ECE 310 Recitation 9 Solution Thursday Apr 15, 2021

Concept check

√ DFT spectral analysis

- O Determine the frequency content of a given signal: $x_a(t) = \sum_{i=1}^{M} A_i \cos(\Omega_i t)$, determine $\{\Omega_i, A_i\}_{i=1}^{M}$
- o Ideally

$$X_a(\Omega) = \sum_{i=1}^M \pi A_i [\delta(\Omega + \Omega_i) + \delta(\Omega - \Omega_i)]$$

$$\frac{1}{2} X_a(\frac{\omega}{2}), \quad 0 < \omega < \pi$$

$$X_d(\omega) = \begin{cases} \frac{1}{T} X_a(\frac{\omega}{T}), \ 0 \le \omega \le \pi \\ \frac{1}{T} X_a(\frac{\omega - 2\pi}{T}), \ \pi < \omega \le 2\pi \end{cases}$$

$$X_m = X_d(\frac{2\pi}{N}m)$$

- o Spectral parameters
 - Amplitudes: $\frac{A_i}{T}$
 - Frequencies: $m_i \to \omega = \frac{2\pi}{N} m_i \to \Omega_i = \frac{\omega_i}{T} \text{ or } \Omega_i = \frac{\omega_i}{T} \frac{2\pi}{T}$
- Windowing effect:

$$\hat{x}[n] = x[n]w[n], \widehat{X_d}(\omega) = \frac{1}{2\pi} \int_{-\pi}^{\pi} X_d(u)W_d(\omega - u)du$$

- For sinusoidal input $x[n] = Acos(\Omega_0 nT)$ and rectangular window $w[n] = \begin{cases} 1, & 0 \le n \le N-1 \\ 0, & else \end{cases}$
- $\widehat{X_d}(\omega) = e^{-j(\omega \Omega_0 T)\frac{N-1}{2}} \frac{\frac{A}{2} sin[(\omega \Omega_0 T)\frac{N}{2}]}{\sin[(\omega \Omega_0 T)\frac{1}{2}]} + e^{-j(\omega + \Omega_0 T)\frac{N-1}{2}} \frac{\frac{A}{2} sin[(\omega + \Omega_0 T)\frac{N}{2}]}{\sin[(\omega + T)\frac{1}{2}]}$
- Main lobe height: $\frac{AN}{2}$, main lobe width $\frac{4\pi}{N}$, peak location $\omega = \pm \Omega_0 T$

√ FFT

$$\circ \ \begin{cases} X_m = Y_m + W_N^m Z_m \\ X_{m+N/2} = Y_m - W_N^m Z_m \end{cases} \ \begin{cases} Y_m = DFT\{x[2l]\}_{l=0}^{n/2-1} \ (even) \\ Z_m = DFT\{x[2l+1]\}_{l=0}^{n/2-1} \ (odd) \end{cases}, \ W_N^m = e^{-j\frac{2\pi m}{N}}$$

o Butterfly diagram, bit-reverse indexing

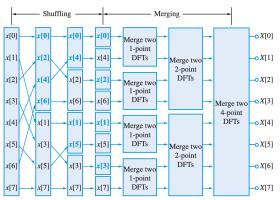


Figure 8.5 The shuffling and merging operations required for recursive computation of the 8-point DFT using the decimation-in-time FFT algorithm.

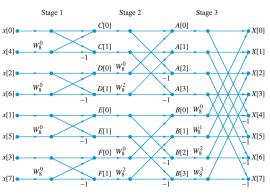


Figure 8.6 Flow graph of 8-point decimation-in-time FFT algorithm using the butterfly computation shown in Figure 8.4. The trivial twiddle factor $W_8^0=1$ is shown for the sake of senerality.

Exercise

- 1. A continuous-time signal $x_c(t) = \cos(\frac{\pi}{3}t)$ is sampled at a rate of 30Hz for 12s to produce a discrete-time signal x[n] with length L = 360.
 - (a) Let X[k] be the length-L DFT of x[n]. At what value(s) of k will X[k] have the greatest magnitude?

$$x[n] = \{x_c(nT)\}_{n=0}^{359} = \left\{\cos(\frac{\pi}{90}n)\right\}_{n=0}^{359}$$

If x[n] was not truncated, the DTFT of a cosine would be a pair of shifted delta functions for every period. Since there is a truncation, the DTFT now has the deltas replaced with sinc-like functions.

Given the resemblance between the DTFT of the cosine and truncated cosine, the greatest magnitudes for both DTFTs are at the same values of ω .

$$\frac{2\pi k}{360} = \frac{\pi}{90} \rightarrow k = 2; \quad 2\pi - \frac{2\pi k}{360} = \frac{\pi}{90} \rightarrow k = 358$$

(b) Suppose that x[n] is zero-padded to a total length of L = 512. At what value(s) of k will X[k] have the greatest magnitude?

Similarly,

$$\frac{2\pi k}{512} = \frac{256}{90} \to k = 2.84 \approx 3; \quad 2\pi - \frac{2\pi k}{512} = \frac{\pi}{90} \to k = 509.15 \approx 509$$

Round to the nearest integer gives the largest X[k].

For the exact derivation of X[k], we can use the above formula:

$$\begin{split} X[k] &= X_d \left(\frac{2\pi}{L} k \right) \\ &= \left(e^{-j(\omega - \Omega_0 T) \frac{N-1}{2}} \frac{\frac{A}{2} \sin \left[(\omega - \Omega_0 T) \frac{N}{2} \right]}{\sin \left[(\omega - \Omega_0 T) \frac{1}{2} \right]} + e^{-j(\omega + \Omega_0 T) \frac{N-1}{2}} \frac{\frac{A}{2} \sin \left[(\omega + \Omega_0 T) \frac{N}{2} \right]}{\sin \left[(\omega + T) \frac{1}{2} \right]} \right) \big|_{\omega = \frac{2\pi}{L} k} \\ &= e^{-j\left(\frac{\pi k}{256} - \frac{\pi}{90} \right) \frac{360-1}{2}} \frac{\sin \left[\left(\frac{\pi k}{256} - \frac{\pi}{90} \right) \frac{360}{2} \right]}{2 \sin \left[\left(\frac{\pi k}{256} - \frac{\pi}{90} \right) \frac{1}{2} \right]} + e^{-j\left(\frac{\pi k}{256} + \frac{\pi}{90} \right) \frac{360-1}{2}} \frac{\sin \left[\left(\frac{\pi k}{256} + \frac{\pi}{90} \right) \frac{360}{2} \right]}{2 \sin \left[\left(\frac{\pi k}{256} + \frac{\pi}{90} \right) \frac{1}{2} \right]} \end{split}$$

Then we can check the values near peak to determine the exact largest X[k].

(c) Suppose that x(t) is only sampled for 2s, so the length of X[k] is L = 60. At what value(s) of k will X[k] have the greatest magnitude?

$$x[n] = \{x_c(nT)\}_{n=0}^{59} = \left\{\cos(\frac{\pi}{90}n)\right\}_{n=0}^{59}$$

$$\frac{2\pi k}{60} = \frac{\pi}{90} \to k = \frac{1}{3} \approx 0; \quad 2\pi - \frac{2\pi k}{60} = \frac{\pi}{90} \to k = \frac{179}{3} \approx 60 \text{ (in next period)}$$

What happens in time domain?

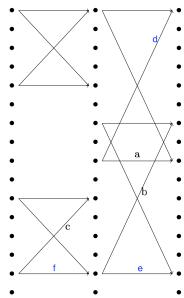
Period of cosine $T_0 = \frac{2\pi}{\pi/3} = 6s$, sampling time $T_s = 2s < T_0$. Not even sampled an entire period! (Information loss!)

(d) Suppose that the x[n] from (c) is zero-padded to a total length of L=64. At what value(s) of k will X[k] have the greatest magnitude?

$$\frac{2\pi k}{64} = \frac{\pi}{90} \to k = \frac{32}{90} \approx 0; 2\pi - \frac{2\pi k}{64} = \frac{\pi}{90} \to k = \frac{179 \times 32}{90} \approx 64 \ (in \ next \ period)$$

Zero padding cannot save an "incomplete" sample!

2. The diagram below represents a part of the computation in a 16-point decimation-in-time radix-2 FFT. Indicate the values of the three branch weights, d, e, and f.



- Branches from the top half are always 1 (coefficient of Y_m is 1).
- Branches going up from the bottom half are W_N^m (coefficient of Z_m).
- Branches going straight across in the bottom half are $-W_N^m$ (coefficient of Z_m).
- N is the current DFT size, m is index of current DFT

$$\begin{aligned} d &= W_{16}^0 = 1 \\ e &= -W_{16}^6 = -e^{-j\frac{2\pi 6}{16}} = -e^{-j\frac{3\pi}{4}} = \frac{\sqrt{2}}{2} + j\frac{\sqrt{2}}{2} \\ f &= -W_{8}^2 = -e^{-j\frac{2\pi 2}{8}} = -e^{-j\frac{\pi}{2}} = j \end{aligned}$$