1. Union Find Disjoint Set & Kruskal

```
struct DisJointSet{
 vector<int> par,rnk,cnt; int numOfsets;
 DisJointSet(int n){
   par.assign(n,-1); rnk.assign(n,0); cnt.assign(n,1); //par==parent
  numOfsets=n; // if wanna count number of disjoint sets
 int Find(int a){
  int i=a,j=a,tmp;
   while(par[i]!=-1){ i=par[i]; }
   while(par[j]!=-1){ tmp=par[j]; par[j]=i; j=tmp; } //path compression
   return i;
 int Uni(int a, int b){
   int A=Find(a),B=Find(b);
   if(A!=B){
    if(rnk[A]<rnk[B]) swap(A,B); // union using rank</pre>
    if(rnk[A]==rnk[B]) rnk[A]++;
    par[B]=A;
    cnt[A]+=cnt[B]; // if we wanna count each set size
    numOfsets--; // if wanna count number of disjoint sets
   return cnt[A]; // if we wanna count each set size
 }
};
struct edge{ int u, v, w;
 edge(int u=0, int v=0, int w=0):u(u), v(v), w(w){};
 bool operator<(const edge& b) const {</pre>
  if(w == b.w \&\& v == b.w) return u < b.u;
   if(w == b.w ) return v < b.v; return w < b.w;</pre>
};
int n, m, bit; vector<edge> e; vi marked;
int Kruskal(){
 DisJointSet djst(n); marked.clear();
 sort(e.begin(), e.end()); int ans=0; int j=0;
 for(int i=0; i<e.size() && j<n-1; i++){</pre>
  if(djst.Find(e[i].u) != djst.Find(e[i].v)){
    djst.Uni(e[i].u, e[i].v); ans+=e[i].w; j++; marked.push back(i);
```

```
return ans;
```

2. Segment Tree

```
class SegmentTree { // the segment tree is stored like a heap array
private: vi st, A; // recall that vi is: typedef vector<int> vi;
    int n:
    int left(int p) { return p << 1; } // same as binary heap operations</pre>
    int right(int p) { return (p << 1) + 1; }</pre>
    void build(int p, int L, int R) { // O(n)
     if (L == R) // as L == R, either one is fine
      st[p] = L; // store the index
     else { // recursively compute the values
      build(left(p), L, (L + R) / 2);
      build(right(p), (L + R) / 2 + 1, R);
      int p1 = st[left(p)], p2 = st[right(p)];
      st[p] = (A[p1] \leftarrow A[p2]) ? p1 : p2;
    int rmg(int p, int L, int R, int i, int j) { // O(log n)
     if (i > R | | j < L) return -1; // current segment outside query
     if (L >= i && R <= j) return st[p]; // inside query range
                  // compute the min position in the left and right part
of the interval
     int p1 = rmq(left(p), L, (L + R) / 2, i, j);
     int p2 = rmq(right(p), (L + R) / 2 + 1, R, i, j);
     if (p1 == -1) return p2; // if we try to access segment outside
query
     if (p2 == -1) return p1; // same as above
     return (A[p1] <= A[p2]) ? p1 : p2; // as in build routine</pre>
public:
 SegmentTree(const vi & A) {
  A = _A; n = (int)A.size(); // copy content for local usage
  st.assign(4 * n, 0); // create large enough vector of zeroes
  build(1, 0, n - 1); // recursive build
 int rmq(int i, int j) { return rmq(1, 0, n - 1, i, j); } // overloading
int main() {
 int arr[] = { 18, 17, 13, 19, 15, 11, 20 }; // the original array
 vi A(arr, arr + 7);
 SegmentTree st(A);
 printf("RMQ(1, 3) = %d\n", st.rmq(1, 3)); // answer = index 2
 printf("RMQ(4, 6) = %d\n", st.rmq(4, 6)); // answer = index 5
```

3. Counting inversions vector<int> tree,a,b; int n;

```
int64 read(int idx){
 int64 sum=0;
 while(idx>0){
   sum+=tree[idx]; idx-=(idx & -idx);
 return sum;
void update(int idx, int val){
 int64 sum=0;
 while(idx<n){</pre>
   tree[idx]+=val; idx+=(idx & -idx);
}
// get largest value with cumulative sum less than or equal to x;
// for smallest, pass x-1 and add 1 to result
int getind(int x) {// ***Change Needed***
  int idx = 0, mask = TREE SIZE; //(must be a power of 2)
  while(mask && idx < TREE SIZE) {</pre>
    int t = idx + mask;
    if(x >= tree[t]) \{idx = t; x -= tree[t]; \}
    mask >>= 1;
  }
  return idx;
int main(){
 while(cin >> n){
   a.assign(n,0); b.assign(n,0); tree.assign(n,0);
   for(int i=0 ; i<n; i++){</pre>
     cin >> a[i]; b[i]=a[i];
   sort(b.begin(),b.end());
   for(int i=0 ; i<n ; i++){</pre>
    int rank=(int)(lower bound(b.begin(),b.end(),a[i])-b.begin());
     a[i]=rank+1;
   int64 invs=0;//num of inversions
   for(int i=n-1; i>=0; i--){
    invs+=read(a[i]-1);
     update(a[i],1);
   cout << invs << endl;</pre>
 return 0;
```

4. FenwickTree

```
class FenwickTree {
private: vi ft; // recall that vi is: typedef vector<int> vi;
public: FenwickTree(int n) { ft.assign(n + 1, 0); } // init n + 1
zeroes
  int rsq(int b) { // returns RSQ(1, b)
    int sum = 0; for (; b; b -= LSOne(b)) sum += ft[b];
    return sum;
  } // note: LSOne(S) (S & (-S))
  int rsq(int a, int b) { // returns RSQ(a, b)
    return rsq(b) - (a == 1 ? 0 : rsq(a - 1));
  // adjusts value of the k-th element by v (v can be +ve/inc or -
ve/dec)
  void adjust(int k, int v) { // note: n = ft.size() - 1
    for (; k < (int)ft.size(); k += LSOne(k)) ft[k] += v;
};
int main() {
 int f[] = \{ 2,4,5,5,6,6,6,7,7,8,9 \}; // m = 11 scores
 FenwickTree ft(10); // declare a Fenwick Tree for range [1..10]
        // insert these scores manually one by one into an empty
Fenwick Tree
 for (int i = 0; i < 11; i++) ft.adjust(f[i], 1); // this is 0(k \log n)
 printf("%d\n", ft.rsq(1, 1)); // 0 => ft[1] = 0
 printf("%d\n", ft.rsq(1, 2)); // 1 => ft[2] = 1
 printf("%d\n", ft.rsq(1, 6)); // 7 => ft[6] + ft[4] = 5 + 2 = 7
 printf("%d\n", ft.rsq(1, 10)); // 11 => ft[10] + ft[8] = 1 + 10 = 11
 printf("%d\n", ft.rsq(3, 6)); // 6 => rsq(1, 6) - rsq(1, 2) = 7 - 1
 ft.adjust(5, 2); // update demo
 printf("%d\n", ft.rsq(1, 10)); // now 13
} // return 0;
```

5. Programming Tips

Lower bound and upper bound for binary Search

- lower bound Returns an iterator pointing to the first element in the range[first, last) which does not compare less than val.

- upper bound Returns an iterator pointing to the first element in the range[first, last) which compares greater than val.

```
map<string, int> dict;
class cmp {
public:
   bool operator()(const string& a, const string& b) const {
    return dict[a] < dict[b];
   }
};
typedef set<string, cmp> sset;
```

6. Maximum Subrectangle Sum

```
for(int i=1; i<n; i++)//preprocess</pre>
     for(int j=0 ; j<n ; j++)</pre>
       a[i][j]+=a[i-1][j];
int Max=0, ans=0;
for(int k=0 ; k<n ; k++){//calc</pre>
 for(int i=0 ; i<n-k ; i++){ Max=0;</pre>
   for(int j=0 ; j<n ; j++){</pre>
     if(Max<0) Max=a[i+k][j]-a[i][j];</pre>
     else Max+=a[i+k][j]-a[i][j];
     if(Max>ans) ans=Max;
} } }
//sub array, finsh and start point p=(val, startidx, finishidx)
p ans=p(-1,0,0); int sum=0,id=1;
for(int i=1 ; i<n ; i++){</pre>
     if(sum<0){sum=0; id=i;}</pre>
     sum+=a[i];
     p tmp=p(sum,id,i+1); ans=Max(ans,tmp);
}
```

7. Optimal Array Multiplication Sequence (Print Path)

```
int n,a[10+5],p[10+5][10+5],dp[10+5][10+5];
int solve(int L, int R){
 if(L==R){ return 0; }
 if(dp[L][R]!=-1) return dp[L][R];
 int Min=INF;
 for(int i=L ; i<R ; i++){</pre>
   int slv=solve(L,i)+solve(i+1,R)+a[(L-1)]*a[i]*a[R];
   if(Min>slv) Min=slv; p[L][R]=i;
 return dp[L][R]=Min;
//prints like this => (A1 x (A2 x A3))
void print(int L, int R){
 if(L==R){ cout << "A" <<L; return; }</pre>
 cout << "("; print(L,p[L][R]);</pre>
 cout << " x ";
 print(p[L][R]+1,R); cout << ")";</pre>
int main(){ int t=1;
 while(cin >> n && n){
   for(int i=1; i<=n; i++)cin >> a[i-1] >> a[i];
   memset(dp,-1,sizeof dp);
   solve(1,n);//cout << solve(1,n) << endl;</pre>
   printf("Case %d: ",t++); print(1,n); printf("\n");
 return 0;
```

8. LIS

```
vector<int> v;
v.push back(inf);
for (int i = 0; i<n; i++) {</pre>
 int x = dolls[i].w; // array element
 int id = lower bound(v.begin(), v.end(), x + 1) - v.begin();
 if (id == v.size() - 1) v.push back(inf); v[id] = x;
cout << v.size() - 1 << endl;</pre>
9. LCS
dp[MAX][MAX] = \{0\};
for(int i=1; i<=n; i++){</pre>
     for(int j=1 ; j<=n ; j++){</pre>
      if(a[i-1]==b[j-1]) dp[i][j]=dp[i-1][j-1]+1;
      else dp[i][j]=max(dp[i-1][j],dp[i][j-1]);
}}
cout << dp[n][n] << endl;</pre>
10.TSP
p a[15]; int n, dp[15][1<<15];
int solve(int pos, int bitset){
 int& dpp=dp[pos][bitset]; //dpp = dp poniter
 if(bitset==(1<<n)-1) return dist(a[pos],a[0]);</pre>
 if(dpp!=-1) return dpp;
 dpp=INF;
 for(int i=0 ; i<n ; i++){</pre>
  if(!(bitset&(1<<i))) dpp=min(dpp,solve(i,bitset|(1<<i))+dist(a[pos],a[i]));</pre>
 return dpp;
int main(){
 int tc; cin >> tc;
 while(tc--){
   cin >> a[0].X >> a[0].Y; cin >> n; n++;
   for(int i=1; i<n; i++) cin >> a[i].X >> a[i].Y;
   memset(dp, -1, sizeof dp);
   cout << solve(0,1) << endl;</pre>
 return 0;
```

11. Articulation Points & Bridges

```
int n, lev, dfsRoot, rootChilds;
int dfsLow[MAX], dfsNum[MAX], parent[MAX];
vvi adj; set<pii> bridges; set<int> artPoints;
void dfs(int u) {
  dfsLow[u] = dfsNum[u] = lev++;
  for (int i = 0; i<adj[u].size(); i++) {</pre>
   int v = adj[u][i];
   if (dfsNum[v] == 0) {
     if (u == dfsRoot) rootChilds++;
     parent[v] = u; dfs(v);
     if (dfsLow[v] >= dfsNum[u] && u != dfsRoot)//u is articulation point
      artPoints.insert(u);
     if (dfsLow[v] > dfsNum[u]) {
       bridges.insert(pii(v, u));
       bridges.insert(pii(u, v));
     dfsLow[u] = min(dfsLow[u], dfsLow[v]);
   else if (parent[u] != v)
     dfsLow[u] = min(dfsLow[u], dfsNum[v]);
}
int main() {
 while (cin >> n) {
   adj.assign(n, vi()); //initialization
   memset(dfsLow, 0, sizeof dfsLow);
   memset(dfsNum, 0, sizeof dfsNum);
   memset(parent, 0, sizeof parent);
   bridges.clear(); artPoints.clear();
   lev = 1; int tmp, u, m;
   for (int i = 0; i<n; i++) { // construct the graph</pre>
     scanf("%d (%d", &u, &m); cin.ignore();
     for (int i = 0; i<m; i++) {
       cin >> tmp; adj[u].push back(tmp);
   for (int i = 0; i<n; i++) {
     if (dfsNum[i] == 0) {
       dfsRoot = i; rootChilds = 0; dfs(i);
      if (rootChilds >= 2) artPoints.insert(dfsRoot);
   printf("%d critical links\n", bridges.size());
   set<pii>::iterator itr; // print answer
   for (itr = bridges.begin(); itr != bridges.end(); itr++)
     printf("%d - %d\n", itr->first, itr->second);
   cout << endl;</pre>
  return 0;
```

12. Finding Strongly Connected Components

```
#define MAX 100000
using namespace std;
int dfsNum[MAX+10],dfsLow[MAX+10],vis[MAX+10],in[MAX+10],n,m,lev,ans;
vector<int> SCC,adj[MAX+10];
void dfs(int u){
 dfsLow[u]=dfsNum[u]=lev++; vis[u]=1; SCC.push_back(u);
 for(int i=0 ; i<adj[u].size() ; i++){</pre>
  int v=adj[u][i];
  if(dfsNum[v]==0) dfs(v);
  if(vis[v]) dfsLow[u]=min(dfsLow[u], dfsLow[v]), in[v]--;
 if(dfsLow[u]==dfsNum[u]){
// this prints all vertices v blong to SCC with dfsLow[v] == dfsLow[u]
   bool flag=true;
  for(int i=0, v ; !SCC.empty() ; i++){
    v=SCC.back(); SCC.pop back(); vis[v]=0;
    printf("%d ", v);
    if(in[v]) flag=false;
    if(v==u) break;
  printf("\n");
  if(flag) ans++;
// counts number of SCCs without indegree outside of other SCCs
}
int main(){
 int tc; scanf("%d", &tc);; int x,y;
 while(tc--){
  scanf("%d %d", &n, &m);
  memset(dfsNum,0,sizeof dfsNum); // memset(adj,0,sizeof adj);
  memset(dfsLow,0,sizeof dfsLow); memset(vis,0,sizeof vis);
   memset(in,0,sizeof in); lev=1; ans=0;
  for(int i=0, j=0; i<m; i++){
    scanf("%d %d", &x, &y); x--; y--;
    adi[x].push back(y); in[y]++;
  for(int i=0 ; i<n ; i++){</pre>
    if(dfsNum[i]==0) dfs(i);
   cout << ans << endl;</pre>
 return 0;
```

13. Graphic Sequence

```
// given a sequence of integers see if it's a sequence of degrees of graph or not.
int a[10010]; long long sum, Min;;
int main(){
 int n;
  while(cin >> n && n){
   for(int i=0 ; i<n ; i++) scanf("%d",&a[i]);</pre>
   sort(a,a+n, ::greater<int>() );
   bool possible=true; sum=0;
   for(int i=0 ; i<n ; i++){</pre>
     sum+=a[i]; Min=0;
     for(int j=i+1; j<n; j++) Min+=min(a[j],i+1);</pre>
     if(sum>i*(i+1)+Min){
       possible=false;
       break;
   if(!possible || sum%2) cout << "Not possible" << endl;</pre>
   else cout << "Possible" << endl;</pre>
 return 0;}
```

14. BFS Topological Sort

```
//store indegree of vertice u in indegree[u]
fr(i,n) if(!indegree[i]) q.push(i);
while(!q.empty()){
  int v = q.front(); q.pop();
  cout << v + 1 << " ";
  int s = adjlist[v].size();
  fr(i,s){
   if(!(--indegree[ adjlist[v][i] ])) q.push(adjlist[v][i]);
  }
}</pre>
```

15.Floyd Warshal (Print Path)

```
#define MAX (100+10)
int adj[MAX][MAX],path[MAX][MAX]; int n;
void print(int i,int j){
 if(i!=j){
  printf(" %d",i );
  print(path[i][j],j);
int main(){
 int tc; cin >> tc;
 while(tc--){
  cin >> n;
  for(int i=0 ; i<n ; i++){</pre>
    for(int j=0 ; j<n ; j++){</pre>
      adj[i][j]=1e9; if(i==j) adj[i][j]=0;
      path[i][j]=j;//initial parent
  for(int k=0 ; k<n ; k++){</pre>
    for(int i=0 ; i<n ; i++){</pre>
      for(int j=0 ; j<n ; j++){</pre>
       if(adj[i][j]>adj[i][k]+adj[k][j]){
         adj[i][j]=adj[i][k]+adj[k][j];
         path[i][j]=path[i][k];//set parent
    }
  int s,d;
  cin >> s >> d;
  printf("%d euros\n",adj[s][d]);
  //this prints the path even if source and distinaion are same
  printf("%d",s); print(path[s][d],d); printf(" %d\n",d);
 return 0;
```

16. Edmonds Karp's //UVa 820 - Internet Bandwidth #define INF (int)1e9 #define MAX 100+10 using namespace std; int res[MAX][MAX],mf,f,s,t,n,m,par[MAX]; vector<int> dist,adj[MAX]; void agument(int v, int minEdge){ if(v==s) f=minEdge; else if(par[v]!=-1){ agument(par[v],min(minEdge,res[par[v]][v])); res[par[v]][v]-=f; res[v][par[v]]+=f; } int main(){ int tc=1: while(cin >> n && n){ mf=0; memset(res,0,sizeof res); for(int i=0; i<n; i++)</pre> adj[i].clear(); cin >> s >> t >> m; s--; t--; int u,v,c; while(m--){ cin >> u >> v >> c; u--; v--; res[u][v]+=c; res[v][u]+=c; adj[u].push back(v); adj[v].push back(u); while(1){ f=0; memset(par,-1,sizeof par); dist.assign(n,INF); dist[s]=0; queue<int> q; q.push(s); while(!q.empty()){ int u=q.front(); q.pop(); if(u==t) break; for(int i=0 ; i<adj[u].size(); i++){</pre> int v=adj[u][i]; if(res[u][v]>0 && dist[v]==INF){ dist[v]=dist[u]+1; q.push(v); par[v]=u; agument(t,INF); if(f==0) break; mf+=f;printf("Network %d\n", tc++); printf("The bandwidth is %d.\n\n", mf); return 0;}

17. Dinic

```
// Adjacency list implementation of Dinic's blocking flow algorithm.
// This is very fast in practice, and only loses to push-relabel flow.
// Running time: 0(|V|^2 |E|)
// INPUT:
     - graph, constructed using AddEdge() - source - sink
// OUTPUT:
      - maximum flow value
       - To obtain the actual flow values, look at all edges with
        capacity > 0 (zero capacity edges are residual edges).
using namespace std:
const int INF = 20000000000:
struct Edge {
 int from, to, cap, flow, index;
 Edge(int from, int to, int cap, int flow, int index) :
   from(from), to(to), cap(cap), flow(flow), index(index) {}
struct Dinic {
 int N; vector<vector<Edge> > G;
  vector<Edge *> dad; vector<int> 0;
  Dinic(int N) : N(N), G(N), dad(N), Q(N) {}
  void AddEdge(int from, int to, int cap) {
   G[from].push back(Edge(from, to, cap, 0, G[to].size()));
   if (from == to) G[from].back().index++;
   G[to].push back(Edge(to, from, 0, 0, G[from].size() - 1));
  long long BlockingFlow(int s, int t) {
   fill(dad.begin(), dad.end(), (Edge *)NULL);
   dad[s] = &G[0][0] - 1;
   int head = 0, tail = 0;
   Q[tail++] = s;
   while (head < tail) {</pre>
     int x = Q[head++];
     for (int i = 0; i < G[x].size(); i++) {</pre>
       Edge &e = G[x][i];
       if (!dad[e.to] && e.cap - e.flow > 0) {
         dad[e.to] = &G[x][i];
         Q[tail++] = e.to;
   if (!dad[t]) return 0;
   long long totflow = 0:
   for (int i = 0; i < G[t].size(); i++) {</pre>
     Edge *start = &G[G[t][i].to][G[t][i].index];
     int amt = INF:
     for (Edge *e = start; amt && e != dad[s]; e = dad[e->from]) {
       if (!e) { amt = 0; break; }
       amt = min(amt, e->cap - e->flow);
     if (amt == 0) continue;
     for (Edge *e = start; amt && e != dad[s]; e = dad[e->from]) {
       e->flow += amt;
       G[e->to][e->index].flow -= amt;
     totflow += amt;
   return totflow:
  long long GetMaxFlow(int s, int t) {
   long long totflow = 0;
   while (long long flow = BlockingFlow(s, t))
     totflow += flow:
   return totflow;
};
```

18. Min Cut // Adjacency matrix implementation of Stoer-Wagner min cut algorithm. // Running time:O(|V|^3) // INPUT: // graph, constructed using AddEdge() // OUTPUT: //(min cut value, nodes in half of min cut) #include <cmath> #include <vector> #include <iostream> using namespace std; typedef vector<int> VI; typedef vector<VI> VVI; const int INF = 10000000000; pair<int, VI> GetMinCut(VVI &weights) { int N = weights.size(); VI used(N), cut, best_cut; int best_weight = -1; for (int phase = N - 1; phase >= 0; phase--) { VI w = weights[0]; VI added = used; int prev, last = 0; for (int i = 0; i < phase; i++) {</pre> prev = last: last = -1;for (int j = 1; j < N; j++)if (!added[j] && (last == -1 || w[j] > w[last])) last = j; if (i == phase - 1) { for (int j = 0; j < N; j++) weights[prev][j] += weights[last][j];</pre> for (int j = 0; j < N; j++) weights[j][prev] = weights[prev][j];</pre> used[last] = true; cut.push back(last): if (best weight == -1 || w[last] < best weight) {</pre> best cut = cut; best weight = w[last]; else { for (int i = 0; i < N; i++) w[j] += weights[last][j]; added[last] = true; return make pair(best weight, best cut);

19. Alternating Path Algorithm for Max Bipartite Matching

```
//UVa 11138 - Nuts and Bolts // O(V^2 + VE)
#define vi vector<int>
using namespace std;
vector< vi > adj; vector<int> owner, vis; int n,b;
int altpath(int u){
 if(vis[u]) return 0; vis[u]=1;
 for(int i=0 ; i<adj[u].size() ; i++){</pre>
   int v=adj[u][i];
   if(owner[v]==-1 || altpath(owner[v])){
    owner[v]=u; return 1;
 return 0;
int main(){
 int tmp,tc,t=1; cin >> tc;
 while(tc--){
   cin >> n >> b; adj.assign(n+b,vi());
   for(int i=0 ; i<n ; i++){</pre>
    for(int j=0 ; j<b ; j++){</pre>
      cin >> tmp; if(tmp==1) adj[i].push back(j+n);
    }
   int ans=0; owner.assign(n+b,-1);
   for(int u=0 ; u<n ; u++){</pre>
    vis.assign(n,0); ans+=altpath(u);
   printf("Case %d: a maximum of %d matched\n", t++, ans);
 return 0;
20. Bitmask
bit&(1<<i) // bit i is 0 or 1
(bit>>j)&1// bit i is 0 or 1 // use this & multiplication to avoid TLE
bit | (1<<i) // set bit i to 1
bit^(1<<i) // toggle bit i
x & (x - 1) // \text{ check if } x \text{ is a power of } 2
string stmp; bitset<12> tmp; //Debuging
tmp=bit; stmp=tmp.to_string();
```

21.MinCost Max Flow dad[k] = make pair(s, dir); width[k] = min(cap, width[s]); // forward and reverse edges separately (so you can set cap[i][j] != } // cap[i][i]). For a regular max flow, set all edge costs to 0. // Running time, $O(|V|^2)$ cost per augmentation L Dijkstra(int s, int t) { max flow: $O(|V|^3)$ augmentations fill(found.begin(), found.end(), false); min cost max flow: $O(|V|^4 * MAX EDGE COST)$ augmentations fill(dist.begin(), dist.end(), INF); // INPUT: fill(width.begin(), width.end(), 0); - graph, constructed using AddEdge() - source - sink dist[s] = 0;// OUTPUT: width[s] = INF; // - (maximum flow value, minimum cost value) // - To obtain the actual flow, look at positive values only. while (s != -1) { int best = -1; #include <cmath> found[s] = true; #include <vector> for (int k = 0; k < N; k++) { #include <iostream> if (found[k]) continue; using namespace std; Relax(s, k, cap[s][k] - flow[s][k], cost[s][k], 1); Relax(s, k, flow[k][s], -cost[k][s], -1); typedef vector<int> VI; if (best == -1 || dist[k] < dist[best]) best = k;</pre> typedef vector<VI> VVI; typedef long long L; s = best; typedef vector<L> VL; typedef vector<VL> VVL; typedef pair<int, int> PII; for (int k = 0; k < N; k++) typedef vector<PII> VPII; pi[k] = min(pi[k] + dist[k], INF);return width[t]; const L INF = numeric limits<L>::max() / 4; struct MinCostMaxFlow { pair<L, L> GetMaxFlow(int s, int t) { int N; VPII dad; L totflow = 0, totcost = 0; VVL cap, flow, cost; while (L amt = Dijkstra(s, t)) { VI found; VL dist, pi, width; totflow += amt; for (int x = t; x != s; x = dad[x].first) { if (dad[x].second == 1) { MinCostMaxFlow(int N) : flow[dad[x].first][x] += amt; N(N), cap(N, VL(N)), flow(N, VL(N)), cost(N, VL(N)), totcost += amt * cost[dad[x].first][x]; found(N), dist(N), pi(N), width(N), dad(N) {} else { void AddEdge(int from, int to, L cap, L cost) { flow[x][dad[x].first] -= amt; this->cap[from][to] = cap; totcost -= amt * cost[x][dad[x].first]; this->cost[from][to] = cost; } void Relax(int s, int k, L cap, L cost, int dir) { return make pair(totflow, totcost); L val = dist[s] + pi[s] - pi[k] + cost; if (cap && val < dist[k]) {</pre> **}**; dist[k] = val;

22.Dijkstra struct ToNode { int v, w; ToNode(int v, int w) $:v(v), w(w) {}$ struct QEntry { int node, cost; QEntry(int node, int cost):node(node), cost(cost) {} bool operator<(const QEntry& op) const {</pre> return cost < op.cost;</pre> }; int n, m; vvtn adj; int dijkstra(int s, int t, vi& dist) { dist.assign(n, INF); priority queue<QEntry> q; q.push(QEntry(s, 0)); dist[s] = 0;while (!q.empty()) { QEntry u = q.top(); q.pop(); if (u.node == t) return u.cost; if (u.cost > dist[u.node]) continue; for (int i = 0; i < adj[u.node].size(); i++) {</pre> QEntry v(adj[u.node][i].v, u.cost + adj[u.node][i].w); if (dist[v.node] > v.cost) { dist[v.node] = v.cost; q.push(v); } return INF; 23.Catalan Catalan(n+1)=(catalan(n)*(2n+2)*(2n+1))/((n+1)*(n+2)) /*n!*/Catalan(n)=2n!/(n!*n!*(n+1)), Catalan(1)=1; 000 001 002 003 004 005 006 001 001 002 005 014 042 132

24. Strongly Connected Componnent - Kosaraju

```
// Doesn't run properly
vvi adjOrg, adjRev; vi vis, ord, col;
void dfsOrg(int u) {
 if (vis[u]) return; vis[u] = true;
 for (int i = 0; i < adjOrg[u].size(); i++) {</pre>
   dfsOrg(adjOrg[u][i]);
 ord.push_back(u);
int dfsRev(int u, int color) {
 if (col[u]) return 0; col[u] = color;
 int ret = 1;
 for (int i = 0; i < adjRev[u].size(); i++) {</pre>
   ret += dfsRev(adjRev[u][i], color);
 return ret;
int main() {
 while (cin >> n && n) {
   int u, v; string line;
   adjOrg.assign(n, vi()); adjRev.assign(n, vi());
   for (int i = 0; i < n; i++) {
     stringstream sstr(line); sstr >> u;
     while (sstr >> v) {
      adjOrg[u].push_back(v); adjRev[v].push_back(u);
   ord.clear();
   vis.assign(n, 0);
   for (int u = 0; u < n; u++) {
    if (!vis[u]) dfsOrg(u);
   int color = 1;
   col.assign(n, 0);
   while (!ord.empty()) {
    int u = ord.back();
     if (!col[u]) {
      int size = dfsRev(u, color); // SCC Size
      if (size > 1) {
        for (int v = 0; v < n; v++) {
         if (col[v] == color); //inSame SCC;
      color++;
     ord.pop_back();
```

25.Primes 26. Extended Euclid ax + by = c, d = GCD(a, b), d | c == 0: const int64 MAX = 1e6 + 100; int x, y, d; bitset<MAX> isp;// isprime void extendedEuclid(int a, int b) { vector<int64> primes, pfs, pws; //pfs = prime factors, pws = prime powers if (b == 0) { x = 1; y = 0; d = a; return; } extendedEuclid(b, a%b); void genprime() { int x1 = y; int y1 = x - (a / b)*y; isp.set(); isp[0] = isp[1] = 0;x = x1; y = y1;for (int64 i = 2; i<MAX; i++) { if (isp[i]) { primes.push back(i); 27. Geometry 1 for (int64 j = i*i; j<MAX; j += i) isp[j] = 0; const double eps = 1e-8; const double PI = acos(-1.0); } struct Point bool isprime(int n) { double x, y; if (n<MAX) return isp[n];</pre> Point(double x = 0, double y = 0) : x(x), y(y) { } bool operator < (const Point& a) const</pre> for (int i = 0; i<primes.size() && primes[i] * primes[i] <= n; i++) {</pre> if (n%primes[i] == 0) return 0; if (a.x != x) return x < a.x; return y < a.y;</pre> return 1; }; //generation prime factors of a number int main() { typedef Point Vector: int64 n; genprime(); while (cin >> n) { struct Line int64 tmp = n, cnt = 0, cop = n, div = 1; Point P; // cop = euler Phi funcion Vector v; // cop = coprimes = all m (m<n && gcd(m,n)==1) double ang; // div = divisors = all m (m<n && gcd(m,n)==m)Line() {} for (int i = 0, pf = 2; pf*pf <= n; i++, $pf = primes[i]) {$ Line(Point P, Vector v) : P(P), v(v) { ang = atan2(v.y, v.x); } int pow = 0: bool operator < (const Line& L) const</pre> while (tmp%pf == 0) { tmp /= pf; pow++; return ang < L.ang;</pre> }; if (pow) { pfs.push back(pf), pws.push back(pow); Vector operator + (Vector A, Vector B) { return Vector(A.x + B.x, A.y + B.y); } cop -= cop / pf; div *= (pow + 1);Vector operator - (Point A, Point B) { return Vector(A.x - B.x, A.y - B.y); } if (tmp>1) cop -= cop / tmp, div *= (1 + 1); // Keep Attention to this Vector operator * (Vector A, double p) { return Vector(A.x*p, A.y*p); } cout << cop + div + 1 << endl;</pre> Vector operator / (Vector A, double p) { return Vector(A.x / p, A.y / p); } } int dcmp(double x) if (fabs(x) < eps) return 0; else return x < 0? -1 : 1; bool operator == (const Point& a, const Point &b)

```
else if (dcmp(Dot(v1, v3) > 0)) return Length(v3);
 return dcmp(a.x - b.x) == 0 && dcmp(a.y - b.y) == 0;
                                                                                              else return fabs(Cross(v1, v2)) / Length(v1);
double Dot(Vector A, Vector B) { return A.x*B.x + A.y*B.y; }
                                                                                            int isPointInPolygon(Point p, Point *poly, int n)
double Length(Vector A) { return sqrt(Dot(A, A)); }
                                                                                              int wn = 0;
                                                                                              for (int i = 0; i < n; i++)
double Angle(Vector A, Vector B) { return acos(Dot(A, B) / Length(A) / Length(B));
                                                                                               const Point& p1 = poly[i], p2 = poly[(i + 1) % n];
                                                                                               if (p == p1 || p == p2 || OnSegment(p, p1, p2)) return -1;
double Cross(Vector A, Vector B) { return A.x*B.y - A.y*B.x; }
                                                                                               int k = dcmp(Cross(p2 - p1, p - p1));
double Area2(Point A, Point B, Point C) { return fabs(Cross(B - A, C - A)) / 2; }
                                                                                               int d1 = dcmp(p1.y - p.y);
                                                                                               int d2 = dcmp(p2.y - p.y);
Vector Rotate(Vector A, double rad)
                                                                                               if (k > 0 & d1 <= 0 & d2 > 0) wn++;
                                                                                               if (k < 0 \&\& d2 <= 0 \&\& d1 > 0) wn--;
 return Vector(A.x*cos(rad) - A.y*sin(rad), A.x*sin(rad) + A.y*cos(rad));
                                                                                              if (wn != 0) return 1;
                                                                                              return 0;
Point GetLineIntersection(Point P, Vector v, Point Q, Vector w)
 Vector u = P - Q;
                                                                                            Vector Normal(Vector A)
 double t = Cross(w, u) / Cross(v, w);
 return P + v*t;
                                                                                              double L = Length(A);
                                                                                              return Vector(-A.y / L, A.x / L);
bool SegmentProperIntersection(Point a1, Point a2, Point b1, Point b2)
                                                                                            double Dist2(Point p1, Point p2)
 double c1 = Cross(a2 - a1, b1 - a1), c2 = Cross(a2 - a1, b2 - a1);
 double c3 = Cross(b2 - b1, a1 - b1), c4 = Cross(b2 - b1, a2 - b1);
                                                                                              return (p1.x - p2.x)*(p1.x - p2.x) + (p1.y - p2.y)*(p1.y - p2.y);
 return dcmp(c1) * dcmp(c2) < 0 && dcmp(c3) * dcmp(c4) < 0;
                                                                                            double RotatingCalipers(Point *P, int n)
bool OnSegment(Point p, Point a1, Point a2)
                                                                                              if (n == 1) return 0;
 return dcmp(Cross(a1 - p, a2 - p)) == 0 && dcmp(Dot(a1 - p, a2 - p)) < 0;
                                                                                              if (n == 2) return Dist2(P[0], P[1]);
                                                                                              P[n] = P[0];
                                                                                              double ans = 0;
double PolygonArea(Point* p, int n)
                                                                                              for (int u = 0, v = 1; u < n; u++)
 double area = 0;
                                                                                               for (;;)
 for (int i = 1; i < n - 1; i++)
   area += Cross(p[i] - p[0], p[i + 1] - p[0]);
                                                                                                 double diff = Cross(P[u + 1] - P[u], P[v + 1] - P[v]);
                                                                                                 if (diff <= 0)
 return area / 2;
                                                                                                   ans = max(ans, Dist2(P[u], P[v]));
double PointDistanceToLine(Point P, Point A, Point B)
                                                                                                   if (diff == 0) ans = max(ans, Dist2(P[u], P[v + 1]));
                                                                                                   break;
 Vector v1 = B - A, v2 = P - A;
 return fabs(Cross(v1, v2)) / Length(v1);
                                                                                                 v = (v + 1) \% n;
double PointDistanceToSegment(Point P, Point A, Point B)
                                                                                              return ans;
 if (A == B) return Length(P - A);
 Vector v1 = B - A, v2 = P - A, v3 = P - B;
                                                                                            bool OnLeft(Line L, Point p)
 if (dcmp(Dot(v1, v2) < 0)) return Length(v2);</pre>
```

return Cross(L.v, p - L.P) > 0; Point GetLineIntersection2(const Line &a, const Line &b) Vector u = a.P - b.P; double t = Cross(b.v, u) / Cross(a.v, b.v); return a.P + a.v*t; int HalfPlaneIntersection(Line* L, int n, Point* poly) sort(L, L + n);int first, last; Point *p = new Point[n]; Line* q = new Line[n]; q[first = last = 0] = L[0];for (int i = 1; i < n; i++) while (first < last && !OnLeft(└[i], p[last - 1])) last--;</pre> while (first < last && !OnLeft(L[i], p[first])) first++;</pre> q[++last] = L[i];if (fabs(Cross(q[last].v, q[last - 1].v)) < eps)</pre> last--: if (OnLeft(q[last], L[i].P)) q[last] = L[i]; if (first < last) p[last - 1] = GetLineIntersection2(q[last - 1], q[last]);</pre> while (first < last && !OnLeft(q[first], p[last - 1])) last--;</pre> if (last - first <= 1) return 0;</pre> p[last] = GetLineIntersection2(q[last], q[first]); for (int i = first; i <= last; i++) poly[m++] = p[i]; return m; vector<Point> CutPolygon(const vector<Point> &poly, Point A, Point B) vector<Point> newpoly; int n = poly.size(); for (int i = 0; i < n; i++)</pre> Point C = polv[i], D = polv[(i + 1) % n]; if (dcmp(Cross(B - A, C - A)) >= 0) newpoly.push back(C); if (dcmp(Cross(B - A, C - D)) != 0) Point ip = GetLineIntersection(A, B - A, C, D - C); if (OnSegment(ip, C, D)) newpoly.push back(ip); } return newpoly;

28. Geometry 2

```
#include <iostream>
#include <vector>
#include <cmath>
#include <cassert>
using namespace std;
double INF = 1e100;
double EPS = 1e-12;
#define M PI acos(-1)
struct PT {
 double x, y;
 PT() {}
 PT(double x, double y) : x(x), y(y) {}
 PT(const PT \&p) : x(p.x), y(p.y) {}
 PT operator + (const PT &p) const { return PT(x + p.x, y + p.y); }
 PT operator - (const PT &p) const { return PT(x - p.x, y - p.y); }
 bool operator<(const PT &p) const { return (x != p.x ? x<p.x : y<p.y); }</pre>
 PT operator * (double c)
                              const { return PT(x*c, y*c); }
                              const { return PT(x / c, y / c); }
 PT operator / (double c)
double dot(PT p, PT q) { return p.x*q.x + p.y*q.y; }
double dist2(PT p, PT q) { return dot(p - q, p - q); }
double cross(PT p, PT q) { return p.x*q.y - p.y*q.x; }
// rotate a point CCW or CW around the origin
PT RotateCCW90(PT p) { return PT(-p.y, p.x); }
PT RotateCW90(PT p) { return PT(p.y, -p.x); }
PT RotateCCW(PT p, double t) {
 return PT(p.x*cos(t) - p.y*sin(t), p.x*sin(t) + p.y*cos(t));
// project point c onto line through a and b
// assuming a != b
PT ProjectPointLine(PT a, PT b, PT c) {
 return a + (b - a)*dot(c - a, b - a) / dot(b - a, b - a);
// project point c onto line segment through a and b
PT ProjectPointSegment(PT a, PT b, PT c) {
 double r = dot(b - a, b - a):
 if (fabs(r) < EPS) return a;</pre>
 r = dot(c - a, b - a) / r;
 if (r < 0) return a:
 if (r > 1) return b;
 return a + (b - a)*r;
// compute distance from c to segment between a and b
double DistancePointSegment(PT a, PT b, PT c) {
 return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
```

```
// Note that it is possible to convert this into an *exact* test using
 // compute distance between point (x,y,z) and plane ax+by+cz=d
                                                                                              // integer arithmetic by taking care of the division appropriately
 double DistancePointPlane(double x, double y, double z,
                                                                                              // (making sure to deal with signs properly) and then by writing exact
                                                                                              // tests for checking point on polygon boundary
  double a, double b, double c, double d)
                                                                                              bool PointInPolygon(const vector<PT> &p, PT q) {
  return fabs(a*x + b*y + c*z - d) / sqrt(a*a + b*b + c*c);
                                                                                               bool c = 0;
                                                                                               for (int i = 0; i < p.size(); i++) {</pre>
                                                                                                 int j = (i + 1) % p.size();
 // determine if lines from a to b and c to d are parallel or collinear
                                                                                                 if ((p[i].y <= q.y && q.y < p[j].y ||</pre>
 bool LinesParallel(PT a, PT b, PT c, PT d) {
                                                                                                   p[j].y \le q.y && q.y < p[i].y) &&
  return fabs(cross(b - a, c - d)) < EPS;</pre>
                                                                                                   q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) / (p[j].y - p[i].y))
                                                                                                   c = !c:
                                                                                               }
 bool LinesCollinear(PT a, PT b, PT c, PT d) {
                                                                                               return c;
  return LinesParallel(a, b, c, d)
    && fabs(cross(a - b, a - c)) < EPS
    && fabs(cross(c - d, c - a)) < EPS;
                                                                                              // determine if point is on the boundary of a polygon
 }
                                                                                              bool PointOnPolygon(const vector<PT> &p, PT q) {
                                                                                               for (int i = 0; i < p.size(); i++)</pre>
                                                                                                 if (dist2(ProjectPointSegment(p[i], p[(i + 1) % p.size()], q), q) < EPS)</pre>
 // determine if line segment from a to b intersects with
 // line segment from c to d
                                                                                                   return true;
 bool SegmentsIntersect(PT a, PT b, PT c, PT d) {
                                                                                               return false;
  if (LinesCollinear(a, b, c, d)) {
    if (dist2(a, c) < EPS || dist2(a, d) < EPS ||</pre>
      dist2(b, c) < EPS || dist2(b, d) < EPS) return true;</pre>
                                                                                              // compute intersection of line through points a and b with
    if (dot(c - a, c - b) > 0 && dot(d - a, d - b) > 0 && dot(c - b, d - b) > 0)
                                                                                              // circle centered at c with radius r > 0
                                                                                              vector<PT> CircleLineIntersection(PT a, PT b, PT c, double r) {
      return false:
    return true;
                                                                                               vector<PT> ret:
                                                                                               b = b - a;
  if (cross(d - a, b - a) * cross(c - a, b - a) > 0) return false;
                                                                                                a = a - c:
  if (cross(a - c, d - c) * cross(b - c, d - c) > 0) return false;
                                                                                               double A = dot(b, b);
  return true;
                                                                                               double B = dot(a, b);
 }
                                                                                               double C = dot(a, a) - r*r;
                                                                                               double D = B*B - A*C;
 // compute intersection of line passing through a and b
                                                                                               if (D < -EPS) return ret;</pre>
 // with line passing through c and d, assuming that unique
                                                                                                ret.push back(c + a + b*(-B + sqrt(D + EPS)) / A);
 // intersection exists; for segment intersection, check if
                                                                                               if (D > EPS)
 // segments intersect first
                                                                                                 ret.push back(c + a + b*(-B - sqrt(D)) / A);
 PT ComputeLineIntersection(PT a, PT b, PT c, PT d) {
                                                                                               return ret;
  b = b - a; d = c - d; c = c - a;
  assert(dot(b, b) > EPS && dot(d, d) > EPS);
  return a + b*cross(c, d) / cross(b, d);
                                                                                              // compute intersection of circle centered at a with radius r
 }
                                                                                              // with circle centered at b with radius R
                                                                                              vector<PT> CircleCircleIntersection(PT a, PT b, double r, double R) {
                                                                                               vector<PT> ret;
                                                                                               double d = sqrt(dist2(a, b));
 // compute center of circle given three points
                                                                                               if (d > r + R \mid | d + min(r, R) < max(r, R)) return ret;
                                                                                                double x = (d*d - R*R + r*r) / (2 * d);
 PT ComputeCircleCenter(PT a, PT b, PT c) {
  b = (a + b) / 2;
                                                                                                double y = sqrt(r*r - x*x);
  c = (a + c) / 2:
                                                                                               PT v = (b - a) / d:
  return ComputeLineIntersection(b, b + RotateCW90(a - b), c, c + RotateCW90(a -
                                                                                                ret.push_back(a + v*x + RotateCCW90(v)*y);
                                                                                                if (y > 0)
 c));
 }
                                                                                                 ret.push_back(a + v*x - RotateCCW90(v)*y);
                                                                                                return ret;
 // determine if point is in a possibly non-convex polygon (by William
 // Randolph Franklin); returns 1 for strictly interior points, 0 for
// strictly exterior points, and 0 or 1 for the remaining points.
                                                                                              double ComputeSignedArea(const vector<PT> &p) {
```

```
double area = 0:
 for (int i = 0; i < p.size(); i++) {</pre>
   int j = (i + 1) % p.size();
   area += p[i].x*p[j].y - p[j].x*p[i].y;
 return area / 2.0;
double ComputeArea(const vector<PT> &p) {
 return fabs(ComputeSignedArea(p));
PT ComputeCentroid(const vector<PT> &p) {
 PT c(0, 0);
 double scale = 6.0 * ComputeSignedArea(p);
 for (int i = 0; i < p.size(); i++) {</pre>
   int j = (i + 1) % p.size();
   c = c + (p[i] + p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
 return c / scale;
// tests whether or not a given polygon (in CW or CCW order) is simple
bool IsSimple(const vector<PT> &p) {
 for (int i = 0; i < p.size(); i++) {</pre>
   for (int k = i + 1; k < p.size(); k++) {</pre>
    int j = (i + 1) % p.size();
     int 1 = (k + 1) % p.size();
     if (i == 1 || j == k) continue;
     if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
       return false;
  }
 }
 return true;
29. Great Circle Distance
struct PT {
 double lat, lon; PT() {}
 PT(double lat, double lon) : lat(lat), lon(lon) {}
 PT operator * (double c) const { return PT(lat*c, lon *c); }
}pts[1000 + 10];
const double eps = 1e-9;
const double PI = 3.141592653589793;
const double R = 6378.00; // radius of earth
double GCDist(PT p1, PT p2) {
 p1 = p1*(PI / 180.); p2 = p2*(PI / 180.);
 double dlon = p2.lon - p1.lon;
 double dlat = p2.lat - p1.lat;
 double a = pow((sin(dlat / 2)), 2)
    + cos(p1.lat) * cos(p2.lat) * pow(sin(dlon / 2), 2);
 double c = 2 * atan2(sqrt(a), sqrt(1 - a));
 double d = R * c;
 return d + eps;}
```

30. Convex Hull

```
#include <algorithm>
#include <vector>
using namespace std;
typedef int coord t; // coordinate type
typedef long long coord2_t; // must be big enough to hold 2*max(|coordinate|)^2
struct Point {
 coord_t x, y;
 bool operator <(const Point &p) const {</pre>
   return x < p.x | | (x == p.x && y < p.y);
};
coord2_t cross(const Point &O, const Point &A, const Point &B)
 return (A.x - 0.x) * (B.y - 0.y) - (A.y - 0.y) * (B.x - 0.x);
// Returns a list of points on the convex hull in counter-clockwise order.
// Note: the last point in the returned list is the same as the first one.
vector<Point> convex hull(vector<Point> P)
 int n = P.size(), k = 0;
 vector<Point> H(2 * n);
 // Sort points lexicographically
 sort(P.begin(), P.end());
 // Build lower hull
 for (int i = 0; i < n; i++) {
   while (k \ge 2 \&\& cross(H[k - 2], H[k - 1], P[i]) <= 0) k--;
   H[k++] = P[i];
 // Build upper hull
 for (int i = n - 2, t = k + 1; i >= 0; i --) {
   while (k \ge t \& cross(H[k - 2], H[k - 1], P[i]) \le 0) k--;
   H[k++] = P[i];
 H.resize(k);
 return H;
int main()
 vector<Point> h, p(1000000);//,p(6);
 srand(time(nullptr));
 for (int i = 0; i < 100000; ++i) {
   p[i] = Point(rand(), rand());
   if (!(i % 1000)) srand(time(nullptr));
 clock t start = clock();
 h = convex hull(p);
 for (int i = 0; i < h.size(); ++i) {</pre>
   cout << "(" << h[i].x << "," << h[i].y << ")" << endl;</pre>
}
```

31. Number Theory #include <iostream> #include <vector> #include <algorithm> using namespace std; typedef vector<int> VI; typedef pair<int, int> PII; // return a % b (positive value) int mod(int a, int b) { return ((a%b) + b) % b; } // computes gcd(a,b) int gcd(int a, int b) { int tmp: while (b) { a %= b; tmp = a; a = b; b = tmp; } return a; } // computes lcm(a,b) int lcm(int a, int b) { return a / gcd(a, b)*b; // returns d = gcd(a,b); finds x,y such that d = ax + byint extended_euclid(int a, int b, int &x, int &y) { int xx = y = 0; int yy = x = 1; while (b) { int q = a / b; int t = b; b = a%b; a = t; t = xx; xx = x - q*xx; x = t; t = yy; yy = y - q*yy; y = t;return a; } // finds all solutions to ax = b (mod n) VI modular linear equation solver(int a, int b, int n) { int x, y; VI solutions; int d = extended_euclid(a, n, x, y); **if** (!(b%d)) { x = mod(x*(b / d), n);for (int i = 0; i < d; i++) solutions.push back(mod(x + i*(n / d), n)); return solutions; // computes b such that ab = 1 (mod n), returns -1 on failure int mod inverse(int a, int n) { int x, y;

```
int d = extended euclid(a, n, x, y);
 if (d > 1) return -1;
 return mod(x, n);
// Chinese remainder theorem (special case): find z such that
// z % x = a, z % y = b. Here, z is unique modulo M = lcm(x,y).
// Return (z.M). On failure, M = -1.
PII chinese remainder theorem(int x, int a, int y, int b) {
 int s, t;
 int d = extended euclid(x, y, s, t);
 if (a%d != b%d) return make pair(0, -1);
 return make pair(mod(s*b*x + t*a*y, x*y) / d, x*y / d);
// Chinese remainder theorem: find z such that
// z % x[i] = a[i] for all i. Note that the solution is
// unique modulo M = lcm i (x[i]). Return (z,M). On
// failure, M = -1. Note that we do not require the a[i]'s
// to be relatively prime.
PII chinese remainder theorem(const VI &x, const VI &a) {
 PII ret = make_pair(a[0], x[0]);
 for (int i = 1; i < x.size(); i++) {</pre>
  ret = chinese_remainder_theorem(ret.second, ret.first, x[i], a[i]);
  if (ret.second == -1) break;
 return ret;
// computes x and y such that ax + by = c; on failure, x = y = -1
void linear diophantine(int a, int b, int c, int &x, int &y) {
 int d = gcd(a, b);
 if (c%d) {
  \times = \vee = -1;
 }
 else {
  x = c / d * mod inverse(a / d, b / d);
  y = (c - a*x) / b;
long conquer fibonacci lgN(long n) {
 long i, h, j, k, t;
 i = h = 1;
 i = k = 0;
 while (n > 0) {
  if (n % 2 == 1) {
    t = i * h:
    j = i * h + j * k + t;
    i = i * k + t;
   t = h * h;
   h = 2 * k * h + t;
   k = k * k + t;
   n = (long)n / 2;
 return j;}
```

32. Gauss - Jordan elimination for (int i = 0; i < n; i++) { a[i] = VT(A[i], A[i] + n);// Uses: b[i] = VT(B[i], B[i] + m);// (1) solving systems of linear equations (AX=B) (2) inverting matrices (AX=I) double det = GaussJordan(a, b); (3) computing determinants of square matrices // expected: 60 // Running time: O(n^3) cout << "Determinant: " << det << endl;</pre> // INPUT: a[][] = an nxn matrix // expected: -0.233333 0.166667 0.133333 0.0666667 b[][] = an nxm matrix 0.166667 0.166667 0.333333 -0.333333 // OUTPUT: X = an nxm matrix (stored in b[][]) // 0.233333 0.833333 -0.133333 -0.0666667 // A^{-1} = an nxn matrix (stored in a[][]) // 0.05 -0.75 -0.1 0.2 11 returns determinant of a[][] cout << "Inverse: " << endl;</pre> using namespace std; for (int i = 0; i < n; i++) { const double EPS = 1e-10; for (int j = 0; j < n; j++) typedef vector<int> VI; cout << a[i][j] << ' '; typedef double T: cout << endl;</pre> typedef vector<T> VT; typedef vector<VT> VVT; // expected: 1.63333 1.3 T GaussJordan(VVT &a, VVT &b) { // -0.166667 0.5 const int n = a.size(); 11 2.36667 1.7 const int m = b[0].size(); -1.85 -1.35 VI irow(n), icol(n), ipiv(n); cout << "Solution: " << endl;</pre> T det = 1;for (int i = 0; i < n; i++) { for (int i = 0; i < n; i++) { for (int j = 0; j < m; j++) int pj = -1, pk = -1; cout << b[i][j] << ' '; for (int j = 0; j < n; j++) if (!ipiv[j]) cout << endl;</pre> for (int k = 0; k < n; k++) if (!ipiv[k])</pre> }} if $(pj == -1 \mid | fabs(a[j][k]) > fabs(a[pj][pk])) { pj = j; pk = k; }$ if (fabs(a[pj][pk]) < EPS) { cerr << "Matrix is singular." << endl; exit(0); }</pre> 33.Big Integer Square ipiv[pk]++; swap(a[pj], a[pk]); import java.math.BigInteger; swap(b[pj], b[pk]); import java.util.Scanner; if (pj != pk) det *= -1; //import java.util. irow[i] = pj;public class Main { icol[i] = pk; // https://en.wikipedia.org/wiki/Integer_square_root T c = 1.0 / a[pk][pk];public static BigInteger sqrt(BigInteger n) { det *= a[pk][pk]; BigInteger cur = null; // X(k) a[pk][pk] = 1.0;BigInteger nxt = n; // X(k+1) for (int p = 0; p < n; p++) a[pk][p] *= c; while (true) { for (int p = 0; p < m; p++) b[pk][p] *= c; cur = nxt; for (int p = 0; p < n; p++) if (p != pk) { nxt = cur.add(n.divide(cur)).divide(BigInteger.valueOf(2)); c = a[p][pk]; a[p][pk] = 0;if (nxt.equals(cur)) break; for (int q = 0; q < n; q++) a[p][q] -= a[pk][q] * c;for (int q = 0; q < m; q++) b[p][q] -= b[pk][q] * c; if (cur.multiply(cur).equals(n)) return cur; else return null; for (int p = n - 1; p >= 0; p --) if (irow[p] != icol[p]) { public static void main(String[] args) { for (int k = 0; k < n; k++) swap(a[k][irow[p]], a[k][icol[p]]); Scanner sc = new Scanner(System.in); int tc = Integer.parseInt(sc.nextLine()); return det: while (tc-- > 0) { sc.nextLine(); int main() { BigInteger y = new BigInteger(sc.nextLine()); const int n = 4; if (y.equals(BigInteger.ZERO)) System.out.println(0); const int m = 2; System.out.println(sqrt(y)); double $A[n][n] = \{ \{ 1,2,3,4 \}, \{ 1,0,1,0 \}, \{ 5,3,2,4 \}, \{ 6,1,4,6 \} \}; \}$ if (tc>0) System.out.println(); double $B[n][m] = \{ \{ 1,2 \}, \{ 4,3 \}, \{ 5,6 \}, \{ 8,7 \} \};$ VVT a(n), b(n); }}

34. Convex Hull Diameter typedef pair<double, double> point; bool cw(const point &a, const point &b, const point &c) { return (b.first - a.first) * (c.second - a.second) - (b.second - a.second) * (c.first a.first) < 0;} vector<point> convexHull(vector<point> p) { int n = p.size(); **if** (n <= 1) return p: int k = 0: sort(p.begin(), p.end()); vector<point> q(n * 2); for (int i = 0; i < n; q[k++] = p[i++]) for (; $k \ge 2 \&\& !cw(q[k - 2], q[k - 1], p[i]); --k)$ for (int i = n - 2, t = k; i >= 0; q[k++] = p[i--]) for (; k > t && !cw(q[k - 2], q[k - 1], p[i]); --k) q.resize(k - 1 - (q[0] == q[1]));return q; double area(const point &a, const point &b, const point &c) { return abs((b.first - a.first) * (c.second - a.second) - (b.second - a.second) * (c.first a.first)); } double dist(const point &a, const point &b) { return hypot(a.first - b.first, a.second - b.second); double diameter(const vector<point> &p) { vector<point> h = convexHull(p); int m = h.size(); **if** (m == 1) return 0; if (m == 2)return dist(h[0], h[1]); while (area(h[m - 1], h[0], h[(k + 1) % m]) > area(h[m - 1], h[0], h[k]))++k; double res = 0; for (int i = 0, j = k; i <= k && j < m; i++) { res = max(res, dist(h[i], h[j])); while (j < m && area(h[i], h[(i + 1) % m], h[(j + 1) % m]) > area(h[i], h[(i + 1) % m], h[(i + 1) % m])res = max(res, dist(h[i], h[(j + 1) % m]));++j; return res; int main() { vector<point> points(4): points[0] = point(0, 0);points[1] = point(3, 0);points[2] = point(0, 3);points[3] = point(1, 1);double d = diameter(points); cout << d << endl;}</pre>

35.2D Fenwick

```
public class FenwickTree2D {
 public static void add(int[][] t, int r, int c, int value) {
  for (int i = r; i < t.length; i = i + 1)
    for (int j = c; j < t[0].length; j = j + 1)
      t[i][i] += value;
 // sum[(0, 0), (r, c)]
 public static int sum(int[][] t, int r, int c) {
  int res = 0;
  for (int i = r; i >= 0; i = (i & (i + 1)) - 1)
    for (int j = c; j >= 0; j = (j & (j + 1)) - 1)
      res += t[i][i];
   return res:
 // sum[(r1, c1), (r2, c2)]
 public static int sum(int[][] t, int r1, int c1, int r2, int c2) {
   return sum(t, r2, c2) - sum(t, r1 - 1, c2) - sum(t, r2, c1 - 1) + sum(t, r1 - 1,
c1 - 1);
 public static int get(int[][] t, int r, int c) {
  return sum(t, r, c, r, c);
 public static void set(int[][] t, int r, int c, int value) {
   add(t, r, c, -get(t, r, c) + value);
 // Usage example
 public static void main(String[] args) {
  int[][] t = new int[10][20];
   add(t, 0, 0, 1);
   add(t, 9, 19, -2);
   System.out.println(-1 == sum(t, 0, 0, 9, 19));
```

36.Data Structure Ideas

```
-Hash Table + Lookup
- Sparse Table
- SQRT Decomposition
- Bucketing
- Interger Arrays as matrices
- Recursive Tree Building
- Shortest Cycles
- Problem DAG
```

37.Extended Fenwick public class FenwickTreeExtended { // T[i] += value public static void add(int[] t, int i, int value) { for (; i < t.length; i |= i + 1) t[i] += value; // sum[0..i] public static int sum(int[] t, int i) { int res = 0;for $(; i \ge 0; i = (i \& (i + 1)) - 1)$ res += t[i]; return res; public static int[] createTreeFromArray(int[] a) { int[] res = new int[a.length]; for (int i = 0; i < a.length; i++) {</pre> res[i] += a[i]; int j = i | (i + 1);if (i < a.length)</pre> res[j] += res[i]; return res; // sum[a..b] public static int sum(int[] t, int a, int b) { return sum(t, b) - sum(t, a - 1); public static int get(int[] t, int i) { int res = t[i]; if (i > 0) { int lca = (i & (i + 1)) - 1; for (--i; i != lca; i = (i & (i + 1)) - 1)res -= t[i]; return res; public static void set(int[] t, int i, int value) { add(t, i, -get(t, i) + value); // interval add public static void add(int[] t, int a, int b, int value) { add(t, a, value); add(t, b + 1, -value);// point query public static int get1(int[] t, int i) {

```
return sum(t, i);
 // interval add
 public static void add(int[] t1, int[] t2, int a, int b, int value) {
  add(t1, a, value);
  add(t1, b, -value);
  add(t2, a, -value * (a - 1));
  add(t2, b, value * b);
 // interval query
 public static int sum(int[] t1, int[] t2, int i) {
  return sum(t1, i) * i + sum(t2, i);
 // Returns min(p|sum[0,p]>=sum)
 public static int lower_bound(int[] t, int sum) {
  --sum:
  int pos = -1;
  for (int blockSize = Integer.highestOneBit(t.length); blockSize != 0; blockSize
>>= 1) {
    int nextPos = pos + blockSize;
    if (nextPos < t.length && sum >= t[nextPos]) {
     sum -= t[nextPos];
     pos = nextPos;
    }
  return pos + 1;
 // Usage example
 public static void main(String[] args) {
  int[] t = new int[10];
  set(t, 0, 1);
  add(t, 9, -2);
  System.out.println(-1 == sum(t, 0, 9));
  t = createTreeFromArray(new int[] {1, 2, 3, 4, 5, 6});
  for (int i = 0; i < t.length; i++)</pre>
    System.out.print(get(t, i) + " ");
  System.out.println();
  t = createTreeFromArray(new int[] {0, 0, 1, 0, 0, 1, 0, 0});
  System.out.println(5 == lower_bound(t, 2));
  int[] t1 = new int[10];
  int[] t2 = new int[10];
  add(t1, t2, 0, 9, 1);
  add(t1, t2, 0, 0, -2);
  System.out.println(sum(t1, t2, 9));
}
```

38.KD Tree if (divX ? a[j].x <= v : a[j].y <= v)</pre> swap(a, ++i, j); import java.util.*; return i; public class KdTreePointQuery { static void swap(Point[] a, int i, int j) { public static class Point { Point t = a[i]; int x, y; a[i] = a[j];a[j] = t;public Point(int x, int y) { this.x = x;this.y = y; long bestDist; int bestNode; } public int findNearestNeighbour(int x, int y) { int[] tx; bestDist = Long.MAX VALUE; int[] ty; findNearestNeighbour(0, tx.length, x, y, true); return bestNode; public KdTreePointQuery(Point[] points) { } int n = points.length; tx = new int[n]; void findNearestNeighbour(int low, int high, int x, int y, boolean divX) { ty = new int[n]; if (low >= high) build(0, n, true, points); return; int mid = (low + high) >> > 1; long dx = x - tx[mid];void build(int low, int high, boolean divX, Point[] points) { long dy = y - ty[mid]; if (low >= high) long dist = dx * dx + dy * dy; return; if (bestDist > dist) { int mid = (low + high) >> > 1; bestDist = dist; nth_element(points, low, high, mid, divX); bestNode = mid: tx[mid] = points[mid].x; long delta = divX ? dx : dy; ty[mid] = points[mid].y; long delta2 = delta * delta; build(low, mid, !divX, points); if (delta <= 0) {</pre> build(mid + 1, high, !divX, points); findNearestNeighbour(low, mid, x, y, !divX); } if (delta2 < bestDist)</pre> findNearestNeighbour(mid + 1, high, x, y, !divX); static void nth_element(Point[] a, int low, int high, int n, boolean divX) { while (true) { else { int k = randomizedPartition(a, low, high, divX); findNearestNeighbour(mid + 1, high, x, y, !divX); if (n < k)if (delta2 < bestDist)</pre> high = k;findNearestNeighbour(low, mid, x, y, !divX); else if (n > k)low = k + 1;} else return; public static void main(String[] args) { Point[] points = new Point[n]; //fill points static final Random rnd = new Random(); //build tree KdTreePointQuery kdTree = new KdTreePointQuery(points); static int randomizedPartition(Point[] a, int low, int high, boolean divX) { swap(a, low + rnd.nextInt(high - low), high - 1); int index = kdTree.findNearestNeighbour(qx, qy); int v = divX ? a[high - 1].x : a[high - 1].y; Point p = points[index]; int i = low - 1;for (int j = low; j < high; j++)</pre>

```
39.Lazy Segment Tree
int64 t arr[100006];
int64_t t[262200];
int64_t lazy[262200];
void build(int64_t node, int64_t a, int64_t b)
  if (a>b) return;
  if (a == b)
    t[node] = arr[a];
   return:
  build(node * 2, a, (a + b) / 2);
  build(node * 2 + 1, (a + b) / 2 + 1, b);
  t[node] = t[node * 2] + t[node * 2 + 1];
int64_t query(int64_t node, int64_t a, int64_t b, int64_t i, int64_t j)
  if (a>b || a>j || b<i) return 0;
  if (lazy[node] != 0)
    t[node] += lazy[node] * (b - a + 1);
    if (a != b)
     lazy[node * 2] += lazy[node];
     lazy[node * 2 + 1] += lazy[node];
   lazy[node] = 0;
  if (a >= i && b <= j) return t[node];</pre>
  int64_t q1 = query(node * 2, a, (a + b) / 2, i, j);
 int64_t q2 = query(node * 2 + 1, (a + b) / 2 + 1, b, i, j);
  return q1 + q2;
void update(int64 t node, int64 t a, int64 t b, int64 t i, int64 t j, int64 t inc)
  if (a>b) return;
  if (lazy[node] != 0)
    t[node] += lazy[node] * (b - a + 1);
   if (a != b)
     lazy[node * 2] += lazy[node];
     lazy[node * 2 + 1] += lazy[node];
   lazy[node] = 0;
  if (a>b || a>j || b<i) return;</pre>
  if (a >= i && b <= j)
    t[node] += inc*(b - a + 1);
   if (a != b)
     lazy[node * 2] += inc;
     lazy[node * 2 + 1] += inc;
```

```
return;
  update(node * 2, a, (a + b) / 2, i, j, inc);
 update(node * 2 + 1, (a + b) / 2 + 1, b, i, j, inc);
 t[node] = t[node * 2] + t[node * 2 + 1];
int main(int argc, char const *argv[])
 int64_t t, n, qu, q, p, a; cin >> t;
 int64 t inc;
  while (t--)
   cin >> n >> qu;
   build(1, 0, n - 1);
   for (int i = 0; i < 262200; ++i) lazy[i] = 0; //CAREFUL
   while (qu--)
     cin >> a:
     if (a == 0)
       cin >> p >> q >> inc;
       update(1, 0, n - 1, p - 1, q - 1, inc);
     else
       cin >> p >> q;
       cout << query(1, 0, n - 1, p - 1, q - 1) << endl;
 return 0;
40.LCA Tree Dist
const int max_nodes, log_max_nodes;
int num_nodes, log_num_nodes, root;
vector<int> children[max nodes]; // children[i] contains the children of node i
int A[max nodes][log max nodes + 1]; // A[i][j] is the 2^j-th ancestor of node i, or -1 if
that ancestor does not exist
int L[max nodes];
                     // L[i] is the distance between node i and the root
             // floor of the binary logarithm of n
int lb(unsigned int n)
 if (n == 0)
   return -1;
  int p = 0;
 if (n >= 1 << 16) { n >>= 16; p += 16; }
 if (n >= 1 << 8) \{ n >>= 8; p += 8; \}
  if (n >= 1 << 4) { n >>= 4; p += 4; }
  if (n >= 1 << 2) \{ n >>= 2; p += 2; \}
 if (n >= 1 << 1) { p += 1; }
  return p;
void DFS(int i, int 1)
 L[i] = 1;
 for (int j = 0; j < children[i].size(); j++)</pre>
```

```
DFS(children[i][j], l + 1);
int LCA(int p, int q)
  // ensure node p is at least as deep as node q
 if (L[p] < L[q])</pre>
   swap(p, q);
  // "binary search" for the ancestor of node p situated on the same level as q
  for (int i = log num nodes; i >= 0; i--)
   if (L[p] - (1 \leftrightarrow i) >= L[q])
     p = A[p][i];
  if (p == q)
   return p;
  // "binary search" for the LCA
  for (int i = log_num_nodes; i >= 0; i--)
   if (A[p][i] != -1 && A[p][i] != A[q][i])
     p = A[p][i];
     q = A[q][i];
  return A[p][0];
int main(int argc, char* argv[])
  // read num nodes, the total number of nodes
  log num nodes = lb(num nodes);
  for (int i = 0; i < num nodes; i++)
    int p;
   // read p, the parent of node i or -1 if node i is the root
   A[i][0] = p;
    if (p != -1)
     children[p].push back(i);
    else
     root = i;
  // precompute A using dynamic programming
  for (int j = 1; j <= log_num_nodes; j++)</pre>
   for (int i = 0; i < num nodes; i++)</pre>
     if (A[i][j - 1] != -1)
       A[i][j] = A[A[i][j - 1]][j - 1];
     else
       A[i][j] = -1;
  // precompute L
  DFS(root, 0);
  return 0;
```

```
41.SegmentTree2D
import java.util.*;
public class SegmentTree2D {
 public static int max(int[][] t, int x1, int y1, int x2, int y2) {
   int n = t.length >> 1;
   x1 += n:
   x2 += n;
   int m = t[0].length >> 1;
   v1 += m;
   v2 += m:
   int res = Integer.MIN VALUE;
   for (int 1x = x1, rx = x2; 1x <= rx; 1x = (1x + 1) >> 1, rx = (rx - 1) >> 1)
    for (int 1y = y1, ry = y2; 1y <= ry; 1y = (1y + 1) >> 1, ry = (ry - 1) >> 1) {
      if ((lx \& 1) != 0 \&\& (ly \& 1) != 0) res = Math.max(res, t[lx][ly]);
      if ((lx \& 1) != 0 \&\& (ry \& 1) == 0) res = Math.max(res, t[lx][ry]);
      if ((rx \& 1) == 0 \& \& (ly \& 1) != 0) res = Math.max(res, t[rx][ly]);
      if ((rx \& 1) == 0 \& (ry \& 1) == 0) res = Math.max(res, t[rx][ry]);
   return res;
 public static void add(int[][] t, int x, int y, int value) {
   x += t.length >> 1;
   y += t[0].length >> 1;
   t[x][v] += value;
   for (int tx = x; tx > 0; tx >>= 1)
    for (int ty = y; ty > 0; ty >>= 1) {
      if (tx > 1) t[tx >> 1][ty] = Math.max(t[tx][ty], t[tx ^ 1][ty]);
      if (ty > 1) t[tx][ty \Rightarrow 1] = Math.max(t[tx][ty], t[tx][ty ^1]);
 }
 public static void main(String[] args) {
   int[][] t = new int[sx * 2][sy * 2];
   add(t, x, y, v);//tree-x-y-value
   int res1 = \max(t, x1, y1, x2, y2); //t-[x1,y1]*[x2,y2]
42.Static RMO
// keep code simple.
int lookup[MAX][LOGMAX];
struct Ouery
 int L, R;
void preprocess(int arr[], int n)
 // Initialize M for the intervals with length 1
 for (int i = 0; i < n; i++)
   lookup[i][0] = i;
```

```
for (int j = 1; (1 << j) <= n; j++)
   for (int i = 0; (i + (1 << j) - 1) < n; i++)
     if (arr[lookup[i][j - 1]] < arr[lookup[i + (1 << (j - 1))][j - 1]])</pre>
      lookup[i][j] = lookup[i][j - 1];
       lookup[i][j] = lookup[i + (1 << (j - 1))][j - 1];
 }
// Returns minimum of arr[L..R]
int query(int arr[], int L, int R)
 int j = (int)\log 2(R - L + 1);
 if (arr[lookup[L][j]] <= arr[lookup[R - (int)pow(2, j) + 1][j]])</pre>
   return arr[lookup[L][j]];
 else return arr[lookup[R - (int)pow(2, j) + 1][j]];
void RMQ(int arr[], int n, Query q[], int m)
 // Fills table lookup[n][Log n]
 preprocess(arr, n);
  for (int i = 0; i<m; i++)</pre>
   // Left and right boundaries of current range
   int L = q[i].L, R = q[i].R;
   // Print sum of current query range
   cout << "Minimum of [" << L << ", '</pre>
     << R << "] is " << query(arr, L, R) << endl;
}
int main()
 int a[] = \{ 7, 2, 3, 0, 5, 10, 3, 12, 18 \};
 int n = sizeof(a) / sizeof(a[0]);
 Query q[] = \{ \{ 0, 4 \}, \{ 4, 7 \}, \{ 7, 8 \} \};
 int m = sizeof(q) / sizeof(q[0]);
  RMQ(a, n, q, m);
 return 0;
```

43.Infix to postfix

```
import java.util.Stack;
public class ShuntingYard {
 public static void main(String[] args) {
   String infix = "3 + 4 * 2 / (1 - 5) ^ 2 ^ 3";
   System.out.printf("infix: %s%n", infix);
   System.out.printf("postfix: %s%n", infixToPostfix(infix));
 static String infixToPostfix(String infix) {
   final String ops = "-+/*^";
   StringBuilder sb = new StringBuilder();
   Stack<Integer> s = new Stack<>();
   for (String token : infix.split("\\s")) {
    if (token.isEmpty())
      continue;
    char c = token.charAt(0);
    int idx = ops.indexOf(c);
    // check for operator
    if (idx != -1) {
      if (s.isEmpty())
        s.push(idx);
      else {
        while (!s.isEmpty()) {
         int prec2 = s.peek() / 2;
         int prec1 = idx / 2;
         if (prec2 > prec1 || (prec2 == prec1 && c != '^'))
           sb.append(ops.charAt(s.pop())).append(' ');
         else break:
        s.push(idx);
    else if (c == '(') {
      s.push(-2); // -2 stands for '('
    else if (c == ')') {
      // until '(' on stack, pop operators.
      while (s.peek() != -2)
        sb.append(ops.charAt(s.pop())).append(' ');
      s.pop();
    else {
      sb.append(token).append(' ');
   while (!s.isEmpty())
    sb.append(ops.charAt(s.pop())).append(' ');
   return sb.toString();
```

```
44.KMP
                                                                                          s = k - 2;
                                                                                          e = s - palLen;
public class Kmp {
                                                                                          bool b = true;
                                                                                          for (j = s; j>e; j--)
 public static int[] prefixFunction(String s) {
  int[] p = new int[s.length()];
                                                                                           d = j - e - 1;
   int k = 0;
   for (int i = 1; i < s.length(); i++) {</pre>
                                                                                           if (out[j] == d){
    while (k > 0 \&\& s.charAt(k) != s.charAt(i))
                                                                                             palLen = d;
      k = p[k - 1];
                                                                                             b = false;
    if (s.charAt(k) == s.charAt(i))
                                                                                             break;
     ++k;
    p[i] = k;
                                                                                           out[k++] = min(d, out[j]);
   return p;
                                                                                          if (b)
 public static int kmpMatcher(String s, String pattern) {
                                                                                           pallen = 1;
  int m = pattern.length();
                                                                                           i++;
  if (m == 0)
    return 0;
   int[] p = prefixFunction(pattern);
                                                                                        out[k++] = palLen;
   for (int i = 0, k = 0; i < s.length(); i++)
                                                                                        llen = k:
    for (; ; k = p[k - 1]) {
                                                                                        s = 1Len - 2;
      if (pattern.charAt(k) == s.charAt(i)) {
                                                                                        e = s - (2 * seqLen + 1 - lLen);
       if (++k == m)
         return i + 1 - m;
                                                                                        for (i = s; i>e; i--)
       break;
                                                                                         d = i - e - 1;
      if (k == 0)
                                                                                          out[k++] = min(d, out[i]);
       break;
   return -1;
                                                                                       //Example
                                                                                       //opposes
                                                                                       //[0, 1, 0, 1, 4, 1, 0, 1, 0, 1, 0, 3, 0, 1, 0]
45.Longest Palindrome
                                                                                       //Longest palindrome has length 4
                                                                                       int main()
using namespace std;
template <class RAI1, class RAI2>
                                                                                        string s; cin >> s;
void fastLongestPalindromes(RAI1 seq, RAI1 seqEnd, RAI2 out)
                                                                                        vector<int> V(2 * s.length() + 1);
{
                                                                                        fastLongestPalindromes(s.begin(), s.end(), V.begin());
 int seqLen = seqEnd - seq;
                                                                                        int best = 0;
 int i = 0, j, d, s, e, lLen, k = 0;
                                                                                        cout << "[";
 int palLen = 0;
                                                                                        for (int i = 0; i<V.size(); i++)</pre>
 while (i<seqLen)</pre>
                                                                                         if (i>0) cout << ", ";</pre>
   if (i>palLen && seq[i - palLen - 1] == seq[i])
                                                                                         cout << V[i];
                                                                                          best = max(best, V[i]);
    palLen += 2;
    i++;
                                                                                        cout << "]" << endl << "Longest palindrome has length " << best << endl;</pre>
    continue;
                                                                                        return 0;
   out[k++] = palLen;
```

46.Simple Parser const char * expressionToParse = "3*2+4*1+(4+9)*6"; char peek(){ return *expressionToParse; char get(){ return *expressionToParse++; int expression(); int number(){ int result = get() - '0'; while (peek() >= '0' && peek() <= '9'){</pre> result = 10 * result + get() - '0'; return result; int factor(){ if (peek() >= '0' && peek() <= '9')</pre> return number(); else if (peek() == '('){ get(); // '(' int result = expression(); get(); // ')' return result; else if (peek() == '-'){ get(); return -factor(); return 0; // error int term(){ int result = factor(); while (peek() == '*' || peek() == '/') **if** (get() == '*') result *= factor(); else result /= factor(); return result; } int expression(){ int result = term(); while (peek() == '+' || peek() == '-') **if** (get() == '+') result += term(); result -= term(); return result; int tmain(int argc, TCHAR* argv[]){ int result = expression(); return 0;

```
47.Suffix array
Suffix array O(n lg^2 n)
LCP table O(n)
#include <cstdio>
#include <algorithm>
#include <cstring>
using namespace std;
#define REP(i, n) for (int i = 0; i < (int)(n); ++i)
const int MAXN = 1 << 21;</pre>
char * S;
int N, gap;
int sa[MAXN], pos[MAXN], tmp[MAXN], lcp[MAXN];
bool sufCmp(int i, int j)
 if (pos[i] != pos[j])
   return pos[i] < pos[j];</pre>
 i += gap;
 j += gap;
 return (i < N && j < N) ? pos[i] < pos[j] : i > j;
void buildSA()
 N = strlen(S);
 REP(i, N) sa[i] = i, pos[i] = S[i];
 for (gap = 1;; gap *= 2)
   sort(sa, sa + N, sufCmp);
   REP(i, N - 1) tmp[i + 1] = tmp[i] + sufCmp(sa[i], sa[i + 1]);
   REP(i, N) pos[sa[i]] = tmp[i];
   if (tmp[N - 1] == N - 1) break;
void buildLCP()
 for (int i = 0, k = 0; i < N; ++i) if (pos[i] != N - 1)
   for (int j = sa[pos[i] + 1]; S[i + k] == S[j + k];)
   lcp[pos[i]] = k;
   if (k)--k;
```

48. Prefix Function std::vector<int> prefix function(const std::string& str) { std::vector<int> prefs(str.size(), 0); for (int i = 1; i < str.size(); ++i) {</pre> int pref = prefs[i - 1]; while (pref > 0 && str[i] != str[pref]) { pref = prefs[pref - 1]; if (str[i] == str[pref]) { ++pref; prefs[i] = pref; return prefs; std::vector<int> z_function(const std::string& str) { std::vector<int> zfunc(str.size(), 0); zfunc[0] = str.size(); for (int i = 1, left = 0, right = 0; i < str.size(); ++i) {</pre> if (i <= right) {</pre> zfunc[i] = std::min(right - i + 1, zfunc[i - left]); while (i + zfunc[i] < str.size() && str[zfunc[i]] == str[i + zfunc[i]]) {</pre> ++zfunc[i]; if (i + zfunc[i] - 1 > right) { left = i: right = i + zfunc[i] - 1: return zfunc; std::string from_prefix_function(const std::vector<int>& prefs) { std::string str(prefs.size(), '.'); char current symbol = 'a'; for (int i = 0; i < prefs.size(); ++i) {</pre> if (prefs[i] > 0) { str[i] = str[prefs[i] - 1]; else { str[i] = current symbol++; return str; std::vector<int> prefix to z(const std::vector<int>& prefs) { return z function(from prefix function(prefs)); std::vector<int> z_to_prefix(const std::vector<int>& z_func) { std::vector<int> prefs(z func.size(), 0); for (int i = 1; i < z func.size(); ++i) {</pre> prefs[i + z func[i] - 1] = std::max(prefs[i + z_func[i] - 1], z_func[i]); for (int i = z func.size() - 2; $i \ge 0$; --i) { prefs[i] = std::max(prefs[i + 1] - 1, prefs[i]); return prefs;

49.KMP

```
void computeLPSArray(char *pat, int M, int *lps);
void KMPSearch(char *pat, char *txt){
 int M = strlen(pat);
 int N = strlen(txt);
 // create lps[] that will hold the longest prefix suffix
 // values for pattern
 int *lps = (int *)malloc(sizeof(int)*M);
 int j = 0; // index for pat[]
       // Preprocess the pattern (calculate lps[] array)
  computeLPSArray(pat, M, lps);
 int i = 0; // index for txt[]
  while (i < N)
   if (pat[j] == txt[i]) {
     j++;
     i++;
     printf("Found pattern at index %d \n", i - j);
     j = lps[j - 1];
   // mismatch after i matches
   else if (i < N && pat[i] != txt[i]){</pre>
     // Do not match lps[0..lps[j-1]] characters,
     // they will match anyway
     if (i != 0)
      j = lps[j - 1];
     else
       i = i + 1;
  free(lps); // to avoid memory leak
void computeLPSArray(char *pat, int M, int *lps){
 int len = 0:
 int i;
 lps[0] = 0;
 i = 1:
  while (i < M){
   if (pat[i] == pat[len]){
     len++;
     lps[i] = len;
     i++;
   else{
     if (len != 0){
       len = lps[len - 1];
     else{
       lps[i] = 0;
       i++:
     }}}
// Driver program to test above function
int main(){
  char *txt = "ABABDABACDABABCABAB";
  char *pat = "ABABCABAB";
  KMPSearch(pat, txt);
  return 0; }
```