

# CS 570: Analysis of Algorithms – H8

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## Exercise 15.2-1

Q- Find an optimal parenthesization of a matrix-chain product whose sequence of dimensions is [5,10,3,12,5,50,6].

# Algorithm for computing the optimal costs using Dynamic Programming

```
MATRIX-CHAIN-ORDER(p):
    n = length[p] - 1
    for i = 1 to n
        m[i, i] = 0
    for l = 2 to n
        for i = 1 to n - l + 1
            j = i + l - 1
            m[i, j] = ∞
            for k = i to j - 1
                q = m[i, k] + m[k + 1, j] + p[i-1]*p[k]*p[j]
                if q < m[i, j]
                    m[i, j] = q
                    s[i, j] = k
    return m and s
```

# Algorithm for constructing an optimal solution

```
PRINT-OPTIMAL-PARENS(s, i, j):
    if i = j
        print "A"i
    else
        print "("
        PRINT-OPTIMAL-PARENS(s, i, s[i, j])
        PRINT-OPTIMAL-PARENS(s, s[i, j] + 1, j)
        print ")"
```

# Solution for <5,10,3,12,5,50,6>

```
p = [5, 10, 3, 12, 5, 50, 6]
m, s = MATRIX-CHAIN-ORDER(p)
PRINT-OPTIMAL-PARENS(s, 1, 6)
```

	6	5	4	3	2	1
6	.	.	.	.	.	.
5	1500	.	.	.	.	.
4	1860	3000	.	.	.	.
3	1770	930	180	.	.	.
2	1950	2430	330	360	.	.
1	2010	1655	405	330	150	.

#Output

((A1A2)((A3A4)(A5A6)))

	6	5	4	3	2	1
6	.	.	.	.	.	.
5	5	.	.	.	.	.
4	4	4	.	.	.	.
3	4	4	3	.	.	.
2	2	2	2	2	.	.
1	2	4	2	2	1	.

### Exercise 15.2-2

**Q-** Give a recursive algorithm **MATRIX-CHAIN-MULTIPLY(A, s, i, j)** that actually performs the optimal matrix-chain multiplication, given the sequence of matrices  $\langle A_1, A_2, \dots, A_n \rangle$ , the  $s$  table computed **MATRIX-CHAIN-ORDER**, and the indices  $i$  and  $j$ . (The initial call would be **MATRIX-CHAIN-MULTIPLY(A, s, 1, n)**.)

```
MATRIX-CHAIN-MULTIPLY(A, s, i, j):
    if i == j:
        return A[i - 1]
    else:
        A_left = MATRIX-CHAIN-MULTIPLY(A, s, i, s[i][j])
        A_right = MATRIX-CHAIN-MULTIPLY(A, s, s[i][j] + 1, j)
        return MATRIX-MULTIPLY(A_left, A_right)
```

```
MATRIX-MULTIPLY(A, B):
    m = number of rows in matrix A
    n = number of columns in matrix A
    p = number of columns in matrix B
    Let C be a new m x p matrix
    for i = 1 to m:
        for j = 1 to p:
            C[i][j] = 0
            for k = 1 to n:
                C[i][j] = C[i][j] + A[i][k] * B[k][j]
    return C
```

```
A = [A1, A2, ..., An]
s = computed s table by MATRIX-CHAIN-ORDER
n = length of A
result = MATRIX-CHAIN-MULTIPLY(A, s, 1, n)
```

**MATRIX-CHAIN-MULTIPLY(A, s, i, j):**

- This function recursively multiplies matrices in the given sequence A between indices  $i$  and  $j$ .
- If  $i$  equals  $j$ , it means there's only one matrix left in the chain, so it returns that matrix.
- Otherwise, it splits the chain at index  $s[i][j]$ , recursively multiplies the left and right subchains, and then multiplies the resulting matrices using the helper function **MATRIX-MULTIPLY**.

**MATRIX-MULTIPLY(A, B):**

- This function performs matrix multiplication between matrices A and B.
- It initializes a new matrix C of dimensions  $m \times p$ . It iterates over each element  $C[i][j]$  of the resulting matrix and computes its value by summing up products of corresponding elements from matrices A and B.
- Finally, it returns the resulting matrix C.