

## ● Q1:

- Admissible heuristics:
  - It's a concept in search algorithms where a heuristic is considered admissible if it never overestimates the cost to reach the goal from any node in the search space.
- "h=0" as a heuristics for a problem: it's admissible
  - "h=0" means there is no estimated cost from the current node to the goal because the estimate is always 0. It's admissible since the "h=0" always either underestimates or equals the true cost to reach the goal and never overestimates.
- I can't say the "h=k" heuristics (where k is a constant) is admissible.
  - One requirement of admissibility is the heuristics can't overestimate the true cost from the current node to the goal. If "h=k" and say k=1 and that is fixed and not changing if the cost of reaching the goal from a node is less, then this heuristics will overestimate.
- $h = \min(h_1, h_2, h_3)$  is admissible:
  - If there is one of them that is admissible, then  $h_{\min} \leq$  any of the 3  $h_1, h_2, h_3$ . That means either it will underestimate or calculate the true cost but never overestimate.
- $h = \max(h_1, h_2, h_3)$  is not admissible
  - Unless  $h_{\max} =$  the heuristics that is admissible, it will overestimate the true cost and thus will be inadmissible.

## ● Q2:

- Link:  
<https://github.com/Initiated0/CSCE580-Fall2024-nayeem-Repo/tree/main/Quizzes/Quiz2>

## ● Q3:

- CSP
  - Variables: T, W, O, F, U, R; the characters that need to be mapped.
  - Domains: Each variable (T, W, O, F, U, R) has the domain  $D=\{0,1,2,3,4,5,6,7,8,9\}$
  - Constraints:
    - Uniqueness  $T \neq W \neq O \neq F \neq U \neq R$
    - $TWO + TWO = FOUR$  ; this sum must hold.
      - $2 \times (100T + 10W + 1O) = 1000F + 100O + 10U + 1R \dots \dots$   
 $\dots \dots \dots \# 1$
    - Range constraint: Domains must be between 0-9
    - In the sum leading letters cant be 0 ( $T \neq 0$ ) ( $F \neq 0$ )

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- Solving with node, arc, and path consistency

- Node consistency:

- Initially, all variables had  $D = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$  but we realized that we can't have  $(T \neq 0) \wedge (F \neq 0)$ .
    - So, for  $D_T = \{1, 2, 3, 4, 5, 6, 7, 8, 9\}$  and  $D_F = \{1, 2, 3, 4, 5, 6, 7, 8, 9\}$

- Arc consistency:

- We must ensure there is consistency and no overlapping in variable's domain for every possible pairs of variables.
    - If  $T \neq W \neq O \neq F \neq U \neq R$ , then  $D_T \neq D_W \neq D_O \neq D_F \neq D_U \neq D_R$

- Path consistency:

- Path consistency ensures that any value assignment for one variable is consistent with value assignments for two other variables along the path.

- Solution:

- Node and arc consistency reduces the search space by decreasing the domains of the variables.
    - Next is checking possible combinations of values for the variables T,W,O,F,U,R to satisfy the equation # 1

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