WHAT'S FAIR ABOUT A FAIR GAME?

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Problem 1 (E4.7) Let $X_1, X_2 \cdots$ be independent RVs such that

$$X_n = \begin{cases} n^2 - 1 & \text{with probability } n^{-2} \\ -1 & \text{with probability } 1 - n^{-2} \end{cases}.$$

Prove that $\mathbb{E}(X_n) = 0, \forall n$. But that if $S_n = \sum_{i=1}^n X_i$, then

$$\frac{S_n}{n} \to -1, \quad a.s..$$

Proof: By definition, it's obvious that $\mathbb{E}(X_i) = 0, \forall i$. Notice that

$$\frac{S_n}{n} + 1 = \frac{S_n + n}{n} = \frac{1}{n} \left(\sum_{i=1}^n Y_i \right).$$

In which $Y_i = 1 + X_i$. Then $\mathbb{E}(Y_i) = 1$. We only need to prove

$$\frac{1}{n}\sum_{i=1}^{n}Y_{i} = \frac{1}{n}S'_{n} \to 0$$
 a.s. .

My first attempt is using **Chebyshev's Inequality**, take $\phi = x^{\alpha}$, here α is a constant we may decide later, we have

$$P\left(\frac{1}{n}S'_n > \varepsilon\right) \le \frac{\mathbb{E}\left(\left(S'_n\right)^{\alpha}\right)}{\left(n\,\varepsilon\right)^{\alpha}}.$$

And we try to prove that $\mathbb{E}\left((S'_n)^{\alpha}\right)$ is really small, at least smaller that n^{α} . However, it's impossible to prove the result this way, whatever α we choose.

Which means this bound is too tight to use Chebyshev. Hence I try another method in the proof of Chernoff Bounds¹,

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¹For more details, refer to wiki https://en.wikipedia.org/wiki/Chernoff_bound

Probability Theory Assignment

Theorem 1 (Chernoff bound) Let X_1, \dots be independent Bernoulli random variables, then

$$P\left(\frac{S_n}{n} \ge \mathbb{E}\left[S_n\right] + \varepsilon\right) \le \exp\left(-D\left(p + \varepsilon \parallel p\right)n\right).$$

The bound is exactly what we want to have. However, the method used in the theorem could not be applied to this problem directly.

Finally I tried the technique used in the proof of weak law for triangular arrays, Let $\overline{Y_i} = Y_i \cdot 1_{Y_i \leq n}$. Hence we have

$$\sum_{i=1}^{n} P(Y_i > n) = \sum_{i=1}^{n} \frac{1}{i^2} < \frac{1}{\sqrt{n}} - \frac{1}{n} \to 0.$$

Plus

$$n^{-2}\sum_{i=1}^n \mathbb{E}\overline{Y_i}^2 < \frac{1}{n^2} \cdot \left\lceil \sqrt{n} \right\rceil \to 0.$$

Let $\overline{S}_n = \sum_{i=1}^n \overline{Y_i}$, it follows that $a_n = \mathbb{E}\left(\overline{S}_n\right) = \lfloor \sqrt{n} \rfloor$. We now prove a stronger result:

$$\frac{S_n'-a_n}{n} \to 0.$$

Clearly

$$P\left(\left|\frac{S_n - a_n}{n}\right| > \varepsilon\right) \le P\left(S'_n \ne \overline{S}_n\right) + P\left(\left|\frac{\overline{S}_n - a_n}{n}\right| > \varepsilon\right).$$

To estimate the first term, notice that

$$P\left(S_n \neq \overline{S}_n\right) \leq P\left(\bigcup_{i=1}^n \{\overline{Y}_i \neq Y_i\}\right) \leq \sum_{i=1}^n P\left(Y_i > n\right) \to 0.$$

For the second term, we use Chebyshev's inequality which is a common trick

$$P\left(\left|\frac{\overline{S}_n - a_n}{n}\right| > \varepsilon\right) \le \frac{\operatorname{var}\left(\overline{S}_n\right)}{n^2 \varepsilon^2} \le \frac{1}{n^2 \varepsilon^2} \cdot \sum_{i=1}^n \mathbb{E}\left(\overline{Y}_i\right)^2 \to 0.$$

Therefore

$$\frac{S_n'}{n} - \frac{1}{\sqrt{n}} \to 0.$$

As $\frac{1}{\sqrt{n}} \to 0$, our proof is completed.

Remark:

The construction of \overline{Y}_i might be a bit tricky. But it's also somewhat natural. When I try the first 2 methods, I found that n^2 is just too big for me to use some inequality tricks on the whole $Y_i, 1 \leq i \leq n$. which prevents me estimating the size, that's why we need to do some truncation.

Probability Theory

Assignment

And back to the title, the result of this problem reveals that even if the expectation of profit in every term of a game you play is 0, which seams fair. The average of your profit can be negative if you play it for too much terms. It tells us don't always take risks in such games. \Box