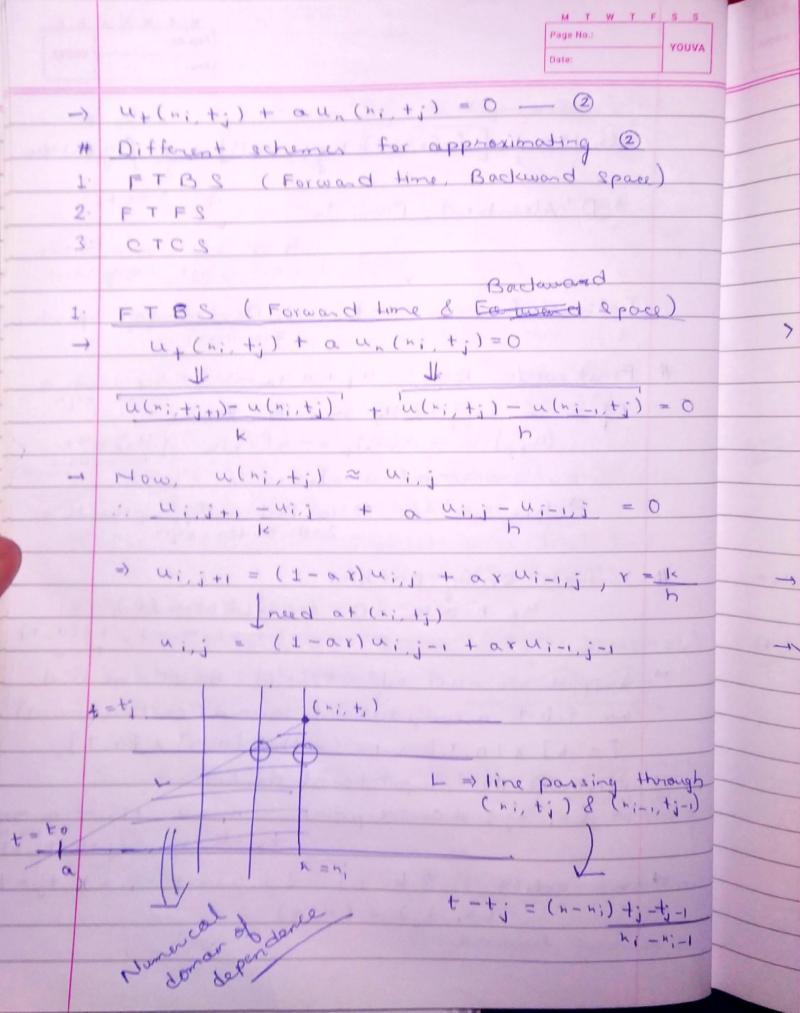
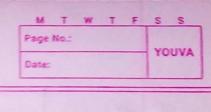


 $u(n,t) = \frac{1}{2} \left[\int (n-ct) + g(n+ct) \right] + \frac{1}{2} \int g(u) du$ > D' Alembert Formula into at whale computation domain is not required. Interval [n-ct, n+ct] => Domain of dependence
(Theoretical) First order Egn: Ut + a Ux = 0 > represents a ware egn. If a is sufficiently smooth, then $(u_{\pm \pm}) = -\alpha(u_n)_{\pm} = -\alpha(u_{\pm})_n \quad \exists u_n = u_{\pm n}$ $(u_{\pm \pm}) = -\alpha(u_n)_{\pm} = -\alpha(u_{\pm})_n \quad \exists u_n = u_{\pm n}$ = -a (-a un)n => ut = a2 unn => so indirectly represents a 2nd order egn. # Infinite String Problem: ut + a un = 0 , u(n, 0) = f(n) $n \in (-\infty, \infty)$, $t \in (0, T]$ suppose we need sol- at (x,t). Based on (x,t) we select a computational domain (con't are 1-0,001) [a,b] x (o,T] s-t- (n,t) e [a,b] x [o,T] - Discretize the computational demain: (n; +;) = cruid points n; = a, xi+= hith, har to = 0, + ++ = + + + 1 + 1 = 1 -> We relect h & k s.t.] i, j s.t n; = x, t; = t $(n_i, t_i) = (n, t)$





of dependence. For (+; +;), themstical domain of dependence

is [n; -ct; n; +ct;

& numerical domain is [a, n;]

Mow, by CFL [n; -ct; n; +ct;] & [a, n,] - Select h & k e.t. h/k/c, so that

CFL cond = is solished.

Ex: u, + u, = 0 , n e (-0, 0), te (0,1]

F.O solo at (4, 4) - Hurarial don as af dependence u(n,0)=\$ 0 a = 0 , b = 1/2

$$a = (n_i - h_i + j_i)$$

$$(n_i, + j_i) = (\frac{1}{4}, \frac{1}{4})$$

$$= 1 - h_i - \frac{1}{4} = 0$$

$$= 1 - k + k$$

