What is the decimal form of the 4-bit, signed, 2's complement number 1101?

A: **-4** 

B: **-2** 

C: 1

D: -3

E: Something else

What is the decimal form of the 8-bit, signed, 2's complement number 4f?

A: 81

B: **79** 

C: **55** 

D: **-77** 

E: Something else

 $\begin{array}{r} 01100101 \\ + 00100010 \\ \hline 10000111 \end{array}$ 

01100101 Operand A

00100010 Operand B

This must run somehow on a physical machine!

10000111 Result

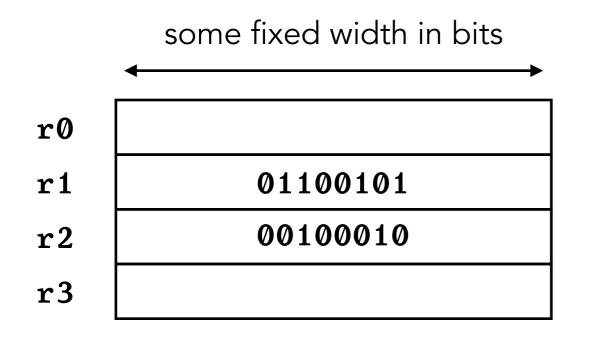
Where (physically) are Op A/B, the Result, and the Instruction?

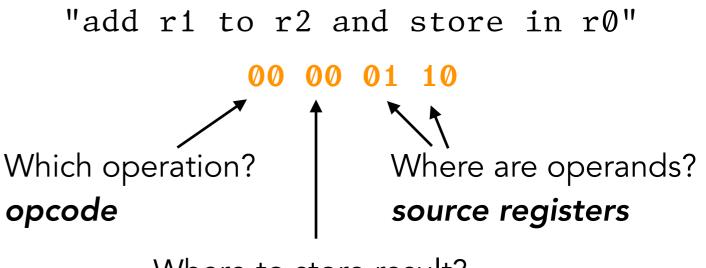
And the *Operation* had better be encoded in binary!

+ Operation

## Registers

#### Instructions





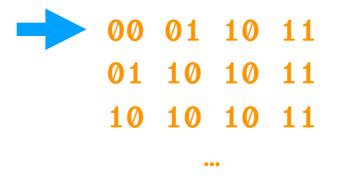
Where to store result? destination register

## To run a program, the machine uses

The content of **Registers** 

r0
r1 01100101
r2 00100010
r3

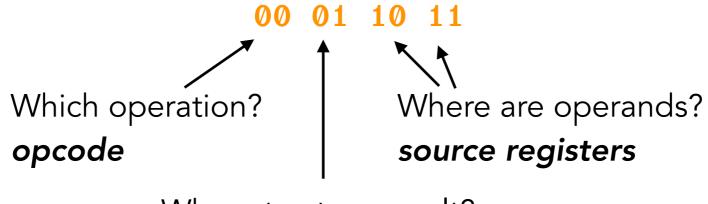
A list of *Instructions*, with a *Current Instruction* 



The rules for the **Instruction Set** 

- 1. If the opcode is 00, add the contents of the operand registers, store the result in the destination register.
- 2. If the opcode is 01, ...

The machine runs each instruction in order



Where to store result? destination register

00001000 (8)	00100000 (32)
00000110 (6)	
00000010 (2)	
00000100 (4)	

```
00 00 01 10 "put r1 + r2 in r0"
01 11 01 10 "put r1 - r2 in r3"
10 00 00 11 "put r0 * r3 in r0"
```

r0

r1

r2

r3

## Instruction Set

- 1. If the opcode is 00, **add** the contents of the source registers, store the result in the destination register.
- 2. If the opcode is 01, **subtract** the contents of the second source from the first, store the result in the destination register.
- 3. If the opcode is 10, **multiply** the contents of the source registers, store the result in the destination register.

# Instructions (The Program!)

```
r0 00001111 (15)
r1 00010100 (20)
r2 00000101 (5)
r3
```

```
00 11 00 01
00 11 11 10
```

#### What is in **r3** after this runs?

A: 00100011 (35)

B: 00010100 (20)

 $\mathbb{C}$ : 00101000 (40)

D: 00011001 (25)

E: Something else

- 1. If the opcode is 00, **add** the contents of the source registers, store the result in the destination register.
- 2. If the opcode is 01, **subtract** the contents of the second source from the first, store the result in the destination register.
- 3. If the opcode is 10, **multiply** the contents of the source registers, store the result in the destination register.

```
r0 00001111 (15)
r1 00010100 (20)
r2 00000101 (5)
r3
```

```
01 11 01 0000 11 11 1010 11 11 10
```

#### What is in **r3** after this runs?

A: 01100100 (100)

B: 00010100 (20)

 $\mathbb{C}$ : 00101000 (40)

D: 00110010 (50)

E: Something else

- 1. If the opcode is 00, **add** the contents of the source registers, store the result in the destination register.
- 2. If the opcode is 01, **subtract** the contents of the second source from the first, store the result in the destination register.
- 3. If the opcode is 10, **multiply** the contents of the source registers, store the result in the destination register.

```
00000010 (2)

00010100 (20)

00000101 (5)
```

r0

r1

r2

r3

# Which of these programs ends with 60 in **r3**?

- If the opcode is 00, add the contents of the source registers, store the result in the destination register.
- 2. If the opcode is 01, **subtract** the contents of the second source from the first, store the result in the destination register.
- 3. If the opcode is 10, **multiply** the contents of the source registers, store the result in the destination register.



```
11 00 00 10
11 01 01 00
00 10 01 00
```

#### What is in **r2** after this runs?

A: 00000110 (6)

B: 00000000 (0)

C: 00001000 (8)

D: 00010110 (22)

E: Something else

- 1. If the opcode is 00, **add** the contents of the source registers, store the result in the destination register.
- 2. If the opcode is 01, **subtract** the contents of the second source from the first, store the result in the destination register.
- 3. If the opcode is 10, **multiply** the contents of the source registers, store the result in the destination register.
- 4. If the opcode is 11, interpret the last 4 bits as a constant, and move them into the destination register

```
00 Rd Rn Rmput Rn + Rm in Rd01 Rd Rn Rmput Rn - Rm in Rd10 Rd Rn Rmput Rn * Rm in Rd11 Rd Imm4put Imm4 in Rd
```

## ARM: (just a few) more opcodes and options

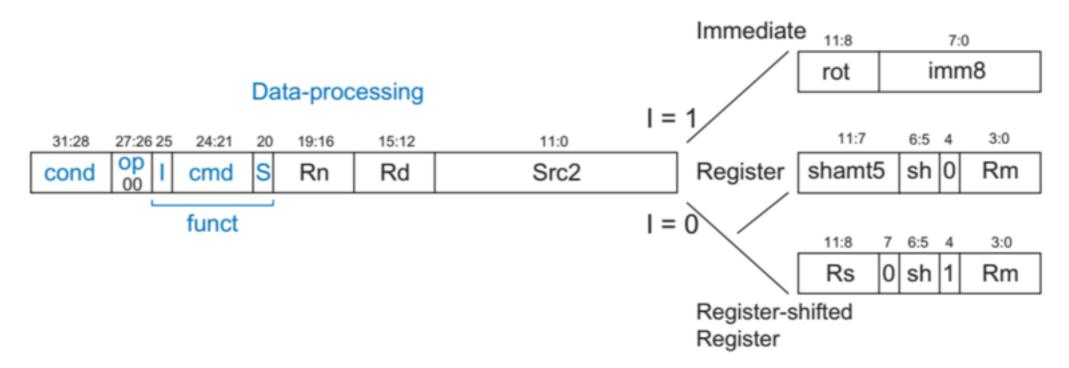


Figure B.1 Data-processing instruction encodings

Some definitions:

word/word size: The number of bits in an instruction/register. In this lecture, we had 8-bit words. Our picluster ARM will use 32-bit words.

**opcode:** The part of the instruction that determines which operation to perform. Our pi-cluster ARM will have many more opcodes (4 bits and more!)

**instruction:** A word that the processor interprets to perform a step of computation. Usually this means changing the value in some register. There are *lots* of instructions we'll talk about this quarter.

**register:** A word of quick-access memory, where most computation happens. Data is stored here briefly before being stored elsewhere. Our pi-cluster ARM uses 16 registers.

**immediate value:** A constant number that appears as part of an instruction.

- 1. If the opcode is 00, **add** the contents of the source registers, store the result in the destination register.
- 2. If the opcode is 01, **subtract** the contents of the second source from the first, store the result in the destination register.
- 3. If the opcode is 10, **multiply** the contents of the source registers, store the result in the destination register.
- 4. If the opcode is 11, interpret the last 4 bits as a constant, and move them into the destination register