj[MAX];
void agument(int v, int minEdge) {
    if(v==s) f=minEdge;
    else if (par[v]!=-1) {
        agument (par[v], min (minEdge, res[par[v]][v]));
        res[par[v]][v]-=f; res[v][par[v]]+=f;
int main(){
    int tc=1;
    while (cin >> n && n) {
        mf=0; memset(res,0,sizeof res); for(int i=0; i<n; i++) adj[i].clear();</pre>
        cin >> s >> t >> m; s--; t--;
        int u, v, c;
        while (m--) {
            cin >> u >> v >> c; u--; v--;
```

```
res[u][v]+=c; res[v][u]+=c;
        adj[u].push_back(v); adj[v].push_back(u);
        f=0; memset(par,-1,sizeof par); dist.assign(n,INF);
        dist[s]=0; queue<int> q; q.push(s);
        while(!q.empty()){
            int u=q.front(); q.pop();
            if(u==t) break;
            for(int i=0 ; i<adj[u].size(); i++){</pre>
                int v=adj[u][i];
                if(res[u][v]>0 && dist[v]==INF) {
                    dist[v]=dist[u]+1; q.push(v); par[v]=u;
        agument (t, INF);
        if(f==0) break;
        mf+=f;
   printf("Network %d\n", tc++);
    printf("The bandwidth is %d.\n\n", mf);
return 0:
```

3.9 Max Flow, Dinic

```
// Adjacency list implementation of Dinic's blocking flow algorithm.
// This is very fast in practice, and only loses to push-relabel flow.
// Running time: O(|V|^2 |E|)
        - graph, constructed using AddEdge() - source - sink
       - maximum flow value
       - To obtain the actual flow values, look at all edges with
        capacity > 0 (zero capacity edges are residual edges).
using namespace std;
const int INF = 2000000000;
struct Edge {
    int from, to, cap, flow, index;
    Edge(int from, int to, int cap, int flow, int index) :
        from(from), to(to), cap(cap), flow(flow), index(index) {}
struct Dinic {
    int N; vector<vector<Edge> > G;
    vector<Edge *> dad; vector<int> Q;
    Dinic(int N) : N(N), G(N), dad(N), Q(N) {}
    void AddEdge(int from, int to, int cap) {
        G[from].push_back(Edge(from, to, cap, 0, G[to].size()));
        if (from == to) G[from].back().index++;
        G[to].push\_back(Edge(to, from, 0, 0, G[from].size() - 1));
    long long BlockingFlow(int s, int t) {
        fill(dad.begin(), dad.end(), (Edge *)NULL);
dad[s] = &G[0][0] - 1;
        int head = 0, tail = 0;
        Q[tail++] = s;
        while (head < tail) {</pre>
            int x = Q[head++];
             for (int i = 0; i < G[x].size(); i++) {</pre>
                 Edge &e = G[x][i];
                 if (!dad[e.to] && e.cap - e.flow > 0) {
                     dad[e.to] = &G[x][i];
                     Q[tail++] = e.to;
        if (!dad[t]) return 0;
        long long totflow = 0;
for (int i = 0; i < G[t].size(); i++) {
    Edge *start = &G[G[t][i].to][G[t][i].index];</pre>
             int amt = INF;
            for (Edge *e = start; amt && e != dad[s]; e = dad[e->from]) {
                 if (!e) { amt = 0; break; }
                 amt = min(amt, e->cap - e->flow);
             if (amt == 0) continue;
             for (Edge *e = start; amt && e != dad[s]; e = dad[e->from]) {
                 e->flow += amt;
                 G[e->to][e->index].flow -= amt;
            totflow += amt:
        return totflow;
    long long GetMaxFlow(int s, int t) {
        long long totflow = 0;
        while (long long flow = BlockingFlow(s, t))
            totflow += flow;
```

```
return totflow;
};
```

3.10 Min Cut

```
// Adjacency matrix implementation of Stoer-Wagner min cut algorithm.
// Running time: O(|V|^3)
// INPUT: graph, constructed using AddEdge()
// OUTPUT: (min cut value, nodes in half of min cut)
#include <cmath>
#include <vector>
#include <iostream>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
const int INF = 10000000000:
pair<int, VI> GetMinCut(VVI &weights) {
  int N = weights.size();
  VI used(N), cut, best_cut;
  int best_weight = -1;
  for (int phase = N-1; phase >= 0; phase--) {
    VI w = weights[0];
    VI added = used;
    int prev, last = 0;
    for (int i = 0; i < phase; i++) {</pre>
      prev = last;
       last = -1;
      for (int j = 1; j < N; j++)
        if (!added[j] && (last == -1 || w[j] > w[last])) last = j;
       if (i == phase-1) {
        for (int j = 0; j < N; j++) weights[prev][j] += weights[last][j];
for (int j = 0; j < N; j++) weights[j][prev] = weights[prev][j];</pre>
        used[last] = true;
         cut.push back(last);
        if (best weight == -1 || w[last] < best weight) {</pre>
          best_cut = cut;
          best_weight = w[last];
      } else {
        for (int j = 0; j < N; j++)
           w[j] += weights[last][j];
        added[last] = true;
  return make_pair(best_weight, best_cut);
// The following code solves UVA problem #10989: Bomb, Divide and Conquer
int main() {
  int N;
  for (int i = 0; i < N; i++) {
    int n, m;
    cin >> n >> m;
    VVI weights(n, VI(n));
    for (int j = 0; j < m; j++) {
      int a, b, c;
      cin >> a >> b >> c;
      weights[a-1][b-1] = weights[b-1][a-1] = c;
    pair<int, VI> res = GetMinCut(weights);
cout << "Case #" << i+1 << ": " << res.first << endl;</pre>
// END CUT
```

3.11 MaxCardinalityBipartiteMatching, Alternating path

```
#include <iostream>
#include <algorithm>
#include <vector>
using namespace std;
```

```
typedef vector<int> vi;
typedef vector<vi > vvi;
vvi adj; vi owner, vis; int n,b;
int altpath(int u){
        if(vis[u]) return 0; vis[u]=1;
        for(int i=0 ; i<adj[u].size() ; i++) {</pre>
                int v=adj[u][i];
                if(owner[v] ==-1 || altpath(owner[v])){
                        owner[v]=u; return 1;
        return 0:
        int tmp,tc,t=1; cin >> tc;
        while (tc--) {
                cin >> n >> b; adj.assign(n+b,vi());
                for(int i=0; i<n; i++) {
                        for (int j=0 ; j<b ; j++) {</pre>
                                // if there is an edge from n group to b group
                                cin >> tmp; if(tmp==1) adj[i].push_back(j+n);
                int ans=0; owner.assign(n+b,-1);
                for(int u=0; u<n; u++){
                        vis.assign(n,0); ans+=altpath(u);
                printf("Case %d: a maximum of %d matched\n", t++, ans);
        return 0;
```

3.12 Min Cost Max Flow

```
// Implementation of min cost max flow algorithm using adjacency // matrix (Edmonds and Karp 1972). This implementation keeps track of
// forward and reverse edges separately (so you can set cap[i][j] !=
// cap[j][i]). For a regular max flow, set all edge costs to 0.
// Running time, O(|V|^2) cost per augmentation
      max flow:
                            O(|V|^3) augmentations
        min cost max flow: O(|V|^4 * MAX\_EDGE\_COST) augmentations
         - graph, constructed using AddEdge(), source, sink
       - (maximum flow value, minimum cost value)
       - To obtain the actual flow, look at positive values only.
#include <cmath>
#include <vector>
#include <iostream>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
typedef long long L;
typedef vector<L> VL;
typedef vector<VL> VVL;
typedef pair<int, int> PII;
typedef vector<PII> VPII;
const L INF = numeric_limits<L>::max() / 4;
struct MinCostMaxFlow {
  int N:
  VVL cap, flow, cost;
  VI found:
  VL dist, pi, width;
  MinCostMaxFlow(int N) :
    N(N), cap(N, VL(N)), flow(N, VL(N)), cost(N, VL(N)),
    found(N), dist(N), pi(N), width(N), dad(N) {}
  void AddEdge(int from, int to, L cap, L cost) {
    this->cap[from][to] = cap;
    this->cost[from][to] = cost;
  void Relax(int s, int k, L cap, L cost, int dir) {
  L val = dist[s] + pi[s] - pi[k] + cost;
  if (cap && val < dist[k]) {</pre>
      dist[k] = val;
```

```
dad[k] = make_pair(s, dir);
      width[k] = min(cap, width[s]);
  L Dijkstra(int s, int t) {
    fill(found.begin(), found.end(), false);
    fill(dist.begin(), dist.end(), INF);
    fill(width.begin(), width.end(), 0);
    dist[s] = 0;
    width[s] = INF;
    while (s != -1) {
     int best = -1;
      found[s] = true;
      for (int k = 0; k < N; k++) {
       if (found[k]) continue;
        Relax(s, k, cap[s][k] - flow[s][k], cost[s][k], 1);
        Relax(s, k, flow[k][s], -cost[k][s], -1);
        if (best == -1 || dist[k] < dist[best]) best = k;
    for (int k = 0; k < N; k++)
     pi[k] = min(pi[k] + dist[k], INF);
    return width[t];
  pair<L, L> GetMaxFlow(int s, int t) {
    L totflow = 0, totcost = 0;
    while (L amt = Dijkstra(s, t)) {
      totflow += amt;
      for (int x = t; x != s; x = dad[x].first) {
       if (dad[x].second == 1) {
          flow[dad[x].first][x] += amt;
          totcost += amt * cost[dad[x].first][x];
        | else {
          flow[x][dad[x].first] -= amt;
         totcost -= amt * cost[x][dad[x].first];
    return make pair (totflow, totcost);
};
// The following code solves UVA problem #10594: Data Flow
int main() {
 int N, M;
  while (scanf("%d%d", &N, &M) == 2) {
    VVL v(M, VL(3));
    for (int i = 0; i < M; i++)
     scanf("%Ld%Ld%Ld", &v[i][0], &v[i][1], &v[i][2]);
   L D, K;
    scanf("%Ld%Ld", &D, &K);
    MinCostMaxFlow mcmf(N+1);
    for (int i = 0; i < M; i++) {
     mcmf.AddEdge(int(v[i][0]), int(v[i][1]), K, v[i][2]);
      mcmf.AddEdge(int(v[i][1]), int(v[i][0]), K, v[i][2]);
    mcmf.AddEdge(0, 1, D, 0);
    pair<L, L> res = mcmf.GetMaxFlow(0, N);
    if (res.first == D) {
      printf("%Ld\n", res.second);
    else {
     printf("Impossible.\n");
  return 0;
// END CUT
```

3.13 Min Cost Matching

```
// algorithm for finding min cost perfect matchings in dense
// graphs. In practice, it solves 1000x1000 problems in around 1
// second.
     cost[i][j] = cost for pairing left node i with right node j
    Lmate[i] = index of right node that left node i pairs with
    Rmate[j] = index of left node that right node j pairs with
// The values in cost[i][j] may be positive or negative. To perform
// maximization, simply negate the cost[][] matrix.
#include <algorithm>
#include <cstdio>
#include <cmath>
#include <vector>
using namespace std;
typedef vector<double> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
double MinCostMatching(const VVD &cost, VI &Lmate, VI &Rmate) {
  int n = int(cost.size());
  // construct dual feasible solution
  VD u(n);
  VD v(n);
  for (int i = 0; i < n; i++) {
    u[i] = cost[i][0];
    for (int j = 1; j < n; j++) u[i] = min(u[i], cost[i][j]);</pre>
  for (int j = 0; j < n; j++) {
    v[j] = cost[0][j] - u[0];
    for (int i = 1; i < n; i++) v[j] = min(v[j], cost[i][j] - u[i]);
  // construct primal solution satisfying complementary slackness
  Lmate = VI(n, -1);
  Rmate = VI(n, -1);
 for (int j = 0; j < n; j++) {
    if (Rmate[j] != -1) continue;</pre>
      if (fabs(cost[i][j] - u[i] - v[j]) < 1e-10) {</pre>
        Lmate[i] = j;
        Rmate[j] = i;
        mated++;
        break;
  VD dist(n):
  VI dad(n):
  VI seen(n):
  // repeat until primal solution is feasible
  while (mated < n) {
    // find an unmatched left node
    while (Lmate[s] != -1) s++;
    // initialize Dijkstra
    fill(dad.begin(), dad.end(), -1);
    fill(seen.begin(), seen.end(), 0);
for (int k = 0; k < n; k++)</pre>
      dist[k] = cost[s][k] - u[s] - v[k];
    int i = 0;
    while (true) {
      // find closest
      for (int k = 0; k < n; k++) {
        if (seen[k]) continue;
        if (j == -1 || dist[k] < dist[j]) j = k;</pre>
      seen[j] = 1;
      // termination condition
      if (Rmate[j] == -1) break;
      // relax neighbors
      const int i = Rmate[j];
      for (int k = 0; k < n; k++) {
        if (seen[k]) continue;
        const double new_dist = dist[j] + cost[i][k] - u[i] - v[k];
        if (dist[k] > new_dist) {
          dist[k] = new_dist;
```

```
dad[k] = j;
  // update dual variables
  for (int k = 0; k < n; k++)
    if (k == j || !seen[k]) continue;
   const int i = Rmate[k];
v[k] += dist[k] - dist[j];
    u[i] -= dist[k] - dist[j];
  u[s] += dist[j];
  // augment along path
  while (dad[j] >= 0) {
   const int d = dad[j];
    Rmate[j] = Rmate[d];
    Lmate[Rmate[j]] = j;
  Rmate[j] = s;
  Lmate[s] = j;
  mated++:
double value = 0;
for (int i = 0; i < n; i++)
 value += cost[i][Lmate[i]];
return value;
```

4 Mathematics and Geometry

4.1 Prime Numbers, Factoring

```
#include <iostream>
#include <cstdio>
#include <vector>
#include <bitset>
using namespace std;
typedef long long int64;
const int64 MAX = 1e6 + 100;
bitset<MAX> isp;// isprime
vector<int64> primes, pfs, pws; //pfs = prime factors, pws = prime powers
void genprime(){
    isp.set(); isp[0]=isp[1]=0;
   for (int64 i=2; i<MAX; i++) {
    if (isp[i]) { primes.push_back(i);</pre>
            for(int64 j=i*i ; j<MAX ; j+=i) isp[j]=0;</pre>
bool isprime(int n) {
    if(n<MAX) return isp[n];</pre>
    if(n%primes[i]==0) return 0;
    return 1:
//generation prime factors of a number
int main(){
    int64 n; genprime();
                                          // cop = euler Phi funcion
    while (cin >> n) {
        int64 tmp=n, cnt=0, cop=n, div=1;// cop = coprimes = all m (m<n && gcd(m,n)==1)
                                          // div = divisors = all m (m<n && gcd(m,n) == m)
        for(int i=0, pf=2; pf*pf<=n; i++, pf=primes[i]){</pre>
            int pow=0;
            while (tmp%pf==0) {
                tmp/=pf; pow++;
            if(pow) {
                pfs.push_back(pf), pws.push_back(pow);
cop -= cop/pf; div *= (pow+1);
        if(tmp>1) cop -= cop/tmp, div *= (1+1);
        cout << cop + div + 1 << endl;
```

1

4.2 Extended Euclid

```
import java.util.Scanner;
public class Main {
    public static class ExtendedEuclid {
        public static int x0;
        public static int v0;
        public static int c;
        public static int d:
        public static int gcd(int a, int b) {
             if (b == 0) return a;
             return gcd(b, a % b);
        public static void calculate(int a, int b) {
             if (b == 0) {
                 ×0 = 1:
                 y0 = 0;
                 d = a:
                 return:
             calculate(b, a % b);
             int x1 = y0;
             int y1 = x0 - (a / b) * y0;
             x0 = x1;
             y0 = y1;
        public static int howManyPositiveSolutions(int a, int b, int c) {
   if (c % gcd(a, b) != 0) return 0; // no solution even negatives
             x0 \neq c / gcd(a, b);

y0 \neq c / gcd(a, b); // x = x0 + (b/d)n, y = y0 - (a/d)n
             double lowerBoundForN = (double) (-x0 + 0.5) / (b / d); // for x>0
             double upperBoundForN = (double) (+y0 - 0.5) / (a / d); // for y>0
             return (int) Math.max(0, Math.ceil(upperBoundForN) - Math.floor(lowerBoundForN) - 1); //
                   how many int between
    public static void main(String[] args) {
        ExtendedEuclid.calculate(25, 18); // Copetitive Programming 2 Example
```

4.3 Number Theory General

```
#include <iostream>
#include <vector>
#include <algorithm>
using namespace std:
typedef vector<int> VI;
typedef pair<int, int> PII;
// return a % b (positive value)
int mod(int a, int b) {
 return ((a%b) + b) % b;
// computes gcd(a,b)
int gcd(int a, int b) {
  while (b) { a %= b; tmp = a; a = b; b = tmp; }
  return a;
// computes lcm(a.b)
int lcm(int a, int b) {
 return a / gcd(a, b) *b;
```

```
// returns d = gcd(a,b); finds x,y such that d = ax + by
int extended_euclid(int a, int b, int &x, int &y) {
  int xx = y = 0;
  int yy = x = 1;
  while (b) {
    int q = a / b;
    int t = b; b = a%b; a = t;
    t = xx; xx = x - q*xx; x = t;
    t = yy; yy = y - q*yy; y = t;
  return a;
// finds all solutions to ax = b \pmod{n}
VI modular_linear_equation_solver(int a, int b, int n) {
  VI solutions:
  int d = extended_euclid(a, n, x, y);
  if (!(b%d)) {
    x = mod(x*(b / d), n);
    for (int i = 0; i < d; i++)
      solutions.push_back(mod(x + i*(n / d), n));
 return solutions:
// computes b such that ab = 1 \pmod{n}, returns -1 on failure
int mod inverse(int a, int n) {
 int x, y;
int d = extended_euclid(a, n, x, y);
  if (d > 1) return -1;
 return mod(x, n);
// Chinese remainder theorem (special case): find z such that
// z % x = a, z % y = b. Here, z is unique modulo M = lcm(x,y). // Return (z,M). On failure, M = -1.
PII chinese_remainder_theorem(int x, int a, int y, int b) {
 int s, t;
  int d = extended_euclid(x, y, s, t);
 if (a%d != b%d) return make_pair(0, -1);
return make_pair(mod(s*b*x + t*a*y, x*y) / d, x*y / d);
// Chinese remainder theorem: find z such that
// z % x[i] = a[i] for all i. Note that the solution is
// unique modulo M = lcm_i (x[i]). Return (z, M). On
// failure, M = -1. Note that we do not require the a[i]'s
// to be relatively prime.
PII chinese_remainder_theorem(const VI &x, const VI &a) {
 PII ret = make_pair(a[0], x[0]);
  for (int i = 1; i < x.size(); i++) {</pre>
    ret = chinese_remainder_theorem(ret.second, ret.first, x[i], a[i]);
    if (ret.second == -1) break;
 return ret:
// computes x and y such that ax + by = c; on failure, x = y = -1
void linear_diophantine(int a, int b, int c, int &x, int &y) {
  int d = gcd(a, b);
  if (c%d) {
   x = y = -1;
  else {
   x = c / d * mod_inverse(a / d, b / d);
    y = (c - a*x) / b;
long conquer_fibonacci_lgN(long n) {
 long i, h, j, k, t;
i = h = 1;
j = k = 0;
  while (n > 0) {
   if (n % 2 == 1) {
     t = j * h;
      j = i * h + j * k + t;
      i = i * k + t;
    h = 2 * k * h + t;

k = k * k + t;
    n = (long) n / 2;
  return i:}
```

4.4 Gauss-Jordan Elimination

```
// Gauss-Jordan elimination with full pivoting.
// Uses:
     (1) solving systems of linear equations (AX=B)
     (2) inverting matrices (AX=I)
// (3) computing determinants of square matrices
// Running time: O(n^3)
// INPUT:
              a[][] = an nxn matrix
                b[][] = an nxm matrix
                       = an nxm matrix (stored in b[][])
                A^{-1} = an \ nxn \ matrix \ (stored in a[][])
                returns determinant of a[][]
#include <iostream>
#include <vector>
#include <cmath>
using namespace std;
const double EPS = 1e-10;
typedef vector<int> VI;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
T GaussJordan (VVT &a, VVT &b) {
  const int n = a.size();
  const int m = b[0].size();
  VI irow(n), icol(n), ipiv(n);
  T det = 1:
   for (int i = 0; i < n; i++) {</pre>
     for (int j = 0; j < n; j++) if (!ipiv[j])
for (int j = 0; j < n; j++) if (!ipiv[j])
for (int k = 0; k < n; k++) if (!ipiv[k])</pre>
         if (pj == -1 || fabs(a[j][k]) > fabs(a[pj][pk])) { pj = j; pk = k; }
     if (fabs(a[pj][pk]) < EPS) { cerr << "Matrix is singular." << endl; exit(0); }</pre>
     ipiv[pk]++;
     swap(a[pj], a[pk]);
     swap(b[pj], b[pk]);
    if (pj != pk) det *= -1;
irow[i] = pj;
icol[i] = pk;
    T c = 1.0 / a[pk][pk];
     det *= a[pk][pk];
     a[pk][pk] = 1.0;
     for (int p = 0; p < n; p++) a[pk][p] *= c;
for (int p = 0; p < m; p++) b[pk][p] *= c;
     for (int p = 0; p < n; p++) if (p != pk) {
       c = a[p][pk];
       for (int q = 0; q < n; q++) a[p][q] -= a[pk][q] * c;
       for (int q = 0; q < m; q++) b[p][q] -= b[pk][q] * c;
  for (int p = n-1; p >= 0; p--) if (irow[p] != icol[p]) {
   for (int k = 0; k < n; k++) swap(a[k][irow[p]], a[k][icol[p]]);</pre>
  return det:
int main() {
  const int n = 4;
  const int m = 2;
  double A[n][n] = \{ \{1,2,3,4\}, \{1,0,1,0\}, \{5,3,2,4\}, \{6,1,4,6\} \}; double B[n][m] = \{ \{1,2\}, \{4,3\}, \{5,6\}, \{8,7\} \};
  VVT a(n), b(n);
   for (int i = 0; i < n; i++) {
    a[i] = VT(A[i], A[i] + n);
b[i] = VT(B[i], B[i] + m);
  double det = GaussJordan(a, b);
   cout << "Determinant: " << det << endl;</pre>
   // expected: -0.233333 0.166667 0.133333 0.0666667
                  0.166667 0.166667 0.333333 -0.333333 
0.233333 0.833333 -0.133333 -0.0666667
                  0.05 -0.75 -0.1 0.2
   cout << "Inverse: " << endl;</pre>
  for (int i = 0; i < n; i++) {
  for (int j = 0; j < n; j++)
    cout << a[i][j] << ' ';</pre>
     cout << endl;
```

// expected: 1.63333 1.3

```
// -0.166667 0.5

// 2.36667 1.7

// -1.85 -1.35

cout << "Solution: " << end1;

for (int i = 0; i < n; i++) {

  for (int j = 0; j < m; j++) {

    cout << b[i][j] << ' ';

    cout << end1;

}
```

4.5 BigInteger Square in java

```
import java.math.BigInteger;
import java.util.Scanner;
//import java.util.
public class Main {
    // https://en.wikipedia.org/wiki/Integer_square_root
    public static BigInteger sqrt(BigInteger n) {
    BigInteger cur = null; // X(k)
        BigInteger nxt = n; // X(k+1)
        while(true) {
            nxt = cur.add(n.divide(cur)).divide(BigInteger.valueOf(2));
            if(nxt.equals(cur)) break;
        if(cur.multiply(cur).equals(n)) return cur;
        else return null:
    public static void main(String[] args) {
        Scanner sc = new Scanner(System.in);
        int tc = Integer.parseInt(sc.nextLine());
        while(tc-- > 0) {
            sc.nextLine();
            BigInteger y = new BigInteger(sc.nextLine());
            if(y.equals(BigInteger.ZERO)) System.out.println(0);
                                           System.out.println(sqrt(y));
            if(tc>0) System.out.println();
```

4.6 Geometry 1

```
// C++ routines for computational geometry.
#include <iostream>
#include <vector>
#include <cmath>
#include <cassert>
using namespace std;
double INF = 1e100;
double EPS = 1e-12;
  double x, y;
  PT() {}
  \texttt{PT}\left(\texttt{double}\ x,\ \texttt{double}\ y\right)\ :\ x\left(x\right),\ y\left(y\right)\ \{\,\}
  PT(const PT &p) : x(p.x), y(p.y) {}
PT operator + (const PT &p) const { return PT(x+p.x, y+p.y); }
  PT operator - (const PT &p) const { return PT(x-p.x, y-p.y);
                                   const { return PT(x*c, y*c );
  PT operator * (double c)
  PT operator / (double c)
                                   const { return PT(x/c, y/c ); ]
double dot (PT p, PT q)
                               { return p.x*q.x+p.y*q.y; }
double dist2(PT p, PT q) { return dot(p-q,p-q); }
double cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
ostream &operator<<(ostream &os, const PT &p) {
  os << "(" << p.x << ", " << p.y << ")";
// rotate a point CCW or CW around the origin
PT RotateCCW90 (PT p) { return PT (-p.y,p.x); }
PT RotateCW90(PT p)
                          { return PT(p.y,-p.x); }
PT RotateCCW(PT p, double t) {
  return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
// project point c onto line through a and b
```

```
// assuming a != b
PT ProjectPointLine(PT a, PT b, PT c) {
  return a + (b-a) *dot (c-a, b-a) /dot (b-a, b-a);
// project point c onto line segment through a and b
PT ProjectPointSegment (PT a, PT b, PT c) {
  double r = dot(b-a, b-a);
  if (fabs(r) < EPS) return a;</pre>
  r = dot(c-a, b-a)/r;
  if (r < 0) return a;
  if (r > 1) return b;
  return a + (b-a) *r;
// compute distance from c to segment between a and b
double DistancePointSegment (PT a, PT b, PT c) {
  return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
// compute distance between point (x,y,z) and plane ax+by+cz=d
double DistancePointPlane(double x, double y, double z,
                           double a, double b, double c, double d)
  return fabs(a*x+b*y+c*z-d)/sqrt(a*a+b*b+c*c);
// determine if lines from a to b and c to d are parallel or collinear
bool LinesParallel(PT a, PT b, PT c, PT d) {
 return fabs(cross(b-a, c-d)) < EPS;
bool LinesCollinear(PT a, PT b, PT c, PT d) {
  return LinesParallel(a, b, c, d)
      && fabs(cross(a-b, a-c)) < EPS
      && fabs(cross(c-d, c-a)) < EPS;
// determine if line segment from a to b intersects with
// line segment from c to d
bool SegmentsIntersect(PT a, PT b, PT c, PT d) {
 if (LinesCollinear(a, b, c, d)) {
   if (dist2(a, c) < EPS || dist2(a, d) < EPS ||</pre>
    dist2(b, c) < EPS || dist2(b, d) < EPS) return true;
if (dot(c-a, c-b) > 0 && dot(d-a, d-b) > 0 && dot(c-b, d-b) > 0)
      return false;
    return true:
  if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return false;
  if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return false;
  return true:
// compute intersection of line passing through a and b
// with line passing through c and d, assuming that unique
// intersection exists; for segment intersection, check if
// segments intersect first
PT ComputeLineIntersection(PT a, PT b, PT c, PT d) {
 b=b-a; d=c-d; c=c-a;
  assert(dot(b, b) > EPS && dot(d, d) > EPS);
  return a + b*cross(c, d)/cross(b, d);
// compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c) {
 b = (a+b)/2:
  c = (a+c)/2;
  return ComputeLineIntersection(b, b+RotateCW90(a-b), c, c+RotateCW90(a-c));
// determine if point is in a possibly non-convex polygon (by William
// Randolph Franklin); returns 1 for strictly interior points, 0 for
// strictly exterior points, and 0 or 1 for the remaining points.
// Note that it is possible to convert this into an *exact* test using
// integer arithmetic by taking care of the division appropriately
// (making sure to deal with signs properly) and then by writing exact
// tests for checking point on polygon boundary
bool PointInPolygon(const vector<PT> &p, PT q) {
  bool c = 0;
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1)%p.size();
    if ((p[i].y <= q.y && q.y < p[j].y ||
  p[j].y <= q.y && q.y < p[i].y) &&</pre>
      q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) / (p[j].y - p[i].y))
      c = !c;
  return c:
// determine if point is on the boundary of a polygon
bool PointOnPolygon(const vector<PT> &p, PT q) {
  for (int i = 0; i < p.size(); i++)
```

```
if (dist2(ProjectPointSegment(p[i], p[(i+1)%p.size()], q), q) < EPS)</pre>
      return true:
    return false;
// compute intersection of line through points a and b with
// circle centered at c with radius r >
vector<PT> CircleLineIntersection(PT a, PT b, PT c, double r) {
  vector<PT> ret;
  b = b-a;
  a = a-c;
  double A = dot(b, b);
double B = dot(a, b);
  double C = dot(a, a) - r*r;
  double D = B*B - A*C;
  if (D < -EPS) return ret;</pre>
  ret.push_back(c+a+b*(-B+sqrt(D+EPS))/A);
  if (D > EPS)
    ret.push_back(c+a+b*(-B-sqrt(D))/A);
  return ret:
// compute intersection of circle centered at a with radius \boldsymbol{r}
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a, PT b, double r, double R) {
  vector<PT> ret:
  double d = sqrt(dist2(a, b));
  if (d > r+R | | d+min(r, R) < max(r, R)) return ret;</pre>
  double x = (d*d-R*R+r*r)/(2*d);
  double y = sqrt(r*r-x*x);
  PT v = (b-a)/d;
  ret.push_back(a+v*x + RotateCCW90(v)*y);
  if (y > \overline{0})
    ret.push_back(a+v*x - RotateCCW90(v)*y);
  return ret;
// This code computes the area or centroid of a (possibly nonconvex)
// polygon, assuming that the coordinates are listed in a clockwise or
// counterclockwise fashion. Note that the centroid is often known as // the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
  double area = 0:
  for(int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    area += p[i].x*p[j].y - p[j].x*p[i].y;
  return area / 2.0;
double ComputeArea(const vector<PT> &p) {
  return fabs(ComputeSignedArea(p));
PT ComputeCentroid(const vector<PT> &p)
 PT c(0.0);
  double scale = 6.0 * ComputeSignedArea(p);
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
  return c / scale;
 // tests whether or not a given polygon (in CW or CCW order) is simple
bool IsSimple(const vector<PT> &p) {
  for (int i = 0; i < p.size(); i++) {
  for (int k = i+1; k < p.size(); k++) {</pre>
     int j = (i+1) % p.size();
int l = (k+1) % p.size();
if (i = 1 | | j = k) continue;
if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
         return false;
  return true;
int main() {
  // expected: (-5,2)
  cerr << RotateCCW90(PT(2,5)) << endl;</pre>
  // expected: (5,-2)
  cerr << RotateCW90(PT(2,5)) << endl;</pre>
  // expected: (-5.2)
  cerr << RotateCCW(PT(2,5),M_PI/2) << endl;</pre>
  cerr << ProjectPointLine(PT(-5,-2), PT(10,4), PT(3,7)) << endl;</pre>
```

```
// expected: (5,2) (7.5,3) (2.5,1)
cerr << ProjectPointSegment(PT(-5,-2), PT(10,4), PT(3,7)) << " "</pre>
     << ProjectPointSegment(PT(7.5,3), PT(10,4), PT(3,7)) << " "
     << ProjectPointSegment(PT(-5,-2), PT(2.5,1), PT(3,7)) << endl;
// expected: 6.78903
cerr << DistancePointPlane(4,-4,3,2,-2,5,-8) << endl;</pre>
// expected: 1 0 1
cerr << LinesParallel(PT(1,1), PT(3,5), PT(2,1), PT(4,5)) << " "</pre>
     << LinesParallel(PT(1,1), PT(3,5), PT(2,0), PT(4,5)) << " "
     << LinesParallel(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
// expected: 0 0 1
<< LinesCollinear(PT(1,1), PT(3,5), PT(5,9), PT(7,13)) << endl;
cerr << SegmentsIntersect(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << " "
     << SegmentsIntersect(PT(0,0), PT(2,4), PT(4,3), PT(0,5)) << " "
     << SegmentsIntersect(PT(0,0), PT(2,4), PT(2,-1), PT(-2,1)) << " "
     << SegmentsIntersect(PT(0,0), PT(2,4), PT(5,5), PT(1,7)) << endl;
// expected: (1,2)
cerr << ComputeLineIntersection(PT(0,0), PT(2,4), PT(3,1), PT(-1,3)) << endl;</pre>
// expected: (1.1)
cerr << ComputeCircleCenter(PT(-3.4), PT(6.1), PT(4.5)) << endl;</pre>
vector<PT> v:
v.push_back(PT(0,0));
v.push_back(PT(5,0));
v.push_back(PT(5,5));
v.push_back(PT(0,5));
// expected: 1 1 1 0 0
cerr << PointInPolygon(v, PT(2,2)) << " "</pre>
     << PointInPolygon(v, PT(2,0)) << " "
     << PointInPolygon(v, PT(0,2)) << " "
     << PointInPolygon(v, PT(5,2)) << " "
     << PointInPolygon(v, PT(2,5)) << endl;
// expected: 0 1 1 1 1
cerr << PointOnPolygon(v, PT(2,2)) << " "
     << PointOnPolygon(v, PT(2,0)) << " "
     << PointOnPolygon(v, PT(0,2)) << " "
     << PointOnPolygon(v, PT(5,2)) << " "
     << PointOnPolygon(v, PT(2,5)) << endl;
             (5,4) (4,5)
             blank line
             (4.5) (5.4)
             blank line
             (4.5) (5.4)
vector<PT> u = CircleLineIntersection(PT(0,6), PT(2,6), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleLineIntersection(PT(0,9), PT(9,0), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
u = CircleCircleIntersection(PT(1,1), PT(10,10), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;
     CircleCircleIntersection(PT(1,1), PT(8,8), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 10, sqrt(2.0)/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl;</pre>
 u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 5, sqrt(2.0)/2.0); \\  for (int i = 0; i < u.size(); i++) cerr << u[i] << " "; cerr << endl; 
// area should be 5.0
// centroid should be (1.1666666, 1.166666)
PT pa[] = { PT(0,0), PT(5,0), PT(1,1), PT(0,5) };
vector<PT> p(pa, pa+4);
PT c = ComputeCentroid(p);
cerr << "Area: " << ComputeArea(p) << endl;
cerr << "Centroid: " << c << endl;</pre>
return 0;
```

4.7 Geometry 2

```
const double eps = 1e-8;
const double PI = acos(-1.0);
struct Point
```

```
double x, y;
  Point (double x = 0, double y = 0) : x(x), y(y) { }
  bool operator < (const Point& a) const
    if (a.x != x) return x < a.x;</pre>
    return y < a.y;
};
typedef Point Vector;
struct Line
  Point P:
  Vector v:
  double ang;
  Line() {}
  Line (Point P, Vector v) : P(P), v(v) { ang = atan2(v.y, v.x); }
  bool operator < (const Line& L) const
    return ang < L.ang;
};
Vector operator + (Vector A, Vector B) { return Vector(A.x + B.x, A.y + B.y); }
Vector operator - (Point A, Point B) { return Vector(A.x - B.x, A.y - B.y); }
Vector operator * (Vector A, double p) { return Vector(A.x*p, A.y*p); }
Vector operator / (Vector A, double p) { return Vector(A.x / p, A.y / p); }
int dcmp(double x)
  if (fabs(x) < eps) return 0; else return x < 0 ? -1 : 1;
bool operator == (const Point & a, const Point &b)
  return demp(a.x - b.x) == 0 && demp(a.y - b.y) == 0;
double Dot(Vector A, Vector B) { return A.x*B.x + A.y*B.y; }
double Length(Vector A) { return sgrt(Dot(A, A)); }
double Angle(Vector A, Vector B) { return acos(Dot(A, B) / Length(A) / Length(B)); }
double Cross(Vector A, Vector B) { return A.x*B.y - A.y*B.x; }
double Area2(Point A, Point B, Point C) { return fabs(Cross(B - A, C - A)) / 2; }
Vector Rotate (Vector A, double rad)
  return Vector(A.x*cos(rad) - A.y*sin(rad), A.x*sin(rad) + A.y*cos(rad));
Point GetLineIntersection (Point P. Vector v. Point O. Vector w)
  Vector u = P - Q;
  double t = Cross(w, u) / Cross(v, w);
 return P + v*t;
bool SegmentProperIntersection (Point al, Point a2, Point b1, Point b2)
 double c1 = Cross(a2 - a1, b1 - a1), c2 = Cross(a2 - a1, b2 - a1);
double c3 = Cross(b2 - b1, a1 - b1), c4 = Cross(b2 - b1, a2 - b1);
  return dcmp(c1) * dcmp(c2) < 0 && dcmp(c3) * dcmp(c4) < 0;
bool OnSegment (Point p, Point a1, Point a2)
  return dcmp(Cross(a1 - p, a2 - p)) == 0 && dcmp(Dot(a1 - p, a2 - p)) < 0;
double PolygonArea(Point* p, int n)
  double area = 0;
  for (int i = 1; i < n - 1; i++)
    area += Cross(p[i] - p[0], p[i + 1] - p[0]);
  return area / 2;
double PointDistanceToLine (Point P. Point A. Point B)
  Vector v1 = B - A, v2 = P - A;
  return fabs(Cross(v1, v2)) / Length(v1);
double PointDistanceToSegment(Point P, Point A, Point B)
  if (A == B) return Length(P - A);
```

```
Vector v1 = B - A, v2 = P - A, v3 = P - B;
  if (dcmp(Dot(v1, v2) < 0)) return Length(v2);</pre>
  else if (dcmp(Dot(v1, v3) > 0)) return Length(v3);
  else return fabs(Cross(v1, v2)) / Length(v1);
int isPointInPolygon(Point p, Point *poly, int n)
  int wn = 0;
  for (int i = 0; i < n; i++)
    const Point& p1 = poly[i], p2 = poly[(i + 1) % n];
    if (p == p1 || p == p2 || OnSegment(p, p1, p2)) return -1;
int k = dcmp(Cross(p2 - p1, p - p1));
   int d1 = dcmp(p1.y - p.y);
int d2 = dcmp(p2.y - p.y);
    if (k > 0 & d1 <= 0 & d2 > 0) wn++;
    if (k < 0 && d2 <= 0 && d1 > 0) wn--;
  if (wn != 0) return 1;
  return 0;
Vector Normal (Vector A)
  double L = Length(A);
  return Vector(-A.y / L, A.x / L);
double Dist2(Point p1, Point p2)
  return (p1.x - p2.x)*(p1.x - p2.x) + (p1.y - p2.y)*(p1.y - p2.y);
double RotatingCalipers(Point *P, int n)
  if (n == 1) return 0;
  if (n == 2) return Dist2(P[0], P[1]);
 P[n] = P[0];
  double ans = 0;
for (int u = 0, v = 1; u < n; u++)
    for (;;)
      double diff = Cross(P[u + 1] - P[u], P[v + 1] - P[v]);
      if (diff <= 0)
        ans = max(ans, Dist2(P[u], P[v]));
        if (diff == 0) ans = max(ans, Dist2(P[u], P[v + 1]));
        break:
      v = (v + 1) % n;
  return ans:
bool OnLeft (Line L, Point p)
  return Cross(L.v, p - L.P) > 0;
Point GetLineIntersection2(const Line &a, const Line &b)
  double t = Cross(b.v, u) / Cross(a.v, b.v);
  return a.P + a.v*t;
int HalfPlaneIntersection(Line* L, int n, Point* poly)
  sort(L, L + n);
  int first, last;
  Point *p = new Point[n];
  Line* q = new Line[n];
q[first = last = 0] = L[0];
  for (int i = 1; i < n; i++)
    while (first < last && !OnLeft(L[i], p[last - 1])) last--;</pre>
    while (first < last && !OnLeft(L[i], p[first])) first++;</pre>
    q[++last] = L[i];
    if (fabs(Cross(q[last].v, q[last - 1].v)) < eps)</pre>
      if (OnLeft(q[last], L[i].P)) q[last] = L[i];
    if (first < last) p[last - 1] = GetLineIntersection2(q[last - 1], q[last]);</pre>
  while (first < last && !OnLeft(q[first], p[last - 1])) last--;
  if (last - first <= 1) return 0;</pre>
  p[last] = GetLineIntersection2(q[last], q[first]);
```

```
int m = 0;
for (int i = first; i <= last; i++) poly[m++] = p[i];
return m;
}

vector<Point> CutPolygon(const vector<Point> &poly, Point A, Point B)
{
    vector<Point> newpoly;
    int n = poly.size();
    for (int i = 0; i < n; i++)
{
        Point C = poly[i], D = poly[(i + 1) % n];
        if (dcmp(Cross(B - A, C - A)) >= 0) newpoly.push_back(C);
        if (dcmp(Cross(B - A, C - D)) != 0)
        {
              Point ip = GetLineIntersection(A, B - A, C, D - C);
              if (OnSegment(ip, C, D)) newpoly.push_back(ip);
        }
    }
}
return newpoly;
```

4.8 Convex Hull

```
// Compute the 2D convex hull of a set of points using the monotone chain
// algorithm. Eliminate redundant points from the hull if REMOVE_REDUNDANT is
// #defined.
// Running time: O(n log n)
// INPUT: a vector of input points, unordered.
    OUTPUT: a vector of points in the convex hull, counterclockwise, starting
              with bottommost/leftmost point
#include <cstdio>
#include <cassert>
#include <vector>
#include <algorithm>
#include <cmath>
#include <map>
using namespace std:
#define REMOVE_REDUNDANT
typedef double T;
const T EPS = 1e-7;
struct PT {
  T x, y;
  PT() {}
 PT(T x, T y) : x(x), y(y) {}
bool operator<(const PT &rhs) const { return make_pair(y,x) < make_pair(rhs.y,rhs.x); }</pre>
 bool operator==(const PT &rhs) const { return make_pair(y,x) == make_pair(rhs.y,rhs.x); }
T cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
T area2(PT a, PT b, PT c) { return cross(a,b) + cross(b,c) + cross(c,a); }
#ifdef REMOVE REDUNDANT
bool between (const PT &a, const PT &b, const PT &c) {
  return (fabs(area2(a,b,c)) < EPS && (a.x-b.x) *(c.x-b.x) <= 0 && (a.y-b.y) *(c.y-b.y) <= 0);
void ConvexHull(vector<PT> &pts) {
 sort(pts.begin(), pts.end());
  pts.erase(unique(pts.begin(), pts.end()), pts.end());
  vector<PT> up, dn;
for (int i = 0; i < pts.size(); i++) {</pre>
    while (up.size() > 1 && area2(up[up.size()-2], up.back(), pts[i]) >= 0) up.pop_back();
    while (dn.size() > 1 && area2(dn[dn.size()-2], dn.back(), pts[i]) <= 0) dn.pop_back();</pre>
    up.push back(pts[i]);
    dn.push_back(pts[i]);
  for (int i = (int) up.size() - 2; i >= 1; i--) pts.push_back(up[i]);
#ifdef REMOVE_REDUNDANT
  if (pts.size() <= 2) return;</pre>
  dn.clear();
  dn.push_back(pts[0]);
  dn.push_back(pts[1]);
  for (int i = 2; i < pts.size(); i++) {</pre>
    if (between(dn[dn.size()-2], dn[dn.size()-1], pts[i])) dn.pop_back();
    dn.push_back(pts[i]);
  if (dn.size() >= 3 && between(dn.back(), dn[0], dn[1])) {
    dn[0] = dn.back();
    dn.pop_back();
```

```
pts = dn;
#endif
// The following code solves SPOJ problem #26: Build the Fence (BSHEEP)
 int t;
 scanf("%d", &t);
 for (int caseno = 0; caseno < t; caseno++) {</pre>
   int n;
   scanf("%d", &n);
   vector<PT> v(n);
for (int i = 0; i < n; i++) scanf("%lf%lf", &v[i].x, &v[i].y);</pre>
   vector<PT> h(v);
   map<PT, int> index;
   for (int i = n-1; i >= 0; i--) index[v[i]] = i+1;
   double len = 0;
   for (int i = 0; i < h.size(); i++) {</pre>
      double dx = h[i].x - h[(i+1)%h.size()].x;
      double dy = h[i].y - h[(i+1)%h.size()].y;
     len += sqrt (dx*dx+dy*dy);
   if (caseno > 0) printf("\n");
   printf("%.2f\n", len);
   for (int i = 0; i < h.size(); i++) {</pre>
     if (i > 0) printf(" ");
     printf("%d", index[h[i]]);
   printf("\n");
```

4.9 Convex Hull Diameter

```
typedef pair<double, double> point;
bool cw (const point &a, const point &b, const point &c) {
 return (b.first - a.first) * (c.second - a.second) - (b.second - a.second) * (c.first - a.first) <
       0;
vector<point> convexHull(vector<point> p) {
  int n = p.size();
  if (n <= 1)
   return p;
  int k = 0:
  sort(p.begin(), p.end());
  vector<point> q(n * 2);
  for (int i = 0; i < n; q[k++] = p[i++])
   for (; k \ge 2 \&\& !cw(q[k-2], q[k-1], p[i]); --k)
  for (int i = n - 2, t = k; i >= 0; q[k++] = p[i--])
for (; k > t && !cw(q[k-2], q[k-1], p[i]); --k)
  q.resize(k - 1 - (q[0] == q[1]));
  return q;
double area(const point &a, const point &b, const point &c) {
 return abs((b.first - a.first) * (c.second - a.second) - (b.second - a.second) * (c.first - a.first)
double dist(const point &a, const point &b) {
 return hypot (a.first - b.first, a.second - b.second);
double diameter(const vector<point> &p) {
  vector<point> h = convexHull(p);
  int m = h.size();
   return 0:
   return dist(h[0], h[1]);
  while (area(h[m-1], h[0], h[(k+1) % m]) > area(h[m-1], h[0], h[k]))
  double res = 0:
  for (int i = 0, j = k; i \le k \&\& j < m; i++) {
   res = max(res, dist(h[i], h[(j + 1) % m]));
     ++1;
```

```
return res;
}
int main() {
  vector<point> points(4);
  points[0] = point(0, 0);
  points[1] = point(3, 0);
  points[2] = point(0, 3);
  points[3] = point(1, 1);
  double d = diameter(points);
  cout < d << end;
}</pre>
```

4.10 Great Circle Distance

```
struct PT{double lat, lon; PT() {}
PT(double lat, double lon) : lat(lat), lon(lon) {}
PT operator * (double c ) const { return PT(lat*c, lon *c); }
}pts[1000+10];

const double eps = le-9;
const double PI = 3.141592653589793;
const double R = 6378.00 ; // radius of earth

double GCDist(PT pl, PT p2) {
    pl = p1*(PT/180.); p2 = p2*(PI/180.);
    double dlon = p2.lon - p1.lon;
    double dlat = p2.lat - p1.lat;
    double a = pow((sin(dlat/2)),2) + cos(p1.lat) * cos(p2.lat) * pow(sin(dlon/2), 2);
    double d = R * c;
    return d+eps;
}
```

5 String Algorithms

5.1 Suffix Array 1

```
// Suffix array construction in O(L log^2 L) time. Routine for
// computing the length of the longest common prefix of any two
// suffixes in O(log L) time.
// INPUT: string s
// OUTPUT: array suffix[] such that suffix[i] = index (from 0 to L-1)
            of substring s[i...L-1] in the list of sorted suffixes.
            That is, if we take the inverse of the permutation suffix[],
            we get the actual suffix array.
#include <vector>
#include <iostream>
#include <string>
using namespace std;
struct SuffixArray {
  const int L;
  string s;
  vector<vector<int> > P;
  vector<pair<int,int>,int> > M;
  SuffixArray(const string &s) : L(s.length()), s(s), P(1, vector<int>(L, 0)), M(L) {
    for (int i = 0; i < L; i++) P[0][i] = int(s[i]);
for (int skip = 1, level = 1; skip < L; skip *= 2, level++) {
      P.push back(vector<int>(L, 0));
      for (int i = 0; i < L; i++)
        M[i] = make_pair(make_pair(P[level-1][i], i + skip < L ? P[level-1][i + skip] : -1000), i);
      sort (M.begin(), M.end());
      for (int i = 0; i < L; i++)
        P[level][M[i].second] = (i > 0 && M[i].first == M[i-1].first) ? P[level][M[i-1].second] : i;
  vector<int> GetSuffixArray() { return P.back(); }
  // returns the length of the longest common prefix of s[i...L-1] and s[j...L-1]
  int LongestCommonPrefix(int i, int j) {
    int len = 0:
    if (i == j) return L - i;
for (int k = P.size() - 1; k >= 0 && i < L && j < L; k--) {</pre>
      if (P[k][i] == P[k][j]) {
        i += 1 << k;
        j += 1 << k;
```

```
len += 1 << k;
    return len;
// BEGIN CUT
// The following code solves UVA problem 11512: GATTACA.
#define TESTING
#ifdef TESTING
int main() {
 int T;
  cin >> T:
  for (int caseno = 0; caseno < T; caseno++) {</pre>
    string s;
    cin >> s;
    SuffixArray array(s);
    vector<int> v = array.GetSuffixArray();
    int bestlen = -1, bestpos = -1, bestcount = 0;
    for (int i = 0; i < s.length(); i++) {</pre>
      int len = 0, count = 0;
      for (int j = i+1, j < s.length(), j++) {
       int 1 = array.LongestCommonPrefix(i, j);
        if (1 >= len) {
          if (1 > len) count = 2; else count++;
          len = 1:
      if (len > bestlen || len == bestlen && s.substr(bestpos, bestlen) > s.substr(i, len)) {
        bestlen = len;
        bestcount = count;
        bestpos = i;
    if (bestlen == 0) {
      cout << "No repetitions found!" << endl;</pre>
    } else {
      cout << s.substr(bestpos, bestlen) << " " << bestcount << endl;</pre>
#else
// END CUT
int main() {
  // bobocel is the 0'th suffix
  // obocel is the 5'th suffix
       bocel is the 1'st suffix
        ocel is the 6'th suffix
        cel is the 2'nd suffix
         el is the 3'rd suffix
           l is the 4'th suffix
  SuffixArray suffix("bobocel");
vector<int> v = suffix.GetSuffixArray();
  // Expected output: 0 5 1 6 2 3 4
  for (int i = 0; i < v.size(); i++) cout << v[i] << " ";</pre>
  cout << endl;
  cout << suffix.LongestCommonPrefix(0, 2) << endl;</pre>
// BEGIN CUT
#endif
// END CUT
```

5.2 Suffix Array 2

```
if (pos[i] != pos[j])
   return pos[i] < pos[j];</pre>
   += gap;
 return (i < N && j < N) ? pos[i] < pos[j] : i > j;
void buildSA()
 N = strlen(S);
 REP(i, N) sa[i] = i, pos[i] = S[i];
 for (gap = 1;; gap *= 2)
   REP(i, N) pos[sa[i]] = tmp[i];
   if (tmp[N - 1] == N - 1) break;
void buildLCP()
  for (int i = 0, k = 0; i < N; ++i) if (pos[i] != N - 1)
   for (int j = sa[pos[i] + 1]; S[i + k] == S[j + k];)
     ++k;
   lcp[pos[i]] = k;
   if (k)--k;
```

5.3 Prefix Function

```
std::vector<int> prefix_function(const std::string& str) {
  std::vector<int> prefs(str.size(), 0);
  for (int i = 1; i < str.size(); ++i) {
  int pref = prefs[i - 1];</pre>
    while (pref > 0 && str[i] != str[pref]) {
      pref = prefs[pref - 1];
    if (str[i] == str[pref]) {
      ++pref;
    prefs[i] = pref;
  return prefs;
\verb|std::vector<| \textbf{int}>| z_function(| \textbf{const}|| std::string&| str)| | \{
  std::vector<int> zfunc(str.size(), 0);
  zfunc[0] = str.size();
  for (int i = 1, left = 0, right = 0; i < str.size(); ++i) {</pre>
    if (i <= right) {
      zfunc[i] = std::min(right - i + 1, zfunc[i - left]);
    while (i + zfunc[i] < str.size() && str[zfunc[i]] == str[i + zfunc[i]]) {</pre>
      ++zfunc[i];
    if (i + zfunc[i] - 1 > right) {
      left = i;
      right = i + zfunc[i] - 1;
  return zfunc;
std::string from_prefix_function(const std::vector<int>& prefs) {
  std::string str(prefs.size(), '.');
  char current_symbol = 'a';
for (int i = 0; i < prefs.size(); ++i) {</pre>
    if (prefs[i] > 0) {
      str[i] = str[prefs[i] - 1];
      str[i] = current_symbol++;
  return str;
std::vector<int> prefix_to_z(const std::vector<int>& prefs) {
 return z_function(from_prefix_function(prefs));
std::vector<int> z to prefix(const std::vector<int>& z func) {
  std::vector<int> prefs(z_func.size(), 0);
  for (int i = 1; i < z_func.size(); ++i) {</pre>
```

```
prefs[i + z_func[i] - 1] = std::max(prefs[i + z_func[i] - 1], z_func[i]);
}
for (int i = z_func.size() - 2; i >= 0; --i) {
    prefs[i] = std::max(prefs[i + 1] - 1, prefs[i]);
}
return prefs;
```

5.4 Infix to Postfix

```
import java.util.Stack;
public class ShuntingYard {
  public static void main(String[] args) {
    String infix = "3 + 4 * 2 / (1 - 5) ^{\circ} 2 ^{\circ} System.out.printf("infix: %s%n", infix);
    System.out.printf("postfix: %s%n", infixToPostfix(infix));
  static String infixToPostfix(String infix) {
    final String ops = "-+/*^";
StringBuilder sb = new StringBuilder();
    Stack<Integer> s = new Stack<>();
    for (String token : infix.split("\\s")) {
      if (token.isEmpty())
        continue;
       char c = token.charAt(0);
      int idx = ops.indexOf(c);
       // check for operator
       if (idx != -1) {
        if (s.isEmpty())
           s.push(idx);
         else (
           while (!s.isEmpty()) {
             int prec2 = s.peek() / 2;
int prec1 = idx / 2;
             if (prec2 > prec1 || (prec2 == prec1 && c != '^'))
                sb.append(ops.charAt(s.pop())).append('');
             else break;
           s.push(idx);
      else if (c == '(') {
   s.push(-2); // -2 stands for '('
       else if (c == ')') {
         // until '(' on stack, pop operators.
         while (s.peek() != -2)
          sb.append(ops.charAt(s.pop())).append(' ');
         s.pop();
      else {
         sb.append(token).append('');
    while (!s.isEmpty())
      sb.append(ops.charAt(s.pop())).append(' ');
    return sb.toString();
```

5.5 Knuth-Morris-Pratt (KMP)

```
public class Kmp {

public static int[] prefixFunction(String s) {
   int[] p = new int[s.length()];
   int k = 0;
   for (int i = 1; i < s.length(); i++) {
      while (k > 0 && s.charAt(k) != s.charAt(i))
        k = p[k - 1];
      if (s.charAt(k) == s.charAt(i))
        ++k;
      p[i] = k;
   }
   return p;
}
```

```
public static int kmpMatcher(String s, String pattern) {
     int m = pattern.length();
       return 0;
     int[] p = prefixFunction(pattern);
     for (int i = 0, k = 0; i < s.length(); i++)</pre>
       for (; ; k = p[k - 1]) {
         if (pattern.charAt(k) == s.charAt(i)) {
           if (++k == m)
              return i + 1 - m;
            break:
         if (k == 0)
           break:
     return -1;
// Competitive Programming KMP
#define MAX_N 100010
char T[MAX_N], P[MAX_N]; // T = text, P = pattern
int b[MAX_N], n, m; // b = back table, n = length of T, m = length of P
 roid kmpFreprocess() {
   int i = 0, j = -1; b[0] = -1; // call this before calling kmpSearch()
   while (i < m) { // starting values
   while (j >= 0 && P[i] != P[j]) j = b[j]; // pre-process the pattern string P
   i++; j++; // if different, reset j using b
     // if same, advance both pointers
     b[i] = j; // observe i = 8, 9, 10, 11, 12 with j = 0, 1, 2, 3, 4
// in the example of P = "SEVENTY SEVEN" above
void kmpSearch() {
  // this is similar as kmpPreprocess(), but on string T
  int i = 0, j = 0;
  // starting values while (i < n) { // search through string T
     while (j >= 0 && T[i] != P[j]) j = b[j]; // if different, reset j using b
    if (j == m) { // if same, advance both pointers
    printf("P is found at index %d in T\n", i - j); // a match found when j == m
       j = b[j]; // prepare j for the next possible match
```

5.6 Simple Parser

```
const char * expressionToParse = "3*2+4*1+(4+9)*6";
char peek(){
 return *expressionToParse;
char get(){
 return *expressionToParse++;
int expression();
int number(){
  int result = get() - '0';
  while (peek() >= '0' && peek() <= '9'){
   result = 10 * result + get() - '0';
 return result:
int factor(){
  if (peek() >= '0' && peek() <= '9')</pre>
   return number();
  else if (peek() == '('){
   get(); // '('
    int result = expression();
    get(); // ')'
    return result;
  else if (peek() == '-'){
    get();
    return -factor();
  return 0; // error
int term() {
 int result = factor();
```

```
while (peek() == '*' || peek() == '/')
   if (get() == '*')
      result += factor();
   else
      result /= factor();
   return result;
}

int expression(){
   int result = term();
   while (peek() == '+' || peek() == '-')
      if (get() == '+')
      result += term();
      else
      result -= term();
   return result;
}

int _tmain(int argc, _TCHAR* argv[]){
   int result = expression();
   return 0;
```

5.7 Longest Palindrome

```
using namespace std;
template <class RAI1, class RAI2>
void fastLongestPalindromes(RAI1 seq, RAI1 seqEnd, RAI2 out)
 int seqLen = seqEnd - seq;
int i = 0, j, d, s, e, lLen, k = 0;
int palLen = 0;
  while (i<seqLen)
    if (i>palLen && seq[i - palLen - 1] == seq[i])
      continue;
    out[k++] = palLen;
    s = k - 2;
    e = s - palLen;
    bool b = true;
    for (j = s; j>e; j--)
      d = j - e - 1;
      if (out[j] == d) {
   palLen = d;
         b = false;
        break;
      out[k++] = min(d, out[j]);
    if (b)
      palLen = 1;
      i++;
  out[k++] = palLen;
```

```
lLen = k;
  s = 1Len - 2;
  e = s - (2 \star seqLen + 1 - lLen);
  for (i = s; i>e; i--)
    d = i - e - 1;
    out[k++] = min(d, out[i]);
//Example
//[0, 1, 0, 1, 4, 1, 0, 1, 0, 1, 0, 3, 0, 1, 0]
//Longest palindrome has length 4
int main()
  string s; cin >> s;
  vector<int> V(2 * s.length() + 1);
  fastLongestPalindromes(s.begin(), s.end(), V.begin());
  int best = 0;
  cout << "[";
for (int i = 0; i<V.size(); i++)
    if (i>0) cout << ", ";</pre>
    cout << V[i];
    best = max(best, V[i]);
  cout << "]" << endl << "Longest palindrome has length " << best << endl;
  return 0:
```

6 Miscellaneous

6.1 Notes

```
/* ----- Bitmask
bit&(1<<i) // bit i is 0 or 1
(bit>>j)&1// bit i is 0 or 1 // use this & multiplication to avoid TLE
bit|(1<<i) // set bit i to 1
bit^(1<<i) // toggle bit i
x & ( x 1) // check if x is a power of 2
string stmp; bitset<12> tmp; //Debuging
tmp=bit; stmp=tmp.to_string();
/* ----- Data Structure Ideas
-Hash Table + Lookup
- Sparse Table
- SQRT Decomposition
- Bucketing
- Interger Arrays as matrices
- Recursive Tree Building
- Shortest Cycles
- Problem DAG
```