# ZooKeeper: Wait-free coordination for Internet-scale systems

Paper Presentation for EECS 591

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# "ZooKeeper"?





## ZooKeeper





## ZooKeeper

- Developed at Yahoo! Research
- Started as sub-project of Hadoop, now a top-level Apache project
- Development is driven by application needs



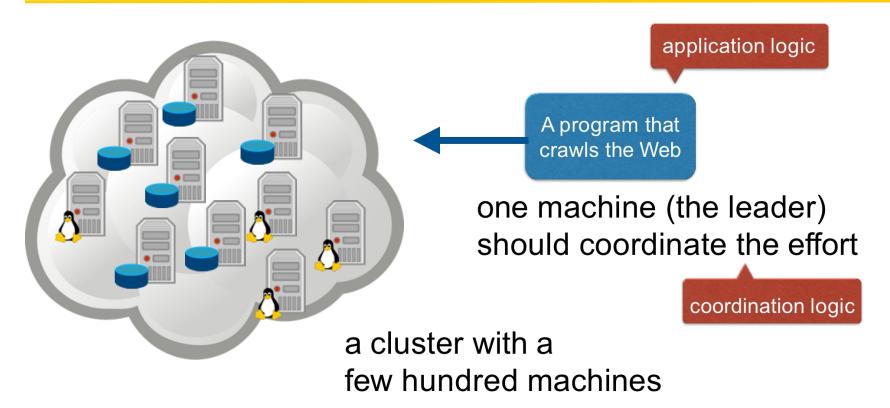
### **Motivation**

- In the past: a single program running on a single computer with a single CPU
- Today: distributed applications consist of independent programs running on a changing set of computers

Developers have to deal with **coordination logic** and **application logic** at the same time!

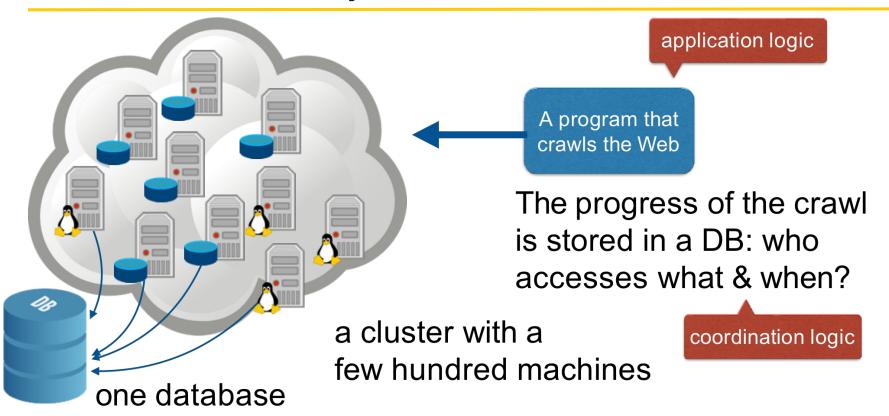


### Question: How do you elect the leader?



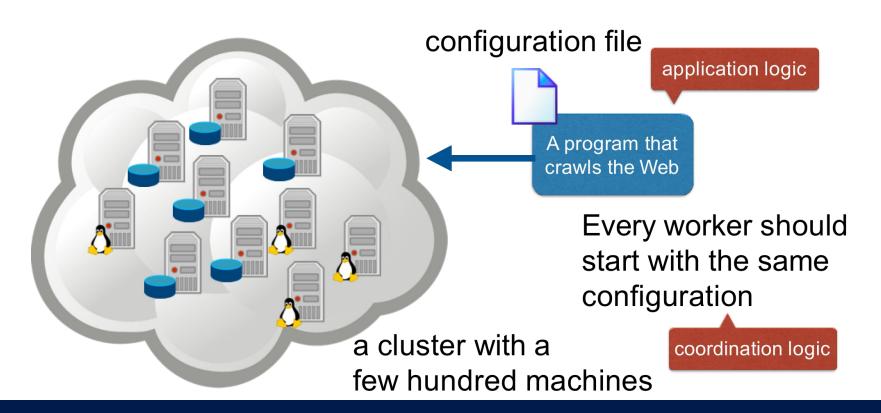


### Question: How do you lock a service?





### Question: How can the configuration be distributed?





## Solution approaches

- Be specific: develop a particular service for each coordination task
  - Leading election
  - Locking service
  - Configuration
- Be general: provide an API to make many services possible

ZooKeeper: exposing an API that enables application developers to implement their own primitives



## **ZooKeeper Terminology**

- Client: user of the ZooKeeper service
- Server: process providing the ZooKeeper service
- znode: in-memory data node in ZooKeeper, organised in a hierarchical name space (the data tree)
- Update/write: any operation which modifies the state of the data tree
- Clients establish a session when connecting to ZooKeeper



## ZooKeeper's data model

- znodes are organized in a hierarchical name space
- use the standard UNIX notation for file system paths
- znodes are not designed for general data storage. Instead,
   they map to abstractions of the client application

/app1/p 1

### **Z**node

Regular: Create and delete by client explicitly

 Ephemeral: Either delete them explicitly, or let the system remove them automatically when the session that creates them terminates (deliberately or due to a failure).



### **Znode**

#### SEQUENTIAL flag:

- monotonically increasing counter appended to a znode's path
- counter value of a new znode under a parent is always larger than value of existing children

#### WATCH flag:

- allow clients to receive timely notifications of changes without requiring polling
- Watches indicate that a change has happened, but do not provide the change



## ZooKeeper API

- String create (path, data, flags)
  - creates a znode with path name path, stores data in it and sets flags (ephemeral, sequential)
- void delete(path, version)
  - deletes the anode if it is at the expected version
- Stat exists(path, watch)
  - watch flag enables the client to set a watch on the znode
- (data, Stat) getData(path, watch)
  - returns the data and meta-data of the znode
- Stat setData(path, data, version)
  - writes data if the version number is the current version of the znode
- String[] getChildren(path, watch)



## **ZooKeeper API**

 ZooKeeper does not use handles to access znodes (no open() or close() methods)

All methods have both a synchronous and an asynchronous version



## **ZooKeeper Guarantees**

- Linearizable writes: all requests that update the state of ZooKeeper are serializable and respect precedence
- FIFO client order: all requests from a given client are executed in the order that they were sent by the client

A-linearizability (asynchronous linearizability)



### Think!

#### When electing a new leader:

- As the new leader starts making changes, we do not want other processes to start using the configuration that is being changed
- If the new leader dies before the configuration has been fully updated, we do not want the processes to use this partial configuration

deleting *ready* znode -> updating the various configuration znodes -> creating *ready* znode



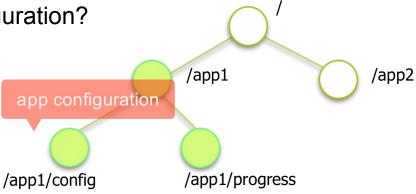
## **Example of primitives: Configuration**

#### **Questions:**

- How does a **new** worker query for a configuration?
- 2. How does an administrator **change** the configuration **on the fly**?
- 3. How do the workers read the **new** configuration?

#### [configuration stored in/app1/config]

- getData(/app1/config,true)
- 2. setData(/app1/config/config\_data,-1)
   [notify watching clients]
- 3. getData(/app1/config,true)





## **Example of primitives: Group Membership**

#### Questions:

 How can all workers (slaves) of an application register themselves?

2. How can a process find out about all active workers of an application?

[a znode is designated to store workers]

- 1. create(/app1/workers/worker, data, EPHEMERAL)
- 2. getChildren(/app1/workers, true)

/app1/workers/worker1

/app1/workers/worker2

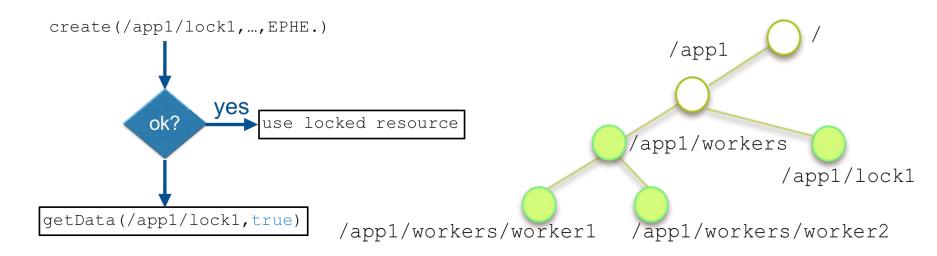
/app1/workers

/app1



## **Example of primitives: Simple Locks**

Questions: How can all workers of an application use a single resource through a lock?





### **Example of primitives: Simple Locks without Herd Effect**

#### Lock 1 n = create(1 + "/lock-", EPHEMERAL|SEQUENTIAL) 2 C = getChildren(l, false) 3 if n is lowest znode in C, exit 4 p = znode in C ordered just before n 5 if exists(p, true) wait for watch event 6 goto 2 /app1 Unlock /app1/locks 1 delete(n) /app1/locks/lock 1 /app1/locks/lock 2

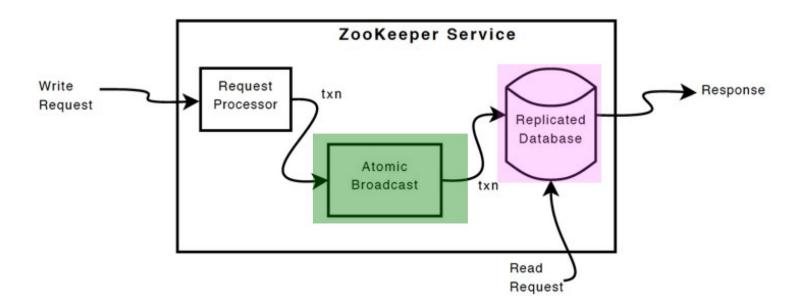


### **Example of primitives: Read/Write Locks**

```
Write Lock
  n = create(1 + "/write-", EPHEMERAL|SEQUENTIAL)
  C = getChildren(l, false)
3 if n is lowest znode in C, exit
4 p = znode in C ordered just before n
  if exists (p, true) wait for event
6 goto 2
Read Lock
1 n = create(1 + "/read-", EPHEMERAL|SEQUENTIAL)
                                                                      /app1
  C = getChildren(l, false)
3 if no write znodes lower than n in C, exit
                                                              /app1/locks
4 p = write znode in C ordered just before n
5 if exists(p, true) wait for event
  goto 3
                                       /app1/locks/lock 1
                                                                 /app1/locks/lock 2
```



# A little about implementation

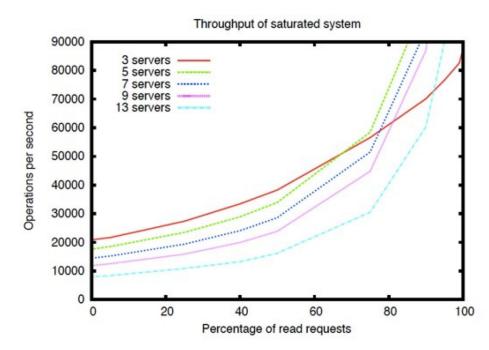




## **Throughput**

Setup: 250 clients, each client has at least 100 outstanding requests (read/write

of 1K data)





# **Summary**

- Wait-free
- Event ordering
- Watch mechanism



## **Thoughts**

- a minimalist and flexible coordination system
- scales well with increase in read operations, but does not with increase in write operations
- punts the ball to the clients



## **Questions?**



### Reference

ZooKeeper: Wait-free coordination for Internet scale system, Hunt al., 2010

https://people.eecs.berkeley.edu/~istoica/classes/cs294/15/notes/18-zookeeper.pdf

http://deptinfo.unice.fr/twiki/pub/Minfo/DistributedAlgo/TUDelft-coordination\_zookeeper.pdf

