CSE3241: Operating System and System Programming

Class-14

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Necessary Conditions for Deadlock

- A **Deadlock** can arise if the following **four** conditions hold **simultaneouly**:
 - 1. **Mutual Exclusion**: At least one R must be held in a nonsharable mode. It would never happen, $P_i \longleftarrow R \longrightarrow P_{i+1}$ when R is non-sharable.
 - 2. Hold and Wait: $R_i \longrightarrow P_i \longrightarrow R_{i+1} \longrightarrow P_{i+1}$.
 - 3. No Preemption: R cannot be taken away from P unless it voluntarily realeases.
 - 4. Circular Wait: $P_i \longrightarrow R_j \longrightarrow P_{i+1} \longrightarrow R_{j+1} \longrightarrow \dots \longrightarrow P_i$

Methods of Handling Deadlock

- Deadlock can be handled by one of three ways:
 - 1. Prevent or Avoid: Take precautions.
 - Detect and Recover: At first let deadlock occurs and then take action.
 - 3. Pretend: Pretend deadlock will never happen.

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How to Pretend: Ostrich Algorithm

■ **Algorithm**: Stick your head in the sand and pretend there is no problem at all.

- It is upto application developers to handle deadlocks.
- Mathematicians found it completely unacceptable.
- Many OS engineers love to embrace this algorithm.
- Most operating systems implement it including: UNIX, Linux, Microsoft Windows.

Ostrich Bird [2]



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Story of an Unlucky Process Rumor

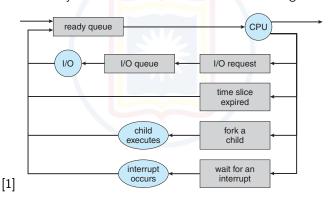
- An unlucky process:
 - was found when MIT's IBM 7094 was shut down.
 - waited from 1967 to 1973.
 - ▶ did not get access of CPU even after waiting for around 7 years.
- The culprit was nobody else but its niceness (i.e., low-priority).
- Moral of the Story: Too much nicesness is dangerous!!!
- Indefinite blocking or starvation: It is a situation, when a process is trapped in the ready queue for the indefinite time.
 - It may badly affect any low-prority process in any system using priority based CPU scheduling.

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CPU Scheduler

CPU scheduler:

- is the part of OS which selects a process from the ready queue and allocates the CPU to that process.
- is known as short-term scheduler.
- works when any of the event occurs as shown in the Figure.



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Coarse Classification of CPU Scheduling Algorithm

CPU scheduling can be:

1. Preemptive Scheduling:

- a running process can be sent back to the ready queue by taking away its allocated CPU.
- incurs an extra cost associated with access to shared data.
- may cause extra burden for the OS because of context switch.

2. Non-Preemptive Scheduling:

- onece the CPU is allocated to a process, it keeps the CPU until it realeases the CPU eitherby terminating or by switching to the waiting state.
- is known as cooperative scheduling.

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Fine Classification of CPU Scheduling Algorithm

- CPU Scheduling can be:
 - First-Come, First-Served (FCFS): Process P that requests the CPU first is allocated the CPU first.
 - 2. Shortest-Job-First (FJS): Process P that requires the least CPU access is allocated the CPU first.
 - Priority Based (PB): Process P with the highest priority is allocated the CPU first.
 - 4. Round-Robin (RR): Process P is allocated the CPU for T time and after that it is put in the last of the ready queue.
- All scheduling algorithms except FCFS, can be preemptive and non-preemptive.

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Negative Points of Scheduling Algorithm

FCFS

- average waiting time of processes is not minimal.
- ► cannot make a balance between CPU burst and I/O burst processes
- ▶ not suitable for time sharing systems, since it is non-preemptive.

SJF:

knowing the CPU usage time is advance is impossible.

PB

indefinite blocking or starvation.

RR:

- average waiting time of processes is often long.
- burden of context switching.

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References



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