

第1章 整除 参考答案

计算证明

1. 计算下面整数对的最大公因子和最小公倍数.

$$(1) (202, 282) \quad 2, 28482$$

$$(2) (-666, 1410) \quad 6, -156510$$

$$(3) (30, 105, 360) \quad 15, 2520$$

$$(4) (8n^2 + 28n + 12, 12n^2 + 30n + 12)$$

$$(8n^2 + 28n + 12, 12n^2 + 30n + 12) = (4(2n+1)(n+3), 6(2n+1)(n+2)) = (4, 6) \cdot ((2n+1)(n+3), (2n+1)(n+2)) \\ = 2(2n+1)$$

$$[8n^2 + 28n + 12, 12n^2 + 30n + 12] = 12(2n+1)(n+2)(n+3)$$

2. 求下面整数的标准分解式.

$$(1) 69 \quad 69 = 3^1 \times 23^1$$

$$(2) 200 \quad 200 = 2^3 \times 5^2$$

$$(3) 3288 \quad 3288 = 2^3 \times 3^1 \times 137^1$$

$$(4) 22345680 \quad 22345680 = 2^4 \times 3^1 \times 5^1 \times 7^1 \times 47^1 \times 283^1$$

3. 若 $a \in \mathbb{Z}^+ \cup \{0\}$, $a^4 - 3a^2 + 9$ 是质数还是合数?

$$\text{解} \quad a^4 - 3a^2 + 9 = (a^2 - 3a + 3)(a^2 + 3a + 3) = [(a-1)(a-2) + 1][(a+1)(a+2) + 1]$$

$$\text{当 } a = 0, \quad a^4 - 3a^2 + 9 = 9 \rightarrow \text{合数}$$

$$\text{当 } a = 1, \quad a^4 - 3a^2 + 9 = 7 \rightarrow \text{质数}$$

$$\text{当 } a = 2, \quad a^4 - 3a^2 + 9 = 13 \rightarrow \text{质数}$$

$$\text{当 } a > 2, \quad (a+1)(a+2) + 1 > (a-1)(a-2) + 1 > 1 \rightarrow \text{合数}$$

综上所述, 当 $a = 1$ 或 $a = 2$ 时为质数, 当 $a = 0$ 或 $a > 2$ 时为合数.

4. 若 $m-p \mid mn+pq$, 求证 $m-p \mid mq+np$.

证明 易知 $m-p \mid (m-p)(n-q)$, 又 $m-p \mid mn+pq$, 而 $-(m-p)(n-q) + (mn+pq) = mq+np$, 则 $m-p \mid mq+np$, 证毕.

5. 设 $3 \mid a^2 + b^2$, 求证: $3 \mid a$ 且 $3 \mid b$.

证明 设 $a = 3m + p$, $b = 3n + q$, ($m, n \in \mathbb{Z}$, $p, q \in \{-1, 0, 1\}$),

则 $a^2 + b^2 = (3m + p)^2 + (3n + q)^2 = 3(3m^2 + 3n^2 + 2mp + 2nq) + p^2 + q^2$.

由 $3 \mid a^2 + b^2$, 得 $3 \mid p^2 + q^2$, 而 $p^2 + q^2 = \begin{cases} 0, & p = q = 0 \\ 1, & p + q = \pm 1 \\ 2, & p = \pm 1, q = \pm 1 \end{cases}$, 故 $p = q = 0$, $3 \mid a$ 且 $3 \mid b$, 证毕.

6. 设 $a = qn - t$, 若 $a \mid pm$, 已知 $p - q = t$ 且 $(a, n + 1) = 1$, 求证: $a \mid tm$.

证明 由 $a \mid pm$ 知 $a \mid pmn$, 又 $a \mid am$, 由整系数线性组合得 $a \mid pmn - am$, 而 $a = qn - t$, 则 $a \mid (p - q)m(n + 1)$, 又有 $p - q = t$ 且 $(a, n + 1) = 1$, 得 $a \mid tm$, 证毕.

7. 求证任意 n 个连续的正整数乘积都被 $n!$ 整除. (写出严谨的证明过程)

证明 数学归纳法.

记任意 n 个连续的正整数分别为 a_0, a_1, \dots, a_{n-1} . 其中 $a_i = m + i$ ($m \in \mathbb{N}$), 乘积记为 T_m .

当 $m = 1$ 时, $T_1 = n!$, $n! \mid T_1$ 成立.

假设当 $m = k$ ($k \in \mathbb{N}$) 时, $n! \mid T_k$, 即 $n! \mid k(k+1) \cdots (k+n-1)$ 成立.

那么当 $m = k + 1$ 时, 有 $T_{k+1} = \prod_{i=0}^{n-1} a_i = (k+1)(k+2) \cdots (k+n) = k(k+1) \cdots (k+n-1) + n(k+1) \cdots (k+n-1)$.

只需证 $n! \mid n(k+1) \cdots (k+n-1)$, 即 $(n-1)! \mid (k+1) \cdots (k+n-1)$.

取 n 为 $n-1$, m 为 k , 重复上述证明过程.

易知经过有限步后, $n = 1$, 而 $\forall m \in \mathbb{N}$, $1 \mid m$, 假设成立, 证毕.

8. 求证 $12 \mid n^4 + 2n^3 + 11n^2 + 10n$, $n \in \mathbb{Z}$.

证明 由 $n^4 + 2n^3 + 11n^2 + 10n = (n-1)n(n+1)(n+2) + 12n(n+1)$, 只需证 $12 \mid (n-1)n(n+1)(n+2)$, 易知 $[1, 2, 3, 4] \mid (n-1)n(n+1)(n+2)$ (严谨证明可用数学归纳法), 证毕.

9. 证明 n 的标准分解式中次数都是偶数当且仅当 n 是完全平方数.

证明 充分性. 设 $n = m^2$, 由算术基本定义有 $\begin{cases} n = p_1^{\alpha_1} \cdots p_s^{\alpha_s}, & p_i \text{ 是素数且 } p_i < p_j (0 < i < j \leq s) \\ m = q_1^{\beta_1} \cdots q_t^{\beta_t}, & q_i \text{ 是素数且 } q_i < q_j (0 < i < j \leq t) \end{cases}$, 由算术分解的唯一性可得 $\begin{cases} s = t \\ p_i = q_i, & 0 < i \leq s. \\ \alpha_i = 2\beta_i, & 0 < i \leq s \end{cases}$. 必要性. 设 $n = p_1^{2\alpha_1} \cdots p_s^{2\alpha_s}$, p_i 是素数且 $p_i < p_j (0 < i < j \leq s)$, 取 $m = p_1^{\alpha_1} \cdots p_s^{\alpha_s}$ 即可满足 $n = m^2$. 证毕.

10. 证明 $\sqrt{5}$ 是无理数. *并将其表示为简单连分数的形式. (*标注表示不作强制要求)

证明 反证法. 反设 $\sqrt{5} = \frac{p}{q}$, 其中 $p, q \in \mathbb{Z}^+$, $(p, q) = 1$. 得 $p^2 = 5q^2$, 则 $5 \mid p$. 可设 $p = 5m$, 代入得 $q^2 = 5m^2$, 同理 $5 \mid q$. 则 $(p, q) = 5 \neq 1$ 矛盾, 故反设不成立, 证毕.

对无理数 $\alpha = \sqrt{5}$, 有:

$$a_0 = [\sqrt{5}] = 2, \alpha_1 = \frac{1}{\sqrt{5} - 2} = \sqrt{5} + 2$$
$$a_1 = [\sqrt{5} + 2] = 4, \alpha_2 = \frac{1}{\sqrt{5} + 2 - 4} = \sqrt{5} + 2 = \alpha_1$$

故 $\sqrt{5} = [2; \overline{4}]$.

编程练习 (基于C/C++)

1. 编写程序使用Eratosthenes筛法打印1 000 000内所有素数及个数, 效果如图所示. (*思考: a.对比筛法与普通算法的性能差异; b.递归调用该算法求更大范围素数进行优化; c.求更大的素数 (如 2^{512} 数量级) 该方法是否适用? 会引入哪些新的问题?)

```

1  #include<iostream>
2  #include<math.h>
3  using namespace std;
4
5  bool is_prime(int n)
6  {
7      if (n < 2)return false;
8      bool flag = true;
9      for (int i = 2; i < n; i++)
10     {
11         if (n % i == 0)
12         {
13             flag = false;
14             break;
15         }
16     }
17     return flag;
18 }
19
20 //based on Eratosthenes once
21 int find_prime(int range)
22 {
23     int count = 0;
24     bool* num = new bool[range + 1];
25     for (int i = 1; i < range + 1; i++)
26         num[i] = true;
27     for (int i = 2; i < sqrt(range); i++)
28     {
29         if(is_prime(i))
30         {
31             for (int temp = i + i; temp <= range; temp += i)
32                 num[temp] = false;
33         }
34     }
35     for (int i = 2; i < range + 1; i++)
36     {
37         if (num[i])
38         {
39             count++;
40             cout << i << ", ";
41         }
42     }
43     cout << endl;
44     return count;
45 }
46
47 int main()
48 {
49     int n;
50     cout << "Please input the range: 1-";
51     cin >> n;
52     int count = find_prime(n);
53     cout << "Total: " << count << endl;
54     system("pause");
55 }

```

2. 编写程序计算最大公因数和最小公倍数，效果如图所示.

Microsoft Visual Studio 调试控制台

```
a=9876
b=6789
gcd(a,b)=3
lcm(a,b)=22349388
```

```
1  #include<iostream>
2  using namespace std;
3
4  int Euclid(int a, int b)
5  {
6
7      if (a < 0)a = -a;
8      if (b < 0)b = -b;
9      int r = a;
10     while (r != 0)
11     {
12         r = a % b;
13         a = b;
14         b = r;
15     }
16     return a;
17 }
18
19 int main()
20 {
21     int a, b;
22     cout << "a=";
23     cin >> a;
24     cout << "b=";
25     cin >> b;
26     int gcd = Euclid(a, b);
27     int lcm = a * b / gcd;
28     cout << "gcd(a,b)=" << gcd << endl;
29     cout << "lcm(a,b)=" << lcm << endl;
30 }
```

3. 编写程序实现算术基本定理，效果如下所示。

Microsoft Visual Studio 调试控制台

```
Please input n(n>0): 888
888=2^3*3^1*37^1
```

```
1  #include<iostream>
2  #include<map>
3  using namespace std;
4  void num_decompose(int n)
5  {
6      cout << n << "=";
7      map<int, int>nums;
8      while (n!=1)
9      {
10         int i = 2;
```

```

11     while (i <= n)
12     {
13         if (n % i == 0)
14         {
15             auto iter = nums.find(i);
16             iter == nums.end() ? nums[i] = 1 : iter->second++;
17             n /= i;
18             break;
19         }
20         i++;
21     }
22 }
23
24 for (auto iter = nums.begin(); iter != nums.end(); )
25 {
26     cout << iter->first << "^" << iter->second;
27     if (++iter == nums.end())
28         cout << endl;
29     else
30         cout << "*";
31 }
32 }
33
34 int main()
35 {
36     int n;
37     cout << "Please input n(n>0): ";
38     cin >> n;
39     if (n == 1)
40         cout << "1=1" << endl;
41     else
42         num_decompose(n);
43     return 0;
44 }

```