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- Direct methods
- Iterative methods

The following general topics will recur all throughout.

Matrix factorizations

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- Speed of algorithms

Most of the methods for solving the problems aim to express *A* as a product of 'simpler' matrices which readily reveal the solution of the problem.

▶ LU decomposition (A = LU).

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- ▶ The Singular Value decomposition (SVD) ( $A = USV^*$ ).

## Rounding: The Silent Killer

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IEEE standard allows to track small errors made when two numbers are added, subtracted, multiplied or divided on a computer.

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On my computer MATLAB produces:

$$\begin{array}{rcl} (\frac{4}{3}-1)*3-1 &=& -2.2204\times 10^{-16} \\ 5\times \frac{(1+\exp(-50))-1}{(1+\exp(-50))-1} &=& \mbox{NaN}. \\ & \frac{\log(\exp(750))}{100} &=& \mbox{Inf}. \end{array}$$

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Prevention is better than cure!

## Stability of algorithms

Analysing the errors caused by the algorithm itself requires knowing the effect of rounding errors during the execution of the algorithm. A desirable property of algorithms is *backward stability:* 

If an algorithm alg(x) is used to compute f(x), then including the effect of rounding error, alg(x) is said to be backward stable if  $alg(x) = f(x + \delta x)$  for small  $\delta x$ . Here  $\delta x$  is called the backward error.

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Thus a backward stable algorithm provides the exact answer to a slightly perturbed problem.

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Thus if the algorithm is backward stable, then the *forward error* which is the difference between its exact and computed solutions is small if the solution is not too sensitive to perturbation.

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If an algorithm is *iterative*, then it is necessary to know the number of iterations necessary to accept any approximate solution as an answer. This is decided by the quality of the convergence, whether *linear*, *quadratic or cubic*....

# Texts & References and Evaluation Policy

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A detailed evaluation policy has already been posted.