# Functional Programming Epistemic Model Checking

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#### 1 Introduction

We implemented the epistemic model checking for Kripke models and succint Kripke models (mental models). We also implemented translating functions between the two. Note that the translation from Kripke to mental models introduces nominal propositions whose sole purpose is to identify a world each. Given that we represent agents, worlds and propositions of Kripke models as positive integers, we let the nominal propositions be the negations of the worlds they represent. As such translating back and forth from a Kripke model nets us one whose universe of propositions is a strict superset of the original. Still if we restrict it to the original universe, it is isomorphic to the original Kripke model.

### 2 Code

satisfies \_ Top

= True

# 2.1 Propositional Logic

Show instance for propositional formulas (includes the knowledge operator).

```
type Proposition = Integer
data Form = P Proposition | Neg Form | Conj Form Form | Disj Form Form | Top | Bottom | Knows Agent Form deriving (Eq.Ord)
instance Show Form where
  showsPrec _ (P name) = showString $ show name
  showsPrec p (Neg formula) = showParen (p > 3) $
    showString "not " o showsPrec 3 formula
  showsPrec p (Conj lhs rhs) = showParen (p > 2) $
    showsPrec 3 lhs o showString " and " o showsPrec 2 rhs
  showsPrec p (Disj lhs rhs) = showParen (p > 1) $
   showsPrec 2 lhs o showString " or " o showsPrec 1 rhs
  showsPrec p Top = showString "T"
  showsPrec p Bottom = showString "B"
  showsPrec p (Knows agent formula) = showParen (p > 3) $
   showString ("K" ++ (show agent) ++ " " ) \circ showsPrec 3 formula
  The definition of the propositional logic satisfaction function adapted from the slides.
satisfies :: Assignment 	o Form 	o Bool
satisfies v (P k)
                      = k 'elem' v
satisfies v (Neg f) = not (satisfies v f)
satisfies v (Conj f g) = satisfies v f && satisfies v g
satisfies v (Disj f g) = satisfies v f || satisfies v g
```

```
satisfies _ Bottom = False

Function "make_beta" that given AP_M, AP_W and the valuations of all worlds, computes \beta_M which we use to discard valuations that do not correspond to any possible world.

make_beta :: [Proposition] \rightarrow [Proposition] \rightarrow [(World, Valuation)] \rightarrow Form

make_beta ap_M ap_W valuations = make_beta2 ap_M ap_W ap_W valuations

make_beta2 :: [Proposition] \rightarrow [Proposition] \rightarrow [Proposition] \rightarrow [(World, Valuation)] \rightarrow Form

make_beta2 _ [] ap_W _ = uExists ap_W

make_beta2 ap_M (prop_w:restProps) ap_W (val:valuations) = Conj (Disj (Neg (P prop_w)) (desc ap_M (snd val))) (make_beta2 ap_M restProps ap_W valuations)
```

- uExists constructs the formula that corresponds to the exclusive existential !∃. ap is the list of all propositions.
- possibleWorld is the conjunction of implications from nominal propositions to the descriptions of the world that they denote. So that if some  $p_w$  is true for an interpretation I, then if I is to satisfy the conjunct, it must also satisfy V(w)

```
\mathtt{uExists} \; :: \; [\mathtt{Proposition}] \; \to \; \mathtt{Form}
uExists ap = uExists2 [(x, remove x ap) | x \leftarrow ap]
uExists2 :: [(Proposition, [Proposition])] \rightarrow Form
uExists2 [] = Bottom
uExists2 ((active_prop, inactive_props):rest) = Disj (possibleWorld active_prop inactive_props) (uExists2 rest)
\texttt{possibleWorld} \; :: \; \texttt{Proposition} \; \rightarrow \; \texttt{[Proposition]} \; \rightarrow \; \texttt{Form}
possibleWorld active_prop [] = P active_prop
possibleWorld active_prop (inactive_prop:rest) = Conj (Neg (P inactive_prop)) (possibleWorld active_prop rest)
   The description of a valuation V is the conjunction of the propositions true in V along with the negations of those not true in V.
{\tt desc} \; :: \; {\tt [Proposition]} \; \rightarrow \; {\tt [Proposition]} \; \rightarrow \; {\tt Form}
desc ap props = desc2 ((nub ap) \\ props) props
{\tt desc2} \, :: \, {\tt [Proposition]} \, \to \, {\tt [Proposition]} \, \to \, {\tt Form}
desc2 out_props (in_prop:in_props) = Conj (P in_prop) (desc2 out_props in_props)
desc2 (out_prop:out_props) []
                                           = Conj (Neg (P out_prop)) (desc2 out_props [])
desc2 [] []
                                            = Top
```

#### 2.2 Programs

Definition of the program. Note that the Universal Program can reach any world from any world while the

```
data Program = AssignTo Proposition Form

| Question Form
| Semicolon Program Program
| Cap Program Program
| Cup Program Program
| IdleProgram
| UniversalProgram
| deriving (Eq,Show)

remove element list = filter (λe → e/=element) list
```

```
{\tt runProgram} :: {\tt Form} 	o [{\tt Proposition}] 	o {\tt Program} 	o {\tt Assignment} 	o {\tt Bool}
runProgram beta _ (AssignTo prop formula) a_1 a_2
                                                               satisfies a_1 formula = (sort $ nub $ prop:a_1) == (sort $ nub a_2)
                                                                                       = (sort $ nub (remove prop a_1)) == (sort $ nub a_2)
                                                         = ((sort $ nub a_1) == (sort $ nub a_2)) && satisfies a_1 formula
runProgram beta _ (Question formula) a_1 a_2
runProgram beta ap (Semicolon pi_1 pi_2) a_1 a_2
           not (all (\lambda x \to \text{not ((runProgram beta ap pi_1 a_1 x) \&\& (runProgram beta ap pi_2 x a_2)))} [ assignment | assignment <math>\leftarrow (allAssignmentsFor ap), satisfies assignment beta])
runProgram beta ap (Cap pi_1 pi_2) a_1 a_2
                                                         = (runProgram beta ap pi_1 a_1 a_2) && (runProgram beta ap pi_2 a_1 a_2)
runProgram beta ap (Cup pi_1 pi_2) a_1 a_2
                                                         = (runProgram beta ap pi_1 a_1 a_2) || (runProgram beta ap pi_2 a_1 a_2)
runProgram beta _ IdleProgram _ _
                                                         = False
runProgram beta _ UniversalProgram _ _
                                                        = True
\texttt{programSet} \ :: \ [\texttt{Proposition}] \ \to \ \texttt{Program}
programSet [] = IdleProgram
programSet (current_prop:rest) =
   Semicolon (Cup (AssignTo current_prop Top) (AssignTo current_prop Bottom)) (programSet rest)
```

#### 2.3 Mental Model

Here follows the definition of the mental model and its model checking algorithm.

```
{\tt data} \; {\tt SuccEpistM} = {\tt Mo}
        [Proposition]
       Form
        [(Agent, Program)]
        deriving (Eq,Show)
{\tt modelCheck} \; :: \; {\tt SuccEpistM} \; \rightarrow \; {\tt Assignment} \; \rightarrow \; {\tt Form} \; \rightarrow \; {\tt Bool}
modelCheck model world (P k)
                                              = k 'elem' world
modelCheck model world (Neg f)
                                              = not (modelCheck model world f)
modelCheck model world (Conj f g)
                                              = (modelCheck model world f) && (modelCheck model world g)
modelCheck model world (Disj f g)
                                              = (modelCheck model world f) || (modelCheck model world g)
modelCheck (Mo ap beta programs) world (Knows agent f)
                                              = case (lookup agent programs) of
                                                    Just program_a 	o all (\lambda x 	o (not (runProgram beta ap program_a world x))
                                                               || (modelCheck (Mo ap beta programs) x f) ) [ assignment | assignment ← (allAssignmentsFor ap), satisfies assignment beta]
                                                    Nothing \rightarrow False
```

# 2.4 Kripke Model

Definitions for the Kripke alongside its model checking algorithm.

```
type World = Integer
type Valuation = [Proposition]
type Relation = [(World, [World])]
data KripkeM = KMo
        [(Agent, Relation)]
        [(World, Valuation)]
        deriving (Eq,Show)
```

```
{\tt modelCheckKrpk} :: {\tt KripkeM} \to {\tt World} \to {\tt Form} \to {\tt Bool}
modelCheckKrpk (KMo relations valuations) world (P k) = case (lookup world valuations) of
                                                  Just valuation \rightarrow k 'elem' valuation
                                                  Nothing \rightarrow False
modelCheckKrpk model world (Neg f)
                                                = not (modelCheckKrpk model world f)
modelCheckKrpk model world (Conj f g)
                                                = (modelCheckKrpk model world f) && (modelCheckKrpk model world g)
modelCheckKrpk model world (Disj f g)
                                                = (modelCheckKrpk model world f) || (modelCheckKrpk model world g)
modelCheckKrpk (KMo relations valuations) world (Knows agent f)
                                            = case (lookup agent relations) of
                                                  Just relation_a \rightarrow case (lookup world relation_a) of
                                                     Just neighbors 	o all (\lambda x 	o modelCheckKrpk (KMo relations valuations) x (f) ) neighbors
                                                     Nothing \rightarrow True
                                                 Nothing \rightarrow True
kripkeToSuccint :: KripkeM \rightarrow SuccEpistM
kripkeToSuccint (KMo relations valuations) = Mo ap beta programs
                                                  where ap = sort (ap_M ++ ap_W)
                                                        beta = make_beta ap_M ap_W valuations
                                                        programs = [(agent, relationToProgram relation ap beta) | (agent, relation) ← relations]
                                                        ap_M = nub $ concat [snd x | x \leftarrow valuations]
                                                        ap_W = [-(fst x) \mid x \leftarrow valuations]
```

#### 2.5 Model Translation Algorithms

```
{\tt succintToKripke} :: {\tt SuccEpistM} \to {\tt KripkeM}
succintToKripke (Mo ap beta programs) = KMo relations worldsAndAssignments
                                                  where relations = [(agent, programToRelation program ap beta worldsAndAssignments) | (agent, program) ← programs]
                                                         worldsAndAssignments = zip [1..(toInteger $ length assignments)] assignments
                                                         assignments = [ assignment | assignment ← (allAssignmentsFor ap), satisfies assignment beta]
programToRelation :: Program 
ightarrow [Proposition] 
ightarrow Form 
ightarrow [(World, Assignment)] 
ightarrow Relation
programToRelation program ap beta worldsAndAssignments = programToRelation1 program ap beta worldsAndAssignments worldsAndAssignments
\texttt{programToRelation1} :: \texttt{Program} \rightarrow [\texttt{Proposition}] \rightarrow \texttt{Form} \rightarrow [(\texttt{World, Assignment})] \rightarrow [(\texttt{World, Assignment})] \rightarrow \texttt{Relation}
programToRelation1 _ _ _ [] _ = []
programToRelation1 program ap beta ((world,assignment):rest) worldsAndAssignments =
          (world, [vorld | (vorld, ass1) ← worldsAndAssignments, runProgram beta ap program assignment ass1]):(programToRelation1 program ap beta rest worldsAndAssignments)
\texttt{relationToProgram} \; :: \; \texttt{Relation} \; \rightarrow \; \texttt{[Proposition]} \; \rightarrow \; \texttt{Form} \; \rightarrow \; \texttt{Program}
relationToProgram relation ap beta = relationToProgram1 relation ap
\texttt{relationToProgram1} \; :: \; \texttt{Relation} \; \rightarrow \; \texttt{[Proposition]} \; \rightarrow \; \texttt{Program}
relationToProgram1 [] _ = IdleProgram
relationToProgram1 ((world, neighbors):rest) ap = Cup (relationToProgram2 world neighbors ap) (relationToProgram1 rest ap)
\texttt{relationToProgram2} \; :: \; \texttt{World} \; \rightarrow \; \texttt{[World]} \; \rightarrow \; \texttt{[Proposition]} \; \rightarrow \; \texttt{Program}
relationToProgram2 _ [] _ = IdleProgram
relationToProgram2 w_1 (w_2:rest) ap = Cup ( Semicolon (Semicolon (Question (P (-w_1)) ) UniversalProgram) (Question (P (-w_2))) ) (relationToProgram2 w_1 rest ap)
```

# 3 Model Checking Examples

After model checking phi 1,2,3 on my\_kmodel2 on all its worlds, on both the original Kripke model and its translation to a mental model, we have that both the kripke one and the mental one agree as expected (and the result is consistent with the usual semantics). The results are as follows:

```
*Main> modelCheckKrpk my_kmodel2 2 phi2
True
*Main> modelCheckKrpk my_kmodel2 3 phi2
True
*Main> modelCheckKrpk my_kmodel2 4 phi2
True
*Main> modelCheckKrpk my_kmodel2 5 phi2
True
*Main> modelCheckKrpk my_kmodel2 6 phi2
True
*Main> modelCheckKrpk my_kmodel2 7 phi2
*Main> modelCheckKrpk my_kmodel2 1 phi3
[False
*Main> modelCheckKrpk my_kmodel2 2 phi3
*Main> modelCheckKrpk my_kmodel2 3 phi3
[False
*Main> modelCheckKrpk my_kmodel2 4 phi3
True
*Main> modelCheckKrpk my_kmodel2 5 phi3
True
*Main> modelCheckKrpk my_kmodel2 6 phi3
[False
*Main> modelCheck (kripkeToSuccint my_kmodel2 ) [-1,1,2,3] phi
[True
*Main> modelCheck (kripkeToSuccint my_kmodel2 ) [-1,1,2,3] phi2
[False
*Main> modelCheck (kripkeToSuccint my_kmodel2 ) [-1,1,2,3] phi3
[False
*Main> modelCheck (kripkeToSuccint my_kmodel2 ) [-2,2,3] phi2
True
*Main> modelCheck (kripkeToSuccint my_kmodel2 ) [-2,2,3] phi3
True
*Main> modelCheck (kripkeToSuccint my_kmodel2 ) [-3,1,3] phi2
True
*Main> modelCheck (kripkeToSuccint my_kmodel2 ) [-3,1,3] phi3
[False
*Main> modelCheck (kripkeToSuccint my_kmodel2 ) [-4,3] phi2
True
*Main> modelCheck (kripkeToSuccint my_kmodel2 ) [-4,3] phi3
True
*Main> modelCheck (kripkeToSuccint my_kmodel2 ) [-5,1] phi2
True
*Main> modelCheck (kripkeToSuccint my_kmodel2 ) [-5,1] phi3
*Main> modelCheck (kripkeToSuccint my_kmodel2 ) [-6,1,2] phi2
*Main> modelCheck (kripkeToSuccint my_kmodel2 ) [-6,1,2] phi3
[False
*Main>
```