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ECE448: Artificial Intelligence

Lecture 7: Planning and Theorem Proving

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1. **Examples**
2. **Automatic Theorem Proving: forward-chaining, backward-chaining**
3. **Planning: forward-chaining, backward-chaining**
4. **Admissible Heuristics for Planning and Theorem Proving**
 - **Number of Steps**
 - **Planning Graph**
5. **Computational Complexity**

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https://en.wikipedia.org/wiki/River_crossing_puzzle

- A farmer has a fox, a goat, and a bag of beans to get across the river
- His boat will only carry him + one object
- He can't leave the fox with the goat
- He can't leave the goat with the bag of beans



https://en.wikipedia.org/wiki/River_crossing_puzzle

fgb -----(farmer, goat)----> fGb

fGb <-----(farmer)-----

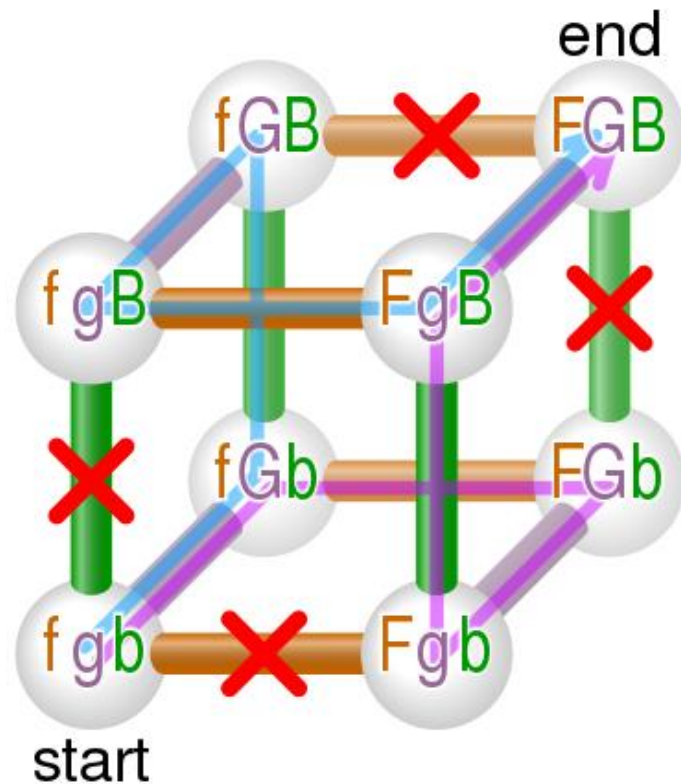
-----(farmer,fox)----> FGb

Fgb <--(farmer,goat)-----

-----(farmer,beans)---> FgB

FgB <------(farmer)-----

-----(farmer,goat)----> FGB



- You have packages waiting for pickup at Atlanta, Boston, Charlotte, Denver, Edmonton, and Fairbanks
- They must be delivered to Albuquerque, Baltimore, Chicago, Des Moines, El Paso, and Frisco
- You have two trucks. Each truck can hold only two packages at a time.

"Simultaneous Selective Disassembly and End-of-Life Decision Making for Multiple Products That Share Disassembly Operations," Sara Behdad, Minjung Kwak, Harrison Kim and Deborah Thurston. *J. Mech. Des* **132**(4), 2010, [doi:10.1115/1.4001207](https://doi.org/10.1115/1.4001207)

- Design decisions limit the sequence in which you can disassemble a product at the end of its life
- Problem statement: design the product in order to make disassembly as cheap as possible

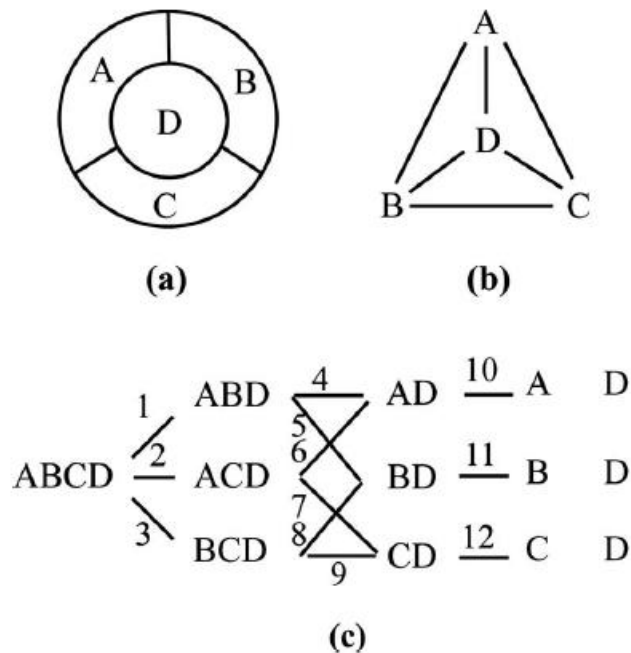


Fig. 1 Simple assembly (a), its connection diagram (b), and its disassembly graph (c) [23]

Application of planning: the Gale-Church alignment algorithm for machine translation

Table 2

Output from alignment program.

English	French
According to our survey, 1988 sales of mineral water and soft drinks were much higher than in 1987, reflecting the growing popularity of these products. Cola drink manufacturers in particular achieved above-average growth rates.	Quant aux eaux minérales et aux limonades, elles rencontrent toujours plus d'adeptes. En effet, notre sondage fait ressortir des ventes nettement supérieures à celles de 1987, pour les boissons à base de cola notamment.
The higher turnover was largely due to an increase in the sales volume.	La progression des chiffres d'affaires résulte en grande partie de l'accroissement du volume des ventes.
Employment and investment levels also climbed.	L'emploi et les investissements ont également augmenté.

Application of planning: the Gale-Church alignment algorithm for machine translation

1. Let $d(x_1, y_1; 0, 0)$ be the cost of substituting x_1 with y_1 ,
2. $d(x_1, 0; 0, 0)$ be the cost of deleting x_1 ,
3. $d(0, y_1; 0, 0)$ be the cost of insertion of y_1 ,
4. $d(x_1, y_1; x_2, 0)$ be the cost of contracting x_1 and x_2 to y_1 ,

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5. $d(x_1, y_1; 0, y_2)$ be the cost of expanding x_1 to y_1 and y_2 , and
6. $d(x_1, y_1; x_2, y_2)$ be the cost of merging x_1 and x_2 and matching with y_1 and y_2 .

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$$\begin{aligned}
 \text{Sentence} \rightarrow & \\
 & \text{Function}(\text{Term}, \dots) \\
 & | \neg \text{Sentence} \\
 & | \text{Sentence} \wedge \text{Sentence} \\
 & | \text{Sentence} \vee \text{Sentence} \\
 & | \text{Sentence} \Rightarrow \text{Sentence} \\
 & | \text{Sentence} \Leftrightarrow \text{Sentence} \\
 & | \text{Quantifier Variable}, \dots \text{Sentence}
 \end{aligned}$$

$$\begin{aligned}
 \text{Term} \rightarrow & \text{Function}(\text{Term}) \\
 & | \text{Variable} | \text{Constant}
 \end{aligned}$$

$$\text{Quantifier} \rightarrow \exists \mid \forall$$

A “sentence” is

- an evaluated function, or
- a negated sentence, or
- the conjunction of 2 sentences, or
- the disjunction of 2 sentences, or
- an implication, or
- an equivalence, or
- a sentence with a quantified variable.

A “term” is an evaluated function, or a variable, or a constant.

A “quantifier” is “there exists,” or “for all.”

English	First-Order Logic Notation
It is a crime for Americans to sell weapons to hostile nations.	$American(x) \wedge Weapon(y) \wedge$ $Sells(x, y, z) \wedge Hostile(z)$ $\Rightarrow Criminal(x)$
Colonel West sold missiles to Ganymede.	$\exists x, Missile(x) \wedge Sells(West, x, Ganymede)$
Colonel West is American.	$American(West)$
Ganymede is an enemy of America.	$Enemy(Ganymede, America)$
Missiles are weapons.	$Missile(x) \Rightarrow Weapon(x)$
An enemy of America is a hostile nation.	$Enemy(x, America) \Rightarrow Hostile(x)$

First-Order Logic Notation

$$\text{American}(x) \wedge \text{Weapon}(y) \wedge \\ \text{Sells}(x, y, z) \wedge \text{Hostile}(z) \\ \Rightarrow \text{Criminal}(x)$$
$$\exists x, \text{Missile}(x) \\ \wedge \text{Sells}(\text{West}, x, \text{Ganymede})$$
$$\text{American}(\text{West})$$
$$\text{Enemy}(\text{Ganymede}, \text{America})$$
$$\text{Missile}(x) \Rightarrow \text{Weapon}(x)$$
$$\text{Enemy}(x, \text{America}) \Rightarrow \text{Hostile}(x)$$

Can we prove the theorem:

Criminal(West)?

- **Universal Instantiation:**

- given the sentence $\forall x, \text{Function}(x)$,
- for any known constant C ,
- it is possible to generate the sentence $\text{Function}(C)$.

- **Existential Instantiation:**

- given the proposition $\exists x, \text{Function}(x)$,
- if no known constant A is known to satisfy $\text{Function}(A)$, then
- it is possible to define a new, otherwise unspecified constant B , and
- to generate the sentence $\text{Function}(B)$.

- **Generalized Modus Ponens:**

- Given the sentence $p_1(x_1) \wedge p_2(x_2) \wedge \dots \wedge p_n(x_n) \Rightarrow q(x_1, \dots, x_n)$, and
- given the sentences $p_1(C_1), \dots, p_n(C_n)$ for any constants C_1, \dots, C_n ,
- it is possible to generate the sentence $q(C_1, \dots, C_n)$

- **Existential Instantiation:**

- Input: $\exists x, \text{Missile}(x) \wedge \text{Sells}(\text{West}, x, \text{Ganymede})$
- Output: $\text{Missile}(M) \wedge \text{Sells}(\text{West}, M, \text{Ganymede})$

- **Generalized Modus Ponens:**

- Input: $\text{Missile}(M)$ and $\text{Missile}(x) \Rightarrow \text{Weapon}(x)$
- Output: $\text{Weapon}(M)$

- **Generalized Modus Ponens:**

- Input: $\text{Enemy}(\text{Ganymede}, \text{America})$ and $\text{Enemy}(x, \text{America}) \Rightarrow \text{Hostile}(x)$
- Output: $\text{Hostile}(\text{Ganymede})$

- **Generalized Modus Ponens:**

- Input: $\text{American}(x) \wedge \text{Weapon}(y) \wedge \text{Sells}(x, y, z) \wedge \text{Hostile}(z) \Rightarrow \text{Criminal}(x)$ and $\text{American}(\text{West}), \text{Weapon}(M), \text{Sells}(\text{West}, M, \text{Ganymede}), \text{Hostile}(\text{Ganymede})$
- Output: $\text{Criminal}(\text{West})$

- State = the set of all currently known sentences
- Action = generate a new sentence
- Goal State = a set of sentences that includes the target sentence

(Question to ponder: how do you disprove a target sentence?)

- **What's Special About Theorem Proving:**
 - A state, at level n , can be generated by the combination of several states at level $n-1$.
- **Definition: Forward Chaining** is a search algorithm in which each action
 - generates a new sentence,
 - by combining as many different preceding states as necessary.

$\{p_1, p_2, p_1 \Rightarrow q_1, p_2 \Rightarrow q_2, q_1 \wedge q_2 \Rightarrow q_3\}$

Initial State

$\{p_1, p_2, p_1 \Rightarrow q_1, p_2 \Rightarrow q_2, q_1 \wedge q_2 \Rightarrow q_3, q_1\}$

Search "Tree" Level 1

$\{p_1, p_2, p_1 \Rightarrow q_1, p_2 \Rightarrow q_2, q_1 \wedge q_2 \Rightarrow q_3, q_2\}$

Search "Tree" Level 2:
Goal Achieved

$\{p_1, p_2, p_1 \Rightarrow q_1, p_2 \Rightarrow q_2, q_1 \wedge q_2 \Rightarrow q_3, q_1, q_2, q_3\}$

- **What Else is Special About Theorem Proving:**
 - The "Goal State" is defined to be any set of sentences that includes the target sentence
- **Definition: Backward Chaining** is a search algorithm in which
 - State = {set of known sentences}, {set of desired sentences}
 - Action = apply a known sentence, backward, to a target sentence, in order to generate a new set of desired sentences
 - Goal = all "desired sentences" are part of the set of "known sentences"

KNOWN: $\{p_1, p_2, p_1 \Rightarrow q_1, p_2 \Rightarrow q_2, q_1 \wedge q_2 \Rightarrow q_3\}$
DESIRED: $\{q_3\}$

Initial State

DESIRED: $\{q_1, q_2\}$

Search Tree Level 1

DESIRED: $\{p_1, q_2\}$

DESIRED: $\{q_1, p_2\}$

Search Tree Level 2

DESIRED: $\{p_1, p_2\}$

DESIRED: $\{p_1, p_2\}$

Search Tree Level 3:
Goal Achieved

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- A search problem is defined by:
 - Initial state
 - Goal state
 - Actions
 - Transition model
 - Cost

- STRIPS (Stanford Research Institute Problem Solver): classical planning framework from the 1970s
- **States** are specified as conjunctions of predicates
 - Start state: $\text{At}(\text{home}) \wedge \text{Sells}(\text{SM}, \text{Milk}) \wedge \text{Sells}(\text{SM}, \text{Bananas}) \wedge \text{Sells}(\text{HW}, \text{drill})$
 - Goal state: $\text{At}(\text{home}) \wedge \text{Have}(\text{Milk}) \wedge \text{Have}(\text{Banana}) \wedge \text{Have}(\text{drill})$
- **Actions** are described in terms of preconditions and effects:
 - $\text{Go}(x, y)$
 - **Precond:** $\text{At}(x)$
 - **Effect:** $\neg \text{At}(x) \wedge \text{At}(y)$
 - $\text{Buy}(x, \text{store})$
 - **Precond:** $\text{At}(\text{store}) \wedge \text{Sells}(\text{store}, x)$
 - **Effect:** $\text{Have}(x)$
- Planning is “just” a search problem

- A planning action is like a “ $p \Rightarrow q$ ” statement.
 - In order to be applied, it requires certain input sentences to be true. For example, the action “put the goat in the boat” requires, as its precondition, that the boat is empty.
 - The result of the action is the generation of an output sentence. For example: “the goat is now in the boat.”
- The initial state is a set of sentences that are initially true.
- The goal state is a set of sentences that we want to “prove.”

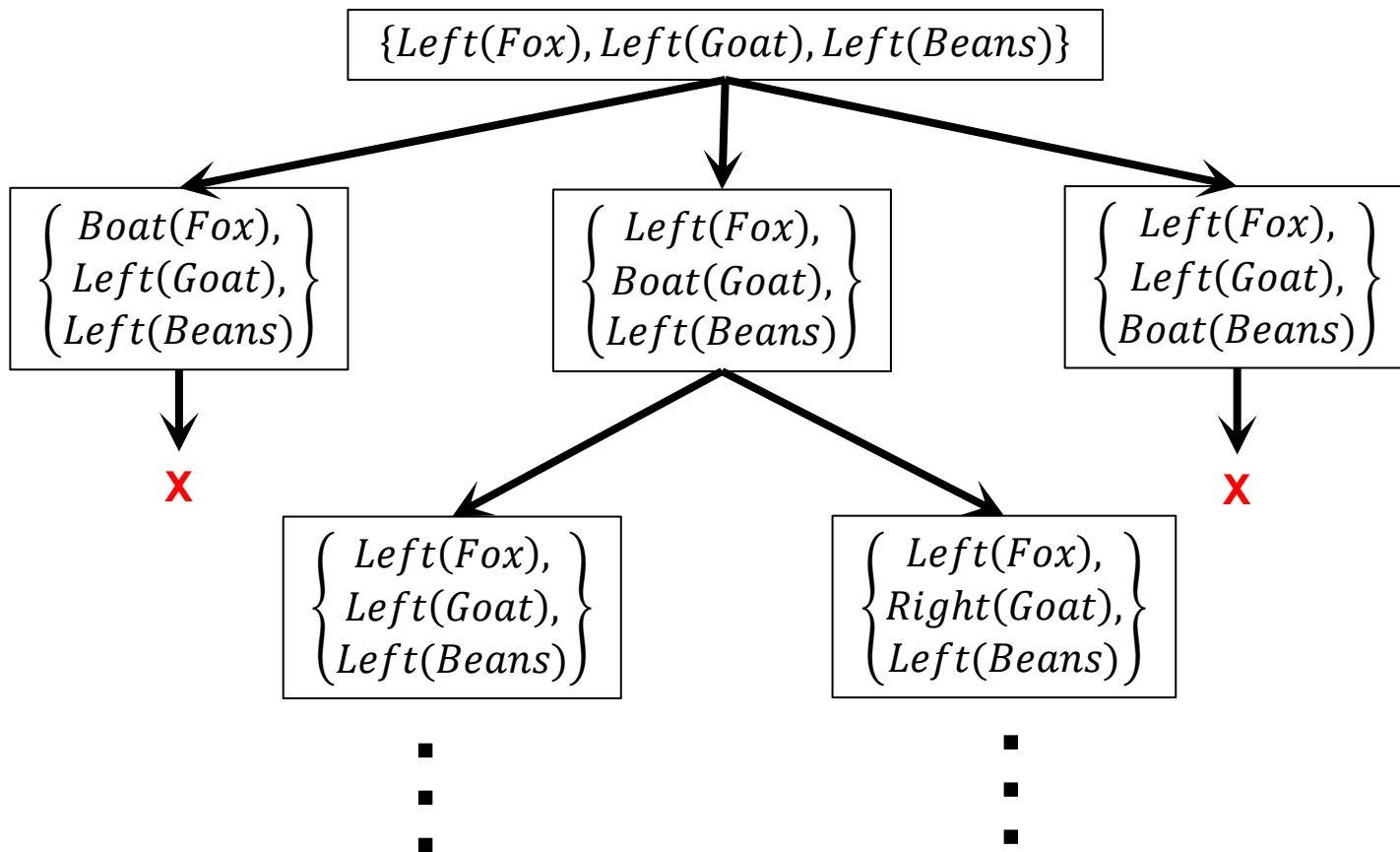
Important differences between Planning and Theorem Proving, #1: Negating your preconditions

- A planning action may **NEGATE** some of its preconditions.
 - Example: the action “put the goat in the boat” requires, as its precondition, the sentence $\neg \text{Boat}(\text{goat})$.
 - It generates, as its output, the sentence: $\text{Boat}(\text{goat})$.
- No action can combine two world states that contain contradictory sentences. For example, you can't combine the states $\{p, q\}$ and $\{p, \neg q\}$ to get the state $\{p, q, \neg q\}$.

Starting with the start state, find all applicable actions (actions for which preconditions are satisfied), compute the successor state based on the effects, keep searching until goals are met

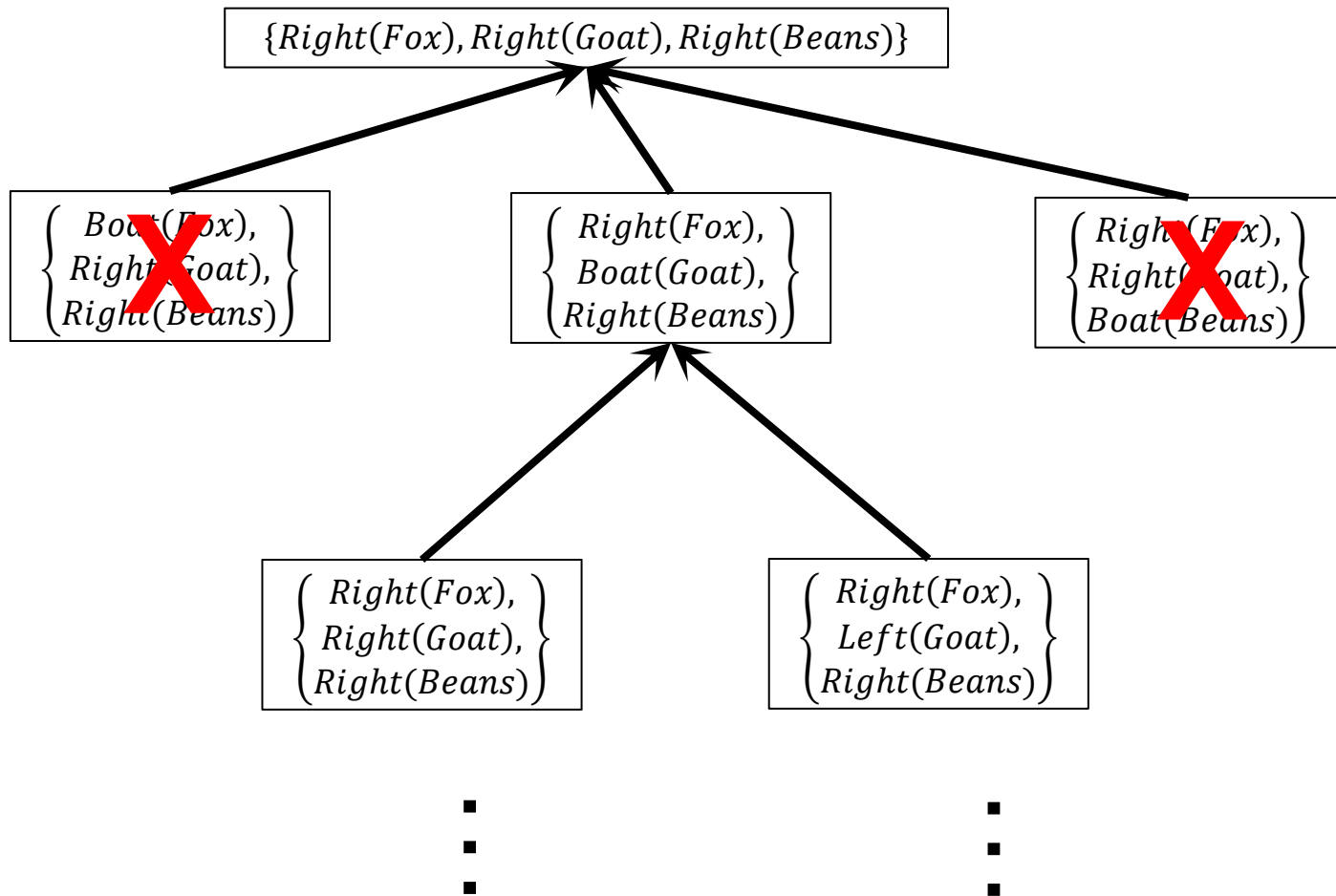
- Can work well with good heuristics





Starting with the goal state (a set of target sentences),

- find all applicable actions (actions that would generate a sentence in the goal state).
- For each, generate the predecessor state as a new set of target sentences.
- Keep searching until all target sentences are in the initial state.



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- Heuristics from Constraint Relaxation: The heuristic $h(n)$ is the number of steps it would take to get from n to G , if problem constraints were relaxed --- this guarantees that $h(n)$ is admissible
- $h_1(n)$ dominates $h_2(n)$ ($h_1(n) \geq h_2(n)$) if $h_1(n)$ is computed by relaxing fewer constraints.

Heuristic #1: Count the number of actions necessary to generate all of the sentences in the goal state that aren't already true.

- What got relaxed: we ignore action pre-requisites.

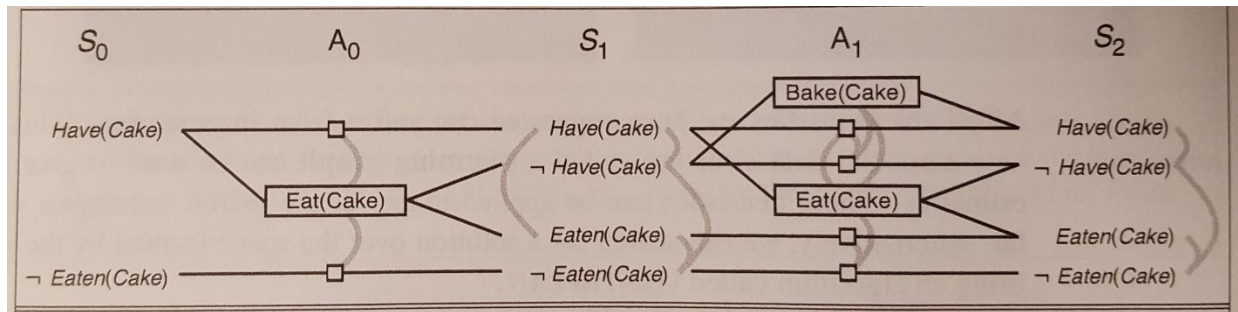
Example: 6 people on left side of the river, we want 6 people on the right side, we have a 2-person boat. Minimum # actions: $h(n) = 3$.

A **planning graph** is a trellis whose stages are:

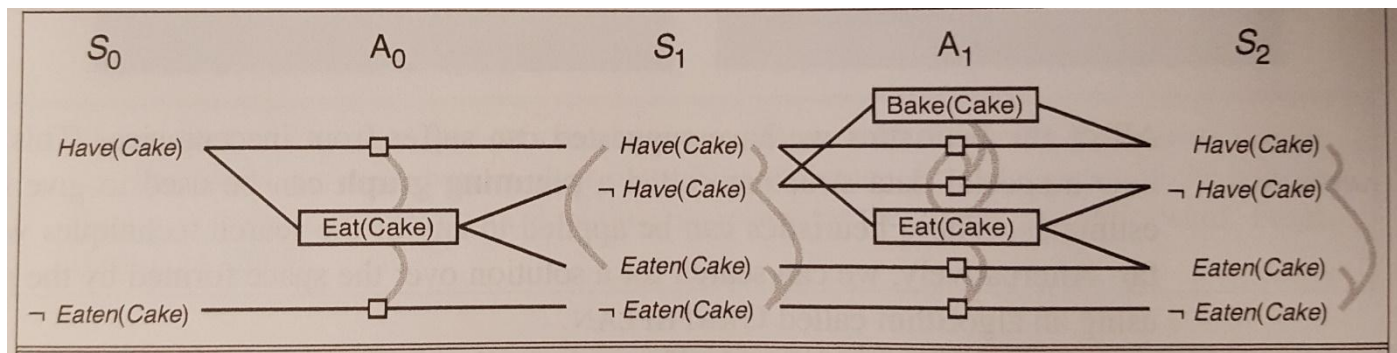
- **Action stages** (A_n): list all of the actions whose pre-requisites are available in “Sentences stage” S_n
- **Sentence stages** (S_{n+1}): list all of the sentences that were available in S_n , plus any new sentences that could have been generated by any action in A_n

And within each stage, we have:

- **Mutex links**: If ALL actions that generate output sentence p also generate $\neg q$, then the sentences p and q become **mutex** (mutually exclusive).

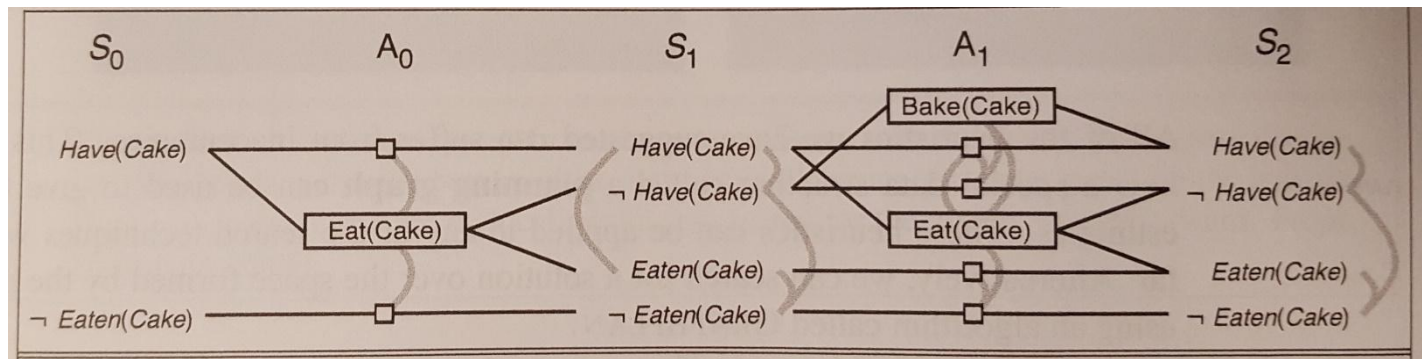


- A_0 has only two possible actions:
 - Do nothing: reproduces the initial state, $\{Have(Cake), \neg Eaten(Cake)\}$
 - Eat(Cake): generates $\{\neg Have(Cake), Eaten(Cake)\}$
- Therefore, at S_1 , $Have(Cake)$ is mutex with $Eaten(Cake)$
- A_1 : $Bake(Cake) \rightarrow Have(Cake)$, without generating $\neg Eaten(Cake)$, so...
- S_1 : $Have(Cake)$ and $Eaten(Cake)$ are no longer mutex.



- **# of mutex links is monotonically non-increasing:** If a pair of sentences are not mutex at stage S_n , then they are also not mutex at S_{n+1}
- **# possible actions is monotonically non-decreasing:** If an action is possible at stage A_n , then it is also possible at A_{n+1}

Heuristic #2: Number of stages until target sentences are non-mutex



Heuristic: # stages between the current stage and the first stage at which all of the goal-state sentences are no longer mutex

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- Planning is PSPACE-complete > NP-complete
 - The computational complexity of finding a plan is exponential
 - The length of the plan is exponential
 - Space necessary to represent it
 - Time necessary to implement it
 - The only thing that's polynomial: the amount of space necessary to represent the world state while finding or implementing a plan
- Example: towers of Hanoi



- Planning is PSPACE-complete
 - The length of a plan can be exponential in the number of “objects” in the problem!
 - So is game search
- Archetypal PSPACE-complete problem: *quantified boolean formula* (QBF)
 - Example: is this formula true?

$$\exists x_1 \forall x_2 \exists x_3 \forall x_4 (x_1 \vee \neg x_3 \vee x_4) \wedge (\neg x_2 \vee x_3 \vee \neg x_4)$$

- Compare to SAT:

$$\exists x_1 \exists x_2 \exists x_3 \exists x_4 (x_1 \vee \neg x_3 \vee x_4) \wedge (\neg x_2 \vee x_3 \vee \neg x_4)$$

- Relationship between SAT and QBF is akin to the relationship between puzzles and games

- Resource constraints
 - Instead of “static,” the world is “semidynamic:” we can’t think forever
- Actions at different levels of granularity: hierarchical planning
 - In order to make the depth of the search smaller, we might convert the world from “fully observable” to “partially observable”
- Contingencies: actions failing
 - Instead of being “deterministic,” maybe the world is “stochastic”
- Incorporating sensing and feedback
 - Possibly necessary to address stochastic or multi-agent environments