

((MARKS)) (1/2/3...)	1
((QUESTION))	Two main measures for the efficiency of an algorithm are
((OPTION_A))	Processor and memory
((OPTION_B))	Complexity and capacity
((OPTION_C))	Time and space
((OPTION_D))	Data and space
((CORRECT_CHOICE))) (A/B/C/D)	C
((EXPLANATION)) (OPTIONAL)	

((MARKS)) (1/2/3...)	1
((QUESTION))	Which of the following case does not exist in complexity theory
((OPTION_A))	Best case
((OPTION_B))	Worst case
((OPTION_C))	Average case
((OPTION_D))	Null case
((CORRECT_CHOICE))) (A/B/C/D)	D
((EXPLANATION)) (OPTIONAL)	

((MARKS)) (1/2/3...)	1
((QUESTION))	The Worst case occur in linear search algorithm when
((OPTION_A))	Item is somewhere in the middle of the array
((OPTION_B))	Item is not in the array at all
((OPTION_C))	Item is the last element in the array
((OPTION_D))	Item is the last element in the array or is not there at all

((CORRECT_CHOICE)) (A/B/C/D)	D
((EXPLANATION)) (OPTIONAL)	.

((MARKS)) (1/2/3...)	1
((QUESTION))	The complexity of merge sort algorithm is
((OPTION_A))	$O(n)$
((OPTION_B))	$O(\log n)$
((OPTION_C))	$O(n^2)$
((OPTION_D))	$O(n \log n)$
((CORRECT_CHOICE)) (A/B/C/D)	D
((EXPLANATION)) (OPTIONAL)	

((MARKS)) (1/2/3...)	1
((QUESTION))	The input to a merge sort is 6,5,4,3,2,1 and the same input is applied to quick sort then which is the best algorithm in this case
((OPTION_A))	Merge sort
((OPTION_B))	Quick sort
((OPTION_C))	Both have same time complexity in this case as they have same running time
((OPTION_D))	Cannot be decided
((CORRECT_CHOICE)) (A/B/C/D)	A
((EXPLANATION)) (OPTIONAL)	

((MARKS)) (1/2/3...)	1
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((QUESTION))	If there exists two functions $f(n)$ and $g(n)$. The constant $c > 0$ and there exists an integer constant $n_0 \geq 1$. If $f(n) \leq c * g(n)$ for every integer $n \geq n_0$ then we say that _____
((OPTION_A))	$f(n) = O(g(n))$
((OPTION_B))	$f(n) = \Theta(g(n))$
((OPTION_C))	$f(n) = \Omega(g(n))$
$f(n) = \Theta(g(n))$	$f(n) = o(g(n))$
((CORRECT_CHOICE))) (A/B/C/D)	A
((EXPLANATION)) (OPTIONAL)	Basic definition of big oh notation

((MARKS)) (1/2/3...)	1
((QUESTION))	In practice _____ is used to define tight upper bound on growth of function $f(n)$
((OPTION_A))	Big oh
((OPTION_B))	Big omega
((OPTION_C))	Big theta
((OPTION_D))	None of these
((CORRECT_CHOICE))) (A/B/C/D)	A
((EXPLANATION)) (OPTIONAL)	The definition of big oh notation is $f(n) \leq c * g(n)$ which defines the upper bound on growth of the function $f(n)$

((MARKS)) (1/2/3...)	1
((QUESTION))	Examples of $O(1)$ are _____
((OPTION_A))	Multiplying two numbers
((OPTION_B))	Assigning some value to a variable
((OPTION_C))	Displaying some integer on console
((OPTION_D))	All of the above

((CORRECT_CHOICE)) (A/B/C/D)	D
((EXPLANATION)) (OPTIONAL)	All these operations are computed by single line expression evaluation

((MARKS)) (1/2/3...)	1
((QUESTION))	Examples of $O(n^2)$ algorithms are
((OPTION_A))	Adding two matrices
((OPTION_B))	Finding transpose of a matrix
((OPTION_C))	Initializing all elements of the matrix by 0
((OPTION_D))	All of the above
((CORRECT_CHOICE)) (A/B/C/D)	D
((EXPLANATION)) (OPTIONAL)	Within two for loops(nested), all these operations are performed.

((MARKS)) (1/2/3...)	1
((QUESTION))	Choose the correct time complexity of following code__ <pre>while(n>0) { n=n/10 }</pre>
((OPTION_A))	$O(1)$
((OPTION_B))	$O(n)$
((OPTION_C))	$O(\log n)$
((OPTION_D))	$O(n^2)$
((CORRECT_CHOICE)) (A/B/C/D)	C
((EXPLANATION)) (OPTIONAL)	

((MARKS)) (1/2/3...)	1
((QUESTION))	The time complexity of binary search is_____
((OPTION_A))	$O(n)$
((OPTION_B))	$O(\log n)$
((OPTION_C))	$O(n \log n)$
((OPTION_D))	$O(n^2)$
((CORRECT_CHOICE)) (A/B/C/D)	B
((EXPLANATION)) (OPTIONAL)	The list is divided at the mid and then the element is searched in either left half or right half.

((MARKS)) (1/2/3...)	1
((QUESTION))	Consider recurrence relation as $T(0)=c1$ $T(n)=T(n-1)+c2$ This can be expressed as
((OPTION_A))	$O(n)$
((OPTION_B))	$O(\log n)$
((OPTION_C))	$O(n \log n)$
((OPTION_D))	$O(n^2)$
((CORRECT_CHOICE)) (A/B/C/D)	A
((EXPLANATION)) (OPTIONAL)	$T(n)=T(n-1)+c2$ $=T(n-2)+2c2$ $=T(n-3)+3c2$ $=T(n-k)+kc2$ If $k=n$ then $T(n)=c1+nc2$ Hence, $T(n)=O(n)$

((MARKS)) (1/2/3...)	1
((QUESTION))	Consider recurrence relation as $T(0)=c1$ and $T(1)=c2$ $T(n)=T(n/2)+c3$ This can be expressed as
((OPTION_A))	$O(n)$
((OPTION_B))	$O(\log n)$
((OPTION_C))	$O(n \log n)$
((OPTION_D))	$O(n^2)$
((CORRECT_CHOICE))) (A/B/C/D)	B
((EXPLANATION)) (OPTIONAL)	

((MARKS)) (1/2/3...)	1
((QUESTION))	Following is the method of solving recurrence relation
((OPTION_A))	Greedy method
((OPTION_B))	Backtracking
((OPTION_C))	Forward substitution method
((OPTION_D))	Divide and Conquer method
((CORRECT_CHOICE))) (A/B/C/D)	C
((EXPLANATION)) (OPTIONAL)	

((MARKS)) (1/2/3...)	1
((QUESTION))	The recurrence relation for factorial function is of the form _____
((OPTION_A))	$T(n)=T(n-1)+c$
((OPTION_B))	$T(n)=T(n-1)+T(n-2)+c$

((OPTION_C))	$T(n/2)+c$
((OPTION_D))	None of these
((CORRECT_CHOICE)) (A/B/C/D)	A
((EXPLANATION)) (OPTIONAL)	<p>The factorial function is as follows-</p> <pre>fact(n) { if n=1 return 1 else return n * fact(n-1) }</pre>

((MARKS)) (1/2/3...)	1
((QUESTION))	The recurrence relation for fibonacci function is of the form _____
((OPTION_A))	$T(n)=T(n-1)+c$
((OPTION_B))	$T(n)=T(n-1)+T(n-2)+c$
((OPTION_C))	$T(n/2)+c$
((OPTION_D))	None of these
((CORRECT_CHOICE)) (A/B/C/D)	B
((EXPLANATION)) (OPTIONAL)	<p>The fibonacci function is as follows-</p> <pre>fibb(n) { if n == 0 return 0 if n == 1 return 1 else return (fibb(n-1) + fibb(n-2)) }</pre>

((MARKS)) (1/2/3...)	1
((QUESTION))	<p>The frequency count of following code is ____</p> <pre>for(i=0;i<m;i++) {</pre>

	<pre> for(j=0;i<n;j++) { C[i][j]=a[i][j]+b[i][j]; } </pre>
((OPTION_A))	$m + mn + mn$
((OPTION_B))	$m + n + mn$
((OPTION_C))	$m + n^2 + mn$
((OPTION_D))	$(m+1) + m(n+1) + mn$
((CORRECT_CHOICE)) (A/B/C/D)	D
((EXPLANATION)) (OPTIONAL)	

((MARKS)) (1/2/3...)	1
((QUESTION))	Consider $T(n)=15n^3 + n^2 + 4$. Select the correct statement
((OPTION_A))	$T(n)=O(n^4)$
((OPTION_B))	$T(n)=\Omega \Omega (n^3)$
((OPTION_C))	$T(n)=\Omega \Omega (n^2)$
((OPTION_D))	All of the above
((CORRECT_CHOICE)) (A/B/C/D)	D
((EXPLANATION)) (OPTIONAL)	

((MARKS)) (1/2/3...)	1
((QUESTION))	<p>Give the frequency count of 3rd Statement</p> <pre> for(i=1;i<=n;i++) for(j=1;j<=i;j++) x=x+1; </pre>

((OPTION_A))	$\frac{1}{2}(n^2+n)$
((OPTION_B))	$\frac{1}{2}(n^2+3n)$
((OPTION_C))	n^2
((OPTION_D))	$(n+1)^2$
((CORRECT_CHOICE)) (A/B/C/D)	A
((EXPLANATION)) (OPTIONAL)	

((MARKS)) (1/2/3...)	1
((QUESTION))	There are four algorithms for solving a problem. Their time complexities are $O(n)$, $O(n^2)$, $O(\log n)$ and $O(n \log n)$. Which is the best algorithm?
((OPTION_A))	$O(n)$
((OPTION_B))	$O(n^2)$
((OPTION_C))	$O(\log n)$
((OPTION_D))	$O(n \log n)$
((CORRECT_CHOICE)) (A/B/C/D)	C
((EXPLANATION)) (OPTIONAL)	

((MARKS)) (1/2/3...)	1
((QUESTION))	The order of the recurrence relation $a_r - 7a_{r-1} + 10a_{r-2} = 0$ is _____.
((OPTION_A))	3
((OPTION_B))	2
((OPTION_C))	1
((OPTION_D))	B
((CORRECT_CHOICE)) (A/B/C/D)	D

((EXPLANATION)) (OPTIONAL)	
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((MARKS)) (1/2/3...)	1
((QUESTION))	Characteristic roots of the recurrence relation $a_r - 2a_{r-1} + a_{r-2} = 0$ are _____
((OPTION_A))	1, -1
((OPTION_B))	-1, -1
((OPTION_C))	1, 1
((OPTION_D))	None of these
((CORRECT_CHOICE)) (A/B/C/D)	C
((EXPLANATION)) (OPTIONAL)	

((MARKS)) (1/2/3...)	1
((QUESTION))	Charactristic polynomial of the recurrence relation $b_n = -3b_{n-1} - b_{n-2}$ is _____.
((OPTION_A))	$Z^2 - 3Z - 2 = 0$
((OPTION_B))	$Z^2 + 3Z - 2 = 0$
((OPTION_C))	$Z^2 + 3Z + 2 = 0$
((OPTION_D))	None of these
((CORRECT_CHOICE)) (A/B/C/D)	C
((EXPLANATION)) (OPTIONAL)	

((MARKS)) (1/2/3...)	1
((QUESTION))	The general solution of the recurrence relation $a_r - 2a_{r-1} = 0$ is _____.
((OPTION_A))	$a^r = c1(-2)^r$

((OPTION_B))	$a^r = c2(2)^r$
((OPTION_C))	$a^r = c1(1)^r$
((OPTION_D))	None of these
((CORRECT_CHOICE)) (A/B/C/D)	B
((EXPLANATION)) (OPTIONAL)	

((MARKS)) (1/2/3...)	1
((QUESTION))	Consider the recurrence relation, $a_n = a_{n-1} + 2a_{n-2}$ with $a_9 = 3$ and $a_{10} = 5$. Find a_7 .
((OPTION_A))	1
((OPTION_B))	3
((OPTION_C))	5
((OPTION_D))	None
((CORRECT_CHOICE)) (A/B/C/D)	A
((EXPLANATION)) (OPTIONAL)	

((MARKS)) (1/2/3...)	1
((QUESTION))	Charactristic polynomial of the recurrence relation $a_{r+2} - a_{r-2} = 0$ is _____.
((OPTION_A))	$Z - 1 = 0$
((OPTION_B))	$Z^2 - 1 = 0$
((OPTION_C))	$(Z - 1)^2 = 0$
((OPTION_D))	None
((CORRECT_CHOICE)) (A/B/C/D)	D

((EXPLANATION)) ((OPTIONAL))	<p>Given homogeneous recurrence relation can be written as</p> $a_{r+2} + 0a_{r+1} + 0a_r + 0a_{r-1} - a_{r-2} = 0$ <p>Order of this recurrence relation is 4.</p> <p>Hence characteristic polynomial is $Z^4 - 1 = 0$</p>
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((MARKS)) (1/2/3...)	1
((QUESTION))	Which of the following is not a homogeneous linear recurrence relation
((OPTION_A))	$a_{r+2} - a_{r-2} = 0$
((OPTION_B))	$a_r = a_{r-1} + a_{r-2}$
((OPTION_C))	$a_r - 2a_{r-1} = -a_{r-2}$
((OPTION_D))	$a_{r+3} + 6a_{r+2} \cdot a_{r+1} - 4a_r = 0$
((CORRECT_CHOICE))) (A/B/C/D)	D
((EXPLANATION)) ((OPTIONAL))	

((MARKS)) (1/2/3...)	1
((QUESTION))	If 4 and -1 are the characteristic roots of the recurrence relation then its homogeneous solution becomes _____
((OPTION_A))	$a^r = c_1(-1)^r + c_2(4)^r$
((OPTION_B))	$a^r = c_0(-1)^r + c_2$
((OPTION_C))	$a^r = (c_1 + c_2 \cdot r)(-1)^r$
((OPTION_D))	None
((CORRECT_CHOICE))) (A/B/C/D)	A
((EXPLANATION)) ((OPTIONAL))	

((MARKS)) (1/2/3...)	1
((QUESTION))	A recurrence relation of degree 1 is called _____.
((OPTION_A))	Linear
((OPTION_B))	Homogeneous
((OPTION_C))	Quadratic
((OPTION_D))	None
((CORRECT_CHOICE)) (A/B/C/D)	A
((EXPLANATION)) (OPTIONAL)	

((MARKS)) (1/2/3...)	1
((QUESTION))	The generating function for the sequence 1, a, a^2 , a^3 , is _____
((OPTION_A))	$1/(1-z)$
((OPTION_B))	$1/(1-az)$
((OPTION_C))	$1/(1+az)$
((OPTION_D))	None
((CORRECT_CHOICE)) (A/B/C/D)	B
((EXPLANATION)) (OPTIONAL)	