

# **FIT2004**

## **Algorithms and Data Structures**

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Referencing materials by  
Nathan Companeze, Aamir Cheema, Arun Konagurthu and Lloyd Allison



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Ready?

# Agenda

- Dynamic Programming

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  - Brute force (aka the starting point)

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  - Overlapping sub-problems sub-problems must be overlapped
  - Backtracking

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- Dynamic Programming
  - Brute force (aka the starting point)
  - Overlapping sub-problems
  - Backtracking
  - Fibonacci
  - Coin change
  - Knapsack
    - Unbounded
    - 0/1
  - Edit distance

# Agenda

- Dynamic Programming
  - Brute force (aka the starting point)
  - Overlapping sub-problems
  - Backtracking for **solution reconstruction**
  - Fibonacci
  - Coin change
  - Knapsack
    - Unbounded
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Let us begin...

# Dynamic Programming

... a smarter way to brute force

# Dynamic Programming

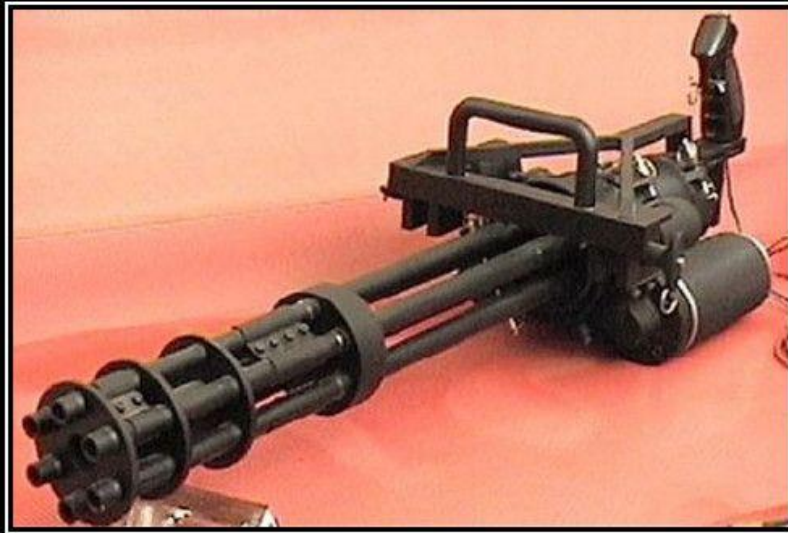
... a smarter way to brute force

- Can we brute force everything?

# Dynamic Programming

... a smarter way to brute force

- Can we brute force everything?



## BRUTE FORCE

If it doesn't work, you're just not using enough.

[themotivationalposters.com](http://themotivationalposters.com)

- Can we brute force everything?
  - So if we can brute force anything, why not just use it?

- Can we brute force everything?
  - So if we can brute force anything, why not just use it?
  - Because brute forcing needs **too much effort**...

- Can we brute force everything?
  - So if we can brute force anything, why not just use it?
  - Because brute forcing needs **too much effort**...
- Consider the Fibonacci problem
  - Find the  $n$ -th Fibonacci number

Questions?

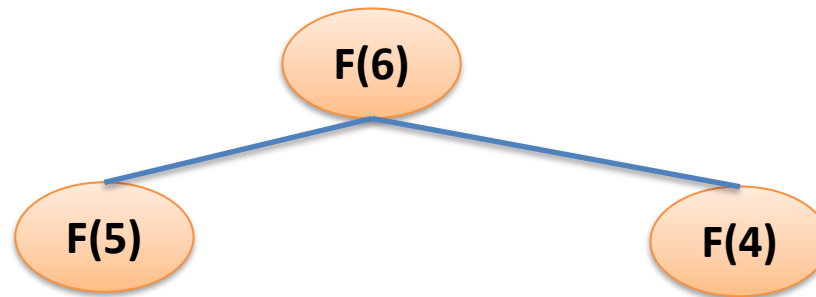


- Now if we want the 6<sup>th</sup> Fibonacci number...
  - How would we get it?



**F(6)**

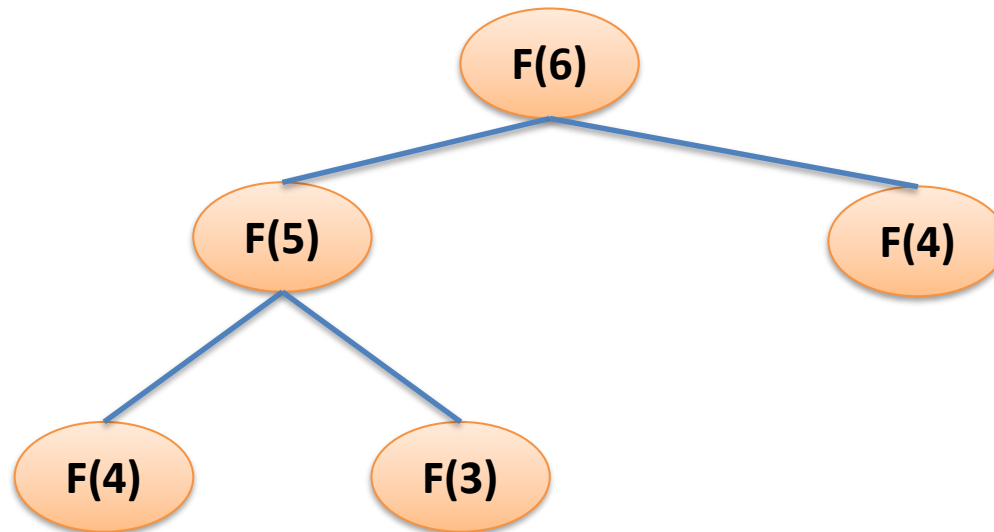
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# Fibonacci

n-th number in the series

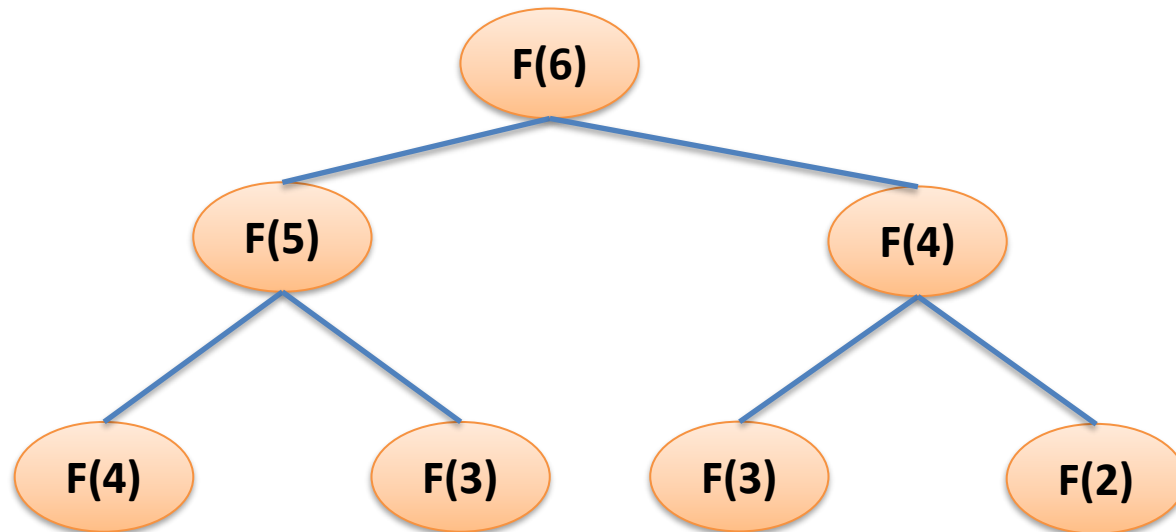
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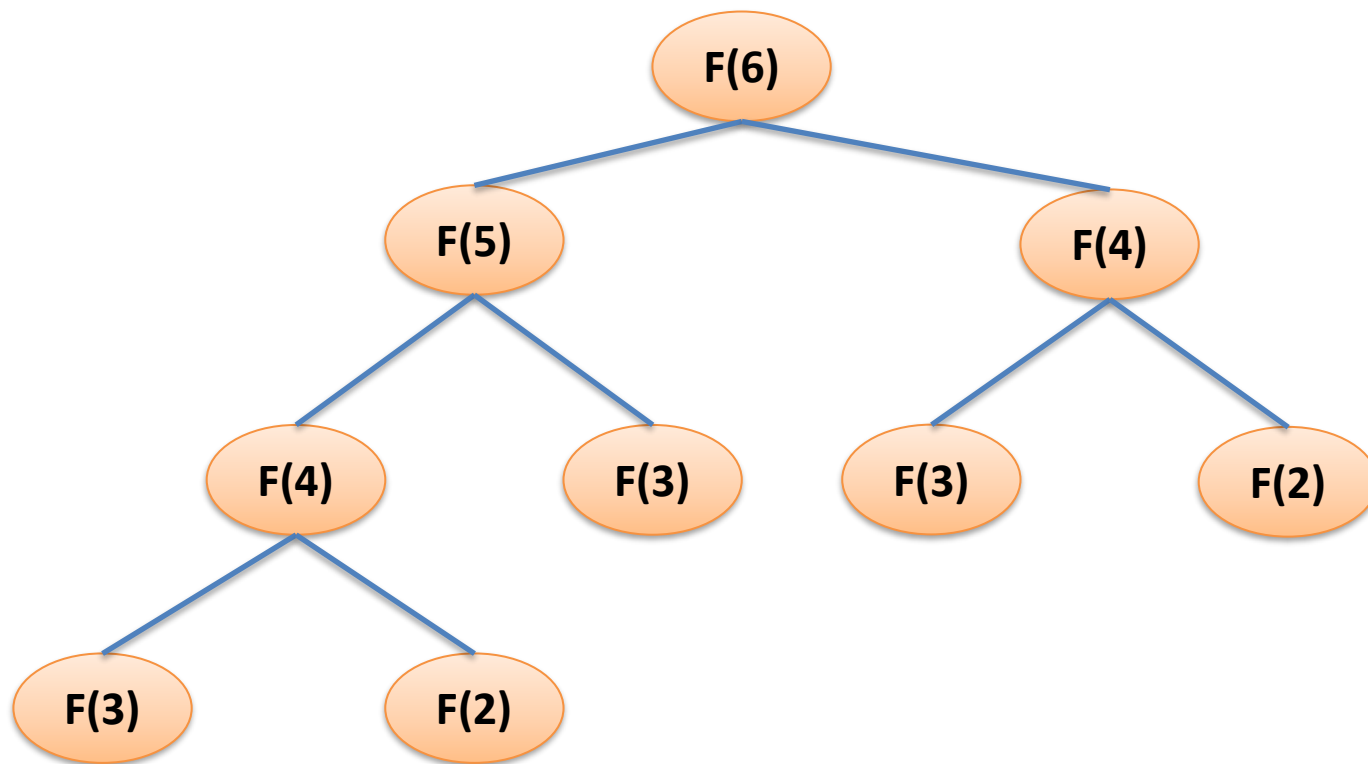
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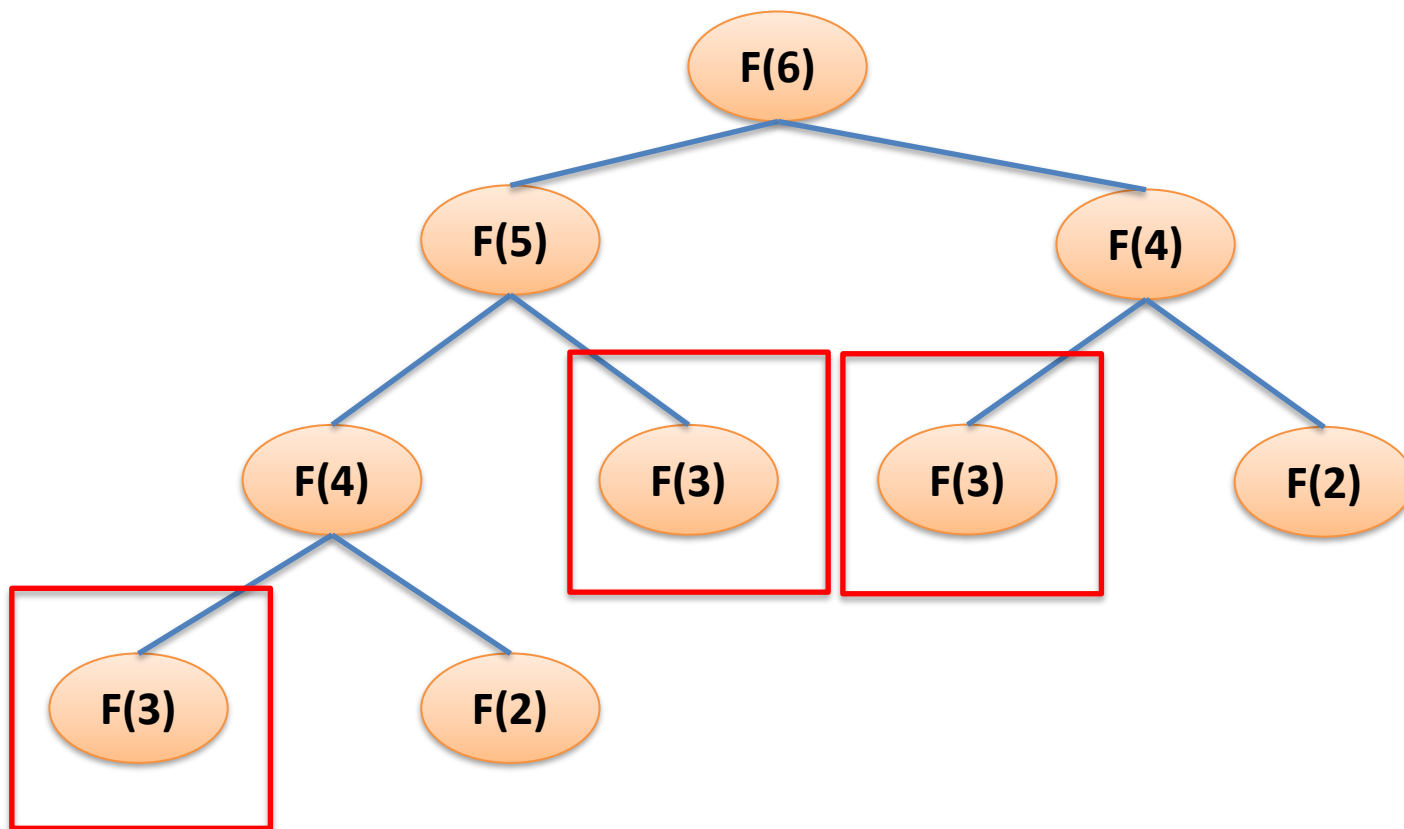
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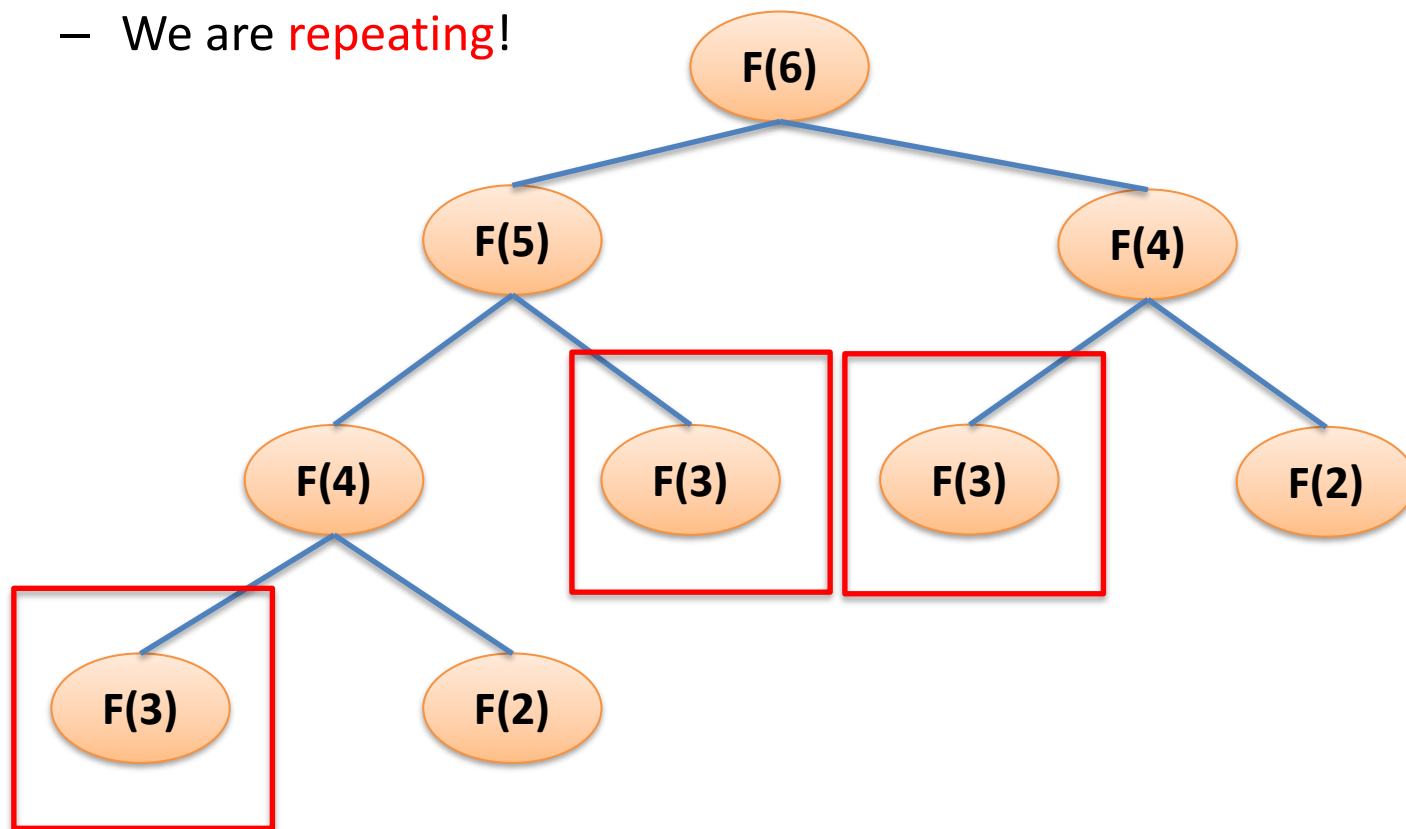
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- Now if we want the 6<sup>th</sup> Fibonacci number...
  - How would we get it?
  - We are **repeating**!

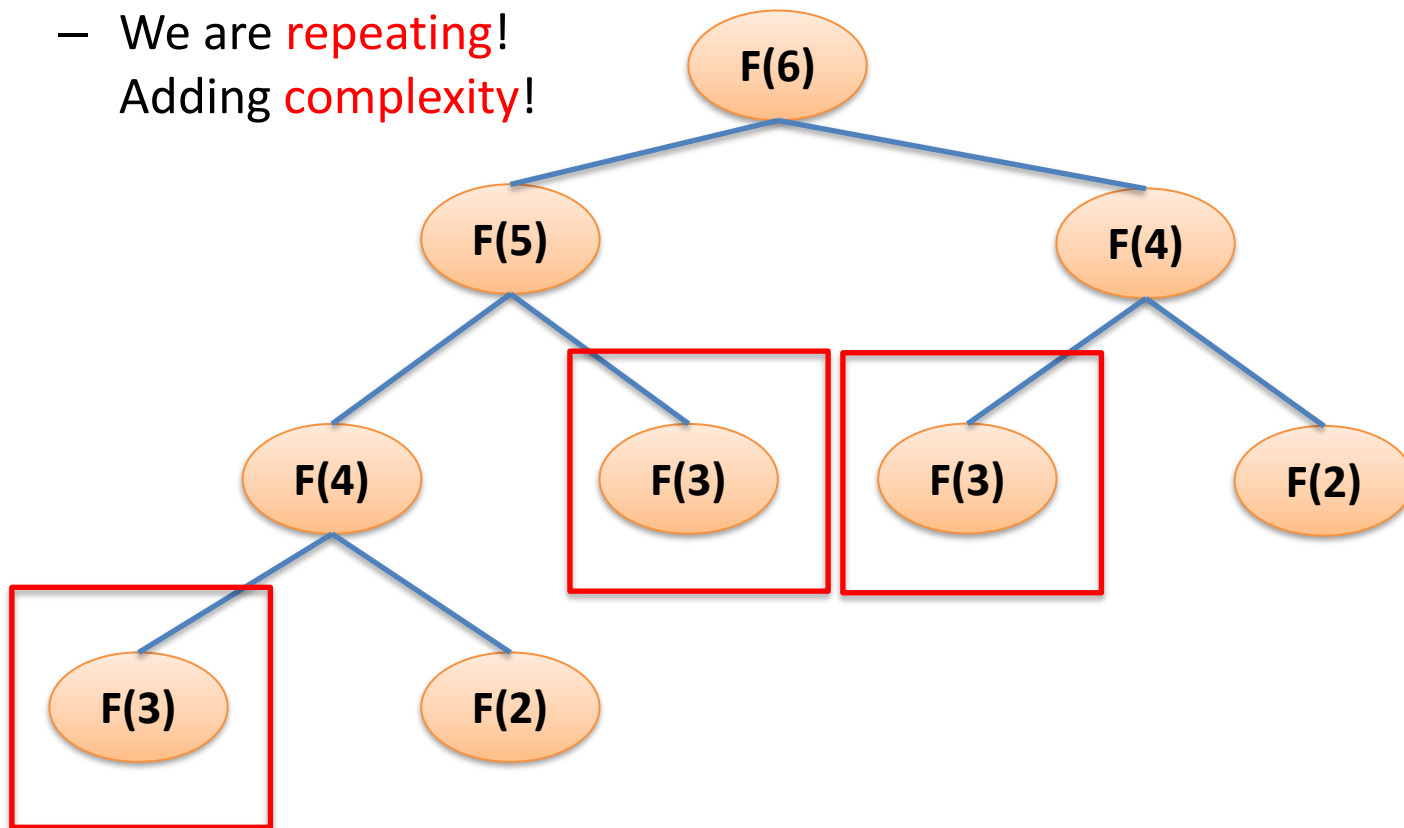


# Fibonacci

n-th number in the series

- Now if we want the 6<sup>th</sup> Fibonacci number...

- How would we get it?
- We are **repeating!**  
Adding **complexity!**









Questions?

- Now if we want the 6<sup>th</sup> Fibonacci number...
  - How would we get it?
  - We are **repeating!**  
Adding **complexity!**
  - So how do we reduce it?
    - **Memoization!**
    - Quick real world examples...

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      - What is 12x12?

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      - What is 12x12?
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    - Quick real world examples...
      - What is 12x12? **Remember this**
      - What is 12x13? Used for this...
      - What is 12x14? Used for this...



- Now if we want the 6<sup>th</sup> Fibonacci number...
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    - Quick real world examples...
      - What is 12x12? **Remember this**
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      - What is 12x14? Used for this...
- Let's apply to Fibonacci

Questions?

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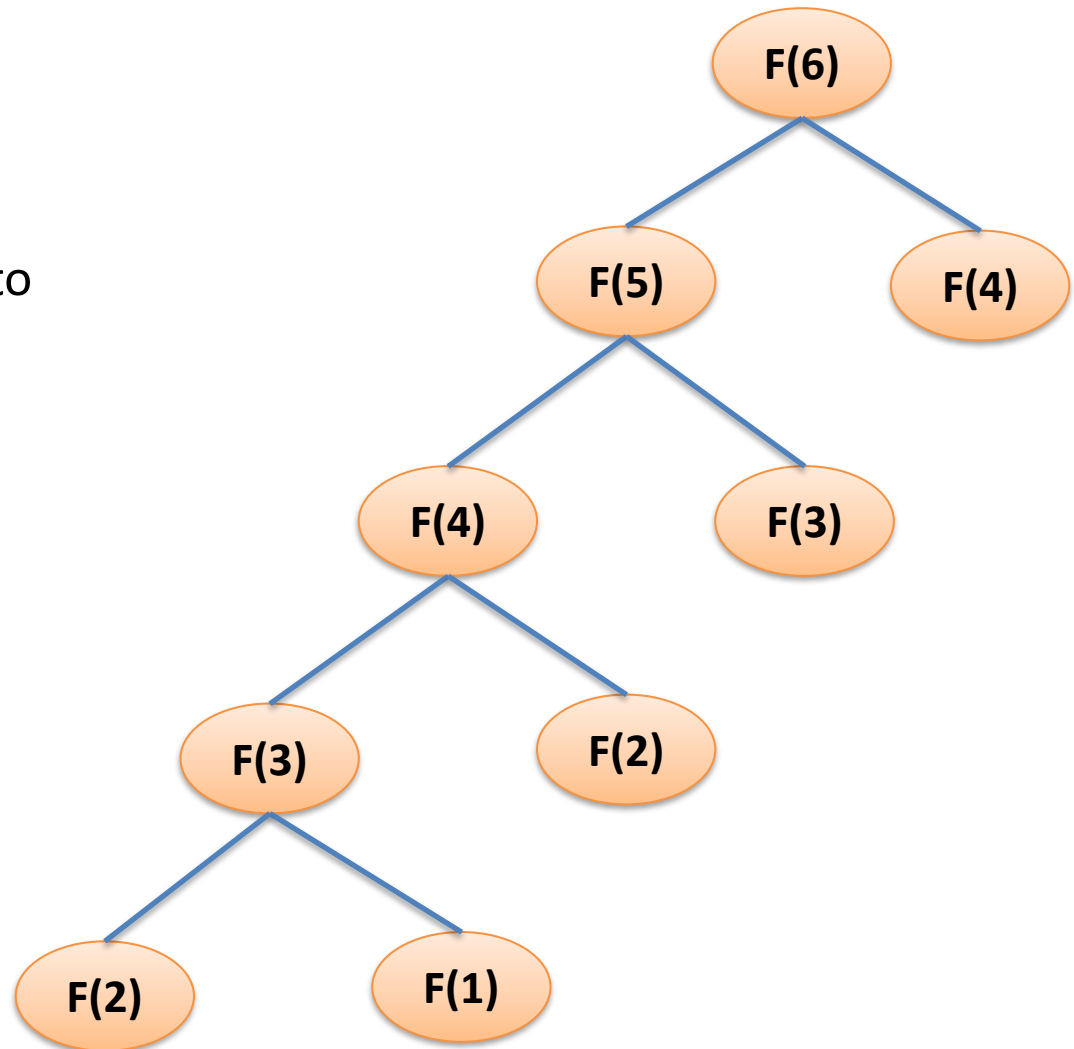
**F(6)**

- Top-down approach
  - Start from top

# Fibonacci

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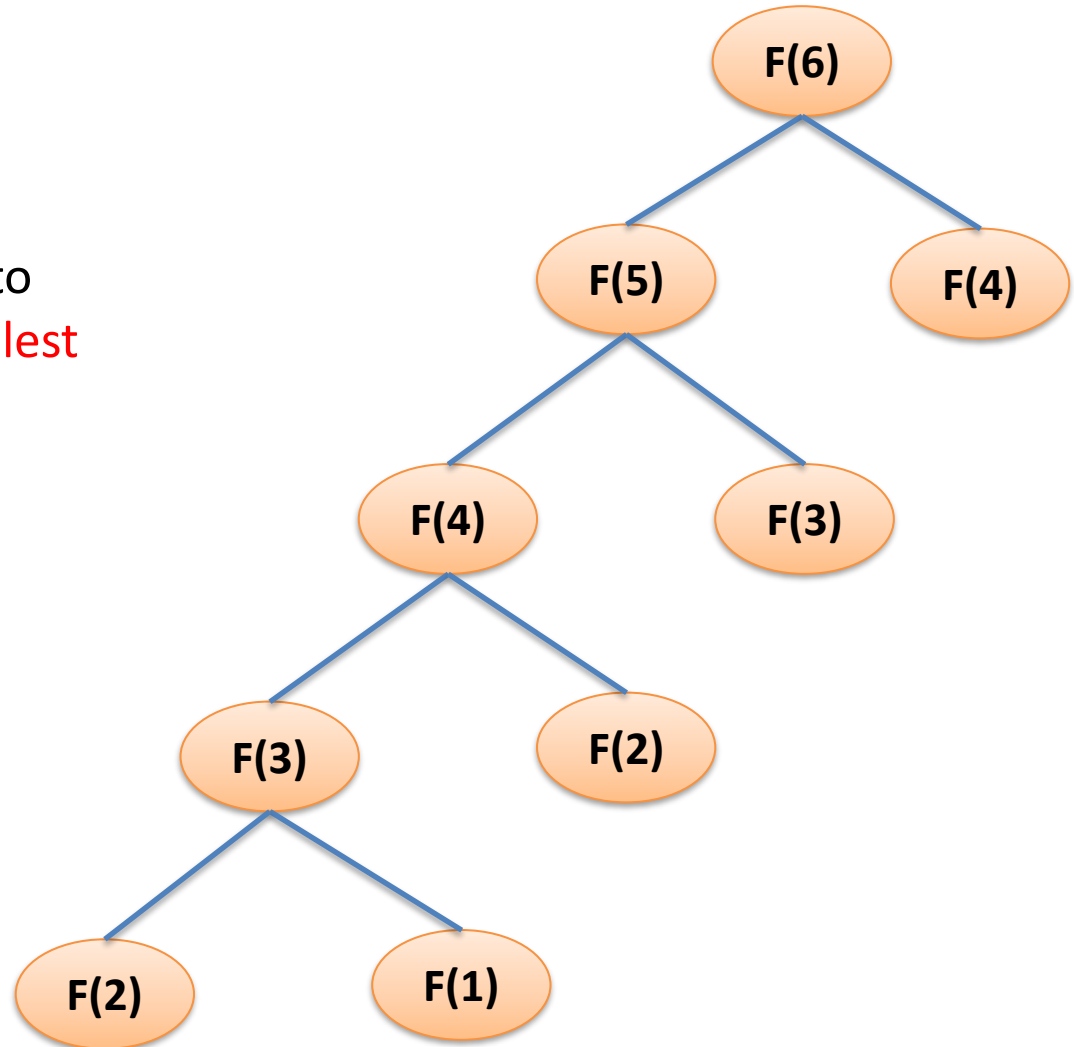
- Top-down approach
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  - Keep on breaking it into smaller problem



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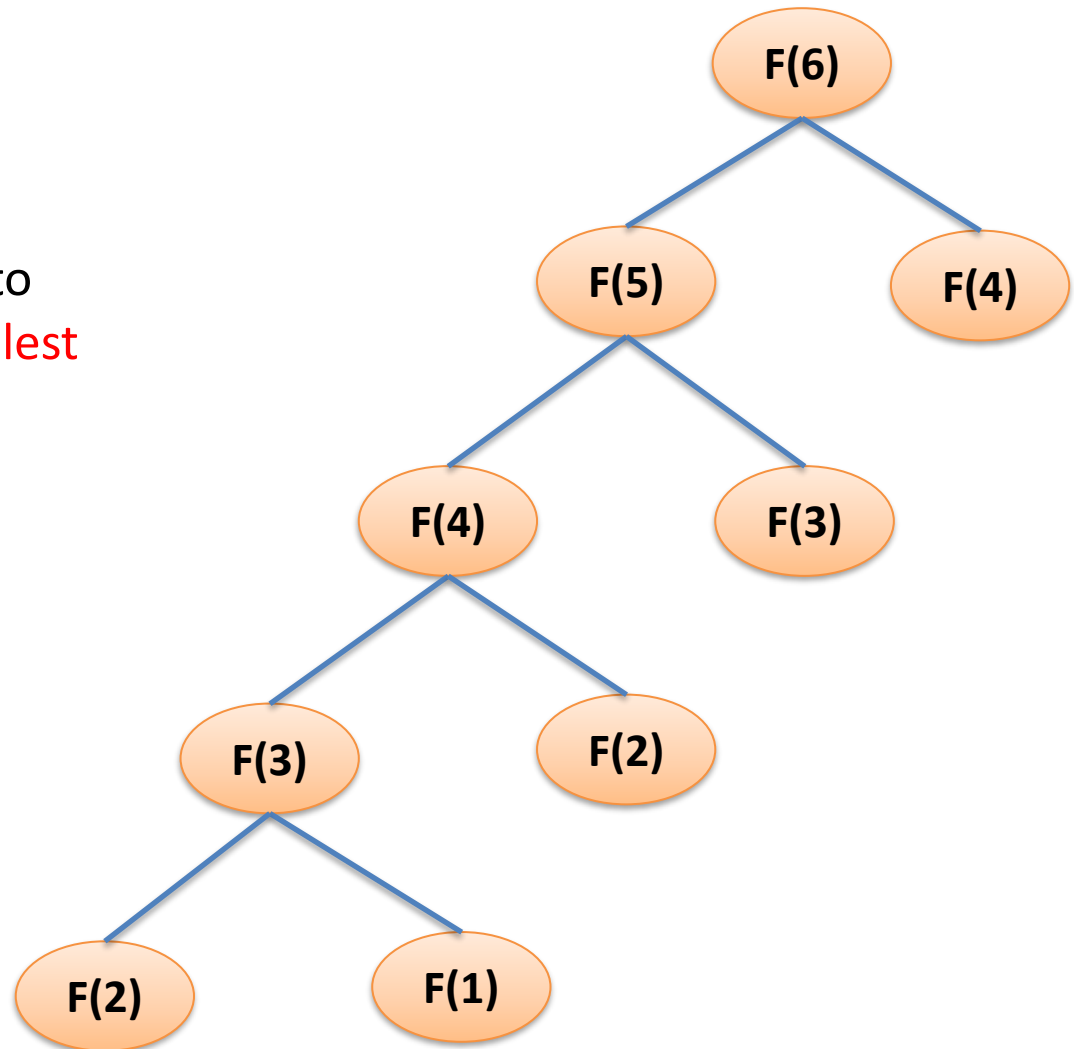
- Top-down approach
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# Fibonacci

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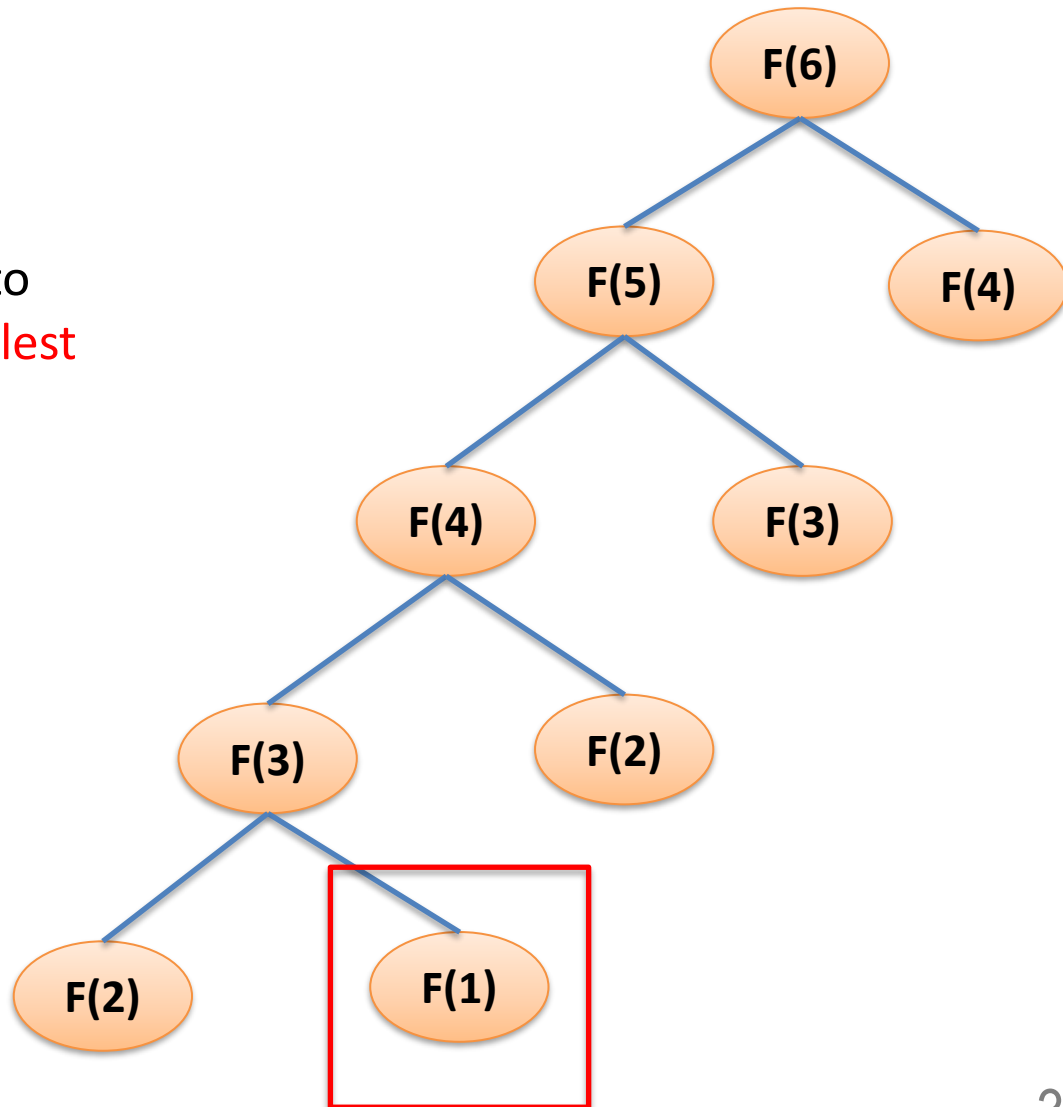
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  - Then solve it, reusing the results



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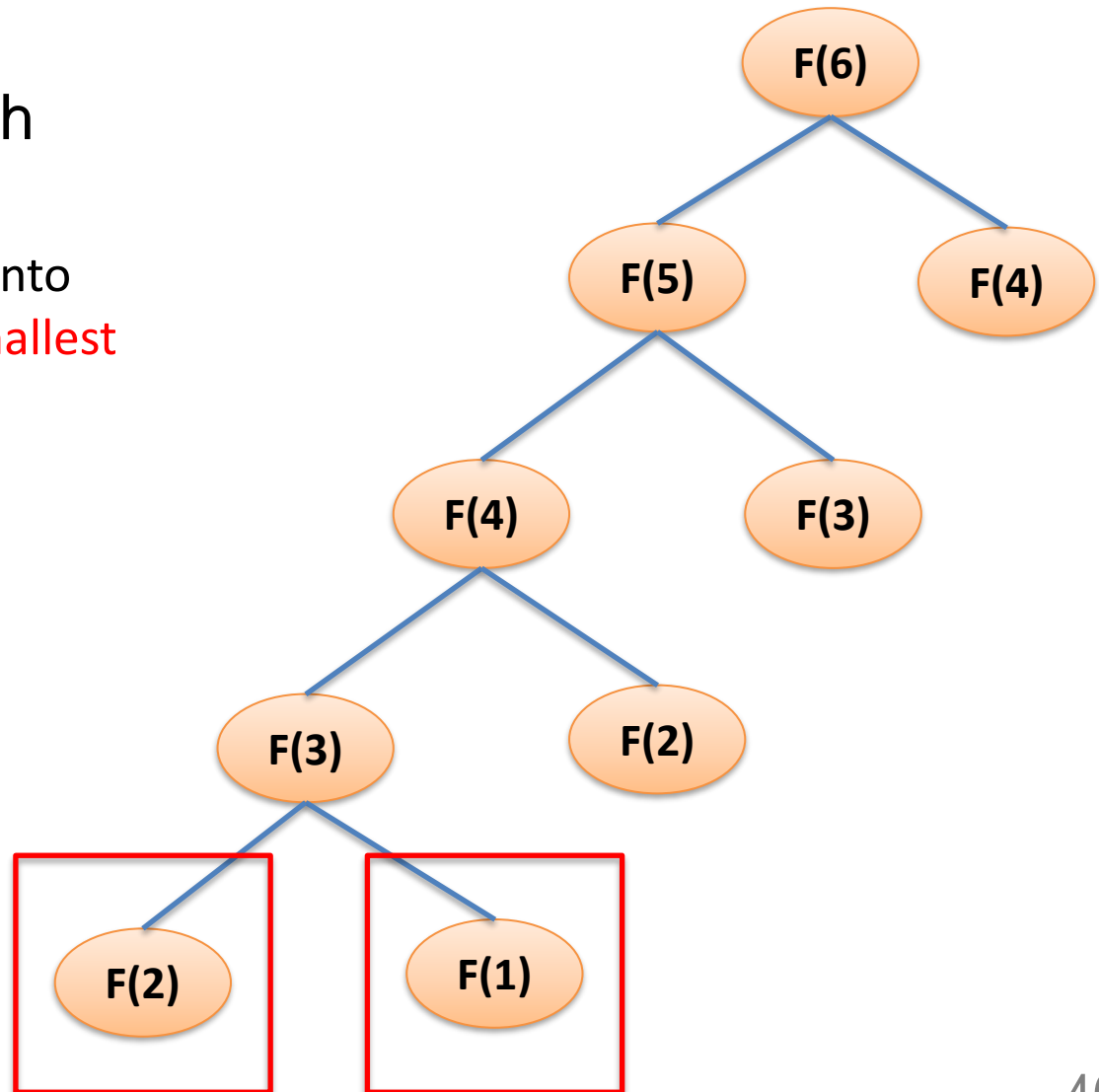
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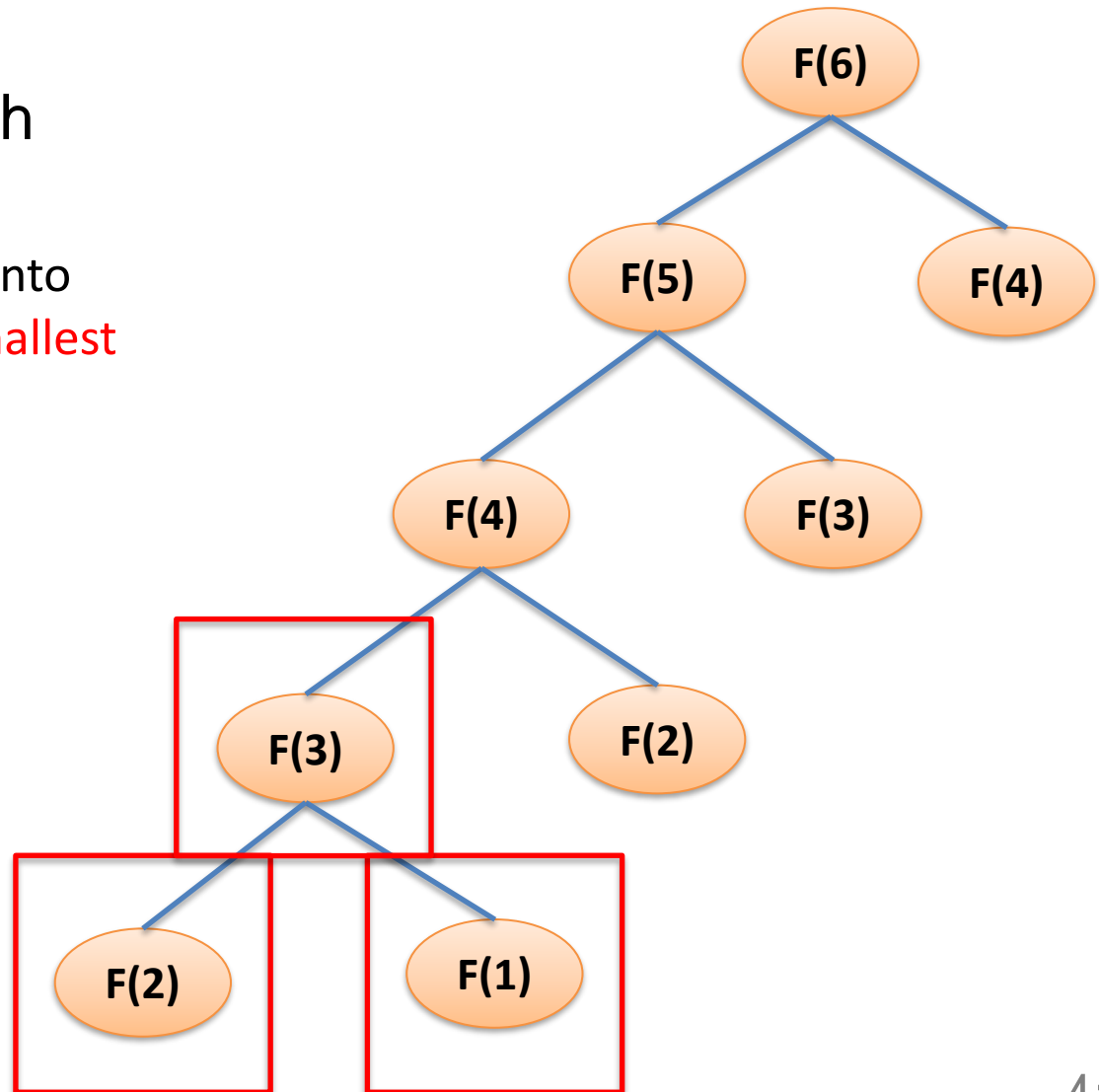




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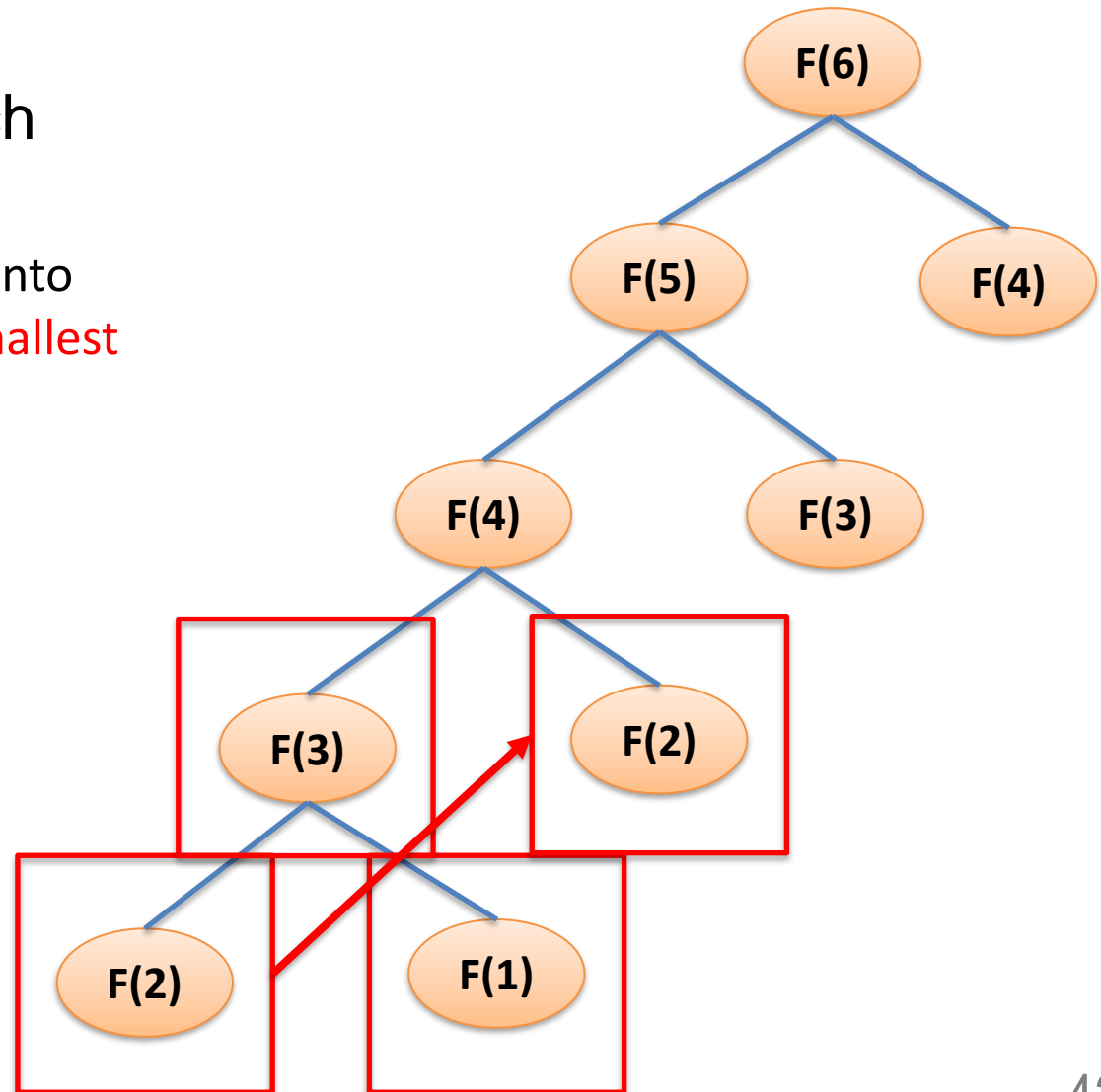
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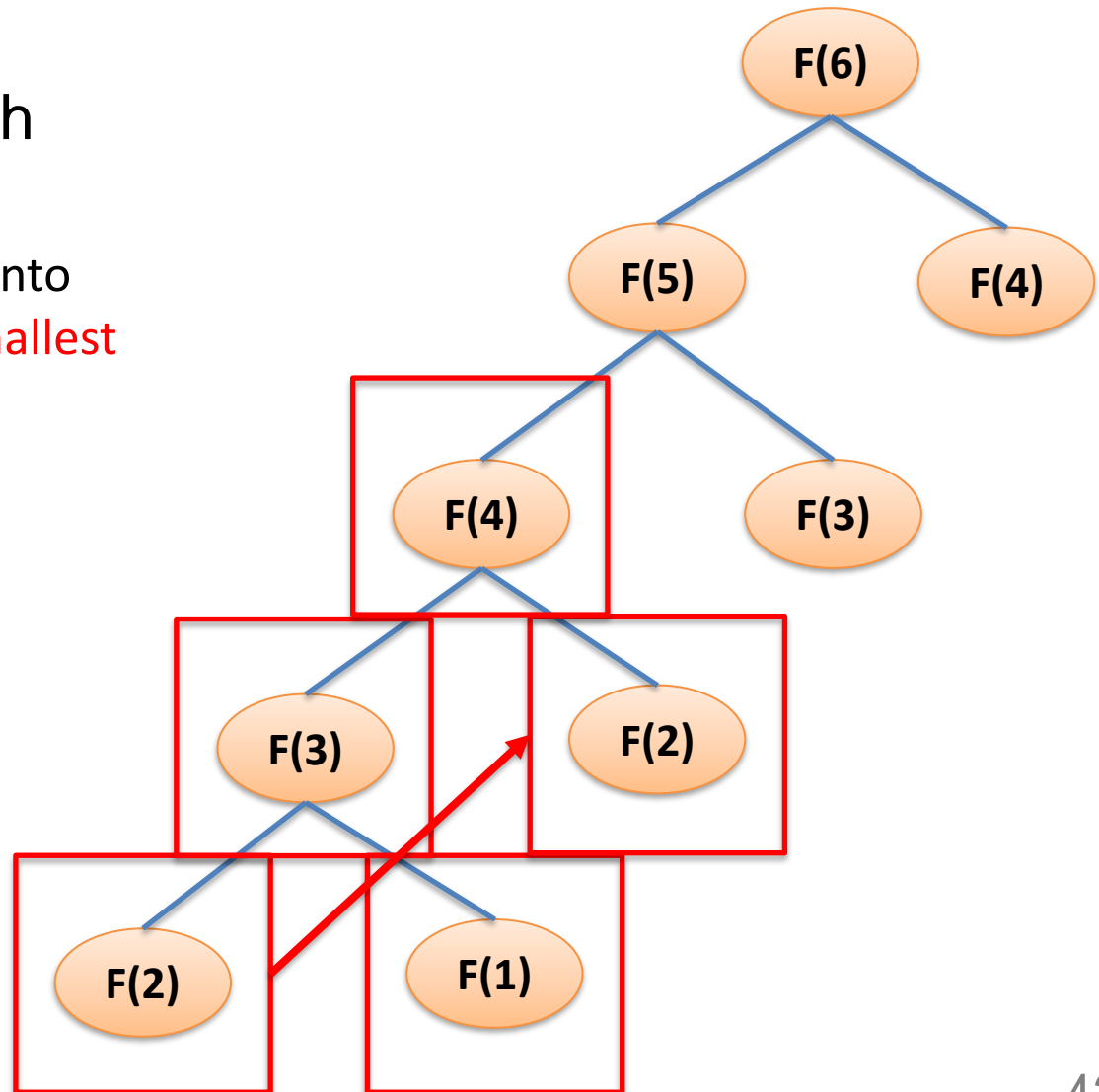
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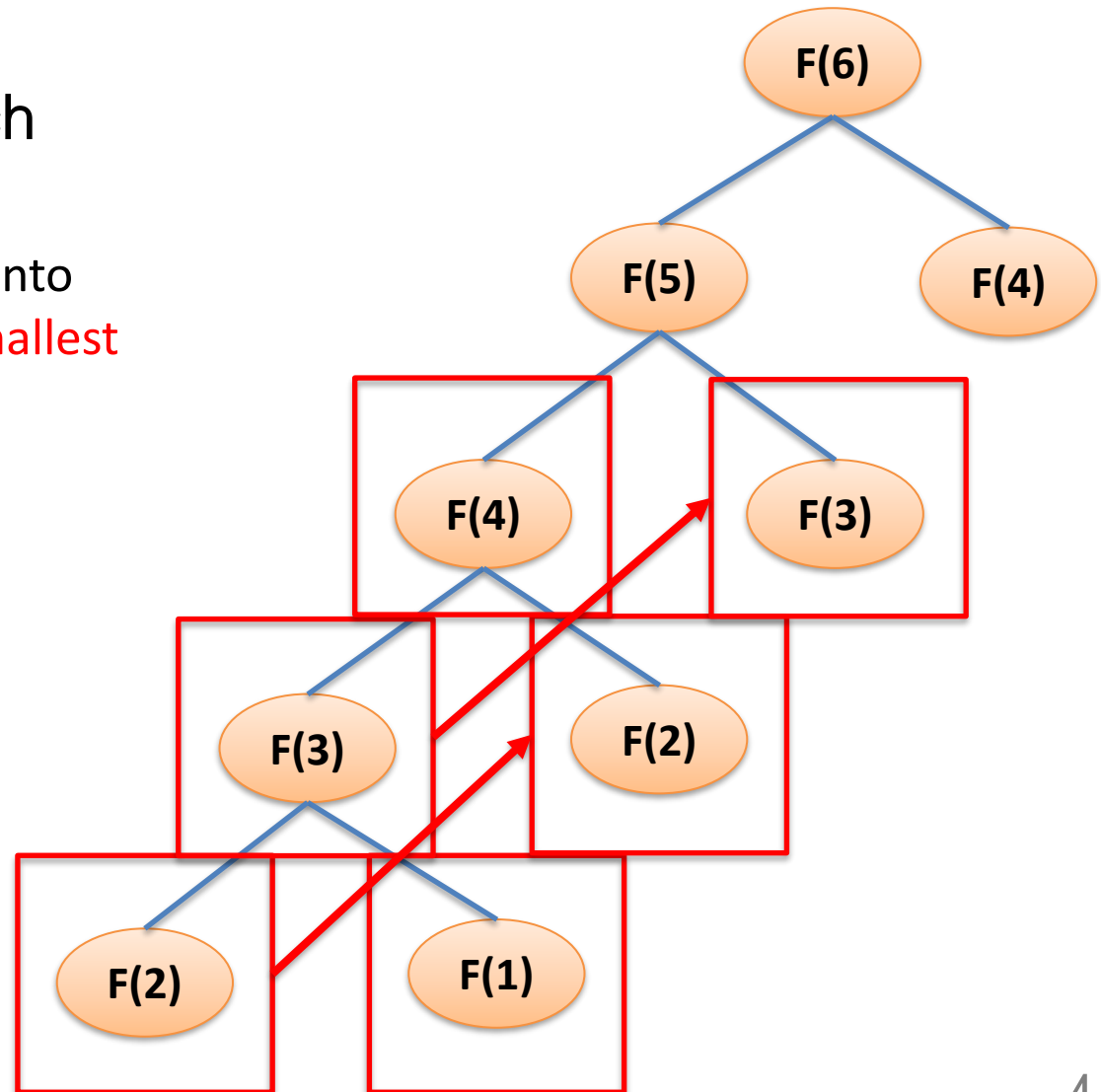
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```
1 memo = [-1] * N
2 memo[0] = 0           # in CompSci, we start from 0
3 memo[1] = 1
4
5 def fib_dpa(N):
6     if memo[N] != -1:  # if we have computed before
7         return memo[N] if the N index inside memo list is not -1 (there is value computed and stored < directly reuse it)
8     # compute the fibonacci
9     memo[N] = fib_dpa(N-1) + fib_dpa(N-2)
10    return memo[N]
```

Questions?

# Dynamic Programming vs Divide and Conquer

- Is it the same?

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These solutions are reusable due to overlapping sub problems
  - Combine solutions up to the big thing...
  
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- Is it the same?
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These solutions are reusable due to overlapping sub problems
  - Combine solutions up to the big thing...
- NO!... But why?
- Let us look at bottom-up to understand better...

# Dynamic Programming vs Divide and Conquer

- Is it the same?

- Take a big thing
- Divide to a smaller one
- Solve the smaller one.

These **OPTIMAL** solutions are reusable due to overlapping subproblems

since only computed once and used forever

- Combine solutions up to the big thing...

differentiate between divide and conquer  
and dynamic programming

merge sort has not repetitive subpart solution  
can not be reused  
so merge sort is divide and conquer(recursion)  
dynamic programming has subpart solution to remember  
to trade for faster computation speed

- NO!... But why?

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Questions?

# Fibonacci

n-th number in the series

- Bottom-up approach
  - Can you explain it?

# Fibonacci

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- Bottom-up approach
  - Start from the base case
  - Solve it
  - Use it to solve bigger case
  - Until we reach the final one...

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```

Questions?

# Fibonacci

n-th number in the series

- What is the complexity?
  - Bottom-up
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# Fibonacci

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5 def fib_dpa(N):
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8     # computer the fibonacci
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10    return memo[N]
```

recursion (Bottom-up)

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```

iterative

# Fibonacci

## n-th number in the series

- What is the complexity?
  - Bottom-up
  - Top-down

brute force  $O(2^n)$



**Suvashish Chakraborty** shared Nondeterministic Memes for NP Complete Teens's photo.

5 hrs · 

[#relatabe](#) [#ADSstuff](#)

When you use Dynamic Programming to solve a naively exponential time problem in polynomial time



Questions?

# Dynamic Programming

## The superpower

- Sounds easy right?



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  - Especially finding combinations

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  - Very popular in programming competitions!

# Dynamic Programming

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  - It isn't that easy however... Why?
- Can be used to solve a lot of problem
  - Especially finding combinations
  - Popular in interviews
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## The superpower

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# Dynamic Programming

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- So how?



# Dynamic Programming

## The superpower

- Sounds easy right?
  - It isn't that easy however... Why? Not easy to break problems down



- So how?
  - Practice



# Dynamic Programming

## The superpower

- Sounds easy right?
  - It isn't that easy however... Why? Not easy to break problems down
  
- So how?
  - Coin change
  - Knapsack
  - Edit-distance



Questions?

# Coin Change problem

The less number of coins...

- Consider the following scenario

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- Consider the following scenario
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    - $2 \times 50 + 1 \times 10$
    - $11 \times 10$
    - ... and many more



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    - $11 \times 10$
    - ... and many more
    - But we want the **smallest** number of coins!



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    - ... and many more
    - But we want the **smallest** number of coins!
- Let us now explore the possible solutions...



Questions?



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  - Try every combination!

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  - Choose the smallest number of coin...

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  - I want 110 doge coin, so what is the possible coin combination?
  
- Brute force?
  - Try every combination!
  - Choose the smallest number of coin...
  - Will it work? Of course!

Questions?

# Coin Change problem

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  - Try with the biggest number of coin
  - Then fill the balance with the rest #ez

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  - Doesn't always work (**greed is not good**)



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 $\{1, 5, 6, 9\}$  and I want 12

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  - Try with the biggest number of coin
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{1,5,6,9} and I want 12
    - $1 \times 9 + 3 \times 1 = 12$  for 4 coins

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- Consider the following scenario
  - Dogecoin currency is  $\{1, 5, 10, 50\}$
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  - Try with the biggest number of coin
  - Then fill the balance with the rest #ez
  - Doesn't always work (**greed is not good**)  
 $\{1, 5, 6, 9\}$  and I want 12
    - $1 \times 9 + 3 \times 1 = 12$  for 4 coins
    - $2 \times 6 = 12$  for 2 coins...

Questions?

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is  $\{1,5,10,50\}$
  - I want 110 doge coin, so what is the possible coin combination?
  
- Dynamic solution?
  - Let us try a smaller problem (easier for me to visualize)

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is  $\{1, 5, 6, 9\}$
  - I want 12 doge coin value, so what is the possible coin combination?

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The less number of coins...

- Consider the following scenario
  - Dogecoin currency is  $\{1, 5, 6, 9\}$
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	base case				N+1								
Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins													

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	0	1	2	3	4	5	6	7	8	9	10	11	12
Value													
Number of coins	0												



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Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1,5,6,9}
    - We will **loop** through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?

Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf

or -1

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Number of coins	0	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf

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$$0+1 = 1$$



Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf


want smallest number of coin then initialise with inf

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- Consider the following scenario

- Dogecoin currency is {1, 5, 6, 9}
  - We will loop through this over and over considering the coins...
- I want 12 doge coin value, so what is the possible coin combination?


$$-5+5 = 1 \text{ coin}$$

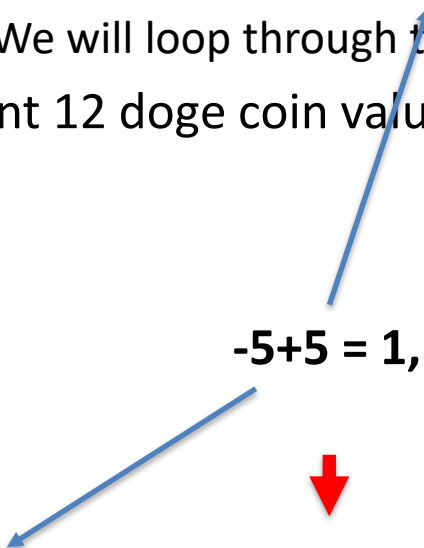
Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1, 5, 6, 9}
  - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?

-5+5 = 1, and so on....



Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1,5,6,9}
    - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?



Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf



# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1,5,6,9}
    - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?

Repeat the process...



Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1,5,6,9}
  - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?

$$1+1=2$$



Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf

optimal

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1,5,6,9}
  - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?

**1+1=2, 2 coins as well**



Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	2	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf

just use result here

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is  $\{1, 5, 6, 9\}$ 
    - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?



Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	2	inf	inf	inf	inf	inf	inf	inf	inf	inf	inf

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1,5,6,9}
    - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?



Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	2	3	4	inf	inf	inf	inf	inf	inf	inf	inf

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1,5,6,9}
    - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?

4+1=5, for 5 coins



Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	2	3	4	inf	inf	inf	inf	inf	inf	inf	inf

4 coins

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1,5,6,9}
    - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?

4+1=5, for 5 coins



Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	2	3	4	5	inf	inf	inf	inf	inf	inf	inf

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1,5,6,9}
    - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?

0+5=5, for 1 coin



Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	2	3	4	5	inf	inf	inf	inf	inf	inf	inf



# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1,5,6,9}
  - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?

0+5=5, for 1 coin

0 + 5

Value  
Number of coins

0	1	2	3	4	5	6	7	8	9	10	11	12
0	1	2	3	4	5	inf	inf	inf	inf	inf	inf	inf

VS  
1

optimal 0 coin

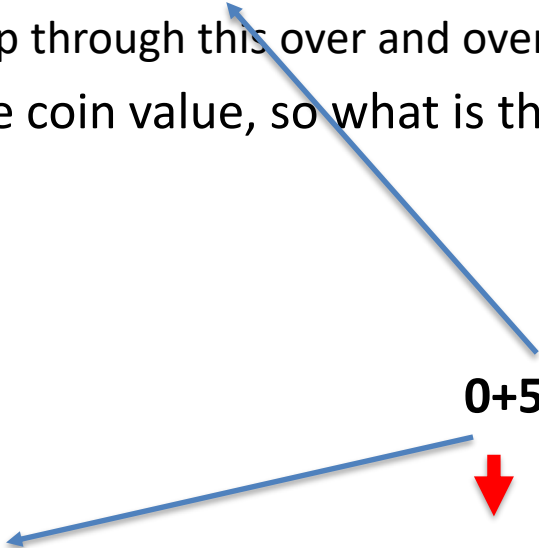
= 1 coin

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1,5,6,9}
    - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?

0+5=5, for 1 coin



Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	2	3	4	1	inf	inf	inf	inf	inf	inf	inf

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1,5,6,9}
    - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?

5+1=6, for 2 coins



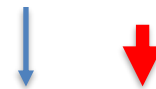
Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	2	3	4	1	inf	inf	inf	inf	inf	inf	inf

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1,5,6,9}
    - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?

5+1=6, for 2 coins



Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	2	3	4	1	2	inf	inf	inf	inf	inf	inf

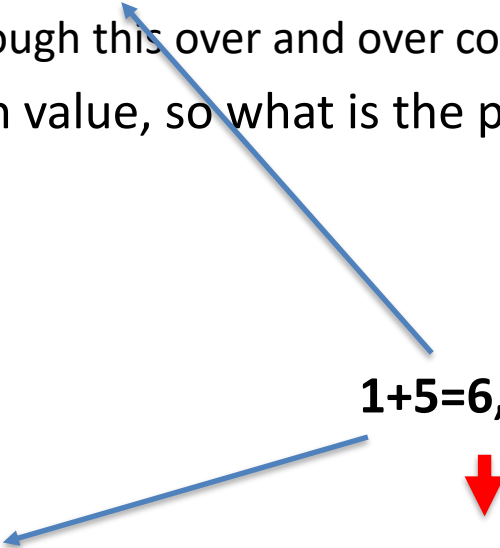
2 < inf

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1,5,6,9}
    - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?

1+5=6, for 2 coins



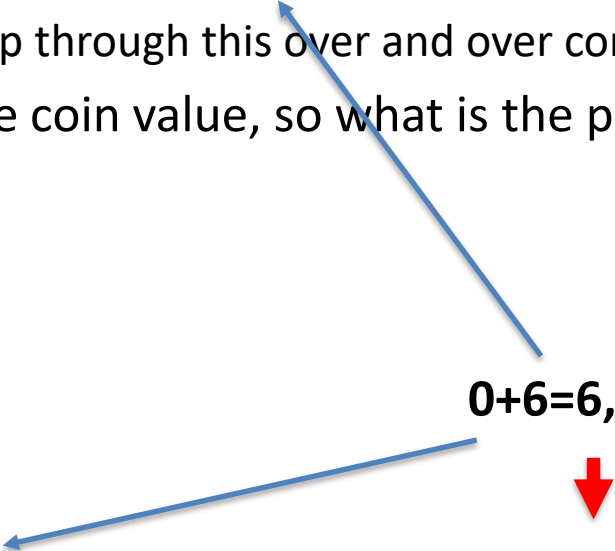
Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	2	3	4	1	2	inf	inf	inf	inf	inf	inf

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1,5,6,9}
    - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?

0+6=6, for 1 coin



Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	2	3	4	1	1	inf	inf	inf	inf	inf	inf

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1,5,6,9}
    - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?
  - So keep on running it and eventually we would be done



Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	2	3	4	1	1	inf	inf	inf	inf	inf	inf

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1,5,6,9}
    - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?
  - So keep on running it and eventually we would be done



Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	2	3	4	1	1	2	inf	inf	inf	inf	inf



# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1,5,6,9}
    - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?
  - So keep on running it and eventually we would be done



Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	2	3	4	1	1	2	3	inf	inf	inf	inf

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1,5,6,9}
  - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?
  - So keep on running it and eventually we would be done



Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	2	3	4	1	1	2	3	4	inf	inf	inf

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1,5,6,9}
    - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?
  - So keep on running it and eventually we would be done




Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	2	3	4	1	1	2	3	1	?	inf	inf

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1,5,6,9}
  - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?
  - So keep on running it and eventually we would be done

optimal case  
base case




Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	2	3	4	1	1	2	3	1	2	inf	inf

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1,5,6,9}
    - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?
  - So keep on running it and eventually we would be done

$7 - 1 = 0$   
 $7 - 5 = 2$   
 $7 - 6 = 1$




	0	1	2	3	4	5	6	7	8	9	10	11	12
Value													
Number of coins	0	1	2	3	4	1	1	2	3	1	2	2	inf

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1,5,6,9}
    - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?
  - So keep on running it and eventually we would be done




	0	1	2	3	4	5	6	7	8	9	10	11	12
Value	0	1	2	3	4	1	1	2	3	1	2	2	2
Number of coins	0	1	2	3	4	1	1	2	3	1	2	2	2

Questions?

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is  $\{1, 5, 6, 9\}$   $O(M)$ 
    - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?
  - So keep on running it and eventually we would be done
    - Complexity?



	$O(N)$ each location on list loop through $\{1,5,6,9\} = O(NM)$												
<b>Value</b>	<b>0</b>	<b>1</b>	<b>2</b>	<b>3</b>	<b>4</b>	<b>5</b>	<b>6</b>	<b>7</b>	<b>8</b>	<b>9</b>	<b>10</b>	<b>11</b>	<b>12</b>
<b>Number of coins</b>	0	1	2	3	4	1	1	2	3	4	2	2	2



# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1,5,6,9}
    - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?
  - So keep on running it and eventually we would be done
    - Complexity?  $O(NM)$



Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	2	3	4	1	1	2	3	4	2	2	2

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1,5,6,9}
    - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?
  - So keep on running it and eventually we would be done
    - Complexity?  $O(NM)$  still much faster than brute force...




Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	2	3	4	1	1	2	3	4	2	2	2

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1,5,6,9}
    - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?
  - So keep on running it and eventually we would be done
    - Complexity?  $O(NM)$  still much faster than brute force...
      - Note: If the list is sorted, we can terminate earlier on the smaller values

3 < 5, skip 6, 9



Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	2	3	4	1	1	2	3	4	2	2	2

Questions?

# Coin Change problem

The less number of coins...

- Consider the following scenario
  - Dogecoin currency is {1,5,6,9}
    - We will loop through this over and over considering the coins...
  - I want 12 doge coin value, so what is the possible coin combination?
  - So keep on running it and eventually we would be done
    - Complexity?  $O(NM)$  still much faster than brute force...
    - Can you **code** it?



Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	2	3	4	1	1	2	3	4	2	2	2

# Coin Change problem

The less number of coins...

- I'll give the algorithm here

```
1 memo = [inf] * (N+1)
2 memo[0] = 0
3 for value in range(1, N+1):
4     for j in range(M):                # this is to go through coins
5         if coin[j] <= value:          if 1 <= 3
6             balance = value - coin[j] 2 = 3 - 1
7             count = 1 + memo[balance] memo[2] = 2, reuse
8             if count < memo[value]:    # if we have new optimal
9                 memo[value] = count
10 return memo[N]
```

# Coin Change problem

The less number of coins...

- I'll give the algorithm here
  - Is this top-down or bottom-up?

```
1 memo = [inf] * (N+1)
2 memo[0] = 0
3 for value in range(1, N+1):
4     for j in range(M):                # this is to go through coins
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6             balance = value - coin[j]
7             count = 1 + memo[balance]
8             if count < memo[value]:    # if we have new optimal
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10 return memo[N]
```

# Coin Change problem

The less number of coins...

- I'll give the algorithm here
  - Is this top-down or **bottom-up**?

```
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2 memo[0] = 0
3 for value in range(1, N+1):
4     for j in range(M):                # this is to go through coins
5         if coin[j] <= value:
6             balance = value - coin[j]
7             count = 1 + memo[balance]
8             if count < memo[value]:    # if we have new optimal
9                 memo[value] = count
10 return memo[N]
```



# Coin Change problem

The less number of coins...

- I'll give the algorithm here
  - Is this top-down or **bottom-up**?
  - The following is the **top-down** from Aamir (usually recursive)

```
1 memo = [inf] * (N+1)
2 memo[0] = 0
3 for value in range(1, N+1):
4     for j in range(M):                # this is to go through coins
5         if coin[j] <= value:
6             balance = value - coin[j]
7             count = 1 + memo[balance]
8             if count < memo[value]:    # if we have new optimal
9                 memo[value] = count
10 return memo[N]
```

# Top-down Solution

Initialize Memo[ ] to contain -1 for all indices # -1 indicates the solution for this index has not been computed yet

Memo[0] = 0

Function CoinChange(value)

watch clayton's

if Memo[value] != -1: \\ DISCUSS infinity is incorrect

return Memo[value]

else:

minCoins = Infinity

for i=1 to N

if Coins[ i ] <= value recursive

c = 1 + CoinChange(value - Coins[ i ])

if c < minCoins

minCoins = c

Memo[value] = minCoins

return Memo[value]

Bottom up solution:

1 + Memo[ value - Coins[i] ]

# Coin Change problem

The less number of coins...

- Top-down vs Bottom up
  - Top-down might save some computations
  - Bottom-up might save space especially since no recursion



# Coin Change problem

The less number of coins...

- Top-down vs Bottom up
  - Top-down might save some computations
  - Bottom-up might save space especially since no recursion
  - I only use bottom-up

# Coin Change problem

The less number of coins...

- Top-down vs Bottom up
  - Top-down might save some computations
  - Bottom-up might save space especially since no recursion
  - I only use bottom-up
    - But some problems could be easier with top-down as it is more intuitive
    - Technically both are interchangeable...

Questions?

# Kapsack Problem

Min-max like a boss

- A problem that can be applicable to a lot of real life scenario

# Kapsack Problem

Min-max like a boss

- A problem that can be applicable to a lot of real life scenario
  - Given a limitation (cost)



# Kapsack Problem

Min-max like a boss

- A problem that can be applicable to a lot of real life scenario
  - Given a limitation (cost)
  - Optimize something (profit)

# Kapsack Problem

## Min-max like a boss

- A problem that can be applicable to a lot of real life scenario
  - Given a limitation (cost)
  - Optimize something (profit)
  - 18<sup>th</sup> most popular algorithmic problem

# Kapsack Problem

## Min-max like a boss

- A problem that can be applicable to a lot of real life scenario
  - Given a limitation (cost)
  - Optimize something (profit)
  - 18<sup>th</sup> most popular algorithmic problem
- Given a capacity  $C$  and a set of items with weights and value... can you find a combination of item such as the total weight  $< C$  but the total value is maximized?

# Kapsack Problem

## Min-max like a boss

- A problem that can be applicable to a lot of real life scenario
  - Given a limitation (cost)
  - Optimize something (profit)
  - 18<sup>th</sup> most popular algorithmic problem
- Given a capacity  $C$  and a set of items with weights and value... can you find a combination of item such as the total weight  $< C$  but the total value is maximized?
  - Unbounded = items are unlimited
  - Bounded = each item can only be taken once

# Kapsack Problem

## Min-max like a boss

- Let say you have these items

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40

# Kapsack Problem

## Min-max like a boss



- Let say you have these items

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40

- And you only have 12 kg... what would you loot?

# Kapsack Problem

## Min-max like a boss



- Let say you have these items

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40

- And you only have 12 kg... what would you loot?
  - $2 \times B + 2 \times D = \$780$

# Kapsack Problem

## Min-max like a boss



- Let say you have these items

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40

- And you only have 12 kg... what would you loot?
  - $2 \times B + 2 \times D = \$780$
  - So how do we code it?  
This is very similar to the coin change!



Questions?

# Kapsack Problem

## Unbounded

- Let us try the bottom-up approach

# Kapsack Problem

## Unbounded

- Let us try the bottom-up approach
  - Start with 0 weight till 12 weight

# Kapsack Problem

## Unbounded

- Let us try the bottom-up approach
  - Start with 0 weight till 12 weight
  - Find the optimal for each weight...

# Kapsack Problem


## Unbounded

- Let us try the bottom-up approach
  - Start with 0 weight till 12 weight
  - Find the optimal for each weight...
  - Use the earlier optimal for the new weight

# Kapsack Problem

## Unbounded

- Let us try the bottom-up approach
  - Start with 0 weight till 12 weight
  - Find the **optimal** for each weight...  
**Maximum profit!**
  - Use the earlier **optimal** for the new weight

 from lowest to maximum


Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0												

# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40




Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	0	0	0	0	0	0	0	0	0	0	0	0

# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40



Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	0	0	0	0	0	0	0	0	0	0	0	0




# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40

9kg, too heavy



Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	0	0	0	0	0	0	0	0	0	0	0	0


# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40

5kg, too heavy



Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	0	0	0	0	0	0	0	0	0	0	0	0


# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40

6kg, too heavy



Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	0	0	0	0	0	0	0	0	0	0	0	0


# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40

1kg, can fit



Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	0	0	0	0	0	0	0	0	0	0	0	0

# Kapsack Problem


## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40



1kg, can fit... so we find the optimal of 0kg + 1kg



Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	0	0	0	0	0	0	0	0	0	0	0	0

# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40

1kg, can fit... so we find the **optimal of 0kg + 1kg**


Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	40	0	0	0	0	0	0	0	0	0	0	0

# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40



	0	1	2	3	4	5	6	7	8	9	10	11	12
Weight													
Profit	0	40	0	0	0	0	0	0	0	0	0	0	0


# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40

9kg too heavy



Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	40	0	0	0	0	0	0	0	0	0	0	0



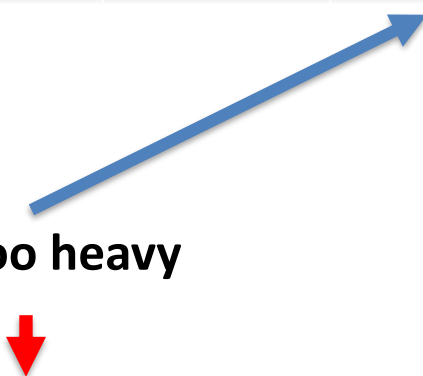
# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40

5kg too heavy



Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	40	0	0	0	0	0	0	0	0	0	0	0

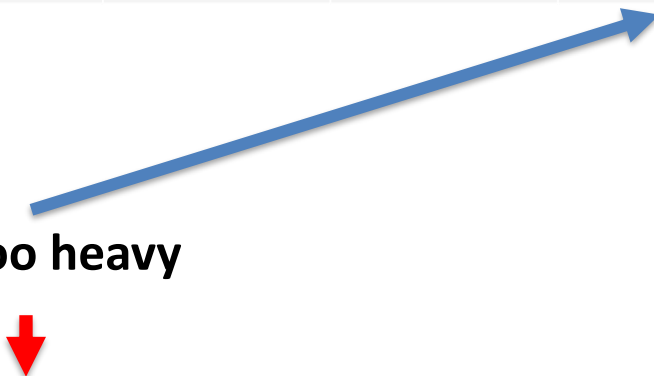
# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40

6kg too heavy



Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	40	0	0	0	0	0	0	0	0	0	0	0

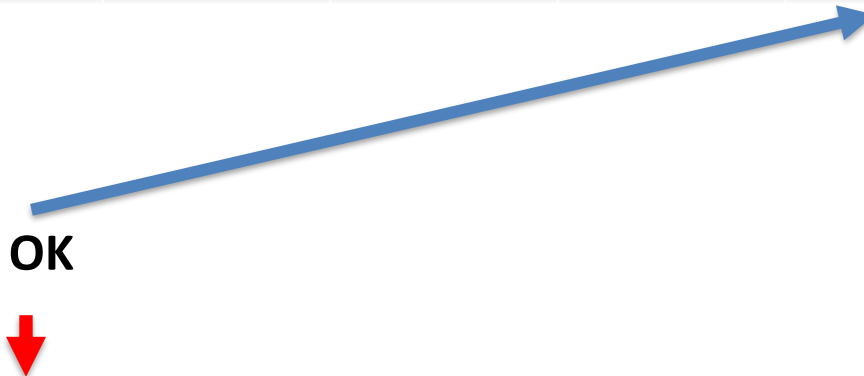
# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40

1kg is OK



Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	40	0	0	0	0	0	0	0	0	0	0	0


# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40

1kg is OK... optimal of 1kg + profit from this item



Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	40	0	0	0	0	0	0	0	0	0	0	0

# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40

1kg is OK... optimal of 1kg + profit from this item

Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	40	0	0	0	0	0	0	0	0	0	0	0

# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40

1kg is OK... optimal of 1kg + profit from this item

Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	40	80	0	0	0	0	0	0	0	0	0	0

# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40

We just repeat the process



Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	40	80	120	0	0	0	0	0	0	0	0	0

# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40

We just repeat the process



Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	40	80	120	160	0	0	0	0	0	0	0	0



# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40

Now let us try here



Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	40	80	120	160	0	0	0	0	0	0	0	0

# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40



9kg too heavy...



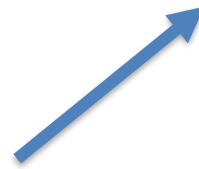
Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	40	80	120	160	0	0	0	0	0	0	0	0

# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40



5kg can fit. Optimal 0kg + current item



Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	40	80	120	160	0	0	0	0	0	0	0	0

# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40



5kg can fit. Optimal 0kg + current item



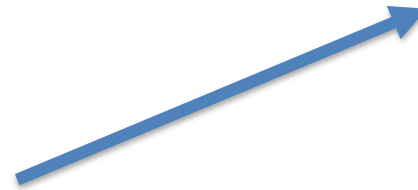
Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	40	80	120	160	350	0	0	0	0	0	0	0

# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40



**6kg too heavy**



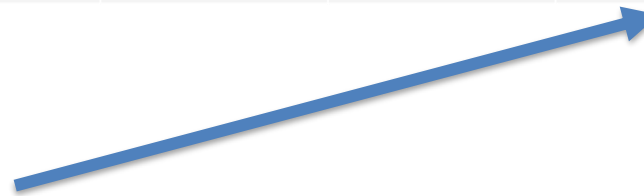
Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	40	80	120	160	350	0	0	0	0	0	0	0

# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40



**1kg is OK. Optimal 4kg + current item**



Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	40	80	120	160	350	0	0	0	0	0	0	0

# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40



1kg is OK. Optimal 4kg + current item but it is only 200

Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	40	80	120	160	350	0	0	0	0	0	0	0

# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40



Moving on....



Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	40	80	120	160	350	0	0	0	0	0	0	0



# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40



heavy



Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	40	80	120	160	350	0	0	0	0	0	0	0

# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40

Optimal 1kg + current item

Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	40	80	120	160	350	0	0	0	0	0	0	0

# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40

Optimal 1kg + current item

Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	40	80	120	160	350	390	0	0	0	0	0	0

# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40

Optimal 0kg + current item = 180 only

Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	40	80	120	160	350	390	0	0	0	0	0	0

# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40

Optimal 5kg + current item = same 390

Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	40	80	120	160	350	390	0	0	0	0	0	0

# Kapsack Problem

## Unbounded

- Let us run through it

Item	A	B	C	D
Weight	9kg	5kg	6kg	1kg
Value	\$550	\$350	\$180	\$40

Eventually....



Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	40	80	120	160	350	390	430	470	550	700	740	780

Questions?

Have a break!



# Kapsack Problem

## Unbounded

- So what is our algorithm?
  - Let us try to produce it now as part of the class activity!

# Kapsack Problem

## Unbounded

- So what is our algorithm?
  - Very similar to the coin change except
    - Finding maximum instead of minimum
    - Initialized to 0

# Kapsack Problem

## Unbounded

- So what is our algorithm?
  - Very similar to the coin change except
    - Finding maximum instead of minimum
    - Initialized to 0
  - You have the 1 array, called items
    - N number of items
    - Items[i].weight for the weight
    - Items[i].profit for the profit

Questions?

# Kapsack Problem

## Unbounded

- So what is our algorithm?
  - Very similar to the coin change except
    - Finding maximum instead of minimum
    - Initialized to 0

```
1 memo = [0] * (N+1) # N is the total weight
2 memo[0] = 0
3 for bag_weight in range(1, N+1):
4     for j in range(M): # this is to go through items
5         if items[j].weight <= bag_weight:
6             balance = bag_weight - items[j].weight
7             profit = item[j].profit + memo[balance]
8             if profit > memo[bag_weight]: # if we have new optimal
9                 memo[bag_weight] = profit
10 return memo[N]
```

# Kapsack Problem

## Unbounded

- So what is our algorithm?
  - Very similar to the coin change except
    - Finding maximum instead of minimum
    - Initialized to 0
  - Complexity?

```
1 memo = [0] * (N+1) # N is the total weight
2 memo[0] = 0
3 for bag_weight in range(1, N+1):
4     for j in range(M): # this is to go through items
5         if items[j].weight <= bag_weight:
6             balance = bag_weight - items[j].weight
7             profit = item[j].profit + memo[balance]
8             if profit > memo[bag_weight]: # if we have new optimal
9                 memo[bag_weight] = profit
10 return memo[N]
```

# Kapsack Problem

## Unbounded

- So what is our algorithm?
  - Very similar to the coin change except
    - Finding maximum instead of minimum
    - Initialized to 0
  - Complexity?  $O(NM)$

see top-down from clayton side

```
1 memo = [0] * (N+1) # N is the total weight
2 memo[0] = 0
3 for bag_weight in range(1, N+1):
4     for j in range(M): # this is to go through items
5         if items[j].weight <= bag_weight:
6             balance = bag_weight - items[j].weight
7             profit = item[j].profit + memo[balance]
8             if profit > memo[bag_weight]: # if we have new optimal
9                 memo[bag_weight] = profit
10 return memo[N]
```

# Kapsack Problem

## Unbounded

- So what is our algorithm?
  - Very similar to the coin change except
    - Finding maximum instead of minimum
    - Initialized to 0
  - Complexity?
  - Top-down? See Nathan's slides



Questions?

# Kapsack Problem

## Unbounded

- But is the code correct?

```
1 memo = [0] * (N+1) # N is the total weight
2 memo[0] = 0
3 for bag_weight in range(1, N+1):
4     for j in range(M): # this is to go through items
5         if items[j].weight <= bag_weight:
6             balance = bag_weight - items[j].weight
7             profit = item[j].profit + memo[balance]
8             if profit > memo[bag_weight]: # if we have new optimal
9                 memo[bag_weight] = profit
10 return memo[N]
```

# Kapsack Problem

## Unbounded

- But is the code correct?
  - Currently, we assume the maximum value is when you find items with a total weight of 12 kg

Eventually....



Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	40	80	120	160	350	390	430	470	550	700	740	780

```
1 memo = [0] * (N+1) # N is the total weight
2 memo[0] = 0
3 for bag_weight in range(1, N+1):
4     for j in range(M): # this is to go through items
5         if items[j].weight <= bag_weight:
6             balance = bag_weight - items[j].weight
7             profit = item[j].profit + memo[balance]
8             if profit > memo[bag_weight]: # if we have new optimal
9                 memo[bag_weight] = profit
10 return memo[N]
```

# Kapsack Problem

## Unbounded

- But is the code correct?
  - Currently, we assume the maximum value is when you find items with a total weight of 12 kg
  - What if we can't reach 12 kg? even bag is not up to 12kg, can still use 12kg list for that bag
  - What if the optimal is at 10 kg instead of 12 kg?

Eventually....



Weight	0	1	2	3	4	5	6	7	8	9	10	11	12
Profit	0	40	80	120	160	350	390	430	470	550	700	740	780

```
1 memo = [0] * (N+1) # N is the total weight
2 memo[0] = 0
3 for bag_weight in range(1,N+1):
4     for j in range(M): # this is to go through items
5         if items[j].weight <= bag_weight:
6             balance = bag_weight - items[j].weight
7             profit = item[j].profit + memo[balance]
8             if profit > memo[bag_weight]: # if we have new optimal
9                 memo[bag_weight] = profit
10 return memo[N]
```

# Kapsack Problem

## Unbounded

- But is the code correct?
  - Currently, we assume the maximum value is when you find items with a total weight of 12 kg
  - What if we can't reach 12 kg?
  - What if the optimal is at 10 kg instead of 12 kg?
  - So what must we change?

```
1 memo = [0] * (N+1) # N is the total weight
2 memo[0] = 0
3 for bag_weight in range(1, N+1):
4     for j in range(M): # this is to go through items
5         if items[j].weight <= bag_weight:
6             balance = bag_weight - items[j].weight
7             profit = item[j].profit + memo[balance]
8             if profit > memo[bag_weight]: # if we have new optimal
9                 memo[bag_weight] = profit
10 return memo[N]
```

# Kapsack Problem

## Unbounded

- But is the code correct?
  - Currently, we assume the maximum value is when you find items with a total weight of 12 kg
  - What if we can't reach 12 kg?
  - What if the optimal is at 10 kg instead of 12 kg?
  - So what must we change?
    - `memo[i] = memo[i-1]`      # copy the previous optimal

```
1 memo = [0] * (N+1) # N is the total weight
2 memo[0] = 0
3 for bag_weight in range(1, N+1):
4     for j in range(M): # this is to go through items
5         if items[j].weight <= bag_weight:
6             balance = bag_weight - items[j].weight
7             profit = item[j].profit + memo[balance]
8             if profit > memo[bag_weight]: # if we have new optimal
9                 memo[bag_weight] = profit
10 return memo[N]
```

# Kapsack Problem

## Unbounded

- But is the code correct?
  - Currently, we assume the maximum value is when you find items with a total weight of 12 kg
  - What if we can't reach 12 kg?
  - What if the optimal is at 10 kg instead of 12 kg?
  - So what must we change?
    - `memo[i] = memo[i-1]` # copy the previous optimal
    - or, linear search through `memo` for the maximum

```
1 memo = [0] * (N+1) # N is the total weight
2 memo[0] = 0
3 for bag_weight in range(1, N+1):
4     for j in range(M): # this is to go through items
5         if items[j].weight <= bag_weight:
6             balance = bag_weight - items[j].weight
7             profit = item[j].profit + memo[balance]
8             if profit > memo[bag_weight]: # if we have new optimal
9                 memo[bag_weight] = profit
10 return memo[N]
```

Questions?



# Kapsack Problem

0/1 items

- Same problem, but you can't repeat the item

# Kapsack Problem

0/1 items

- Same problem, but you can't repeat the item
  - So how would we solve it?

# Kapsack Problem

0/1 items

- Same problem, but you can't repeat the item
  - So how would we solve it?
- This is where we can see the growing of problems...

# Kapsack Problem

0/1 items

- Same problem, but you can't repeat the item
  - So how would we solve it?
- This is where we can see the growing of problems...
  - Grow from 0 weight to N weight

different weight capacity

# Kapsack Problem

0/1 items

- Same problem, but you can't repeat the item
  - So how would we solve it?
- This is where we can see the growing of problems...
  - Grow from 0 weight to N weight
  - Grow from a set of 0 items till M items

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

item 1, either pick or not pick

<b>Weight</b>	0	1	2	3	4	5	6	7	8	9	10	11	12
<b>Profit</b>	0	?	?	?	?	?	?	?	?	?	?	?	?

not decision make

# Kapsack Problem

0/1 items

- We use a **matrix** 2 constraint  
weight, item

Increasing weight





	0	1	2	3	4	5	6	7	8	9	10	11	12
0													
A													
B													
C													
D													

# Kapsack Problem

0/1 items

- We use a matrix

**Increasing weight** 

**Increasing items in set** 


	0	1	2	3	4	5	6	7	8	9	10	11	12
0													
A													
B													
C													
D													

# Kapsack Problem


0/1 items

- We use a matrix

**Increasing weight** weight capacity



**Increasing items in set**



	0	1	2	3	4	5	6	7	8	9	10	11	12
{}													
{A}													
{A,B}													
{A,B, C}													
{A,B, C,D}													



# Kapsack Problem

0/1 items

- We use a matrix
  - So we fill up the matrix

	0	1	2	3	4	5	6	7	8	9	10	11	12
0													
A													
B													
C													
D													

# Kapsack Problem

0/1 items

- We use a matrix
  - Base cases

	0	1	2	3	4	5	6	7	8	9	10	11	12
0													
A													
B													
C													
D													

# Kapsack Problem

0/1 items

- We use a matrix
  - Base cases
    - No item to choose from...

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A													
B													
C													
D													

# Kapsack Problem

0/1 items

- We use a matrix
  - Base cases
    - No item to choose from...
    - Max weight is 0....

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0												
B	0												
C	0												
D	0												

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0												
B	0												
C	0												
D	0												

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0												
B	0												
C	0												
D	0												

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0							
B	0												
C	0												
D	0												

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	?						
B	0												
C	0												
D	0												



# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230						
B	0												
C	0												
D	0												

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	?					
B	0												
C	0												
D	0												

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230					
B	0												
C	0												
D	0												

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230				
B	0												
C	0												
D	0												

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230			
B	0												
C	0												
D	0												

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0												
C	0												
D	0												

reuse from this row

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40											
C	0												
D	0												

# Kapsack Problem

0/1 items

all item just 1

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

just itemA

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40										
C	0												
D	0												



# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40							
C	0												
D	0												

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	40						
C	0												
D	0												

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	40	?					
C	0												
D	0												

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	40	270					
C	0												
D	0												

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	40	270	270				
C	0												
D	0												

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?
    - Is this correct?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	40	270	270	270	270	270	270
C	0												
D	0												

# Kapsack Problem

0/1 items

- We use a matrix

- So row by row...

- Start with item A

- Should we add it?

- Is this correct? Here, we can choose not to include B, having only A in bag

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	40	270	270	270	270	270	270
C	0												
D	0												

# Kapsack Problem

0/1 items

- We use a matrix

- So row by row...

- Start with item A

- Should we add it?

- Is this correct? Here, we can choose not to include B, having only A in bag

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0												
D	0												



# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

exclude item B

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	?											
D	0												

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	?											
D	0												

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40											
D	0												

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40	40	40	40								
D	0												

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40	40	40	40	350							
D	0												

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40	40	40	40	350	390						
D	0												

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40	40	40	40	350	390	390					
D	0												

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40	40	40	40	350	390	390	390	390	390	580	
D	0												



# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40	40	40	40	350	390	390	390	390	390	580	620
D	0												

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40	40	40	40	350	390	390	390	390	390	580	620
D	0	?											

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40	40	40	40	350	390	390	390	390	390	580	620
D	0	40	40	40	40	350	390	390	390				

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40	40	40	40	350	390	390	390	390	390	580	620
D	0	40	40	40	40	350	390	390	390	550			

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40	40	40	40	350	390	390	390	390	390	580	620
D	0	40	40	40	40	350	390	390	390	550	590		

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40	40	40	40	350	390	390	390	390	390	580	620
D	0	40	40	40	40	350	390	390	390	550	590	590	

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40	40	40	40	350	390	390	390	390	390	580	620
D	0	40	40	40	40	350	390	390	390	550	590	590	?

# Kapsack Problem

0/1 items

- We use a matrix
  - So row by row...
    - Start with item A
    - Should we add it?

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40	40	40	40	350	390	390	390	390	390	580	620
D	0	40	40	40	40	350	390	390	390	550	590	590	620



Questions?

# Kapsack Problem

0/1 items

- We use a matrix
  - What is the algorithm/ code?

j

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40	40	40	40	350	390	390	390	390	390	580	620
D	0	40	40	40	40	350	390	390	390	550	590	590	620

i

# Kapsack Problem

0/1 items

- We use a matrix
  - What is the algorithm/ code?

```
1  # for every row (item)
2  for i=1 to M:    if i = 1
3      # for every column (weight)
4      for j=1 to N:    j = 6
5          # get the excluded at current weight (row above)
6          exclude = memo[i-1][j]    current weight from the last row in the same column
7          # calculate the include
8          include = 0                j is total weight
9          if item[i].weight <= j:    available item's value    memo[1,0] = 0
10             include = item[i].value + memo[i-1][j-item[i].weight]
11             memo[i][j] = max(exclude, include)    third row for third item    6 - 6 (item A. weight) = 0
```

# Kapsack Problem

0/1 items

- We use a matrix
  - What is the algorithm/ code?
  - Complexity?

```
1  # for every row (item)
2  for i=1 to M:
3      # for every column (weight)
4      for j=1 to N:
5          # get the excluded at current weight (row above)
6          exclude = memo[i-1][j]
7          # calculate the include
8          include = 0
9          if item[i].weight <= j:
10             include = item[i].value + memo[i-1][j-item[i].weight]
11         memo[i][j] = max(exclude,include)
```

# Kapsack Problem

0/1 items

- We use a matrix
  - What is the algorithm/ code?
  - Complexity?
    - $O(NM)$  time from filling matrix
    - $O(NM)$  space for the matrix

```
1  # for every row (item)
2  for i=1 to M:
3      # for every column (weight)
4      for j=1 to N:
5          # get the excluded at current weight (row above)
6          exclude = memo[i-1][j]
7          # calculate the include
8          include = 0
9          if item[i].weight <= j:
10             include = item[i].value + memo[i-1][j-item[i].weight]
11             memo[i][j] = max(exclude, include)
```

choose the bigger one

# Kapsack Problem

0/1 items

- We use a matrix
  - What is the algorithm/ code?
  - Complexity?
    - $O(NM)$  time from filling matrix
    - $O(NM)$  space for the matrix... we can reduce this however

```
1  # for every row (item)
2  for i=1 to M:
3      # for every column (weight)
4      for j=1 to N:
5          # get the excluded at current weight (row above)
6          exclude = memo[i-1][j]
7          # calculate the include
8          include = 0
9          if item[i].weight <= j:
10             include = item[i].value + memo[i-1][j-item[i].weight]
11         memo[i][j] = max(exclude, include)
```

- We use a matrix
  - What is the algorithm/ code?
  - Complexity?
    - $O(NM)$  time from filling matrix
    - $O(NM)$  space for the matrix... we can reduce this however  
We realize we do not need the entire matrix! We get the current value by looking at the row above only. So we can just store the latest 2 row...

# Kapsack Problem

0/1 items

- We use a matrix
  - What is the algorithm/ code?
  - Complexity?
    - $O(NM)$  time from filling matrix
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We realize we do **not need the entire matrix!** We **get** the **current value** by **looking at the row above only**. So we can just **store the latest 2 row...**  
**Reducing complexity to  $O(2N+M)$**  just a list for items



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  - What is the algorithm/ code?
  - Complexity?
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Reducing complexity to  $O(2N+M)$
    - But in reality, we can't do this space saving... because we need it to reconstruct the solution...

Questions?

Take a break!

# Kapsack Problem

0/1 items

- So what are the items?

# Kapsack Problem

0/1 items

- So what are the items?

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40	40	40	40	350	390	390	390	390	390	580	620
D	0	40	40	40	40	350	390	390	390	550	590	590	620

# Kapsack Problem

0/1 items

- So what are the items?
  - Recall that we compare the current value (inclusive) with the value in the row above (exclusive).

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40	40	40	40	350	390	390	390	390	390	580	620
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A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40	40	40	40	350	390	390	390	390	390	580	620
D	0	40	40	40	40	350	390	390	390	550	590	590	620

# Kapsack Problem

0/1 items

- So what are the items?
  - Recall that we compare the current value (inclusive) with the value in the row above (exclusive). If **same value**, means we **do not include**...

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
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0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40	40	40	40	350	390	390	390	390	390	580	620
D	0	40	40	40	40	350	390	390	390	550	590	590	620

■ {}

# Kapsack Problem

0/1 items

- So what are the items?
  - Recall that we compare the current value (inclusive) with the value in the row above (exclusive). If **different value**, means we **include**...

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40	40	40	40	350	390	390	390	390	390	580	620
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- {C}

# Kapsack Problem

0/1 items

- So what are the items?
  - Recall that we compare the current value (inclusive) with the value in the row above (exclusive). If **different value**, means we **include**...
  - Then we update to the suitable weight of the included item

	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40	40	40	40	350	390	390	390	390	390	580	620
D	0	40	40	40	40	350	390	390	390	550	590	590	620

- {C} just leave item C in the bag if the value is not increased

# Kapsack Problem

0/1 items

Item	A	B	C	D
Weight	6kg	1kg	5kg	9kg
Value	\$230	\$40	\$350	\$550

## ■ So what are the items?

- Recall that we compare the current value (inclusive) with the value in the row above (exclusive). If **different value**, means we **include**...
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	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
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A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40	40	40	40	350	390	390	390	390	390	580	620
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	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40	40	40	40	350	390	390	390	390	390	580	620
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0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40	40	40	40	350	390	390	390	390	390	580	620
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## ■ {C,B}

12 - 5 = 7 (column 7)

since unique item, so no longer can choose C and D, so A and B  
if row A and row B have the same values, then just includes item A

# Kapsack Problem

0/1 items

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  - Recall that we compare the current value (inclusive) with the value in the row above (exclusive). So now do the same...

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A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40	40	40	40	350	390	390	390	390	390	580	620
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- {C,B}



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0/1 items

Item	A	B	C	D
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	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40	40	40	40	350	390	390	390	390	390	580	620
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- {C,B}

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	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40	40	40	40	350	390	390	390	390	390	580	620
D	0	40	40	40	40	350	390	390	390	550	590	590	620

- {C,B,A}

# Kapsack Problem

0/1 items

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	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
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- {C,B,A} makes up the item for total value of 620

# Kapsack Problem

0/1 items

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	0	1	2	3	4	5	6	7	8	9	10	11	12
0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	0	0	0	230	230	230	230	230	230	230
B	0	40	40	40	40	40	230	270	270	270	270	270	270
C	0	40	40	40	40	350	390	390	390	390	390	580	620
D	0	40	40	40	40	350	390	390	390	550	590	590	620

- {C,B,A} makes up the item for total value of 620
- This is what we call **backtracking**!

Questions?

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- We often need to reconstruct solutions
  - Coin change = what are the coins?
  - Knapsack = what are the items?
  - Edit distance = what are the modifications made (insert/delete/replace)
  - When we update the optimal value, we can store the decision we made... But this **waste a lot of memory** as we need to store decisions/combinations at every step!
  - So, we leave **bread crumbs to backtrack** or only the final **decision made!**

Questions?

# Backtracking

## Coin change

- Let say we store our decisions...

Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	2	3	4	1	1	2	3	1	2	2	2

# Backtracking

## Coin change

- Let say we store our decisions...

Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	2	3	4	1	1	2	3	1	2	2	2
The coins	{}	1	1, 1	1, 1, 1	1, 1, 1, 1	5	6	6, 1	6, 1, 1	9	5, 5	6, 5	6, 6



# Backtracking

## Coin change

- Let say we store our decisions...
  - Space complexity?

	0	1	2	3	4	5	6	7	8	9	10	11	12
Value													
Number of coins	0	1	2	3	4	1	1	2	3	1	2	2	2
The coins	{}	1	1, 1	1, 1, 1	1, 1, 1, 1	5	6	6, 1	6, 1, 1	9	5, 5	6, 5	6, 6

# Backtracking

## Coin change

- Let say we store our decisions...
  - Space complexity? Can be  $O(N^2)$

if just 1, 10 coin

$O(N) * O(N) = O(N^2)$

	0	1	2	3	4	5	6	7	8	9	10	11	12
Value	0	1	2	3	4	1	1	2	3	1	2	2	2
Number of coins	0	1	2	3	4	1	1	2	3	1	2	2	2
The coins	{}	1	1, 1	1, 1, 1	1, 1, 1, 1	5	6	6, 1	6, 1, 1	9	5, 5	6, 5	6, 6

$1,1,1,1,1 = O(N)$

# Backtracking

## Coin change

- Let say we store our decisions...
  - Space complexity? Can be  $O(N^2)$
  - Improve it further?

	0	1	2	3	4	5	6	7	8	9	10	11	12
Value	0	1	2	3	4	1	1	2	3	1	2	2	2
Number of coins	0	1	2	3	4	1	1	2	3	1	2	2	2
The coins	{}	1	1, 1	1, 1, 1	1, 1, 1, 1	5	6	6, 1	6, 1, 1	9	5, 5	6, 5	6, 6

# Backtracking

## Coin change

- Let say we store our decisions...
  - Space complexity? Can be  $O(N^2)$
  - Improve it further? Remember the last coin you added

Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	2	3	4	1	1	2	3	1	2	2	2
The coins	{}	1	1	1	1	5	6	1	1	9	5	6	6

# Backtracking

## Coin change

- Let say we store our decisions...
  - Space complexity? Can be  $O(N^2)$
  - Improve it further? Remember the last coin you added  
So space complexity now?

	0	1	2	3	4	5	6	7	8	9	10	11	12
Value													
Number of coins	0	1	2	3	4	1	1	2	3	1	2	2	2
The coins	{}	1	1	1	1	5	6	1	1	9	5	6	6

$5 - 5 = 0$   
 $\{\}$  null set coin

$6 - 6 = 0$

# Backtracking

## Coin change

- Let say we store our decisions...
  - Space complexity? Can be  $O(N^2)$
  - Improve it further? Remember the last coin you added  
So space complexity now?  $O(N)$

Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	2	3	4	1	1	2	3	1	2	2	2
The coins	{}	1	1	1	1	5	6	1	1	9	5	6	6

# Backtracking

## Coin change

- Let say we store our decisions...
  - Space complexity? Can be  $O(N^2)$
  - Improve it further? Remember the last coin you added  
So space complexity now?  $O(N)$   
Coins = {6,6}

backtracking always go back to base case

12 - 6 = 6 (go to column 6)

Value	0	1	2	3	4	5	6	7	8	9	10	11	12
Number of coins	0	1	2	3	4	1	1	2	3	1	2	2	2
The coins	{}	1	1	1	1	5	6	1	1	9	5	6	6

last item added

Questions?



# Backtracking

## vs Decision Array

- Backtracking save space
- Decision array save time

# Backtracking vs Decision Array

- Backtracking save space
  - Less auxiliary space not saving decision made but save the result and backtrack to compare to retrieve back the decision
  - Same space complexity
- Decision array save time

# Backtracking

## vs Decision Array

- Backtracking save space
  - Less auxiliary space
  - Same space complexity
- Decision array save time
  - Faster
  - But time complexity lies in finding the solution still...

Questions?

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- We will skip this since this is similar to the Knapsack really...

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  - Will be covered in the tutorial, linking it up with the longest common subsequence (LCS) problem there

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  - Will be covered in the tutorial, linking it up with the longest common subsequence (LCS) problem there
  - Note: Nathan's slide at the end

watch Ian W5 tutorial

Questions?



# Summary

## Dynamic Programming Algorithm (DPA)

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- Break it down

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- Take problem
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Thank You