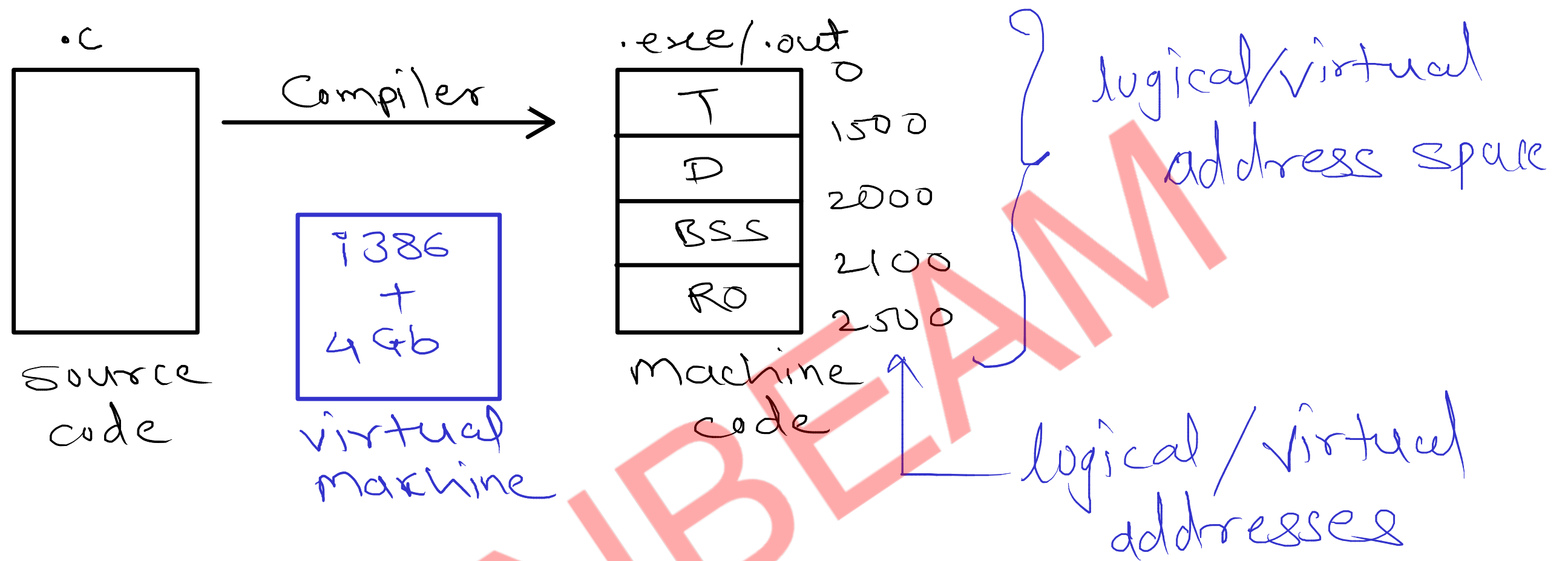


Memory Management



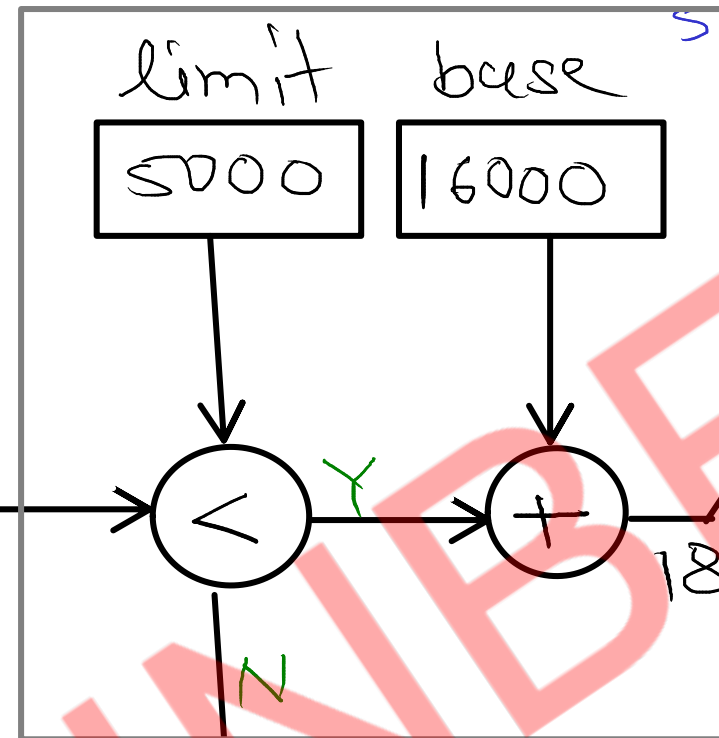
Memory Management

Process

| | |
|---|------|
| T | 0 |
| D | 1500 |
| H | 2000 |
| S | 3000 |
| | 5000 |

CPU 2500

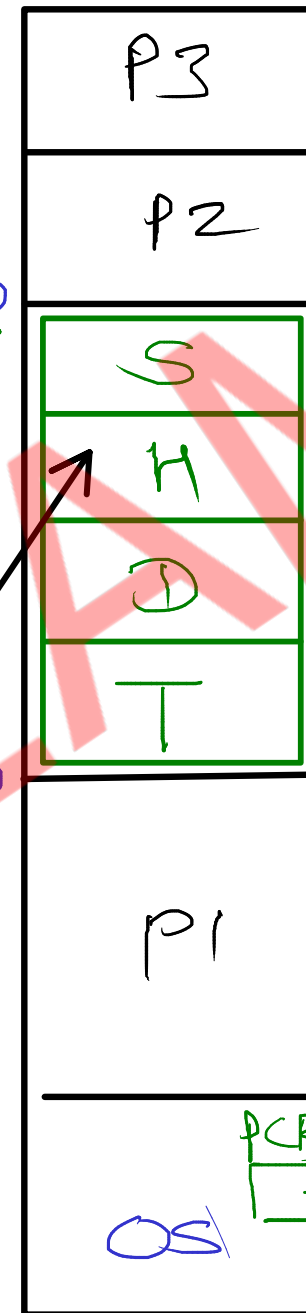
Loader



Simple MMU

Contiguous Memory Allocation

RAM



physical/actual addresses

Physical Address space

limit - 5000
base - 16000

Fixed Partition

Limitations

- 1) Max no. of processes = no. of partitions
- 2) Max size of processes = max. size of partition

| |
|-----------|
| RAM |
| P5 3kb |
| P4 3kb |
| 2kb |
| P3 2kb |
| 1kb |
| P6 4kb |
| P2 2kb |
| P1 4kb |

Internal Fragmentation

↳ if process is not utilizing complete partition which is allocated then some space is wasted

Dynamic Partition

Free slot
table

| | | | |
|---|-----------|-------|-------|
| P | first fit | limit | base |
| | Best fit | 2000 | 3000 |
| | Worst fit | 1000 | 12000 |
| | | 3000 | 7000 |

850 ✓

4000 ✗

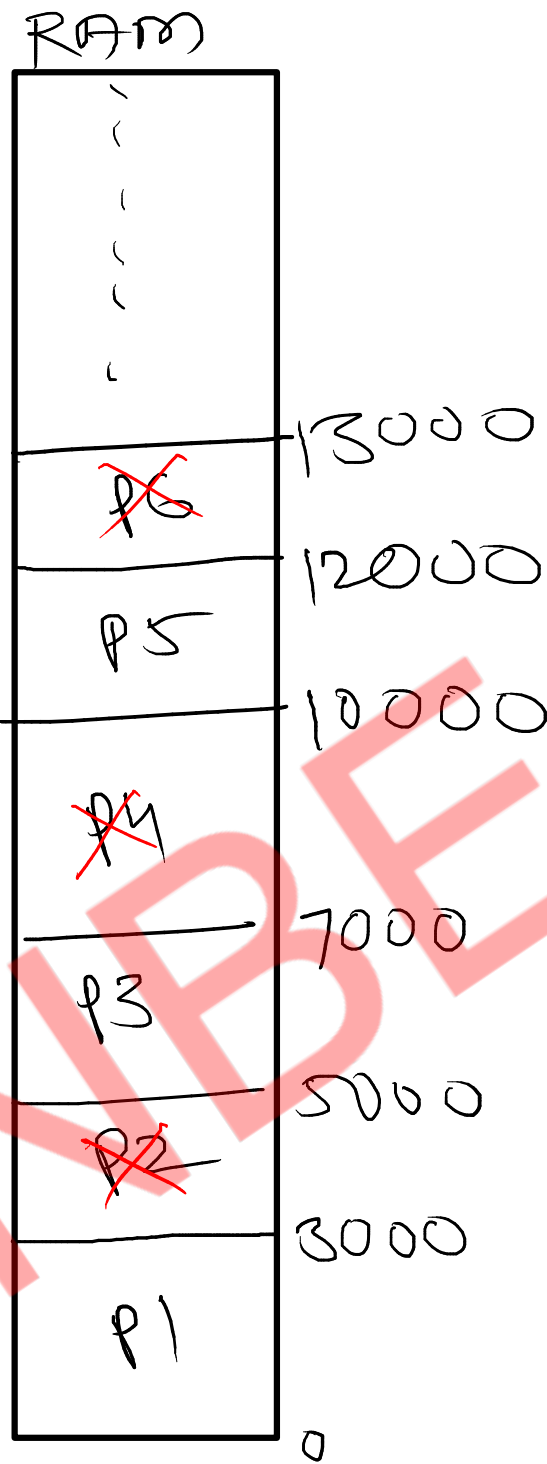
External Fragmentations

↓
due to unavailability

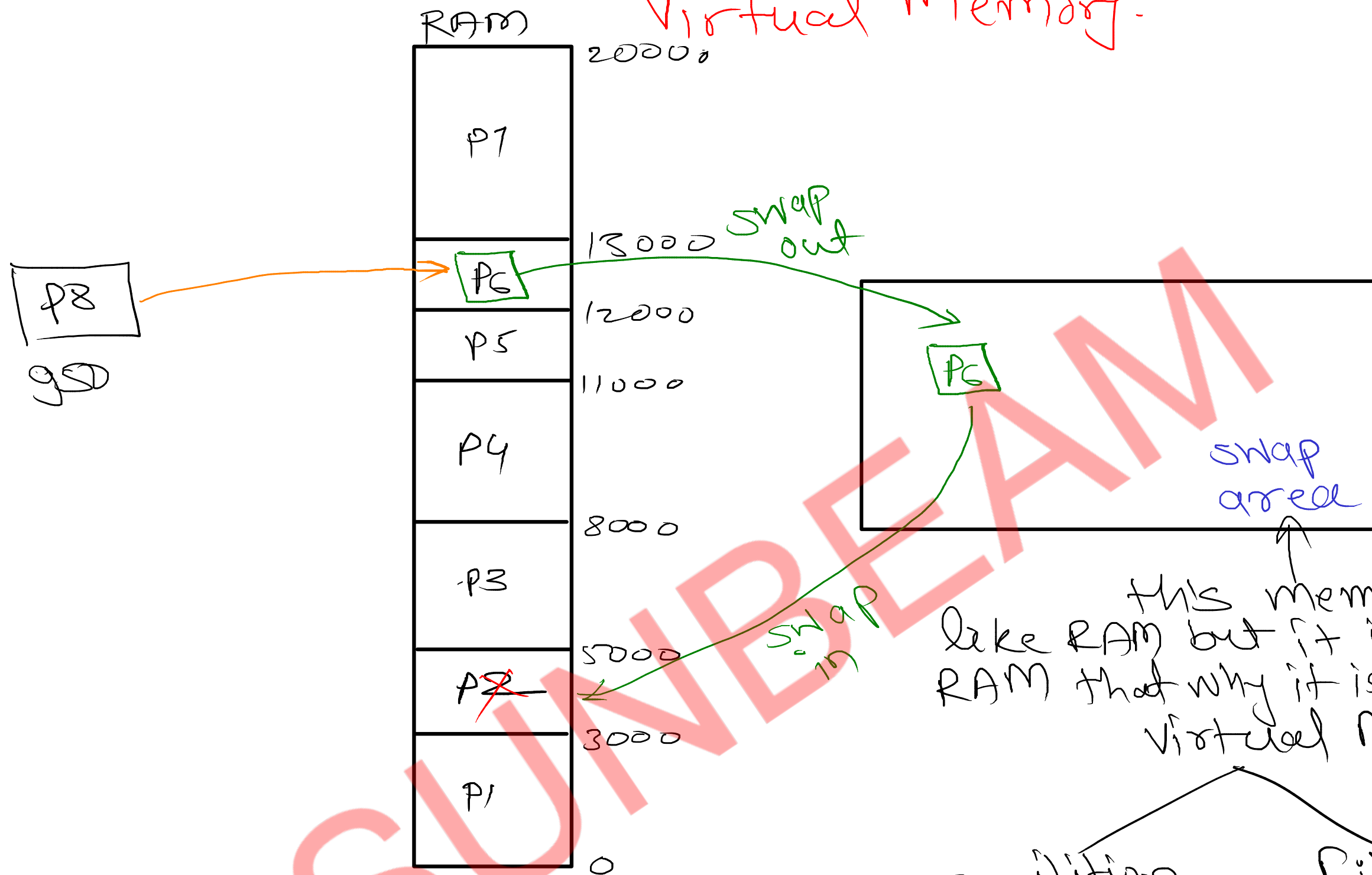
of contiguous free space,
bigger process can not
be loaded into memory

Compaction

↳ processes are moved into RAM to
create large contiguous free space



Virtual Memory.



this memory looks like RAM but it is not actual RAM that why it is known as Virtual Memory

partition
↓
linux

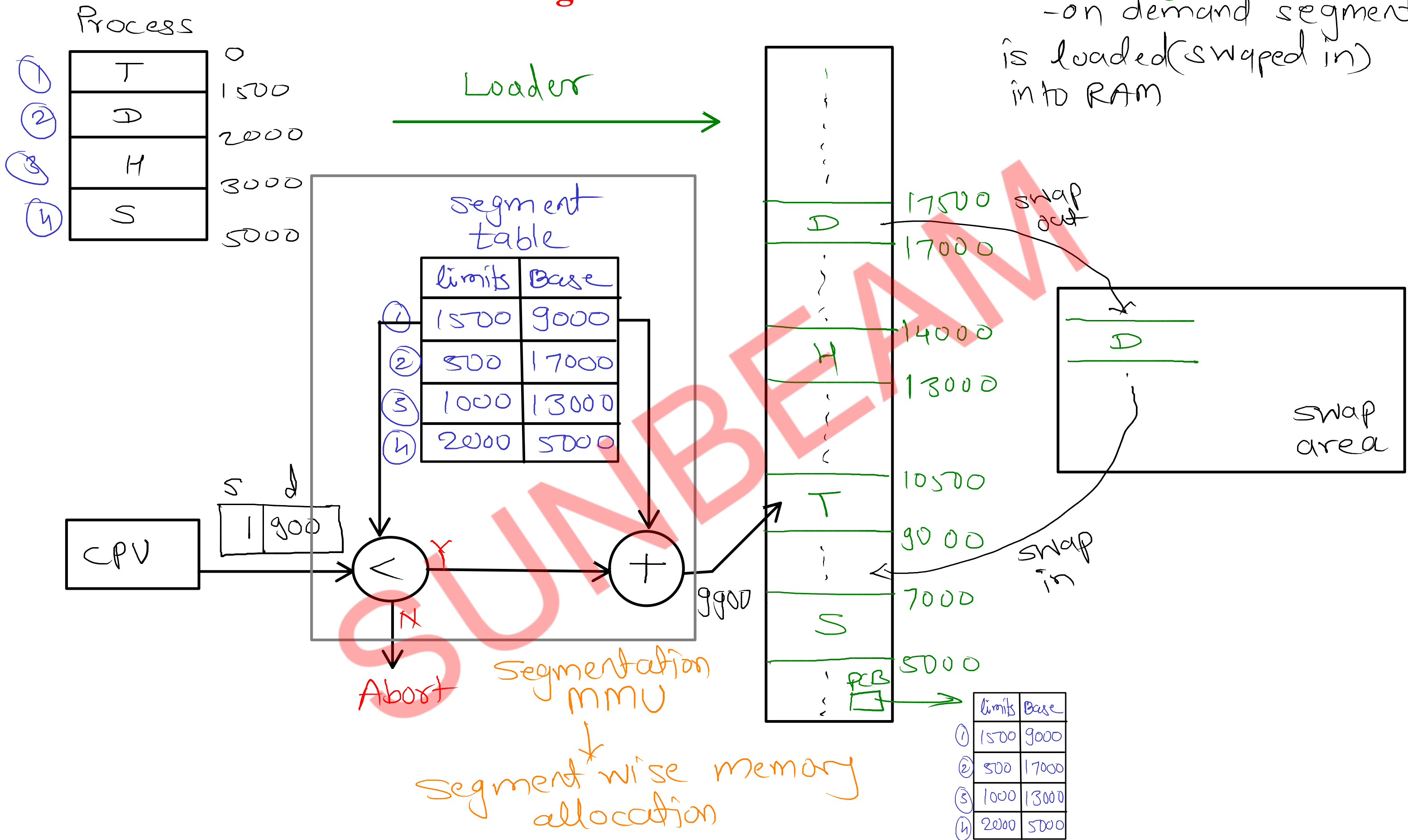
file
↓
windows

size = 2 * RAM size

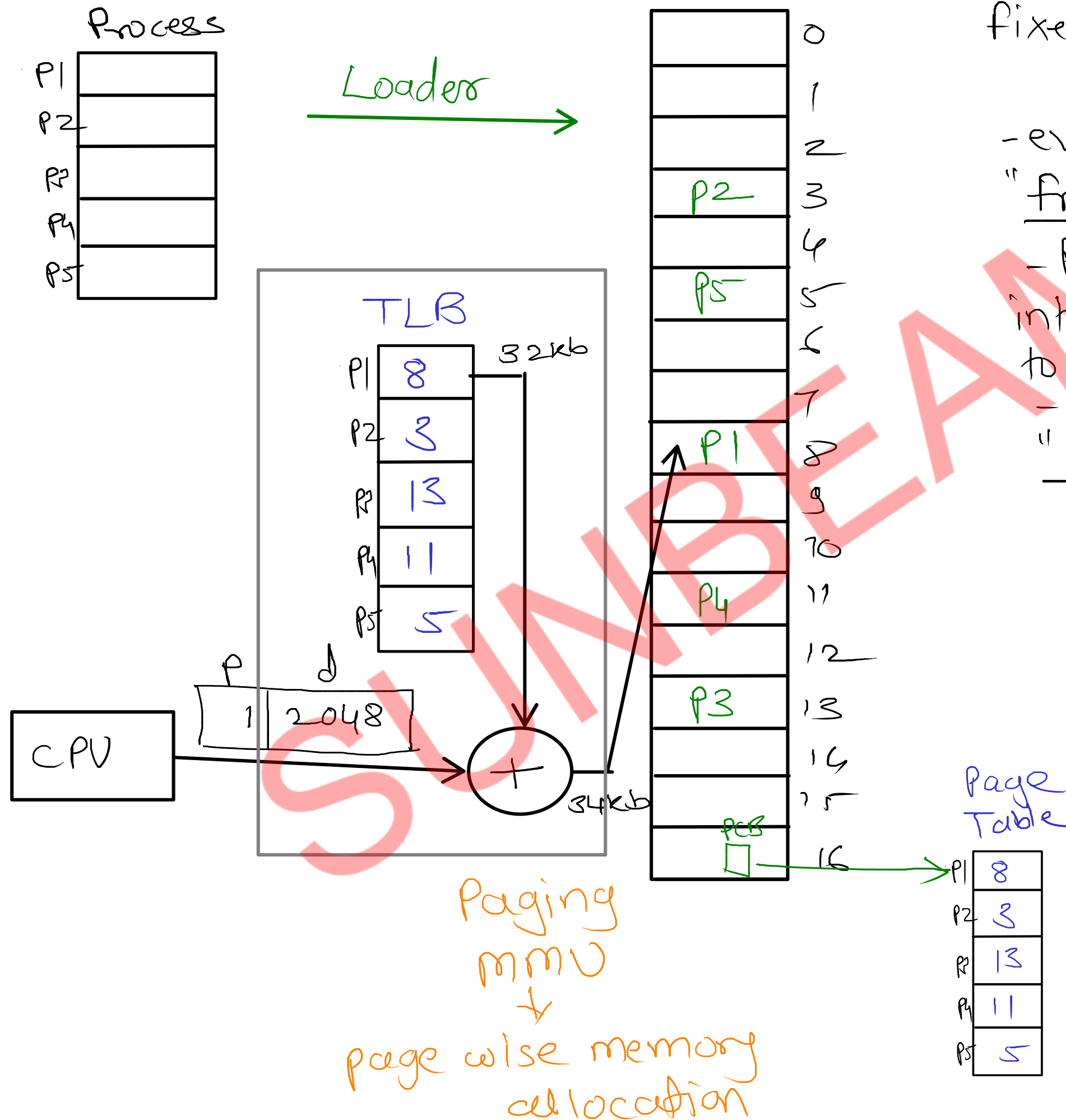
(pagefile.sys)

Segmentation MMU

Demand segmentation
- on demand segment
is loaded (swapped in)
into RAM



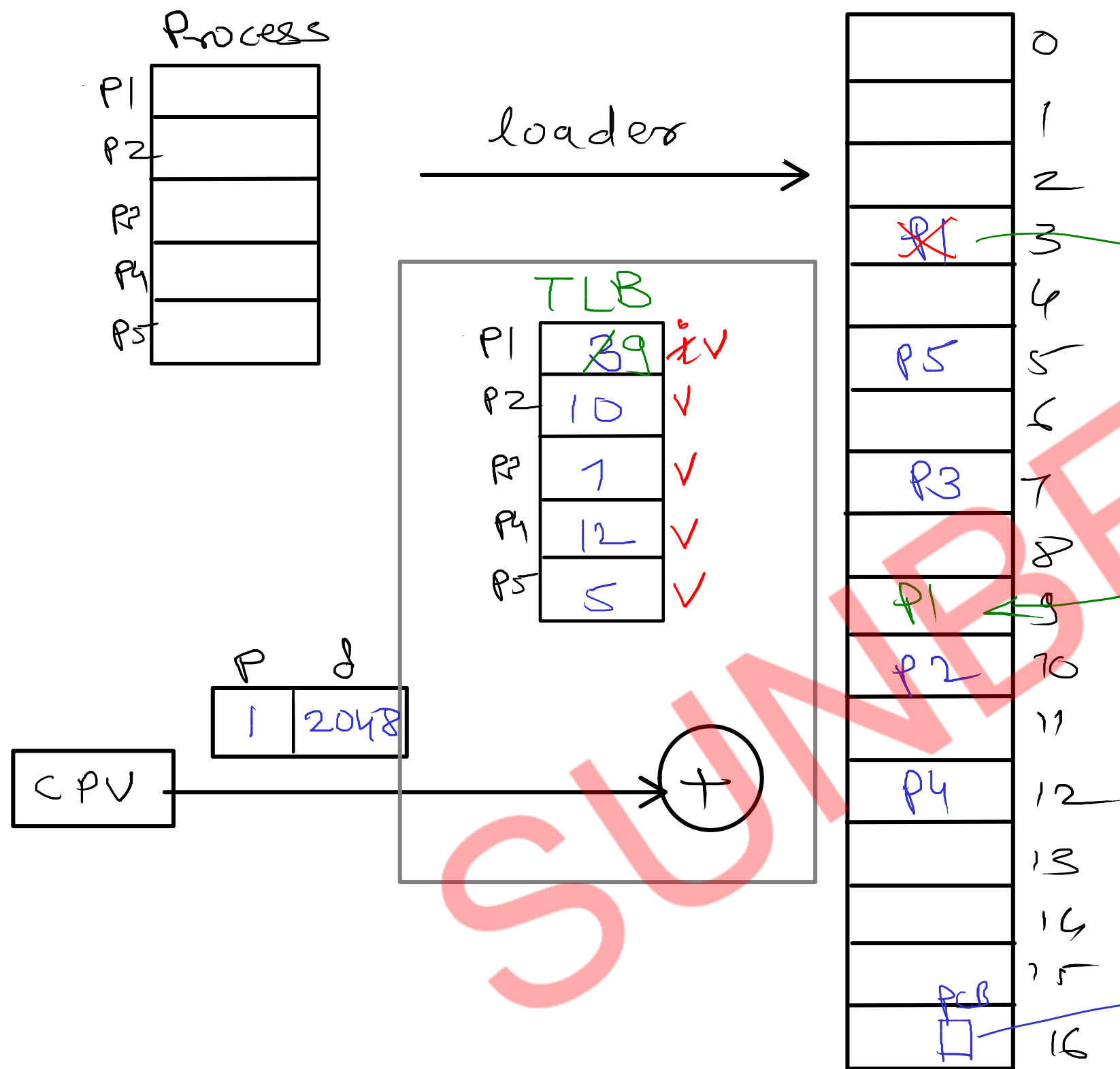
Paging MMU



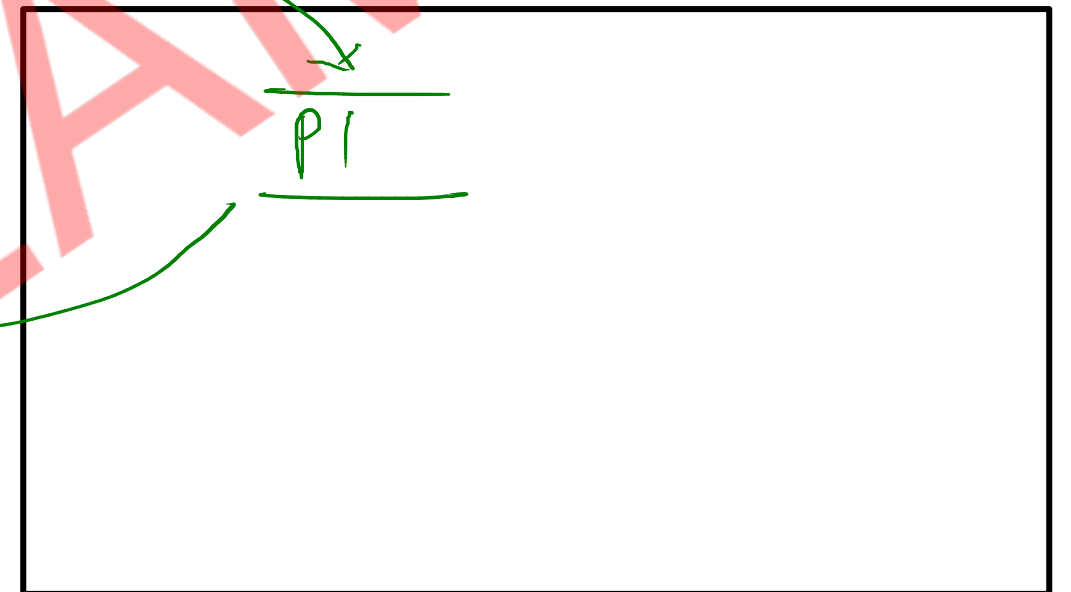
- RAM is divided into fixed & equal size partitions
size = 4096 bytes (4 Kb)
- every partition is known as "frame" / "physical page"
- Process is also divided into partitions of size equal to frame size
- every partition is known as "page" / "logical page"

Paging
MMU
↓
page wise memory
allocation

SUNBEAM



Page Fault:
 - When CPU request for an address of invalid entry of page table, fault is generated into system & known as page fault.



- whenever page fault is generated, page fault handler of OS is called.

Page Table

| | | |
|----|----|---------|
| P1 | 39 | invalid |
| P2 | 10 | valid |
| P3 | 1 | valid |
| P4 | 12 | valid |
| P5 | 5 | valid |

page_fault_handler() {

1) check requested address is valid or not?
if not valid abort the process

2) check for read/write access
if no access then abort the process

3) find free frame to swap in the page which is
on swap area

4) swap in the requested page and update the
entry of page table

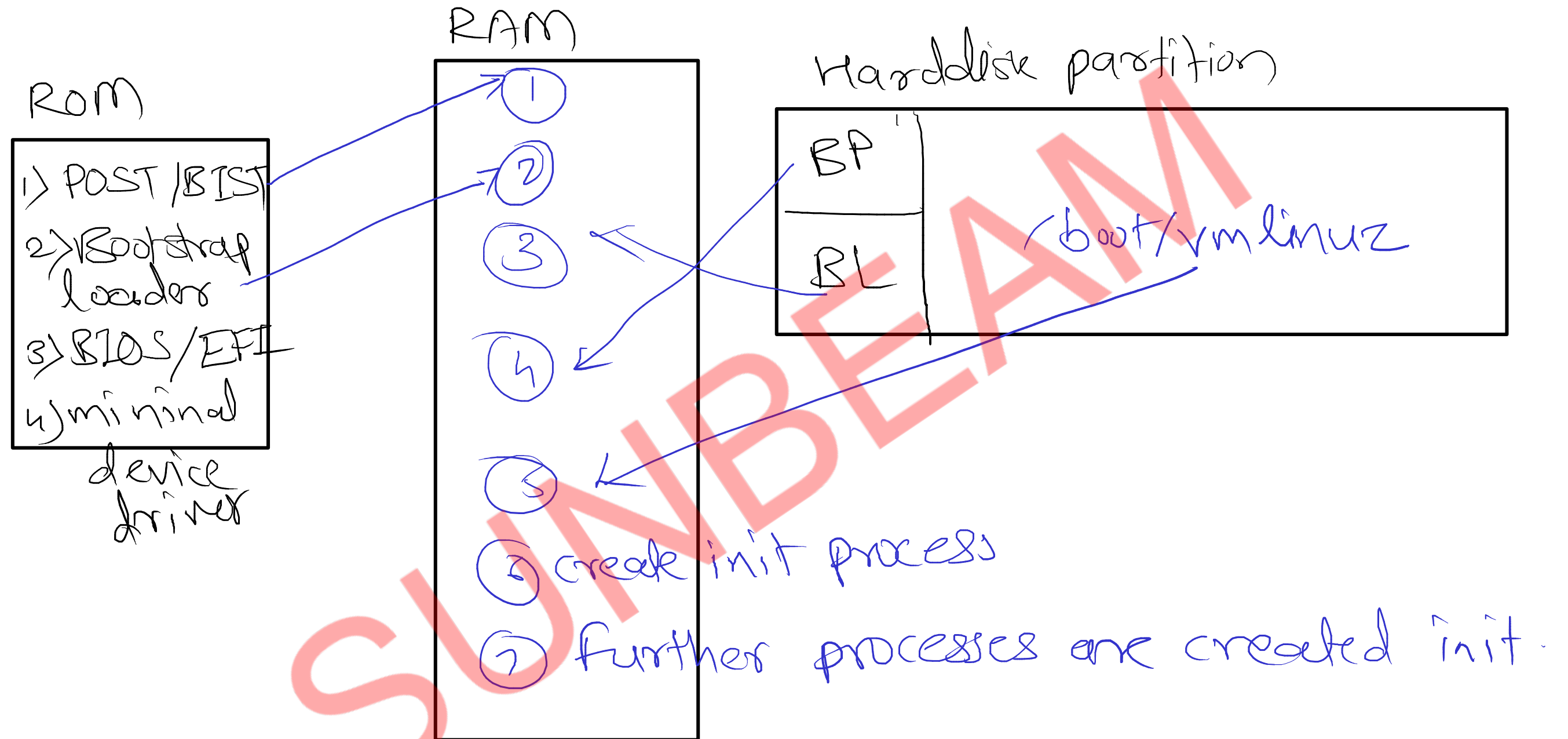
5) request the same address for which page
fault was occurred.

}

OS Booting

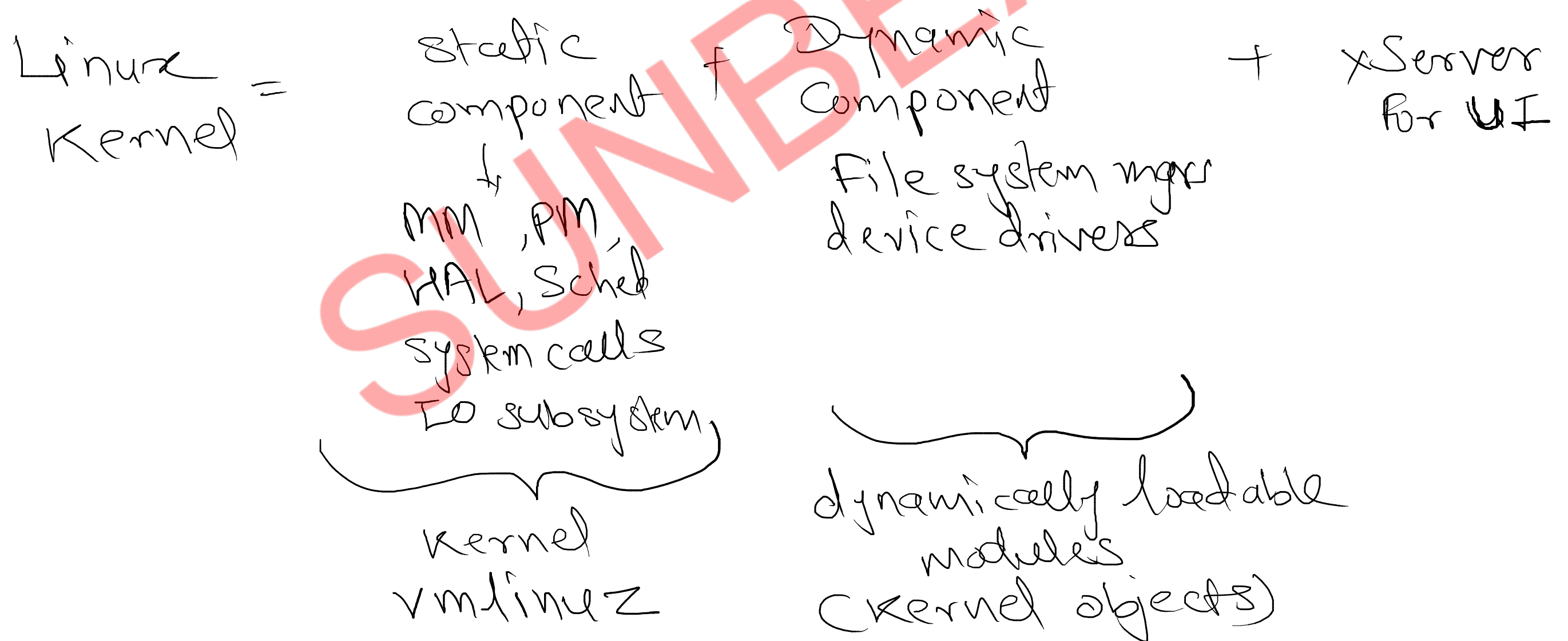
- 1) Power ON
 - 2) POST / BIST — check for all hardware
 - 3) Bootstrap loader — find bootable partition of harddisk
 - 4) load & start boot loader — show menu to the user & depending on user's choice load bootstrap program
 - 5) Bootstrap program → load kernel image into RAM
 - 6) self extraction of kernel
 - 7) start init / systemd process (first process)
 - 8) remaining processes are created by this process and also it starts all the services
- ROM {
- Boot sector {

Booting Process



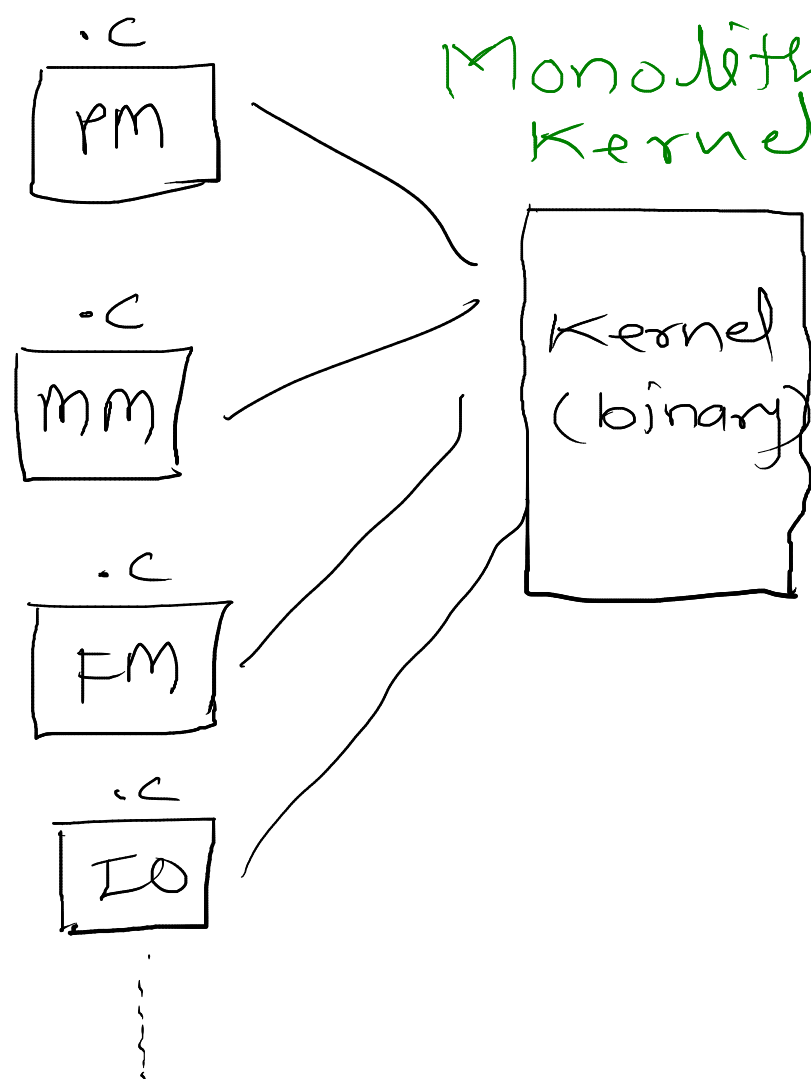
Types of kernel

- 1) Monolithic kernel - BSD UNIX
- 2) Micro kernel - Symbian, MACH
- 3) Modular kernel - Windows
- 4) Hybrid kernel - iOS - Darwin = BSD UNIX + MACH
- 5) Nano kernel - FreeRTOS

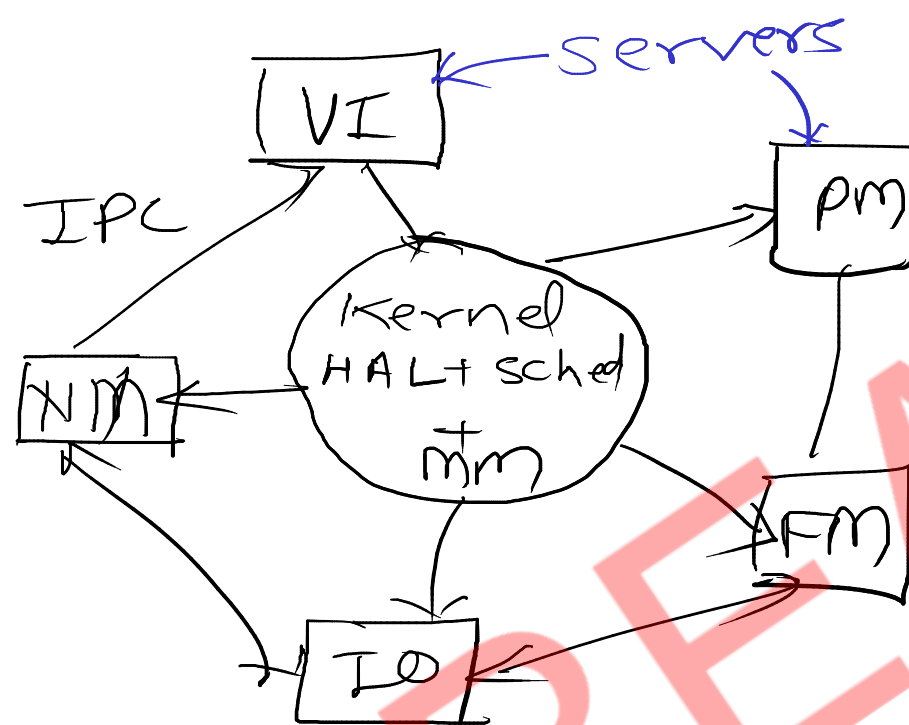


Types of kernel

Monolithic Kernel



Micro Kernel



Modular Kernel dynamically loadable modules

