Name: John Moi

10: 1801187-

Problem 1: Custom Hash

Suppose we created the following object for an application:

```
struct BadMovie {
    int databaseID;
    string name;
    string director;
    int runtimeInSeconds;
    int rating;
};
```

We decided that a Hash Table is the most appropriate data structure for our purposes. However, in our experiments, hashing on the databaseID alone results in many collisions. So, we decided to write a custom hashing function that incorporates all of the member variables to determine the appropriate bucket. Complete such a hash function below, you may assume that the constant ARRAY_SIZE is defined.

```
int hash(const BadMovie &bm) {

int total = 0;

for (int z=0) it bm. name, length(); itt) {

total = total + CT+1) * bm. name[i];

}

for (int z=0; it bm. director.longth(); it+) {

total = total + (T+1) * bm. director[i];

}

total = total + bm. diabase[D;

total = total + bm. nutime In Seconds;

total = total + ton. nating;

beturn total 1/6 bm. APRAT_SIZE;
```

Name: Toto (ho

ID: 18017817

Problem 2: Hash Tables

Consider a Closed Hash Table with HASH SIZE of 10:

```
class ClosedHashTable {
                                                                              63
                                                                          0
public:
       void insert(int key){
                                                                           1
               int bucket = hashFunc(key);
                                                                               ١
                                                                           2
               for (int tries = 0; tries<HASH SIZE; tries++) {</pre>
                       if (hash_array[bucket] == EMPTY) {
                                                                           3
                              hash_array[bucket] = key;
                                                                           4
                              return:
                                                                           5
                                                                              37.
                       bucket = (bucket + 1) % HASH SIZE;
                                                                           6
                                                                              23
               }
                                                                           7
                                                                              53
private:
                                                                           8
                                                                               11/4
       int hashFunc(int x) {
               return (x * 2) % HASH SIZE;
                                                                               19
                                                                           9
       int hash_array[HASH_SIZE];
};
ClosedHashTable ch;
ch.insert(7); 14% (0 = 4
                  1 % 10=2.
ch.insert(1);
ch.insert(23); 46 \epsilon_{10} = 6.
ch.insert(14); 38\% (0 = 8 ch.insert(19); 38\% (0 = 8
ch.insert(53); 106\% 6 = 6.
ch.insert(37); ηψ<sub>70</sub> 10 = Ψ.
ch.insert(83); 166% to = 6
```

a. Show result of the insert commands in the array to the right. Using the state of the hash table after all the inserts, what is the load factor and average number of checks for this hash table?

b. How much larger would we have to make the hash table if we wanted to have average number of checks to be roughly 1.10?

So. Assuming Average Number is 1.10.

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$$3 | 1.20 | = 1.20 | = 1.$$

$$3 | 1.20 | = 0.2$$

$$3 | 1.20 | = 0.2$$

$$4 | 1.20 | = 0.2$$

$$\frac{t \cdot b}{8} = 6$$

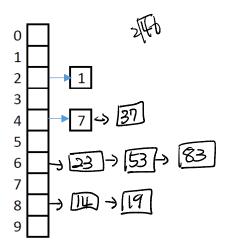
$$= t \cdot b = 48$$

$$total budget = 48$$

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Name: Inho choi

c. Suppose we, instead, inserted those values into an **open hash table** using the same hash function. What would the state of the table look like? (This first two have been done for you.)



d. What is the average number of checks for this open hash table given its state after all the inserts? Open hash table average number of checks = $1+ L_2$.

$$= 1+\frac{2}{5} = \frac{7}{5} = 1.40$$

e. Discuss the quality of this hash function.

This hugh function's load L is 0.8 so this function away has 20% more buckers than you need (You will fill 60% of the buckers).

And open has function is more efficient than closed high table.