

SMART CONTRACT AUDIT REPORT

for

Infusion

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1 Introduction

Given the opportunity to review the design document and related source code of the Infusion protocol, we outline in the report our systematic approach to evaluate potential security issues in the smart contract implementation, expose possible semantic inconsistencies between smart contract code and design document, and provide additional suggestions or recommendations for improvement. Our results show that the given version of smart contracts can be further improved due to the presence of several issues related to either security or performance. This document outlines our audit results.

1.1 About Infusion

Infusion is designed to allow low-cost, low-slippage trades on uncorrelated or tightly correlated assets. It is in essence a DEX that is built starting from Solidly/Velodrome with a unique AMM. The DEX is compatible with all the standard features as popularized by UniswapV2 with a number of novel improvements, including price oracles without upkeeps, a new curve $(x^3y + xy^3 = k)$ for efficient stable swaps. The basic information of audited contracts is as follows:

Item Description

Name Infusion

Type Smart Contract

Language Solidity

Audit Method Whitebox

Latest Audit Report February 19, 2024

Table 1.1: Basic Information of Infusion

In the following, we show the Git repository of reviewed files and the commit hash value used in this audit:

https://github.com/Infusion-Finance/infusion-contracts.git (5b2b8d4)

And this is the commit ID after all fixes for the issues found in the audit have been checked in:

https://github.com/Infusion-Finance/infusion-contracts.git (3ece057)

1.2 About PeckShield

PeckShield Inc. [11] is a leading blockchain security company with the goal of elevating the security, privacy, and usability of current blockchain ecosystems by offering top-notch, industry-leading services and products (including the service of smart contract auditing). We are reachable at Telegram (https://t.me/peckshield), Twitter (http://twitter.com/peckshield), or Email (contact@peckshield.com).

High Critical High Medium

High Medium

Low

Medium Low

High Medium

Low

High Medium

Low

Likelihood

Table 1.2: Vulnerability Severity Classification

1.3 Methodology

To standardize the evaluation, we define the following terminology based on OWASP Risk Rating Methodology [10]:

- <u>Likelihood</u> represents how likely a particular vulnerability is to be uncovered and exploited in the wild;
- Impact measures the technical loss and business damage of a successful attack;
- Severity demonstrates the overall criticality of the risk.

Likelihood and impact are categorized into three ratings: H, M and L, i.e., high, medium and low respectively. Severity is determined by likelihood and impact, and can be accordingly classified into four categories, i.e., Critical, High, Medium, Low shown in Table 1.2.

To evaluate the risk, we go through a list of check items and each would be labeled with a severity category. For one check item, if our tool or analysis does not identify any issue, the contract is considered safe regarding the check item. For any discovered issue, we might further

Table 1.3: The Full List of Check Items

Category	Check Item		
	Constructor Mismatch		
	Ownership Takeover		
	Redundant Fallback Function		
	Overflows & Underflows		
	Reentrancy		
	Money-Giving Bug		
	Blackhole		
	Unauthorized Self-Destruct		
Basic Coding Bugs	Revert DoS		
Dasic Couling Dugs	Unchecked External Call		
	Gasless Send		
	Send Instead Of Transfer		
	Costly Loop		
	(Unsafe) Use Of Untrusted Libraries		
	(Unsafe) Use Of Predictable Variables		
	Transaction Ordering Dependence		
	Deprecated Uses		
Semantic Consistency Checks	Semantic Consistency Checks		
	Business Logics Review		
	Functionality Checks		
	Authentication Management		
	Access Control & Authorization		
	Oracle Security		
Advanced DeFi Scrutiny	Digital Asset Escrow		
, tavanieca Dei i Geraemi,	Kill-Switch Mechanism		
	Operation Trails & Event Generation		
	ERC20 Idiosyncrasies Handling		
	Frontend-Contract Integration		
	Deployment Consistency		
	Holistic Risk Management		
	Avoiding Use of Variadic Byte Array		
	Using Fixed Compiler Version		
Additional Recommendations	Making Visibility Level Explicit		
	Making Type Inference Explicit		
	Adhering To Function Declaration Strictly		
	Following Other Best Practices		

deploy contracts on our private testnet and run tests to confirm the findings. If necessary, we would additionally build a PoC to demonstrate the possibility of exploitation. The concrete list of check items is shown in Table 1.3.

In particular, we perform the audit according to the following procedure:

- <u>Basic Coding Bugs</u>: We first statically analyze given smart contracts with our proprietary static code analyzer for known coding bugs, and then manually verify (reject or confirm) all the issues found by our tool.
- <u>Semantic Consistency Checks</u>: We then manually check the logic of implemented smart contracts and compare with the description in the white paper.
- Advanced DeFi Scrutiny: We further review business logics, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.
- Additional Recommendations: We also provide additional suggestions regarding the coding and development of smart contracts from the perspective of proven programming practices.

To better describe each issue we identified, we categorize the findings with Common Weakness Enumeration (CWE-699) [9], which is a community-developed list of software weakness types to better delineate and organize weaknesses around concepts frequently encountered in software development. Though some categories used in CWE-699 may not be relevant in smart contracts, we use the CWE categories in Table 1.4 to classify our findings. Moreover, in case there is an issue that may affect an active protocol that has been deployed, the public version of this report may omit such issue, but will be amended with full details right after the affected protocol is upgraded with respective fixes.

1.4 Disclaimer

Note that this security audit is not designed to replace functional tests required before any software release, and does not give any warranties on finding all possible security issues of the given smart contract(s) or blockchain software, i.e., the evaluation result does not guarantee the nonexistence of any further findings of security issues. As one audit-based assessment cannot be considered comprehensive, we always recommend proceeding with several independent audits and a public bug bounty program to ensure the security of smart contract(s). Last but not least, this security audit should not be used as investment advice.

Table 1.4: Common Weakness Enumeration (CWE) Classifications Used in This Audit

Category	Summary		
Configuration	Weaknesses in this category are typically introduced during		
	the configuration of the software.		
Data Processing Issues	Weaknesses in this category are typically found in functional-		
	ity that processes data.		
Numeric Errors	Weaknesses in this category are related to improper calcula-		
	tion or conversion of numbers.		
Security Features	Weaknesses in this category are concerned with topics like		
	authentication, access control, confidentiality, cryptography,		
	and privilege management. (Software security is not security		
	software.)		
Time and State	Weaknesses in this category are related to the improper man-		
	agement of time and state in an environment that supports		
	simultaneous or near-simultaneous computation by multiple		
	systems, processes, or threads.		
Error Conditions,	Weaknesses in this category include weaknesses that occur if		
Return Values,	a function does not generate the correct return/status code,		
Status Codes	or if the application does not handle all possible return/status		
	codes that could be generated by a function.		
Resource Management	Weaknesses in this category are related to improper manage-		
	ment of system resources.		
Behavioral Issues	Weaknesses in this category are related to unexpected behav-		
	iors from code that an application uses.		
Business Logics	Weaknesses in this category identify some of the underlying		
	problems that commonly allow attackers to manipulate the		
	business logic of an application. Errors in business logic can		
	be devastating to an entire application.		
Initialization and Cleanup	Weaknesses in this category occur in behaviors that are used		
	for initialization and breakdown.		
Arguments and Parameters	Weaknesses in this category are related to improper use of		
	arguments or parameters within function calls.		
Expression Issues	Weaknesses in this category are related to incorrectly written		
	expressions within code.		
Coding Practices	Weaknesses in this category are related to coding practices		
	that are deemed unsafe and increase the chances that an ex-		
	ploitable vulnerability will be present in the application. They		
	may not directly introduce a vulnerability, but indicate the		
	product has not been carefully developed or maintained.		

2 | Findings

2.1 Summary

Here is a summary of our findings after analyzing the design and implementation of the Infusion protocol smart contracts. During the first phase of our audit, we study the smart contract source code and run our in-house static code analyzer through the codebase. The purpose here is to statically identify known coding bugs, and then manually verify (reject or confirm) issues reported by our tool. We further manually review business logics, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.

Severity	# of Findings		
Critical	0		
High	0		
Medium	2		
Low	3		
Informational	0		
Total	5		

We have so far identified a list of potential issues: some of them involve subtle corner cases that might not be previously thought of, while others refer to unusual interactions among multiple contracts. For each uncovered issue, we have therefore developed test cases for reasoning, reproduction, and/or verification. After further analysis and internal discussion, we determined a few issues of varying severities need to be brought up and paid more attention to, which are categorized in the above table. More information can be found in the next subsection, and the detailed discussions of each of them are in Section 3.

2.2 Key Findings

Overall, these smart contracts are well-designed and engineered, though the implementation can be improved by resolving the identified issues (shown in Table 2.1), including 2 medium-severity vulnerabilities and 3 low-severity vulnerabilities.

Table 2.1: Key Audit Findings

ID	Severity	Title	Category	Status
PVE-001	Medium	Improved Pair Initialization With Mini-	Time And State	Resolved
		mal Liquidity Enforcement		
PVE-002	Low	Incorrect getAmountOut() Logic in Infu-	Business Logic	Resolved
		sionLibrary		
PVE-003	Low	Improved _lockerFeesP Validation in	Coding Practices	Resolved
		Pair Construction		
PVE-004	Low	Inconsistent K Invariants Between Pair	Business Logic	Resolved
		And Router		
PVE-005	Medium	Trust Issue Of Admin Keys	Security Features	Mitigated

Beside the identified issues, we emphasize that for any user-facing applications and services, it is always important to develop necessary risk-control mechanisms and make contingency plans, which may need to be exercised before the mainnet deployment. The risk-control mechanisms should kick in at the very moment when the contracts are being deployed on mainnet. Please refer to Section 3 for details.

3 Detailed Results

3.1 Improved Pair Initialization With Minimal Liquidity Enforcement

• ID: PVE-001

Severity: MediumLikelihood: Low

• Impact: High

• Target: Pair

• Category: Time and State [8]

• CWE subcategory: CWE-663 [3]

Description

The Infusion protocol has the built-in curve $(x^3y + xy^3 = k)$ for efficient stable swaps. While examining the k calculation, we notice a possible denial-of-service issue that may allow the first stable swap LP to brick the (new) pair.

In the following, we show the implementation of the related $_{\tt k}()$ routine, which basically computes the k value from the above curve. However, it comes to our attention that the internal variable $_{\tt a}$ (line 665) may yield 0, which effectively computes the return value to be 0 and cascadingly meets the k invariant maintained in the $_{\tt swap}()$ routine. As a result, the first LP of a stable swap pair may abuse it to initialize the pair with k=0 and then empty the pool by resetting $_{\tt reserve0}$ and/or $_{\tt reserve1}$ to be 0. With that, any later LP may find infeasible to add new liquidity to the pair, hence effectively bricking the pair.

```
661
         function _k(uint x, uint y) internal view returns (uint) {
662
             if (stable) {
663
                 uint _x = (x * 1e18) / decimals0;
664
                 uint _y = (y * 1e18) / decimals1;
665
                 uint _a = (_x * _y) / 1e18;
                 uint _b = ((_x * _x) / 1e18 + (_y * _y) / 1e18);
666
                 return (_a * _b) / 1e18; // x3y+y3x >= k
667
668
             } else {
669
                 return x * y; // xy >= k
670
```

671 }

Listing 3.1: Pair::_k()

Recommendation To fix the above issue, there is a need to ensure the initial k upon the stable swap pair initialization will be larger than MINIMUM_LIQUIDITY*MINIMUM_LIQUIDITY. An example revision is shown as below:

```
460
        function mint(address to) external lock returns (uint liquidity) {
461
            (uint _reserve0, uint _reserve1) = (reserve0, reserve1);
462
            uint _balance0 = IERC20(token0).balanceOf(address(this));
463
            uint _balance1 = IERC20(token1).balanceOf(address(this));
464
            uint _amount0 = _balance0 - _reserve0;
465
            uint _amount1 = _balance1 - _reserve1;
467
            uint _totalSupply = totalSupply; // gas savings, must be defined here since
                totalSupply can update in _mintFee
468
            if (_totalSupply == 0) {
469
                liquidity = Math.sqrt(_k(_amount0, _amount1)) - MINIMUM_LIQUIDITY;
470
                 _mint(address(0), MINIMUM_LIQUIDITY); // permanently lock the first
                     MINIMUM_LIQUIDITY tokens
471
            } else {
472
                liquidity = Math.min(
473
                     (_amount0 * _totalSupply) / _reserve0,
474
                     (_amount1 * _totalSupply) / _reserve1
475
                );
476
477
            require(liquidity > 0, "ILM"); // Pair: INSUFFICIENT_LIQUIDITY_MINTED
478
            _mint(to, liquidity);
480
             _update(_balance0, _balance1, _reserve0, _reserve1);
481
            emit Mint(msg.sender, _amount0, _amount1);
482
```

Listing 3.2: Revised Pair::mint()

Status This issue has been fixed in the following commit: 3ece057.

3.2 Incorrect getAmountOut() Logic in InfusionLibrary

• ID: PVE-002

• Severity: Low

Likelihood: Low

• Impact: Low

• Target: InfusionLibrary

• Category: Business Logic [7]

• CWE subcategory: CWE-841 [4]

Description

To facilitate the user interaction, the Infusion protocol has an InfusionLibrary contract that provides a number of user convenience functions. While examining a specific getAmountOut() helper routine, we notice its implementation needs to be revised.

```
118
         function getAmountOut(
119
             uint amountln,
120
             address tokenIn,
121
             address tokenOut,
122
             bool stable
123
         ) external view returns (uint) {
             (uint dec0, uint dec1, uint r0, uint r1, bool st, address t0, ) = IPair(
124
125
                 router.pairFor(tokenIn, tokenOut, stable)
126
             ).metadata();
127
             return
128
                 ( getAmountOut(amountIn, tokenIn, r0, r1, t0, dec0, dec1, st) *
129
                     1e18) / amountIn;
130
```

Listing 3.3: InfusionLibrary :: getAmountOut()

To elaborate, we show above the implementation of this <code>getAmountOut()</code> routine. As the name indicates, this routine is used to compute the <code>tokenOut</code> amount after swapping the given <code>amountIn</code> of <code>tokenIn</code>. However, the computed amount is equal to the expected <code>tokenOut</code> amount by scaling it to the input amount of <code>1e18</code>, not <code>amountIn</code>.

Recommendation Revise the above routine to compute the intended tokenOut amount based on the given amountIn of tokenIn.

Status This issue has been fixed in the following commit: 3ece057.

3.3 Improved Validation on Protocol Parameters

• ID: PVE-003

Severity: Low

Likelihood: Low

• Impact: Low

• Target: Pair

• Category: Coding Practices [6]

• CWE subcategory: CWE-1126 [1]

Description

DeFi protocols typically have a number of system-wide parameters that can be dynamically configured on demand. The Infusion protocol is no exception. Specifically, if we examine the Pair contract, it has defined a number of protocol-wide risk parameters, such as _lockerFeesP and _feeDistributor. In the following, we show the corresponding constructor routine that initializes their values.

```
constructor() {
114
115
             factory = msg.sender;
116
117
                 address token0,
                 address _token1,
118
119
                 bool _stable,
120
                 uint
                       lockerFeesP,
121
                 address feeDistributor
122
             ) = PairFactory (msg. sender).getInitializable();
123
             require (
                  lockerFeesP === 0
124
125
                      ( lockerFeesP != 0 && feeDistributor != address(0)),
126
                 "MISS_FEE_DIST"
127
             );
128
             (token0, token1, stable, lockerFeesP, feeDistributor) = (
129
130
                  token1,
131
                  stable ,
132
                  lockerFeesP,
133
                 IFeeDistributor ( feeDistributor)
134
             );
135
             fees = address(new PairFees( token0, token1));
136
137
```

Listing 3.4: Pair :: constructor()

These parameters define various aspects of the protocol operation and maintenance and need to exercise extra care when configuring or updating them. Our analysis shows the update logic on these parameters can be improved by applying more rigorous sanity checks. Based on the current implementation, certain corner cases may lead to an undesirable consequence. For example, we can

improve the above constructor by further enforcing the following requirement: require(_lockerFeesP < LOCKER_FEES_SCALE).

Recommendation Validate any changes regarding these system-wide parameters to ensure they fall in an appropriate range. In the meantime, the above constructor in essence allows the first LP can decide _lockerFeesP, which may need to be revisited.

Status This issue has been fixed in the following commit: 3ece057.

3.4 Inconsistent K Invariants Between Pair And Router

ID: PVE-004

Severity: Low

• Likelihood: Low

• Impact: Low

• Target: Router

• Category: Business Logic [7]

• CWE subcategory: CWE-841 [4]

Description

As mentioned earlier, the Infusion protocol has a built-in curve, which is different from xy = k. While analyzing the k usages between Pair and Router routines, we notice certain inconsistency that needs to be resolved before deployment.

To elaborate, we show below the related quoteLiquidity() routine from the Router contract. While it always follows the traditional curve xy=k to compute the output amount, it does not take into account the new curve when the given pair is a stable swap one.

```
77
        function quoteLiquidity(
78
            uint amountA,
79
            uint reserveA.
80
            uint reserveB
81
        ) internal pure returns (uint amountB) {
82
            require(amountA > 0, "Router: INSUFFICIENT_AMOUNT");
83
            require(reserveA > 0 && reserveB > 0, "Router: INSUFFICIENT_LIQUIDITY");
84
            amountB = (amountA * reserveB) / reserveA;
85
```

Listing 3.5: Router::quoteLiquidity()

Recommendation Revise the above logic to properly compute the quote amount for both stable and volatile pairs.

Status This issue has been resolved as it intends to compute the normal quote for liquidity.

3.5 Trust Issue of Admin Keys

• ID: PVE-005

• Severity: Medium

Likelihood: Medium

• Impact: Medium

• Target: Multiple Contracts

• Category: Security Features [5]

• CWE subcategory: CWE-287 [2]

Description

In the Infusion protocol, there is a privileged owner account that plays a critical role in governing and regulating the system-wide operations (e.g., configuring various parameters and pausing swaps). It also has the privilege to control or govern the flow of assets managed by this protocol. Our analysis shows that the privileged account needs to be scrutinized. In the following, we examine the privileged account and the related privileged accesses in current contracts.

```
66
        function setPauser(address _pauser) external {
67
            require(msg.sender == pauser);
68
            pendingPauser = _pauser;
69
70
71
        function acceptPauser() external {
72
            require(msg.sender == pendingPauser);
73
            pauser = pendingPauser;
74
       }
75
76
       function setPause(bool _state) external {
77
            require(msg.sender == pauser);
78
            isPaused = _state;
79
       }
80
81
        function setFeeManager(address _feeManager) external {
82
            require(msg.sender == feeManager, "not fee manager");
83
            pendingFeeManager = _feeManager;
84
85
86
        function acceptFeeManager() external {
87
            require(msg.sender == pendingFeeManager, "not pending fee manager");
88
            feeManager = pendingFeeManager;
89
90
91
        function setFee(bool _stable, uint256 _fee) external {
92
            require(msg.sender == feeManager, "not fee manager");
93
            require(_fee <= MAX_FEE, "fee too high");</pre>
94
            require(_fee != 0, "fee must be nonzero");
95
            if (_stable) {
96
                stableFee = _fee;
97
            } else {
```

```
98 volatileFee = _fee;
99 }
100 }
```

Listing 3.6: Example Privileged Functions in PairFactory

Note that if the privileged owner account is a plain EOA account, this may be worrisome and pose counter-party risk to the exchange users. A multi-sig account could greatly alleviate this concern, though it is still far from perfect. Specifically, a better approach is to eliminate the administration key concern by transferring the role to a community-governed DAO. In the meantime, a timelock-based mechanism can also be considered as mitigation.

Recommendation Promptly transfer the privileged account to the intended DAO-like governance contract. All changed to privileged operations may need to be mediated with necessary timelocks. Eventually, activate the normal on-chain community-based governance life-cycle and ensure the intended trustless nature and high-quality distributed governance.

Status This issue has been mitigated as the team makes use of a multisig to act as the privileged owner.



4 Conclusion

In this audit, we have analyzed the design and implementation of the Infusion protocol, which is designed to allow low-cost, low-slippage trades on uncorrelated or tightly correlated assets. It is in essence a DEX that is built starting from Solidly/Velodrome with a unique AMM. The DEX is compatible with all the standard features as popularized by UniswapV2 with a number of novel improvements, including price oracles without upkeeps, a new curve $(x^3y + xy^3 = k)$ for efficient stable swaps. The current code base is well structured and neatly organized. Those identified issues are promptly confirmed and addressed.

Meanwhile, we need to emphasize that Solidity-based smart contracts as a whole are still in an early, but exciting stage of development. To improve this report, we greatly appreciate any constructive feedbacks or suggestions, on our methodology, audit findings, or potential gaps in scope/coverage.

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