

Exercise 11

11.1 Arithmetic circuits (5pt)

The BGW protocol for secure multiparty computation evaluates an arithmetic circuit among n parties. However, most computer hardware, hardware-layout compilers, and the familiar circuit model in theory operate with binary (Boolean) wires and gates.

Describe how to turn a binary circuit into an equivalent arithmetic circuit over $GF(q)$, for some prime q . To be concrete, represent FALSE and TRUE with values in $GF(q)$ and describe how to implement AND, OR, NOT, and XOR gates. Justify the operations of each gate.

11.2 Accumulator based on the strong RSA assumption (5pt)

Recall the Camenisch-Lysyanskaya (CL) RSA-based accumulator scheme. It uses a hash function $H : \{0, 1\}^* \rightarrow \mathbb{N}$, which maps arbitrary bit strings to primes. How can H be implemented in a deterministic way?

Existing primality tests are probabilistic. Assume we are given an algorithm $isPrime(n, k)$, which outputs a Boolean value indicating whether the integer n is prime; the value k is a security parameter. Let it be implemented by the Miller-Rabin test [KL21, Sec. 8.2], which has access to a random-number generator R . If n is prime, this method always returns TRUE; if n is composite, it returns TRUE with probability at most $O(4^{-k})$ and FALSE otherwise.

Suppose R is a pseudorandom generator with seed s , i.e., there is a method $R.setSeed(s)$ and after every call of this, R outputs a deterministic sequence of values that only depends on s .

- a) Describe an implementation of H with text or using pseudocode.
- b) (*Bonus: +5pt*) Implement the complete CL accumulator including H in python or your favorite programming language.

If you submit a solution to the bonus question, please upload one ZIP file containing the source code and the solution in PDF.

References

- [KL21] Jonathan Katz and Yehuda Lindell. *Introduction to Modern Cryptography: Principles and Protocols*. CRC Press, 3rd edition, 2021.