



Offloading QUIC – AN IMPLEMENTATION GUIDE

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Agenda

- Review the challenges and solution to the proof of concept
- Three types of possible solutions
 - Agree on the possible solutions across multiple implementations.

Outcome ?

- Possible solutions are classified into:
 - Using meta-data to solve the problem.
 - Generic meta-data processing imposed by Linux. Limited by the OS upstreaming barriers.
 - Modification of the implementation guideline
 - With large number of stacks, can we agree on a few implementation rules.
 - Modification of the header.
 - Folks in this room do not like this solution 😊

Connection Id has 16 different sizes in the short header

Connection Id size varies on both Tx and Rx

Solve the problem with some meta data and implementation rule:

- Transmit solution
 - Augment the meta data with the CID size
- Receive solution
 - Header parsing will be programmed with a single size, for a server.

Connection Id has 16 different sizes in the short header

Protocol solution: Encode connection Id with a varint

```

0      1      2      3
0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+
|0|1|S|R|R|K|P P|
+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+
|          Destination Connection ID with a varint encoding (0..144) ...
+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+
|          Packet Number (8/16/24/32)          ...
+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+
|          Protected Payload (*)          ...
+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+

```

Optional connection Id in the short header

- The hardware has a classification schema to identify the crypto key.
- The current spec imposes that the hardware implements multiple schemas
 - Connection id match is higher priority over the outer 4 tuple
- Transmit solution
 - Meta data flowing with the packet identifies the size of CID.
- Receive solution
 - Packets receive at server will always have CID
 - Single size of CID

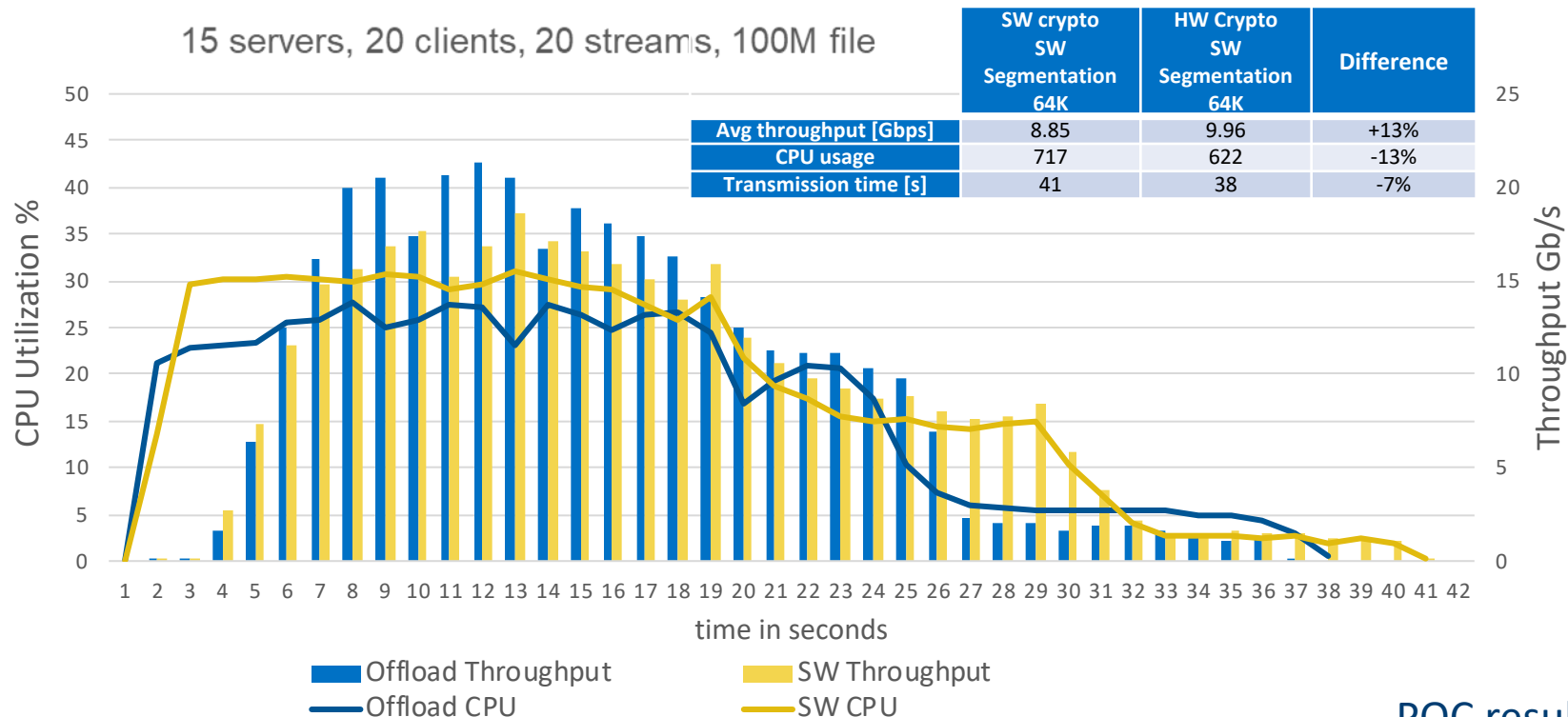
Experiment with Chromium

Chromium serving large number of clients.

- Saturate the network bandwidth with minimal clients.
- Each client does an HTTP request.
- Measurements on Tx side.

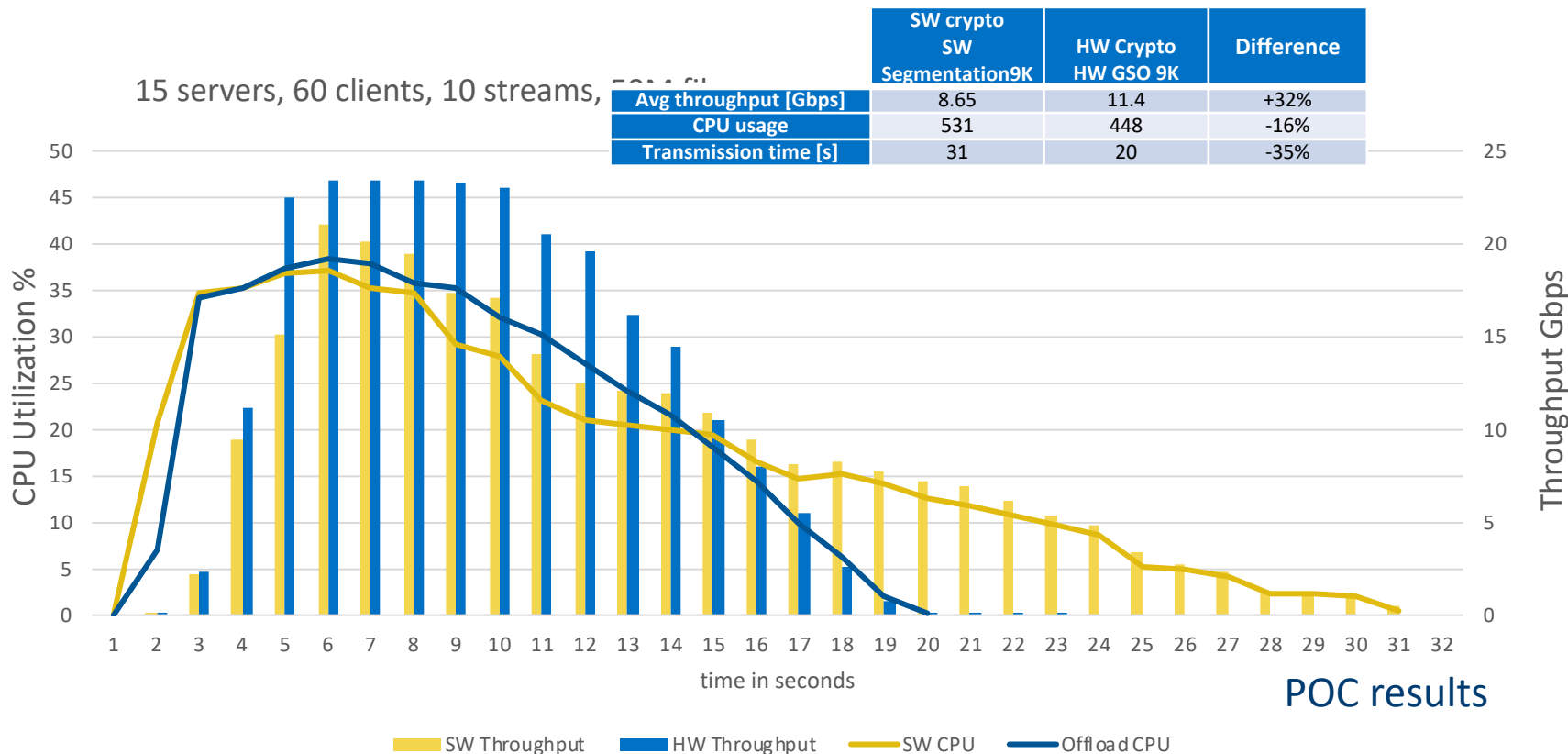
QUIC Crypto Offload Performance

15 servers, 20 clients, 20 streams, 100M file



POC results

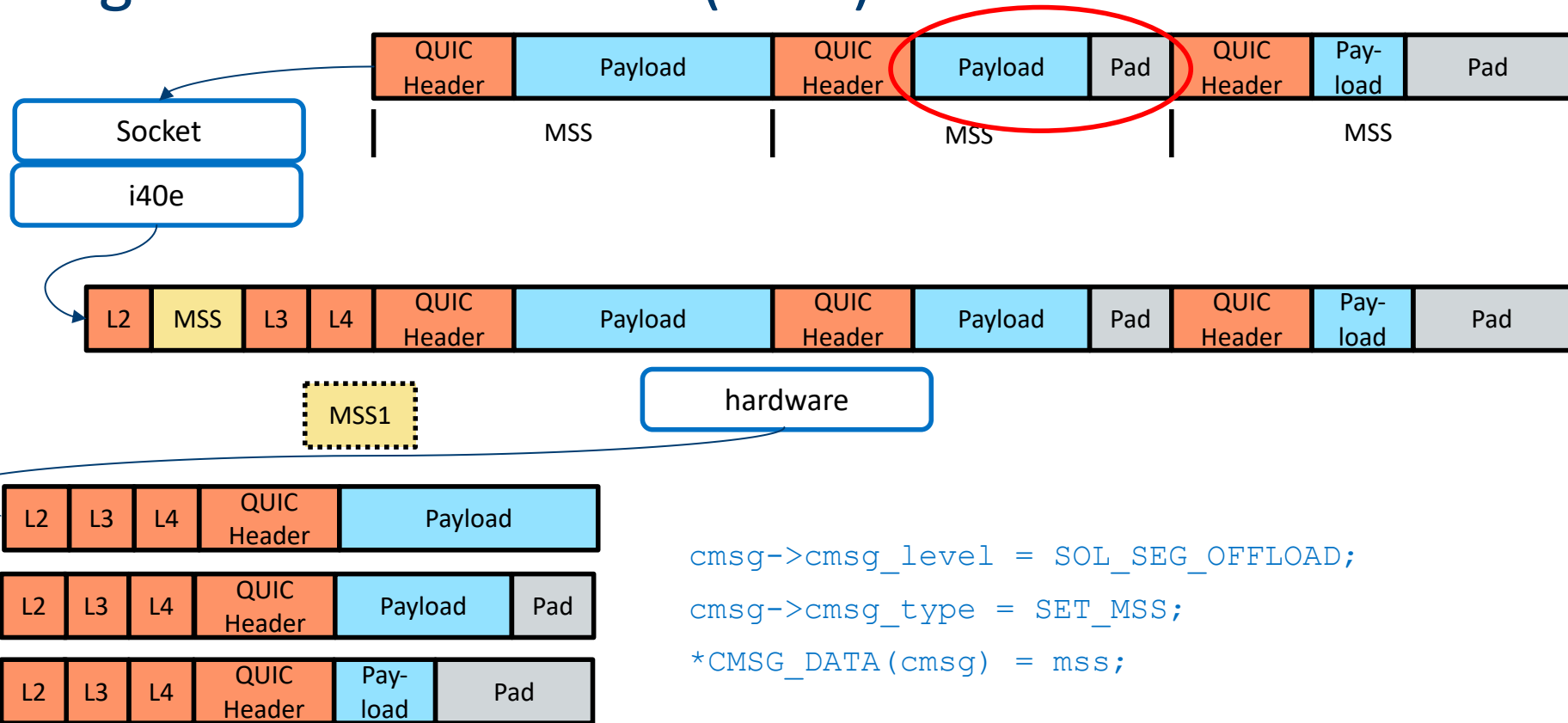
Now let's put it all together



Programming the HW – User space to Driver

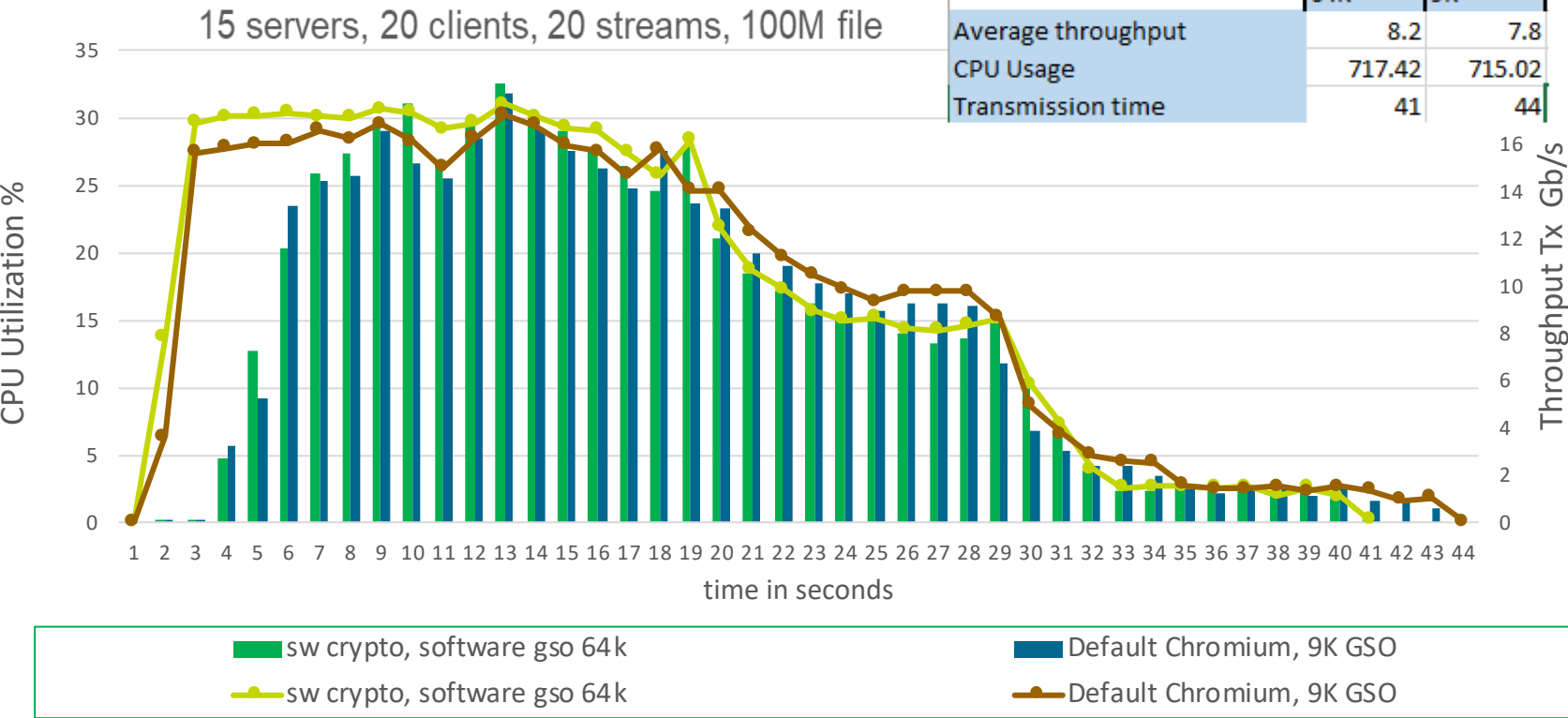
```
struct ulp_offload_devops {  
setsockopt(sk, SOL_OFFLOAD, INIT_DEVICE, &init, ...)    ---->    int (*ulp_offload_init)(struct net_device *netdev, ...);  
  
setsockopt(sk, SOL_OFFLOAD, ADD_SA_{TX|RX}, &addsa, ...) ---->    int (*ulp_add_sa)(struct net_device *netdev, ...);  
  
setsockopt(sk, SOL_OFFLOAD, UPDATE_SA_{TX|RX}, &upsa, ...) ---->    int (*ulp_update_sa)(struct net_device *netdev, ...);  
  
setsockopt(sk, SOL_OFFLOAD, DEL_SA_{TX|RX}, &delsa, ...) ---->    int (*ulp_del_sa)(struct net_device *netdev, ...);  
  
getsockopt(sk, SOL_OFFLOAD, GET_CAPS, &capabilities, ...) ---->    int (*ulp_get_caps)(struct net_device *netdev, ...);  
  
getsockopt(sk, SOL_OFFLOAD, OFFLOAD_OK, &status, ...)  ---->    bool (*ulp_offload_ok)(struct net_device *netdev);  
};
```

Segmentation Offload (USO)

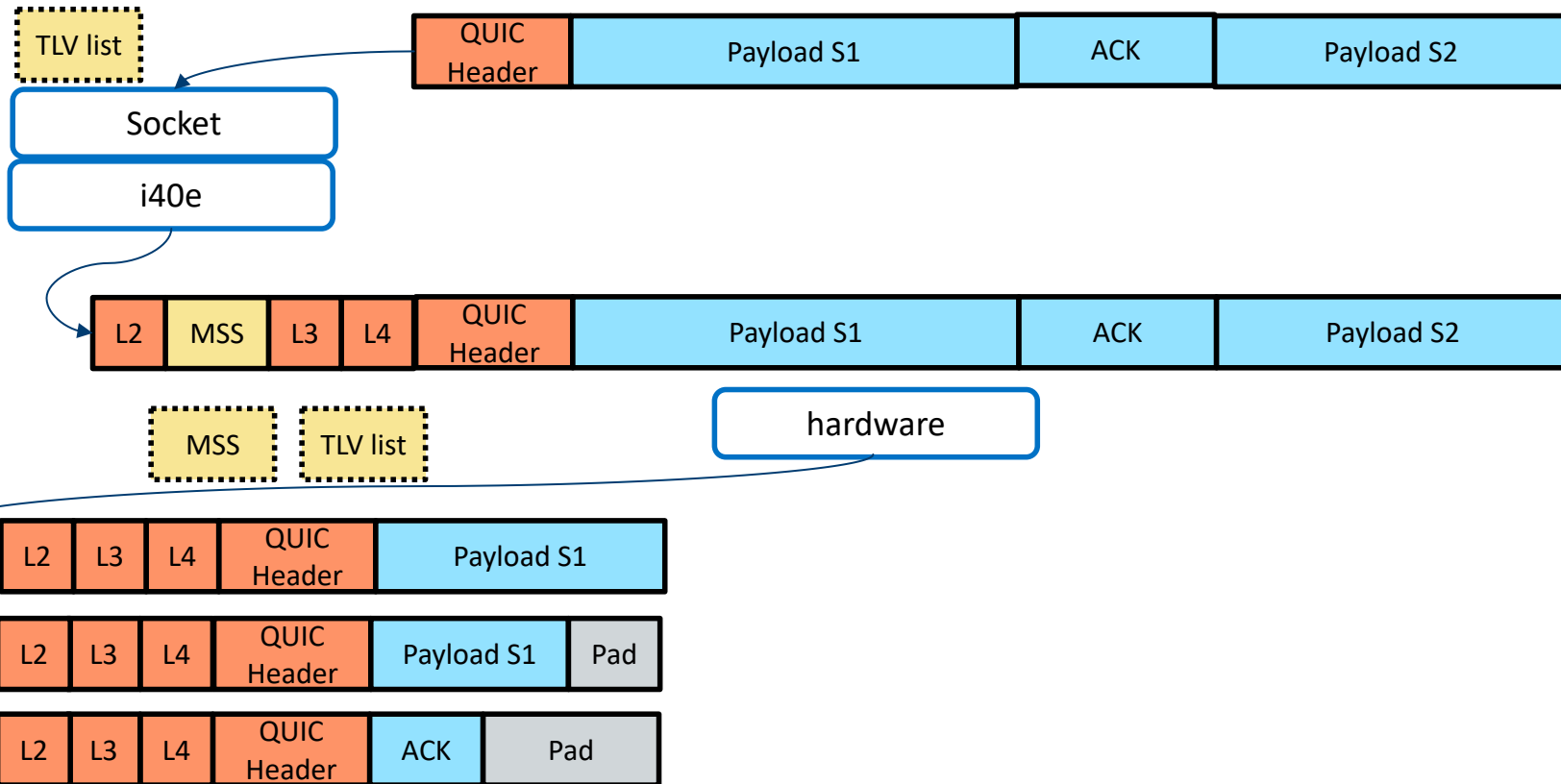


```
cmmsg->cmmsg_level = SOL_SEG_OFFLOAD;  
cmmsg->cmmsg_type = SET_MSS;  
*CMMSG_DATA(cmmsg) = mss;
```

Comparing 9K GSO with 64K GSO



Segmentation Offload like TCP



Segmentation Options

- Option 1: Only segment a single stream in a GSO offload
- Option 2: Limit the segmentation to only use ACK and stream frames
 - Only a single ACK is present at the start in the offloaded buffer
- TLV meta data to present the object list
 - Kernel folks pushed back on sending a large meta – data
- Having a length in the ACK would help
 - Small surgical change to simplify reduce the TLV style meta data

More Discussion?

Detailed discussion today at 4pm @ Congress Hall 3

Ingress Metadata

- Agent passes status to driver: authentication status, decryption status, protocol errors
- How does the driver communicate that status to the stack per packet?
 - Store status in control buffer in skb
 - Create cmsg header once skb hits socket layer
 - Stack uses recvmsg and extracts metadata from cmsg header

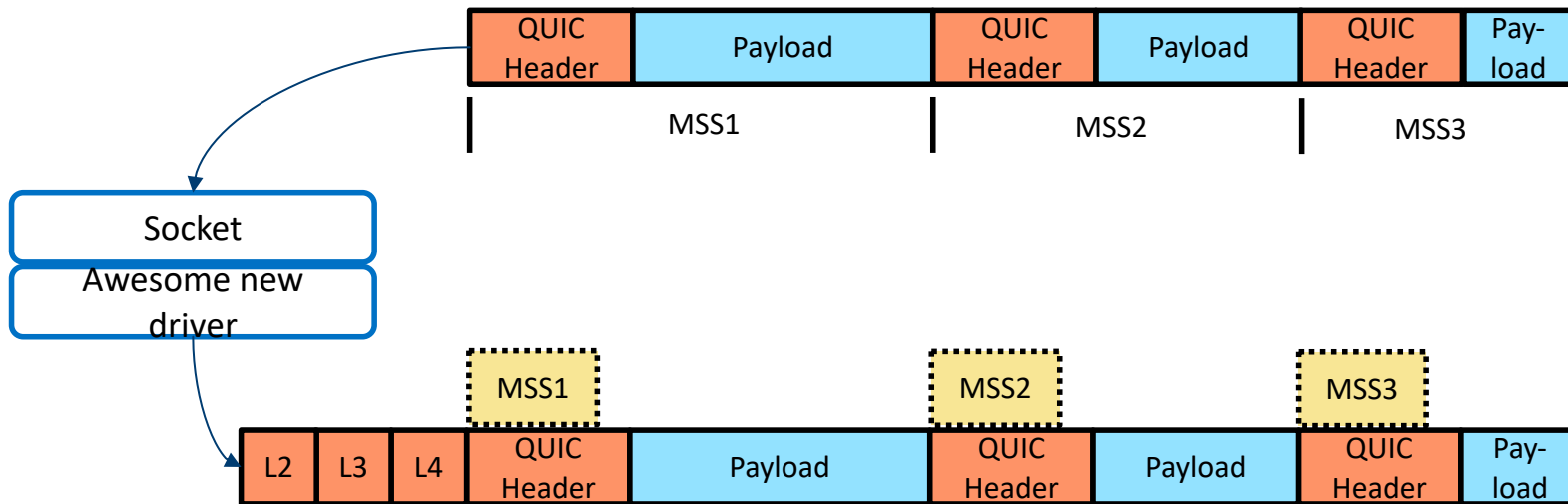


Flexible Interfaces for Flexible HW

- Abstract nonce from HW, send multiple MSS
 - Sent per segment via cmsg headers from the user space stack?

```
int send_quic(socket,  
              quic_header = {flags, CID},  
              nonces = [nonce1, nonce2, nonce3],  
              quic_data = [  
                  {QUIC_segment_1, mss1},  
                  {QUIC_segment_2, mss2},  
                  {QUIC_segment_3, mss3}],  
              npackets = 3);
```

Segmentation Offload – Future Interface



```
for (i = 0; i < numpkts; i++) {  
    cmsg->cmsg_level = SOL_SEG_OFFLOAD;  
    cmsg->cmsg_type = SET_MSS;  
    *CMSG_DATA(cmsg) = mss[i];  
}
```

Takeaways

- Offload saves ~**16%** CPU Usage, improves throughput by ~**32%** in certain test cases
- Segmentation interface and crypto interface are independent
- Crypto offload is impossible without an interface to get crypto parameters from user space to hardware
- A generic interface can enable crypto and segmentation offloads for other protocols
- Opens
 - What is the best way to bind a socket to an interface?
 - Is there a better way to do ingress metadata?
 - Are there other protocols that could make use of such an interface?

THANK YOU!

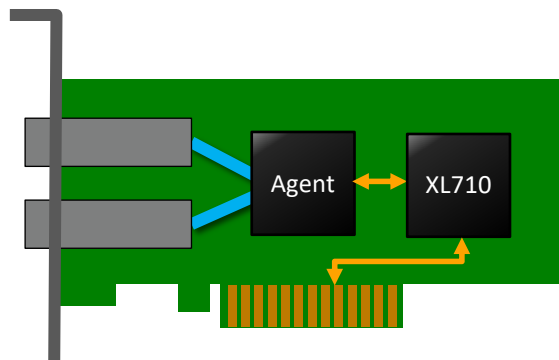
BACKUP

Programming the HW – User space to Driver

```
struct ulp_offload_devops {  
    setsockopt(sk, SOL_OFFLOAD, INIT_DEVICE, &init, ...)    ---->    int (*ulp_ofload_init)(struct net_device *netdev,  
                                                                    struct ulp_offload_init *init);  
    setsockopt(sk, SOL_OFFLOAD, ADD_SA_{TX|RX}, &addsa, ...) ---->    int (*ulp_add_sa)(struct net_device *netdev,  
                                                                    struct ulp_sa_context *sa,  
                                                                    struct ulp_crypto_info *crypto_info);  
    setsockopt(sk, SOL_OFFLOAD, UPDATE_SA_{TX|RX}, &upsa, ...) ---->    int (*ulp_update_sa)(struct net_device *netdev,  
                                                                    struct ulp_sa_context *sa,  
                                                                    struct ulp_sa_update *update);  
    setsockopt(sk, SOL_OFFLOAD, DEL_SA_{TX|RX}, &delsa, ...) ---->    int (*ulp_del_sa)(struct net_device *netdev,  
                                                                    struct ulp_sa_context *sa);  
    getsockopt(sk, SOL_OFFLOAD, GET_CAPS, &capabilities, ...) ---->    int (*ulp_get_caps)(struct net_device *netdev,  
                                                                    struct ulp_offload_caps *caps);  
    getsockopt(sk, SOL_OFFLOAD, OFFLOAD_OK, &status, ...) ---->    bool (*ulp_offload_ok)(struct net_device *netdev);  
};
```

Connectivity with our QUIC Agent

- No separate control plane for Configuration and Metadata
- All control data has to go through the MAC to get to the Agent
- Use one L2 tag to denote Control packets
- Different L2 tag to insert Metadata into a packet



- See https://www.netdevconf.org/2.2/slides/klassert_ipsec_workshop03.pdf for more info about agent connectivity

SW segmentation – 9K vs 64K

