

A hand is shown placing a blue L-shaped block onto a colorful geometric structure made of various blocks. The structure is composed of blocks in shades of blue, orange, yellow, purple, and pink. The background is a solid light blue, and the surface is a light-colored wooden table. Several other blocks are scattered on the table in the foreground, including a green L-shaped block, a blue L-shaped block, a red L-shaped block, and a yellow L-shaped block.

# EMTH403

Mathematical Foundation  
for Computer Science

Nitin K. Mishra (Ph.D.)

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Associate Professor

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# Lecture Outcomes



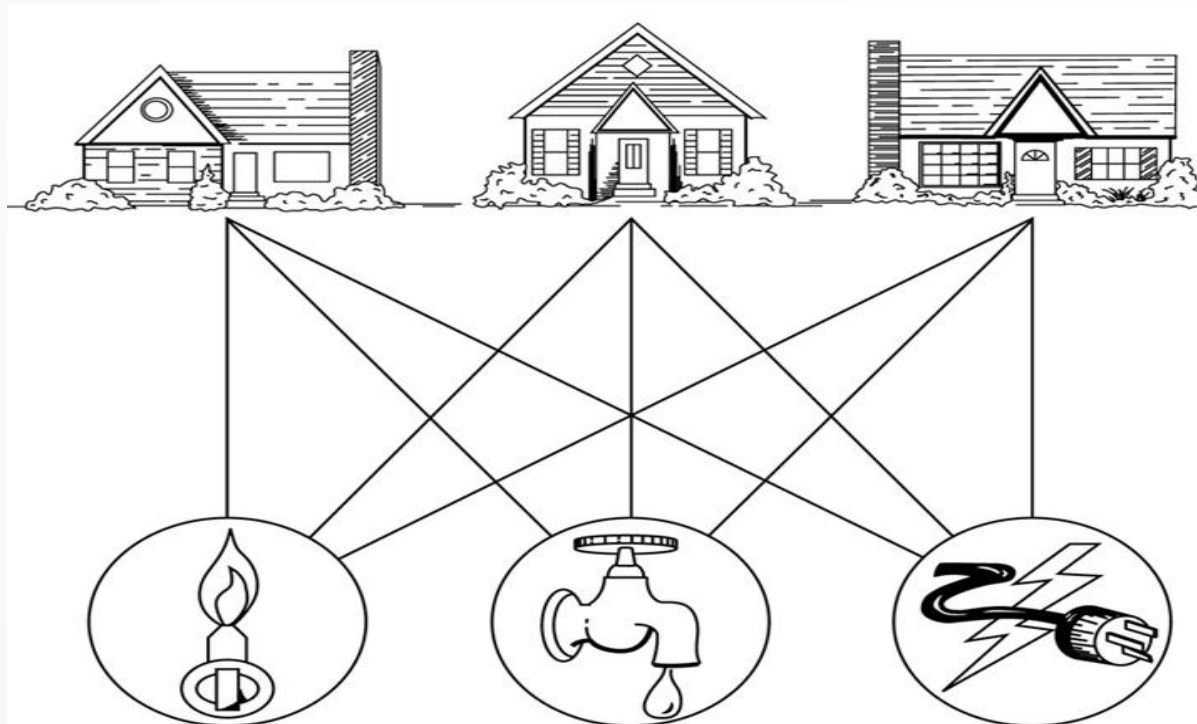
After this lecture, you will be able to

- Understand what are Planar Graphs
- Understand how  $K_{3,3}$  is a non-planar Graph.

# Planar Graphs

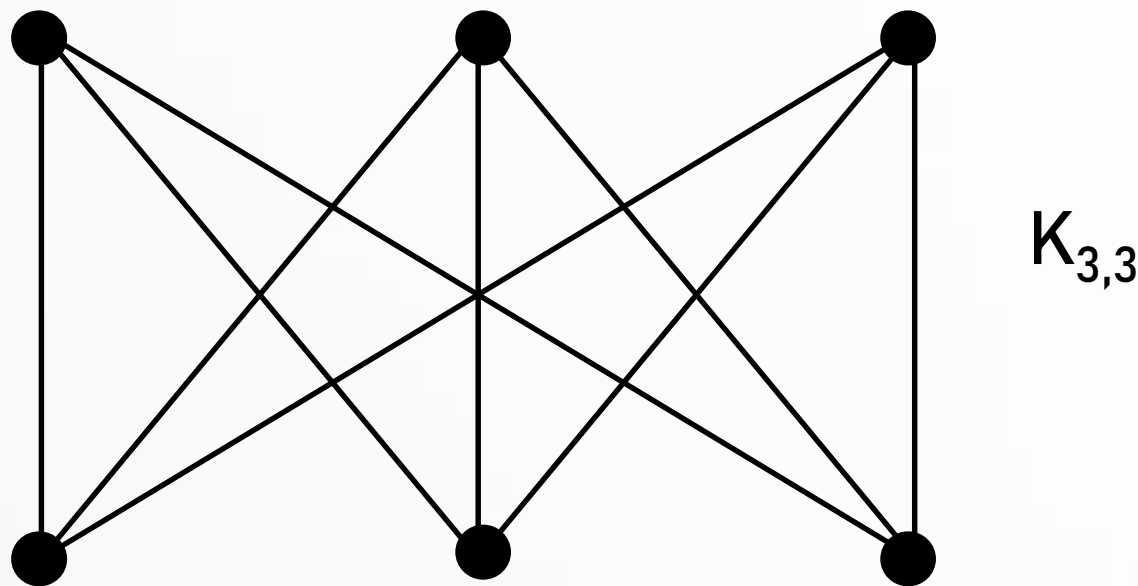
## The House-and-Utilities Problem

Is it possible to join the three houses to the three utilities in such a way that none of the connections cross?



# Planar Graphs

Phrased another way, this question is equivalent to:  
Given the complete bipartite graph  $K_{3,3}$ , can  $K_{3,3}$  be drawn in the plane so that no two of its edges cross?



# Planar Graphs

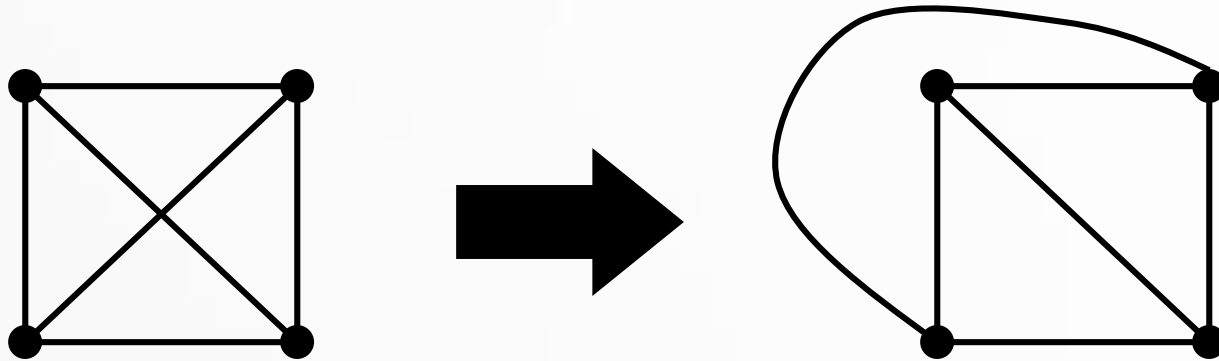
A graph is called planar if it can be drawn in the plane without any edges crossing.

A crossing of edges is the intersection of the lines or arcs representing them at a point other than their common endpoint.

Such a drawing is called a planar representation of the graph.

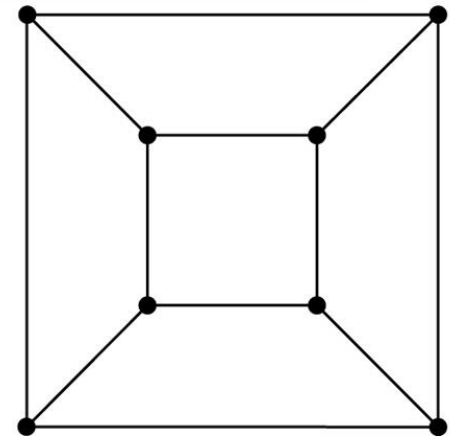
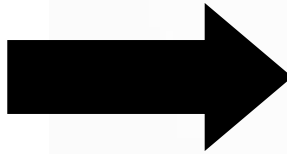
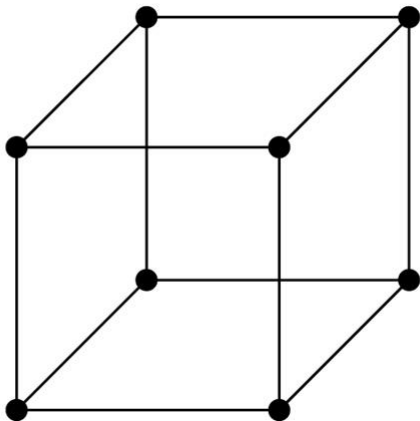
# Planar Graphs

A graph may be planar even if it is usually drawn with crossings, since it may be possible to draw it in another way without crossings.



# Planar Graphs

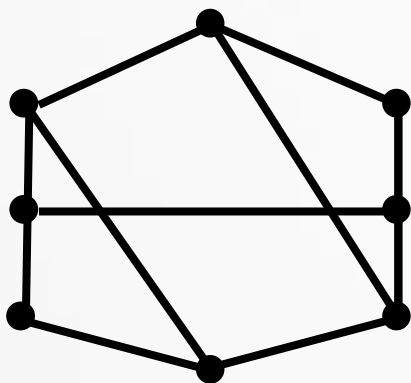
A graph may be planar even if it represents a 3-dimensional object.



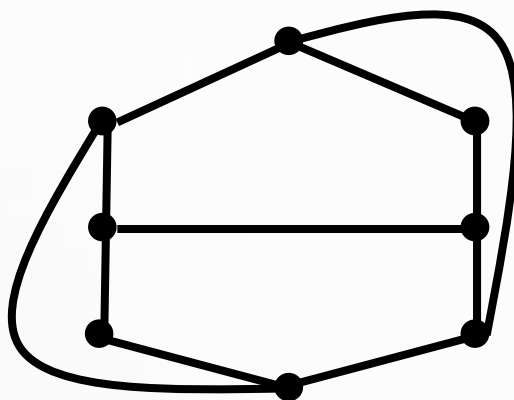
# Planar Graphs

## Definition:

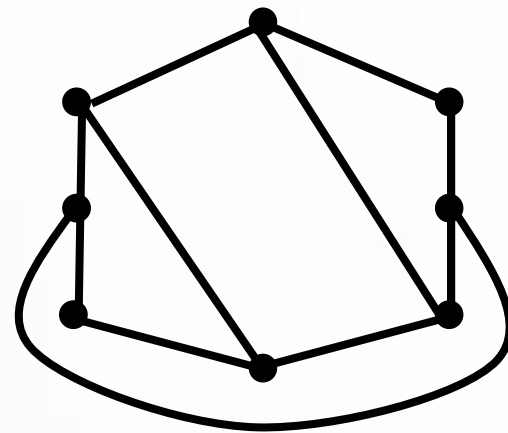
A planar graph  $G$  that is drawn in the plane so that no two edges intersect (that is,  $G$  is embedded in the plane) is called a plane graph.



(a) planar,  
not a plane graph



(b) a plane  
graph

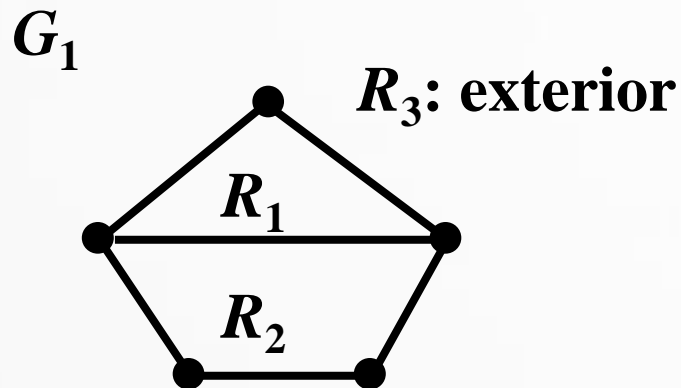


(c) another  
plane graph



# Planar Graphs

Note. A given planar graph can give rise to several different plane graph.

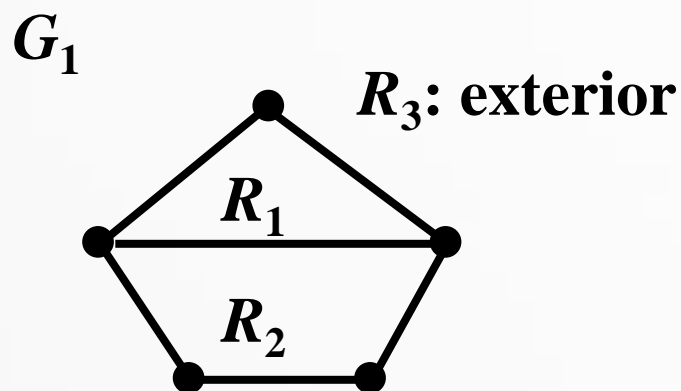


$G_1$  has 3 regions.

# Planar Graphs

## Definition:

Let  $G$  be a plane graph. The connected pieces of the plane that remain when the vertices and edges of  $G$  are removed are called the regions of  $G$ .



$G_1$  has 3 regions.

# Planar Graphs

We can prove that a particular graph is planar by showing how it can be drawn without any crossings.

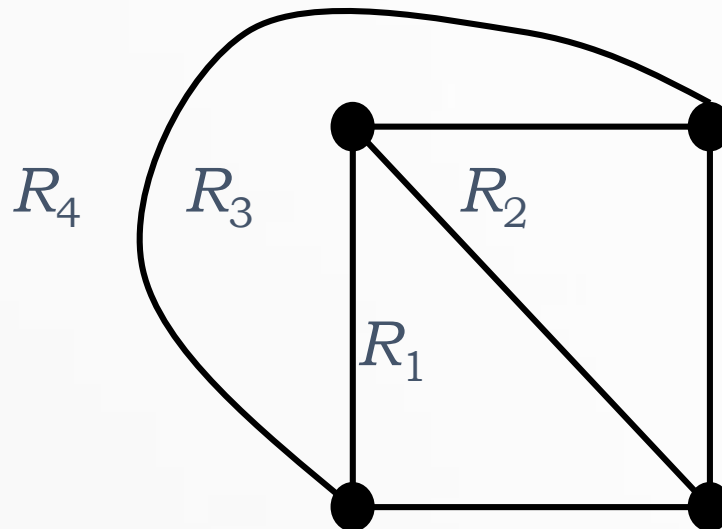
However, not all graphs are planar.

It may be difficult to show that a graph is nonplanar.

We would have to show that there is *no way* to draw the graph without any edges crossing.

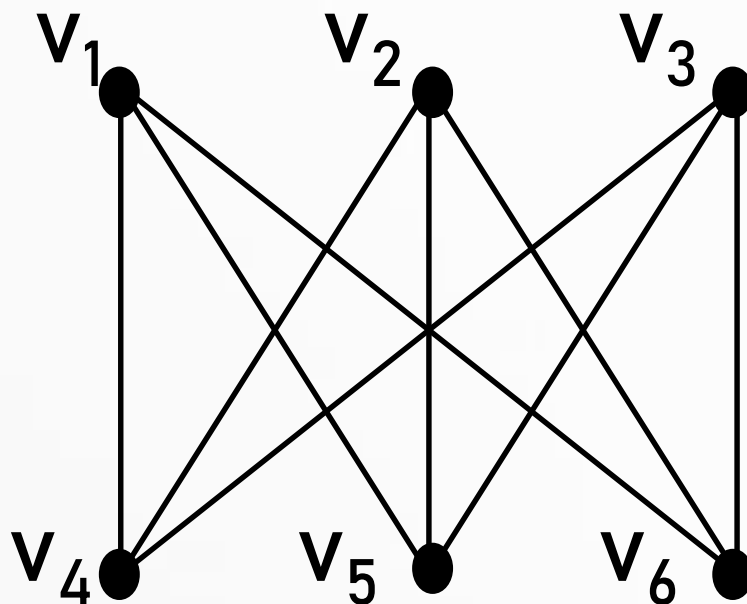
# Regions

Euler showed that all planar representations of a graph split the plane into the same number of *regions*, including an unbounded region.



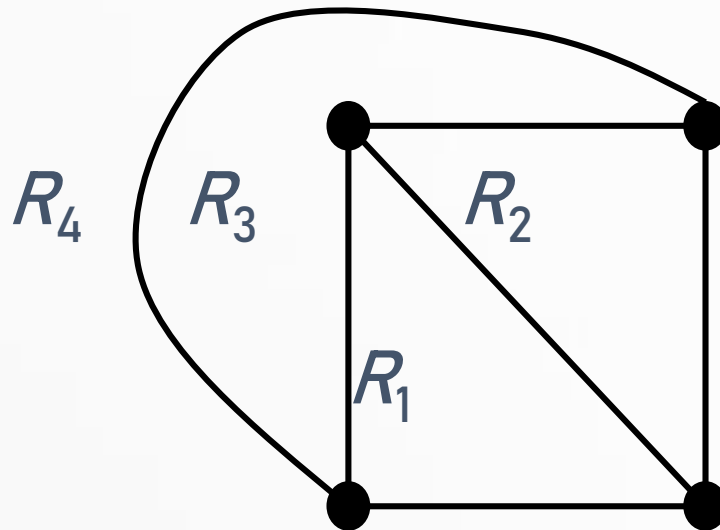
# Planar Graphs

In any planar representation of  $K_{3,3}$ , vertex  $v_1$  must be connected to both  $v_4$  and  $v_5$ , and  $v_2$  also must be connected to both  $v_4$  and  $v_5$ .



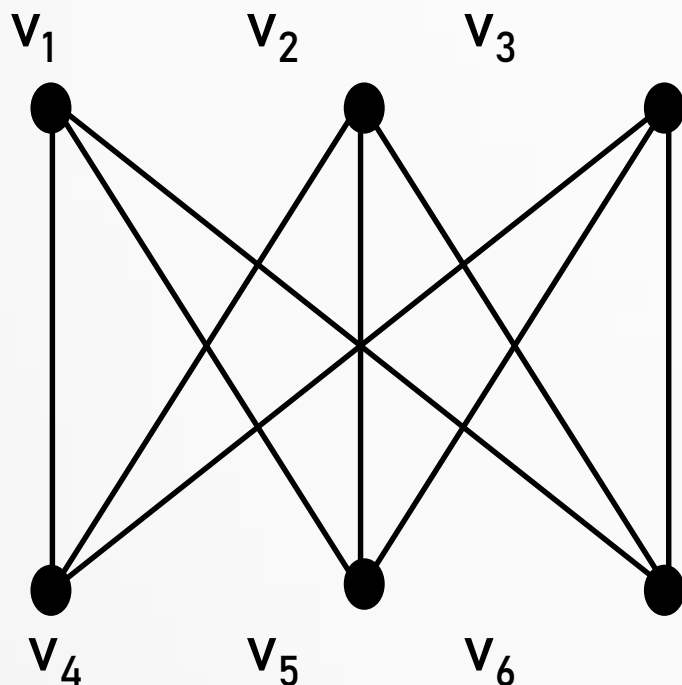
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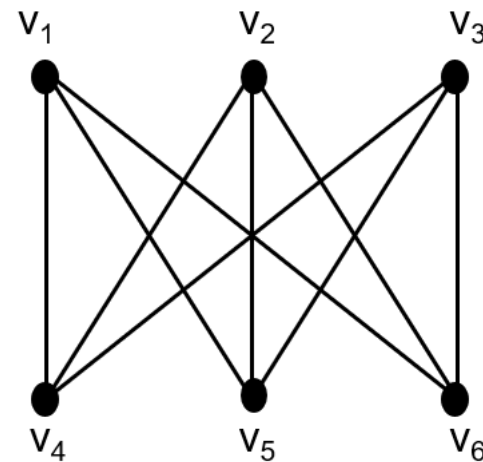
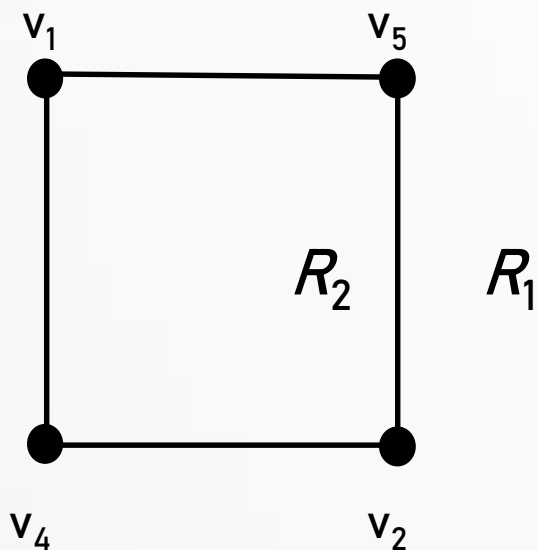
# Regions

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# Regions

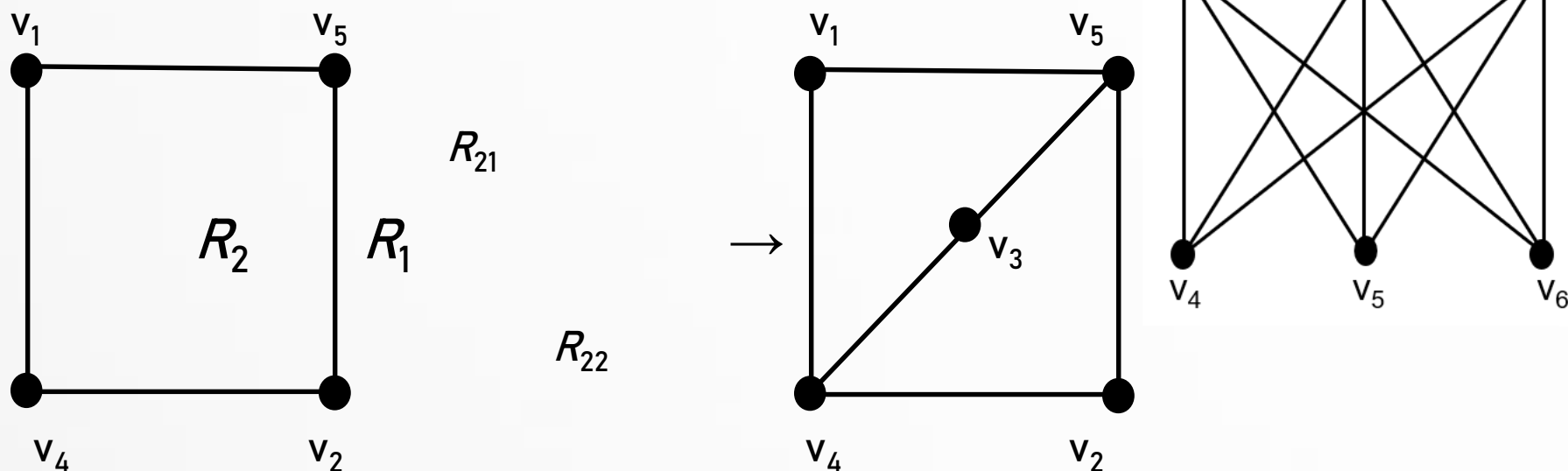
The four edges  $\{v_1, v_4\}$ ,  $\{v_4, v_2\}$ ,  $\{v_2, v_5\}$ ,  $\{v_5, v_1\}$  form a closed curve that splits the plane into two regions,  $R_1$  and  $R_2$ .





# Regions

- Next, we note that  $v_3$  must be in either  $R_1$  or  $R_2$ .
- Assume  $v_3$  is in  $R_2$ . Then the edges  $\{v_3, v_4\}$  and  $\{v_4, v_5\}$  separate  $R_2$  into two subregions,  $R_{21}$  and  $R_{22}$ .



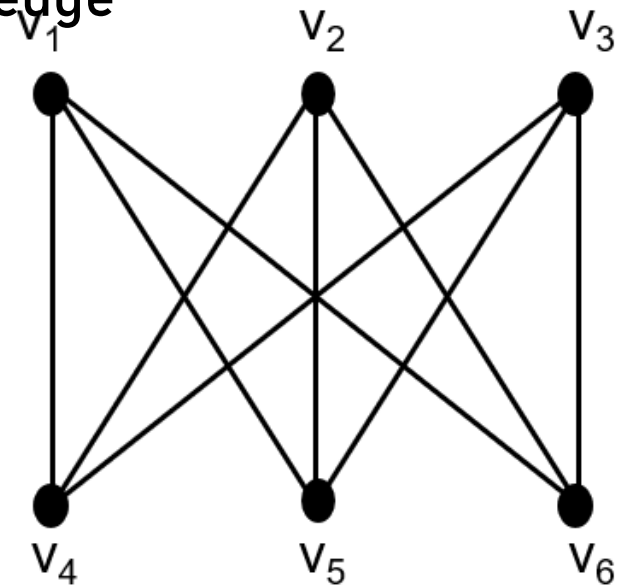
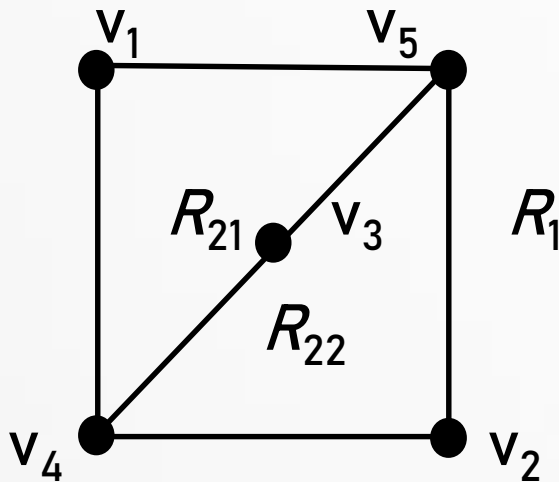
# Regions

Now there is no way to place vertex  $v_6$  without forcing a crossing:

If  $v_6$  is in  $R_1$  then  $\{v_6, v_3\}$  must cross an edge

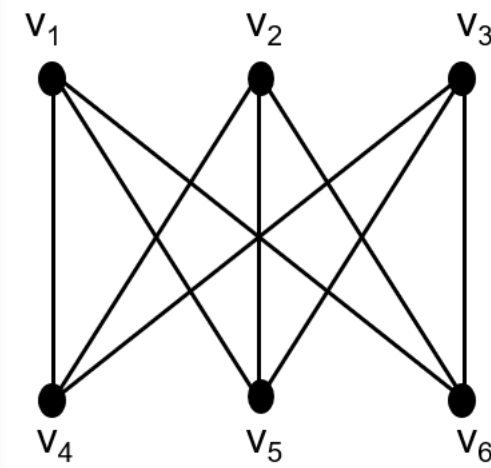
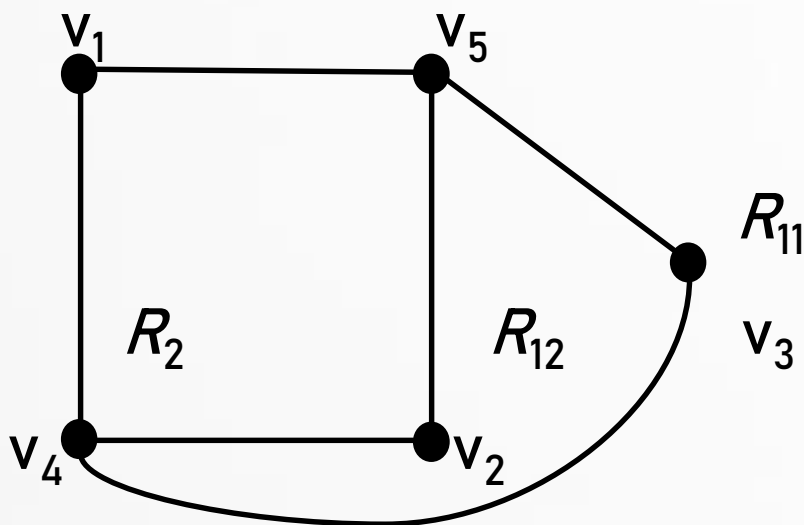
If  $v_6$  is in  $R_{21}$  then  $\{v_6, v_2\}$  must cross an edge

If  $v_6$  is in  $R_{22}$  then  $\{v_6, v_1\}$  must cross an edge



# Regions

Alternatively, assume  $v_3$  is in  $R_1$ . Then the edges  $\{v_3, v_4\}$  and  $\{v_4, v_5\}$  separate  $R_1$  into two subregions,  $R_{11}$  and  $R_{12}$ .



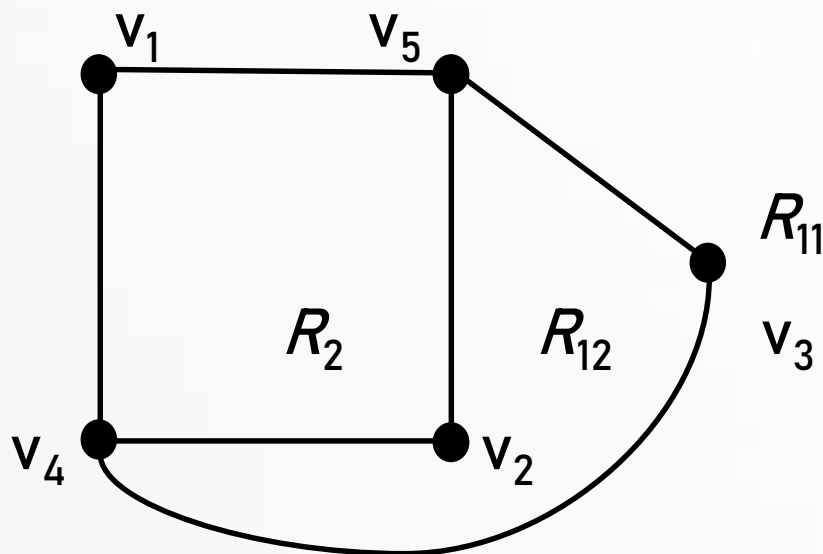
# Regions

Now there is no way to place vertex  $v_6$  without forcing a crossing:

If  $v_6$  is in  $R_2$  then  $\{v_6, v_3\}$  must cross an edge

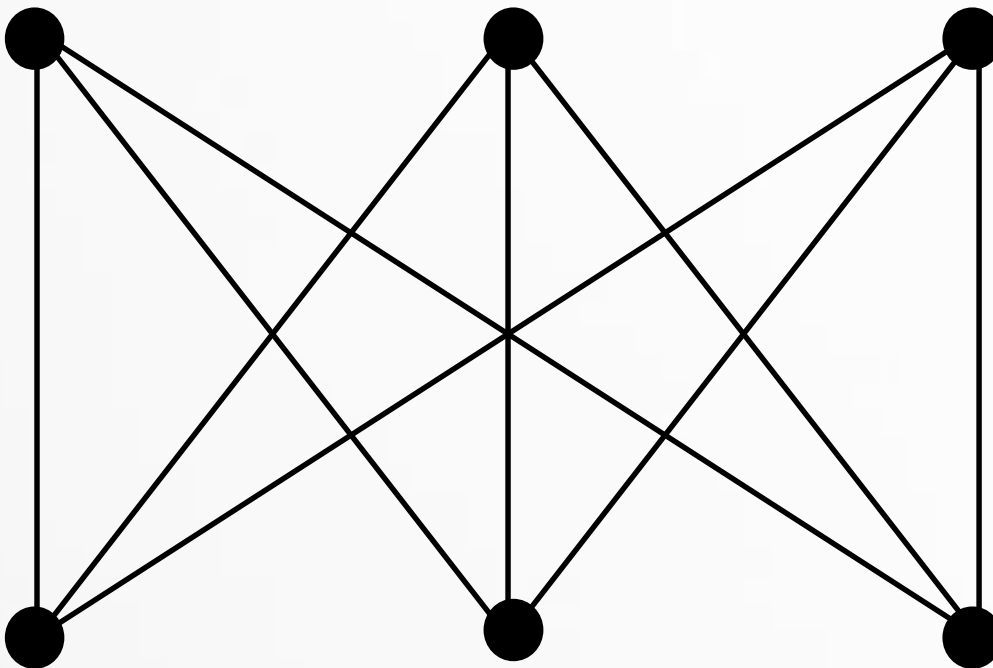
If  $v_6$  is in  $R_{11}$  then  $\{v_6, v_2\}$  must cross an edge

If  $v_6$  is in  $R_{12}$  then  $\{v_6, v_1\}$  must cross an edge



# Planar Graphs

Consequently, the graph  $K_{3,3}$  must be nonplanar.



$K_{3,3}$

# Theorem 1 (Kuratowski, 1930)

A graph is planar if and only if it does not contain a subgraph that is a  $K_5$  or  $K_{3,3}$  configuration.

That's all for now...