

Custom Search

Courses

Login

Write an Article

# Implementation of Diffie-Hellman Algorithm





#### **Background**

Elliptic Curve Cryptography (ECC) is an approach to public-key cryptography, based on the algebraic structure of elliptic curves over finite fields. ECC requires a smaller key as compared to non-ECC cryptography to provide equivalent security (a 256-bit ECC security have an equivalent security attained by 3072-bit RSA cryptography).

For a better understanding of Elliptic Curve Cryptography, it is very important to understand the basics of Elliptic Curve. An elliptic curve is a planar algebraic curve defined by an equation of the form

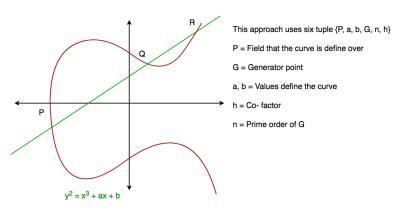
$$y^2 = x^3 + ax + b$$

where 'a' is the co-efficient of x and 'b' is the constant of the equation

## Most Accurate Email Verifier

The curve is non-singular; that is its graph has no cusps or self-intersections (when the characteristic of the co-efficient field is equal to 2 or 3).

In general, an elliptic curve looks like as shown below. Elliptic curves could intersect atmost 3 points when a straight line is drawn intersecting the curve. As we can see that elliptic curve is symmetric about the x-axis, this property plays a key role in the algorithm.



**A** 

Diffie-Hellman algorithm

The Diffie-Hellman algorithm is being used to establish a shared secret that can be used for secret communications while exchanging data over a public network using the elliptic curve to generate points and get the secret key using the parameters.

- For the sake of simplicity and practical implementation of the algorithm, we will consider only 4 variables one prime P and G (a primitive root of P) and two private values a and b.
- P and G are both publicly available numbers. Users (say Alice and Bob) pick private values a and b and they generate a key and exchange it publicly, the opposite person received the key and from that generates a secret key after which they have the same secret key to encrypt.

### **Step by Step Explanation**

Public Keys available = P, G

Private Key Selected = a

Private Key Selected = b

Key generated =  $x = G^a mod P$ Key generated =  $y = G^b mod P$ 

Exchange of generated keys takes place

Key received = y

key received = x

Generated Secret Key =  $k_a = y^a mod P$ 

Generated Secret Key =  $k_b = x^b mod P$ 

Algebraically it can be shown that  $k_a=k_b$ 

Users now have a symmetric secret key to encrypt

#### Example

```
Step 1: Alice and Bob get public numbers P = 23, G = 9
 Step 2: Alice selected a private key a = 4 and
         Bob selected a private key b = 3
 Step 3: Alice and Bob compute public values
 Alice:
           x = (9^4 \mod 23) = (6561 \mod 23) = 6
                 y = (9^3 \mod 23) = (729 \mod 23) = 16
 Step 4: Alice and Bob exchange public numbers
 Step 5: Alice receives public key y =16 and
         Bob receives public key x = 6
 Step 6: Alice and Bob compute symmetric keys
         Alice: ka = y^a \mod p = 65536 \mod 23 = 9
                 kb = x^b \mod p = 216 \mod 23 = 9
 Step 7: 9 is the shared secret.
Implementation:
/* This program calculates the Key for two persons
using the Diffie-Hellman Key exchange algorithm */
#include<stdio.h>
#include<math.h>
// Power function to return value of a ^{\wedge} b mod P
long long int power(long long int a, long long int b,
```

**A** 

**if** (b == 1)

long long int P)

```
else
        return (((long long int)pow(a, b)) % P);
//Driver program
int main()
    long long int P, G, x, a, y, b, ka, kb;
    // Both the persons will be agreed upon the
        // public keys G and P
    P = 23; // A prime number P is taken
    printf("The value of P : %lld\n", P);
    G = 9; // A primitve root for P, G is taken
    printf("The value of G : %lld\n\n", G);
    // Alice will choose the private key a
    a = 4; // a is the chosen private key
    printf("The private key a for Alice : %lld\n", a);
    x = power(G, a, P); // gets the generated key
    // Bob will choose the private key b
    b = 3; // b is the chosen private key
    printf("The private key b for Bob : %lld\n\n", b);
    y = power(G, b, P); // gets the generated key
    // Generating the secret key after the exchange
        // of keys
    ka = power(y, a, P); // Secret key for Alice
    kb = power(x, b, P); // Secret key for Bob
    printf("Secret key for the Alice is : %lld\n", ka);
    printf("Secret Key for the Bob is : %11d\n", kb);
    return 0;
}
Output
 The value of P : 23
 The value of G: 9
 The private key a for Alice : 4
 The private key b for Bob : 3
```

This article is contributed by **Souvik Nandi**. If you like GeeksforGeeks and would like to contribute, you can also write an article using contribute.GeeksforGeeks.org or mail your article to contribute@GeeksforGeeks.org. See your article appearing on the GeeksforGeeks main page and help other Geeks.

Please write comments if you find anything incorrect, or you want to share more information about the topic discussed above.

## **Recommended Posts:**

#### RSA Algorithm in Cryptography

Computer Network | Leaky bucket algorithm

Computer Network | HMAC Algorithm

Secret key for the Alice is : 9 Secret Key for the Bob is : 9

TCP Server-Client implementation in C

Probabilistic shortest path routing algorithm for optical networks

RSA Algorithm using Multiple Precision Arithmetic Library

UDP Server-Client implementation in C

Computer Network | RC4 Encryption Algorithm

Back-off Algorithm for CSMA/CD

Computer Network | Algorithm for Dynamic Time out timer Calculation