

CSE 31

Computer Organization

Lecture 18 – MIPS Procedures (wrap up)

Announcements

- Labs

- Lab 6 grace period ends this week
 - » No penalty for submission during grace period
 - » Demo is REQUIRED to receive full credit
- Lab 7 due this week (with **7 days grace period** after due date)
 - » Demo is REQUIRED to receive full credit
- Lab 8 out this week
 - » Due at 11:59pm on the same day of your lab after next (with 7 days grace period after due date)
 - » You must demo your submission to your TA within 14 days from posting of lab
 - » Demo is REQUIRED to receive full credit

- Reading assignments

- Reading 05 (zyBooks 1.6 – 1.7, 6.1 – 6.3) due **tonight**, 03-APR and Reading 06 (zyBooks 6.4 – 6.7) due 10-APR
 - » Complete **Participation Activities** in each section to receive grade
 - » IMPORTANT: Make sure to submit score to CatCourses by using the link provided on CatCourses

Announcements

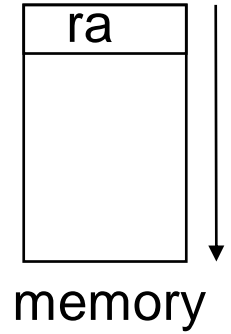
- Homework assignment
 - Homework 04 (zyBooks 4.1 – 4.9) due **tonight**, 03-APR and Homework 05 (zyBooks 1.6 – 1.7) due 10-APR
 - » Complete **Challenge Activities** in each section to receive grade
 - » IMPORTANT: Make sure to submit score to CatCourses by using the link provided on CatCourses
- Project 02
 - Due 05-MAY
 - Can work in teams of 2 students
 - » Each team member must identify teammate in “Comments...” text-box at the submission page
 - » If working in teams, each student must submit code (can be the same as teammate) and demo individually
 - » Grade can vary among teammates depending on demo
 - Demo required for project grade
 - » No partial credit for submission without demo
 - **No grace period**
 - » **Must complete submission and demo by due date.**

Basic Structure of a Function (review)

Prologue

```
entry_label:  
addi $sp, $sp, -framesize  
sw $ra, framesize-4($sp)  # save $ra  
save other regs if need be
```

Body ... (call other functions...)



Epilogue

```
restore other regs if need be  
lw $ra, framesize-4($sp)  # restore $ra  
addi $sp, $sp, framesize  
jr $ra
```

Rules for Procedures (review)

- Called with a `jal` instruction
 - returns with a `jr $ra`
- Accepts up to 4 arguments in `$a0`, `$a1`, `$a2` and `$a3`
- Return value is always in `$v0` (and if necessary, in `$v1`)
- Must follow `register conventions`
 - Covered later

Register Conventions (1/4)

- CalleR: the calling function (where you call a function)
- CalleE: the function being called
- When callee returns from executing, the caller needs to know which registers may have changed and which are guaranteed to be unchanged.
- **Register Conventions**: A set of generally accepted rules as to which registers will be unchanged after a procedure call (**jal**) and which may be changed.

Register Conventions (2/4) – saved

- **\$0**: No Change. Always 0.
- **\$s0-\$s7**: Restore if you change. Very important, that's why they're called saved registers. If the callee changes these in any way, it must restore the original values **before returning**.
- **\$sp**: Restore if you change. The stack pointer must point to the same place before and after the `jal` call, or else the caller won't be able to restore values from the stack.
- HINT -- All saved registers start with **S**!

It's callee's job to restore!

Register Conventions (3/4) – volatile

- **\$ra: Can Change.**
 - The `jal` call itself will change this register. Caller needs to save on stack before next call (nested call).
- **\$v0-\$v1: Can Change.**
 - These will contain the new returned values.
- **\$a0-\$a3: Can change.**
 - These are volatile argument registers. Caller needs to save if they are needed after the call.
- **\$t0-\$t9: Can change.**
 - That's why they're called temporary: any procedure may change them at any time. Caller needs to save if they'll need them afterwards.

It's caller's job to backup!

Register Conventions (4/4)

- What do these conventions mean?
 - If function **R** calls function **E**, then function **R** must save any **V (volatile) registers** that it may be using onto the stack **before making** a **jal** call.
 - Function **E** must save any **S (saved) registers** it intends to use before garbling up their values. It must restore any modified **S** registers **before returning** back to **R**
- Remember: **caller/callee** need to save only **volatile/saved** registers **they are using**, not all registers.

Register Conventions Summary

- If function **R** calls function **E**, then function **R** must save any **V (volatile) registers** that it may be using onto the stack **before making** a **jal** call.
 - Volatile registers: `$ra`, `$v0–$v1`, `$a0–$a3`, `$t0–$t9`
- Function **E** must save any **S (saved) registers** it intends to use before garbling up their values. It must restore any modified **S** registers **before returning** back to **R**
 - Saved registers: `$0`, `$s0–$s7`, `$sp`
- Remember: **caller/callee** need to save only **volatile/saved** registers **they are using**, not all registers.

Example: Fibonacci Numbers 1/7

- The **Fibonacci** numbers are defined as follows:

$$F(n) = F(n - 1) + F(n - 2),$$

$F(0)$ and $F(1)$ are defined to be 1

- In C we have:

```
int fib(int n) {  
    if(n == 0)  
        return 1;  
    if(n == 1)  
        return 1;  
    return (fib(n - 1) + fib(n - 2));  
}
```

Example: Fibonacci Numbers 2/7

- Now, let's translate this to MIPS!
- You will need space for **three words** on the stack
- The function will use one $\$s$ register, $\$s0$
- Write the Prologue:

`fib:`

```
addi $sp, $sp, -12  # Space for 3 words
sw  $ra, 8($sp)      # Save return address
sw  $s0, 4($sp)       # Save s0
```

Example: Fibonacci Numbers 3/7

- Now write the Epilogue:

fin:

```
lw $s0, 4($sp)      # Restore $s0
lw $ra, 8($sp)       # Restore return address
addi $sp, $sp, 12    # Pop the stack frame
jr $ra               # Return to caller
```

Example: Fibonacci Numbers 4/7

```
int fib(int n) {  
    if(n == 0) /* Base case 0*/  
        return 1;  
    if(n == 1) /* Base case 1 */  
        return 1;  
    return (fib(n - 1) + fib(n - 2));  
}
```

```
addi $v0, $zero, 1      # $v0 = 1  
beq  $a0, $zero, fin    # Base case 0  
addi $t0, $zero, 1      # $t0 = 1  
beq  $a0, $t0, fin      # Base case 1  
Continued on next slide. . .
```

Example: Fibonacci Numbers 5/7

```
int fib(int n) {  
    if(n == 0) /* Base case 0*/  
        return 1;  
    if(n == 1) /* Base case 1 */  
        return 1;  
    return (fib(n - 1) + fib(n - 2));  
}
```

Write fib(n-1):

addi \$a0, \$a0, -1	# \$a0 = n - 1
sw \$a0, 0(\$sp)	# Need \$a0 after jal
jal fib	# fib(n - 1)

Example: Fibonacci Numbers 6/7

```
int fib(int n) {  
    if(n == 0) /* Base case 0*/  
        return 1;  
    if(n == 1) /* Base case 1 */  
        return 1;  
    return (fib(n - 1) + fib(n - 2));  
}
```

Write fib(n-2) and + :

```
lw $a0, 0($sp)          # restore $a0 (= n - 1)  
addi $a0, $a0, -1       # $a0 = n - 2  
add $s0, $v0, $zero     # Place fib(n - 1) somewhere  
jal fib                 # fib(n - 2)  
add $v0, $v0, $s0       # $v0 = fib(n-1) + fib(n-2)
```


Example: Fibonacci Numbers 7/7

- Here's the complete code for reference:

```
fib:  addi $sp, $sp, -12
      sw $ra, 8($sp)
      sw $s0, 4($sp)
      addi $v0, $zero, 1
      beq $a0, $zero, fin
      addi $t0, $zero, 1
      beq $a0, $t0, fin
      addi $a0, $a0, -1
      sw $a0, 0($sp)
      jal fib
```

```
      lw $a0, 0($sp)
      addi $a0, $a0, -1
      add $s0, $v0, $zero
      jal fib
      add $v0, $v0, $s0
fin:  lw $s0, 4($sp)
      lw $ra, 8($sp)
      addi $sp, $sp, 12
      jr $ra
```

Quiz

```
int fact(int n) {  
    if(n == 0) return 1; else return(n*fact(n-1));  
}
```

When translating this to MIPS...

- 1) We COULD copy \$a0 to \$a1 (and then not store \$a0 or \$a1 on the stack) to store n across recursive calls.
- 2) We MUST save \$a0 on the stack since it gets changed.
- 3) We MUST save \$ra on the stack since we need to know where to return to...

	1	2	3
a)	F	F	F
b)	F	F	T
c)	F	T	F
d)	F	T	T
e)	T	F	F
f)	T	F	T
g)	T	T	F
h)	T	T	T

Quiz

```
int fact(int n){  
    if(n == 0) return 1; else return(n*fact(n-1));  
}
```

When translating this to MIPS...

- 1) We COULD copy \$a0 to \$a1 (and then not store \$a0 or \$a1 on the stack) to store n across recursive calls.
- 2) We MUST save \$a0 on the stack since it gets changed.
- 3) We MUST save \$ra on the stack since we need to know where to return to...

We can implement it using iterations

	123
a)	FFF
b)	FFT
c)	FTF
d)	FTT
e)	TFF
f)	TFT
g)	TTF
h)	TTT

Summary

- Functions called with **jal**, return with **jr \$ra**.
- The stack is your friend: Use it to save anything you need. Just leave it the way you found it!
- Instructions we know so far...
 - Arithmetic: **add, addi, sub, addu, addiu, subu**
 - Memory: **lw, sw, lb, sb, lbu**
 - Decision: **beq, bne, slt, slti, sltu, sltiu**
 - Unconditional Branches (Jumps): **j, jal, jr**
- Registers we know so far
 - All of them!
 - There are CONVENTIONS when calling procedures!