

Pushpin (maybe I got the game name wrong)

Proofs

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For the following lemmas, I'm considering all games in which left has a winning strategy (values > 0). By swapping \underline{L} and \underline{R} , you can get the same results for \underline{R} and thus for all valid positions of the game.

Notation:

- Each literal symbol will be underlined (\underline{L} or \underline{R})
- \mathbb{X}_n represents, without loss of generality any sequence of characters or spaces of length n . \mathbb{X}'_n represents the sequence \mathbb{X} with the rightmost space removed, if one exists. If a rightmost space does not exist, then $\mathbb{X}_n = \mathbb{X}'_n$
- Upper case symbols (like G , A , and B) represent sequences/game states.
- Greek symbols will be used to denote a single character, that can either be \underline{L} , \underline{R} , $_$.
- The entire game state will be completely underlined (e.g. $\underline{\underline{LLR}}$)
- “+” means concatenation of a sequence

The evaluation of a sequence will be denoted like so, where G is the current game state and the number preceding the colon is the square being pushed:

$$G := \underline{\underline{LLR}}$$

$$G_0 : \underline{\underline{LLR}} \rightarrow \underline{\underline{LR}}$$

$$G_1 : \underline{\underline{LLR}} \rightarrow \underline{\underline{LR}}$$

$$G_2 : \underline{\underline{LLR}} \rightarrow \underline{\underline{LR}}$$

The game is zero-indexed.

Different game positions can be compared like so:

$$G > G_1 \geq G_2$$

Lemmas

Accessibility Lemma (done)

If one game position Q is accessible from another game position P , ending in \underline{L} , by a series of moves, then $P > Q$.

Rationale: Whenever a move is played, by \underline{L} or \underline{R} , the value of the game is always closer to zero.

Reduction Lemma (done)

If \mathbb{X}_n has spaces, then:

$$\mathbb{X}_n > \mathbb{X}'_n$$

If \mathbb{X}_n has no spaces, then:

$$\mathbb{X}_n = \mathbb{X}'_n$$

Same-Length Concat Lemma

If $P > Q$ and P and Q are of the same length and A ends in \underline{L} , then $P + A > Q + A$. Proof.

Starting from the rightmost square, if $P > Q$ then $P \neq Q$ and P and Q must differ at a location n . Let $P = P_1 + \underline{\alpha} + A$ and $Q = Q_1 + \underline{\beta} + A$ where β is the symbol at the location n where P and Q differ. By construction, $\underline{\alpha} \neq \underline{\beta}$. Suppose for contradiction that $\underline{\alpha} < \underline{\beta}$. Then, Suppose P is positive. Q is either positive, zero, or negative. If zero, done .

General Prepend Lemma (done)

If A ends in \underline{L} , then for any characters $\underline{\alpha}$ and $\underline{\beta}$ such that $\underline{\alpha} > \underline{\beta}$, then $P := \underline{\underline{\alpha}} + \underline{A} > Q := \underline{\underline{\beta}} + \underline{A}$. Proof.

There are three cases to examine:

- $\underline{\alpha} = \underline{L}$, $\underline{\beta} = \underline{_}$ In this case, $P_0 : \underline{\underline{\alpha}} + \underline{A} \rightarrow \underline{\underline{\beta}} + \underline{A}$, so $P > Q$ by the accessibility lemma.

- $\underline{\alpha} = \underline{_}$, $\underline{\beta} = \underline{R}$ In this case, Q has more \underline{R} values than P and the same number and location of \underline{L} values. Since right has more moves available in Q than in P , $P > Q$ (TODO: better justification).
- $\underline{\alpha} = \underline{L}$, $\underline{\beta} = \underline{R}$ From the first case and the second case and the transitive property, $P > Q$ follows.

In all three cases, $P > Q$, so necessarily, $P > Q$.

Append Lemma (unproven)

For any sequence of characters (empty or not), \mathbb{X}_n :

$$\underline{\underline{\mathbb{X}_n \underline{L}}} > \underline{\underline{\mathbb{X}_n}}$$

Prepend Lemma

For any sequence of characters (empty or not), \mathbb{X}_n :

$$P := \underline{\underline{\mathbb{X}_n \underline{L}}} > Q := \underline{\underline{\mathbb{X}_n \underline{L}}}$$

Proof.

Suppose that \mathbb{X}_n does not contain spaces. Then,

$$P_n : \underline{\underline{\mathbb{X}_n \underline{L}}} \rightarrow \underline{\underline{\mathbb{X}_n \underline{L}}}$$

So, $P_n = Q$ and Q is accessible from P . Therefore, by the one-step accessibility lemma, $P > Q$.

Suppose that \mathbb{X}_n does contain spaces. Then, $\underline{\underline{\mathbb{X}_n \underline{L}}} \rightarrow \underline{\underline{\mathbb{X}'_n \underline{L}}}$. Without loss of generality, split \mathbb{X}_n by its rightmost space. Define $\mathbb{X}_n = A + _ + B$ where \mathbb{X}_b contains no spaces. Then, $\mathbb{X}'_n = A + B$

\mathbb{X}_n either contains or does not contain spaces, so by disjunction elimination, the lemma holds.

Scratch (problematic because it depends on the concat lemma): Suppose that \mathbb{X}_n does contain spaces. Then, $\underline{\underline{\mathbb{X}_n \underline{L}}} \rightarrow \underline{\underline{\mathbb{X}'_n \underline{L}}}$. By the reduction lemma, since the sequence \mathbb{X}_n has spaces, $\mathbb{X}_n > \mathbb{X}'_n$.

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