

LECTURE-3

Optimization Problems: Some Examples

Eg-1: (Farmer's Problem)

'a' acres available, To cultivate sugarcane & Rice

'r' investment amount

'm' man-days of labour

Return: Sugarcane: Rs. 1700/acre Rice: 800/acre

Cost of seeding & cultivating one acre of land:
Rs. 300 for sugarcane & Rs. 120 for rice.

Labour needed: Sugarcane: 6 man-days/acre
Rice: 3 man-days/acre

Objective: How many acres each of sugarcane and rice he should grow so that his return is maximum.

Decision variables:

x_1 = acres for plantation of sugarcane

x_2 = acres for plantation of rice.

Maximize $Q = 1700x_1 + 800x_2$

subject to $x_1 + x_2 \leq a$

$6x_1 + 3x_2 \leq m$

$300x_1 + 120x_2 \leq r$

$x_1 \geq 0, x_2 \geq 0$

LPP

Eg-2 : (Approximation)

Let $g: [a, b] \rightarrow \mathbb{R}$ be a continuous function whose values $g(x_k)$ at x_k ($k=0, 1, 2, \dots, m$) are known.

To find a polynomial of degree $n < m$ which approximates $g(\cdot)$ in a certain sense.

Let a general polynomial of degree n be of the form

$$p(x) = a_0 + a_1 x + a_2 x^2 + \dots + a_n x^n$$

Define

$$\begin{aligned} \varepsilon_k(a_0, a_1, \dots, a_n) &= g(x_k) - p(x_k) \\ &= g(x_k) - \sum_{i=0}^n a_i x_k^i, \quad k=0, 1, \dots, m. \end{aligned}$$

Objective: To choose a_0, a_1, \dots, a_n such that

$$f(a_0, a_1, \dots, a_n) = \sum_{k=0}^m \varepsilon_k^2(a_0, a_1, \dots, a_n) \text{ is minimized}$$

$$\text{Minimize } f(a_0, a_1, \dots, a_n)$$

$$\text{s.t. } \underline{a} \in \mathbb{R}^{n+1}.$$

Modifications: (eg) $a_i \geq 0$

Eg-3: (Utility Maximization)

Let a consumer consume 'n' commodities with the amount of the i^{th} commodity being $x_i \geq 0$.

The satisfaction level of consuming x_i units of the i^{th} commodity can be described by the value

$u(x) = u(x_1, \dots, x_n)$ of a utility function

$$u: \mathbb{R}_+^n \rightarrow \mathbb{R}.$$

Suppose the price for the i^{th} commodity is $p_i > 0$ and the consumer has a total budget of $a > 0$

$$\begin{array}{ll} \text{Maximize} & z = u(\underline{x}) \\ \text{subject to} & \underline{p}^T \underline{x} \leq a \\ & \underline{x} \geq \underline{0} \end{array}$$

$$\underline{p}^T \underline{x} = \sum_{i=1}^n p_i x_i$$

A concrete example for $u(x)$:

Cobb - Douglas function

$$u(x_1, x_2, \dots, x_n) = x_1^{a_1} x_2^{a_2} \dots x_n^{a_n}, \quad a_i > 0 \quad \forall i$$

$$\sum_{i=1}^n a_i = 1$$

Eg-4: (Expenditure Minimization)

Flip side of utility maximization

Given $p \in \mathbb{R}_+^n$, what is the minimum amount of income needed to give a utility-maximizing consumer a utility level of at least \bar{u} (\bar{u} fixed).

$$\text{Minimize } p^T \underline{x}$$

$$\text{subject to } u(\underline{x}) \geq \bar{u}$$

$$\underline{x} \geq 0$$

Eg-5: (Portfolio Optimization)

Given n securities (eg. stocks) and their characteristics in terms of their expected returns (μ_i) and the variance of their returns (σ_i^2), along with the covariances (σ_{ij}).

$$\mu = \begin{pmatrix} \mu_1 \\ \vdots \\ \mu_n \end{pmatrix}$$

$$C = \begin{bmatrix} \sigma_1^2 & \sigma_{12} & \dots & \sigma_{1n} \\ \sigma_{21} & \sigma_2^2 & \dots & \sigma_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ \sigma_{n1} & \sigma_{n2} & \dots & \sigma_n^2 \end{bmatrix}$$

$$w = \begin{pmatrix} w_1 \\ \vdots \\ w_n \end{pmatrix}$$

$$\sum_{i=1}^n w_i = 1$$

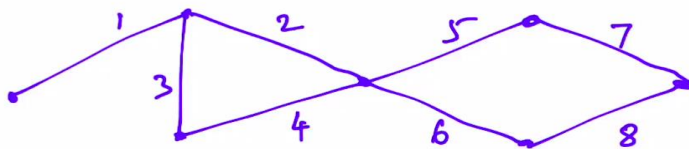
$$\begin{aligned} \mu_V &= w^T m & \sigma_V^2 &= w^T C w \\ \text{portfolio return} & & \text{portfolio risk} & \\ \text{MVP:} & \text{Minimize } \sigma_V^2 & & \\ & \text{Subject to } \sum_i w_i = 1. & & \end{aligned}$$

Variation: (i) add the constraint $\mu_V = w^T m = \mu^0$ (μ^0 fixed)

(ii) $a_i \leq w_i \leq b_i$ $0 \leq a_i \leq b_i \leq 1$

(iii) $w_i \geq 0$

Eg-b: (Matching and Covering Problems)



Maximum Cardinality Matching Problem (MCMF)
Given a graph G with n edges and m nodes, determine a matching containing the largest number of edges

$M = \{1, 2, \dots, m\} \rightarrow$ set of nodes

$N = \{1, 2, \dots, n\} \rightarrow$ set of edges

Define, for $j \in N$
 $x_j = \begin{cases} 1 & \text{if the edge } j \text{ is in the matching} \\ 0 & \text{if the edge } j \text{ is not in the matching} \end{cases}$

MEMP: Maximize $\sum_{j \in N} x_j$

Subject to $x_i + x_j \leq 1$, for i and j sharing a common endpoint and $i, j \in N$

$$x_j = 0 \text{ or } 1, j \in N.$$

Minimum Cardinality Covering Problem (MCCP)

→ determine a covering containing the fewest number of edges.

Define, for each $j \in N$

$$y_j = \begin{cases} 1 & \text{if the edge } j \text{ is in the covering} \\ 0 & \text{if the edge } j \text{ is not in the covering} \end{cases}$$

Also, let S_i be the set of edges that intersect at node $i \in M$.

MCCP: Minimize $\sum_{j \in N} y_j$

Subject to $\sum_{j \in S_i} y_j \geq 1, i \in M$

$$y_j = 0 \text{ or } 1, j \in N.$$

Integer linear programming problems

More examples can be found in almost all areas of science, engineering, economics, etc.

→ Prevalent almost everywhere.
