

A hand is shown placing a blue L-shaped block onto a colorful geometric structure made of various blocks. The structure is composed of blocks in shades of blue, orange, yellow, purple, and pink. The background is a solid light blue, and the surface is a light-colored wooden table. Several other blocks are scattered on the table in the foreground, including a green L-shaped block, a blue L-shaped block, a red L-shaped block, and a yellow L-shaped block.

# EMTH403

## Mathematical Foundation for Computer Science

Nitin K. Mishra (Ph.D.)

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Associate Professor

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# Lecture Outcomes



After this lecture, you will be able to

- understand what are path, cycles and connectivity in a graph.
- understand how to discard an isomorphism with the help of cycles.

# Connectivity – Path

- Many problems can be modeled with paths formed by traveling along the edges of graphs.
- For instance, the problem of determining whether a **message can be sent between two computers** using intermediate links can be studied with a graph model.
- Problems of efficiently planning routes for **mail delivery, garbage pickup, diagnostics in computer networks**, and so on can be solved using models that involve paths in graphs.

# Connectivity – Path

- Informally, a path is a **sequence of edges** that begins at a vertex of a graph and travels from vertex to vertex along edges of the graph.
- As the path travels along its edges, it visits the vertices along this path, that is, the endpoints of these edges.

# Connectivity – Path

A formal definition of paths and related terminology is given below.

# Connectivity – Path

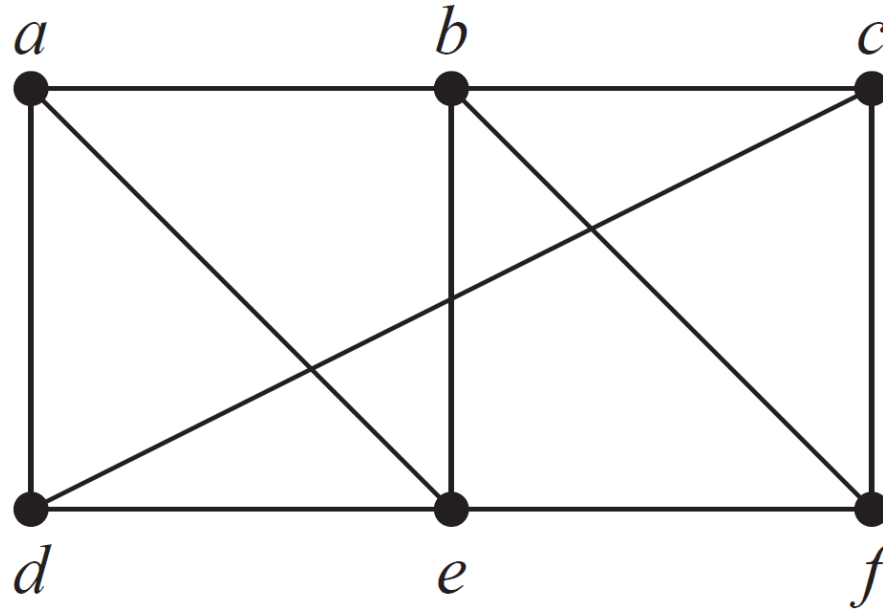
Let  $n$  be a nonnegative integer and  $G$  an undirected graph.

A path of **length  $n$**  from  $u$  to  $v$  in  $G$  is a sequence of  $n$  edges  $e_1, \dots, e_n$  of  $G$  for which there exists a sequence  $x_0 = u, x_1, \dots, x_{n-1}, x_n = v$  of vertices such that  $e_i$  has, for  $i = 1, \dots, n$ , the endpoints  $x_{i-1}$  and  $x_i$ .

# Connectivity – Path

- The path is a **circuit** if it begins and ends at the same vertex, that is, if  $u = v$ , and has length greater than zero.
- The **path or circuit** is said to pass through the vertices  $x_1, x_2, \dots, x_{n-1}$  or traverse the edges  $e_1, e_2, \dots, e_n$ .
- A path or circuit is **simple** if it does not contain the same edge more than once.

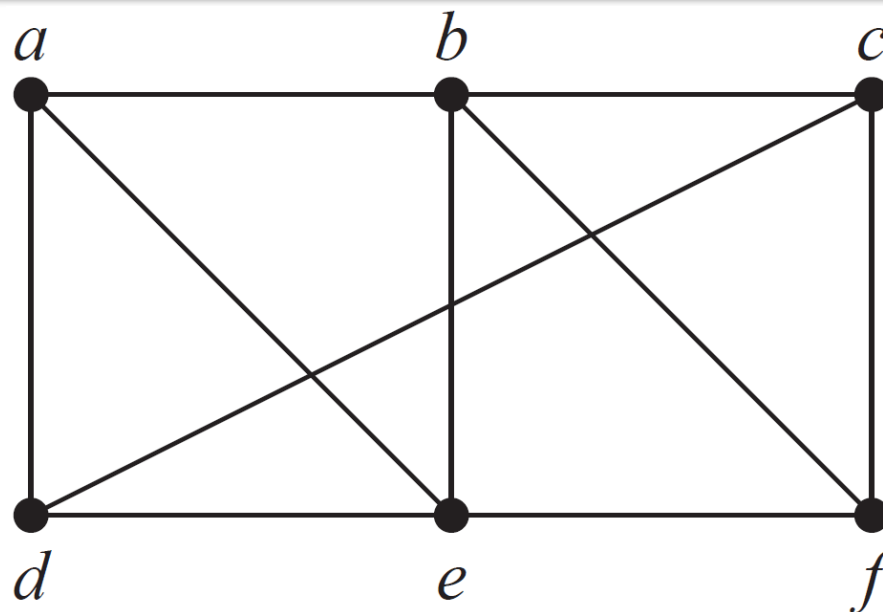
# Connectivity – Path



In the simple graph shown in above,  $a, d, c, f, e$  is a simple path of length 4, because  $\{a, d\}$ ,  $\{d, c\}$ ,  $\{c, f\}$ , and  $\{f, e\}$  are all edges.



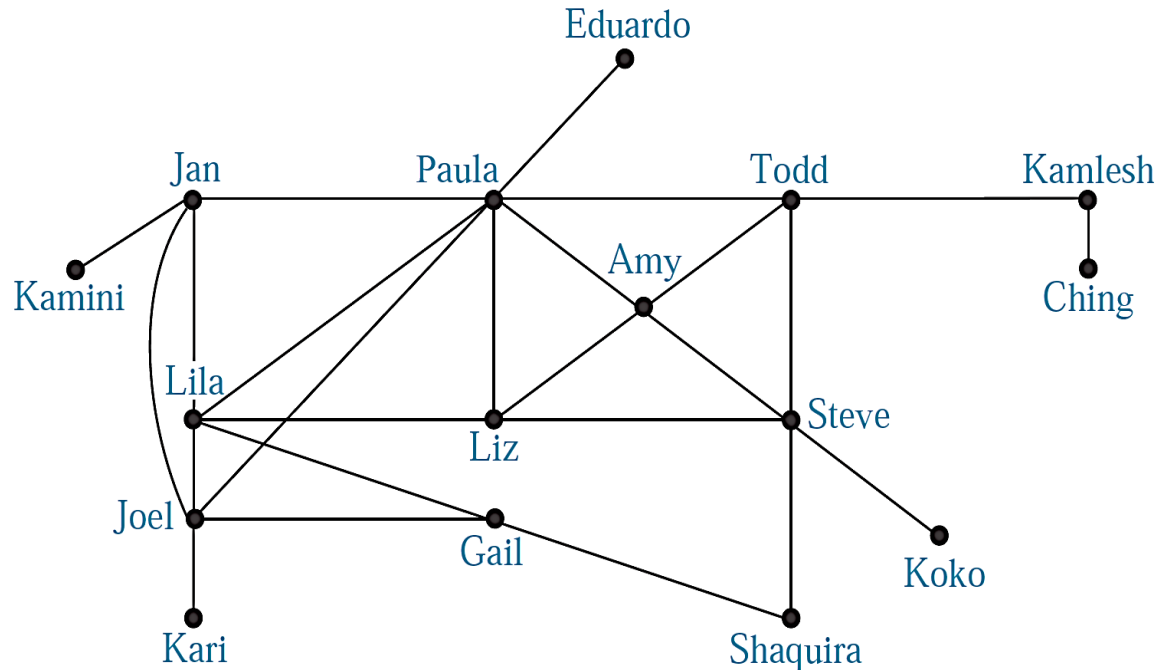
# Connectivity – Path



However,  $d, e, c, a$  is not a path, because  $\{e, c\}$  is not an edge.

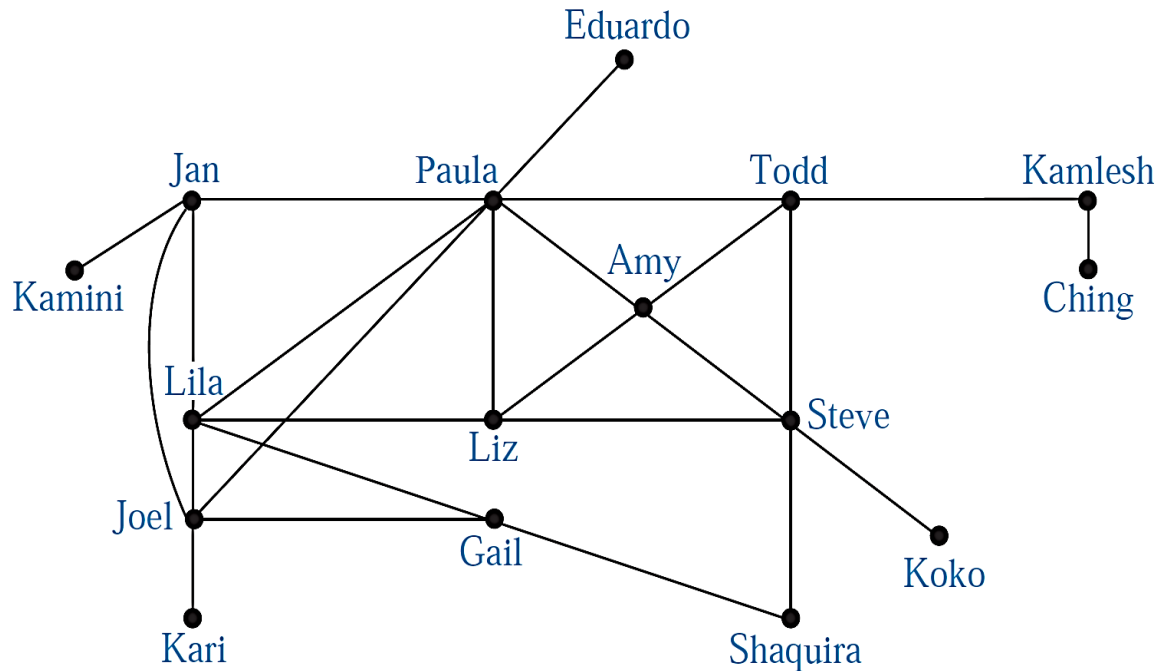
Note that  $b, c, f, e, b$  is a circuit of length 4 because  $\{b, c\}, \{c, f\}, \{f, e\},$  and  $\{e, b\}$  are edges, and this path begins and ends at  $b$ .

# Connectivity – Path – Example



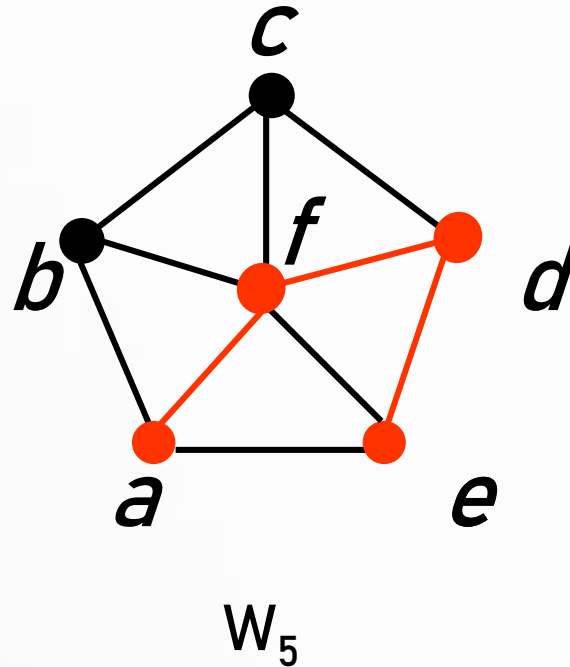
In an acquaintanceship graph there is a path between two people if there is a chain of people linking these people, where two people adjacent in the chain knows one another. For example, there is a chain of six people linking Kamini and Ching

# Connectivity – Path – Example



Many social scientists have conjectured that almost every pair of people in the world are linked by a small chain of people, perhaps containing just five or fewer people.

# Connectivity – Path



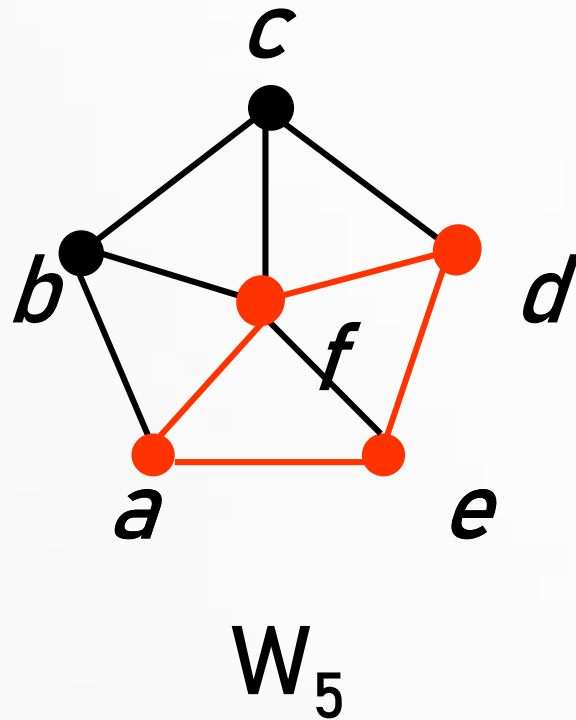
This path passes through vertices  $f$  and  $d$  in that order.

# Connectivity – Path

A **path of length  $n$**  from  $u$  to  $v$  in an undirected graph is a sequence of edges  $e_1, e_2, \dots, e_n$  of the graph such that edge  $e_1$  has endpoints  $x_0$  and  $x_1$ , edge  $e_2$  has endpoints  $x_1$  and  $x_2$ ,  $\dots$  and edge  $e_n$  has endpoints  $x_{n-1}$  and  $x_n$ ,

**where  $x_0 = u$  and  $x_n = v$ .**

# Connectivity – Path

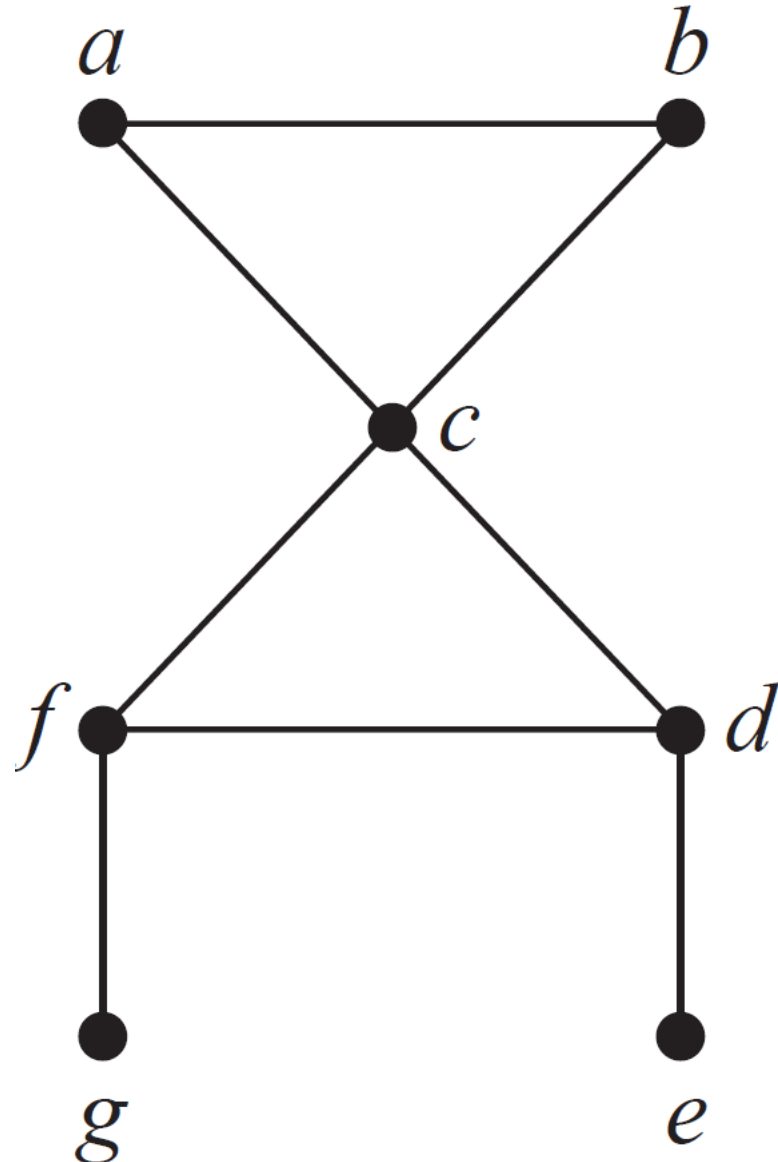


This path passes through vertices  $f, d, e$ , in that order. It has **length** 4.

It is a **circuit** because it begins and ends at the same vertex.

It is called **simple** because it does not contain the same edge more than once.

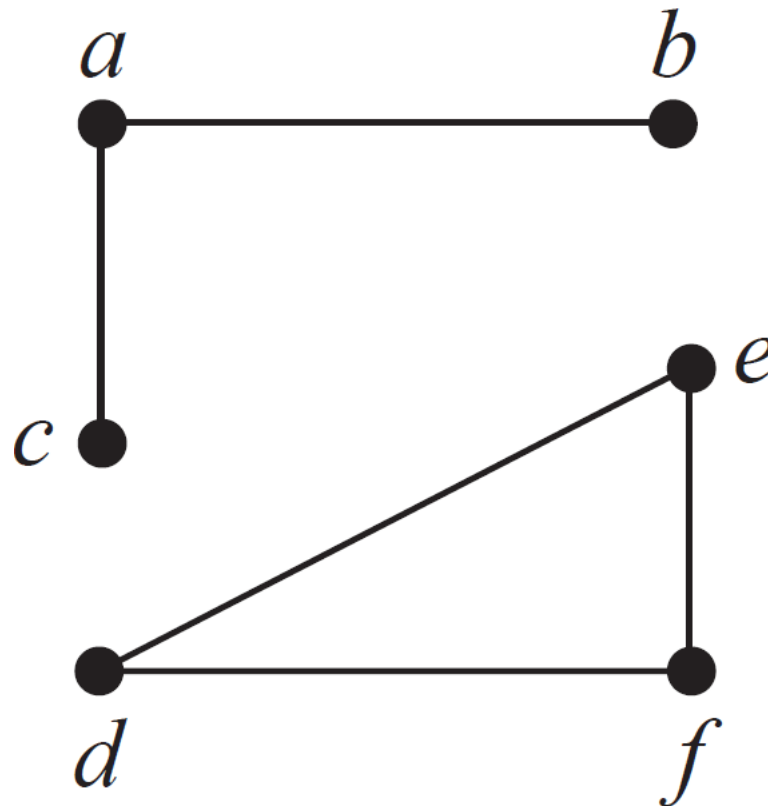
# Connectivity – Path



- An undirected graph is called connected if there is a path between every pair of distinct vertices of the graph.

# Connectivity – Path

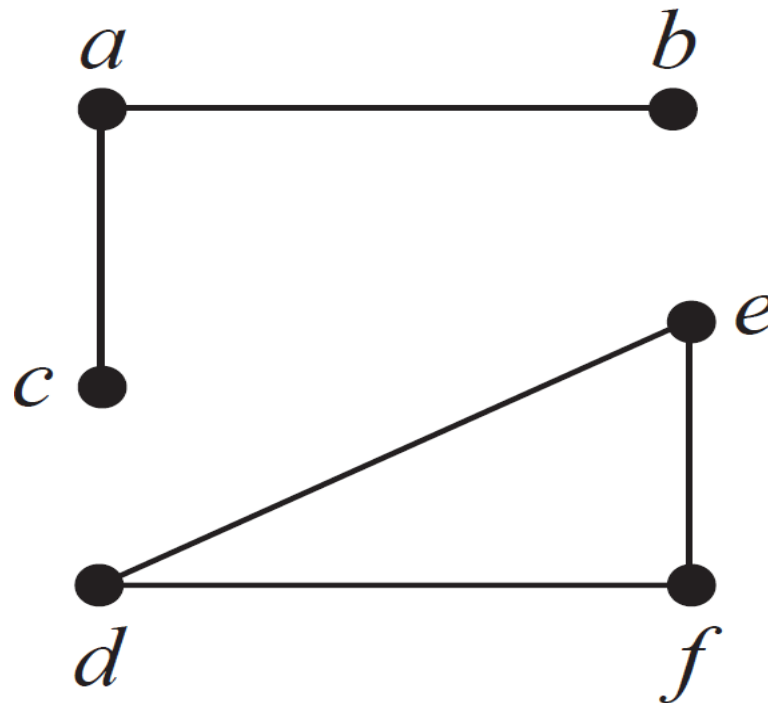
- An undirected graph that is not connected is called disconnected.





# Connectivity – Path

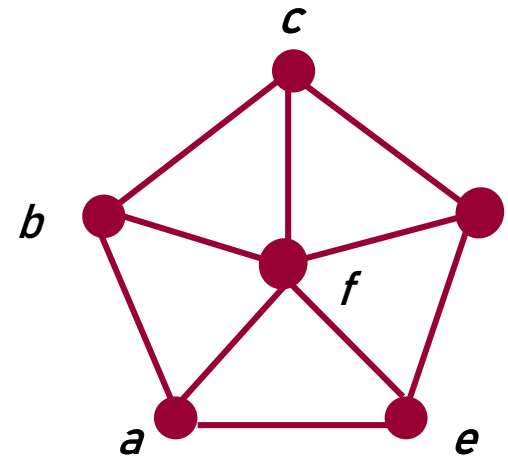
- We say that we disconnect a graph when we remove vertices or edges, or both, to produce a disconnected subgraph.



# Connectedness in Undirected Graphs

**Definition:-** An undirected graph is called connected if there is a path between every pair of distinct vertices of the graph.

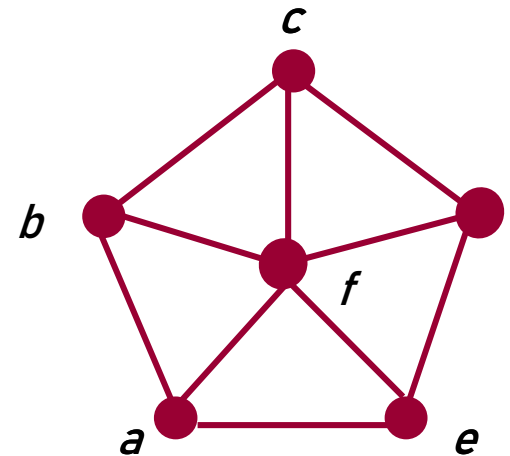
IS THIS GRAPH CONNECTED?



$W_5$

# Connectedness in Undirected Graphs

**Theorem:-** There is a simple path between every pair of distinct vertices of a connected undirected graph.



$W_5$

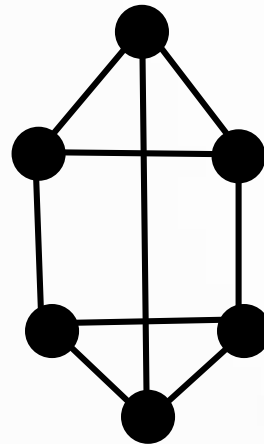
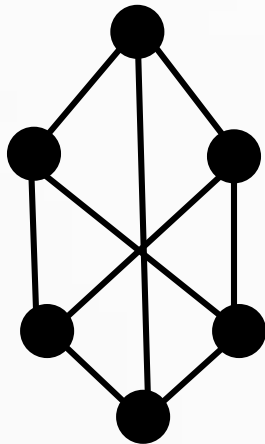
# Graph Isomorphism – Definition

The simple graphs  $G_1 = (V_1, E_1)$  and  $G_2 = (V_2, E_2)$  are *isomorphic* if there exists a one to-one and onto function  $f$  from  $V_1$  to  $V_2$  with the property that  $a$  and  $b$  are adjacent in  $G_1$  if and only if  $f(a)$  and  $f(b)$  are adjacent in  $G_2$ , for all  $a$  and  $b$  in  $V_1$ .

Such a function  $f$  is called an *isomorphism*.\* Two simple graphs that are not isomorphic are called *non-isomorphic*.

# Graph Isomorphism

Ques:- Are the pairs of graphs are isomorphic?

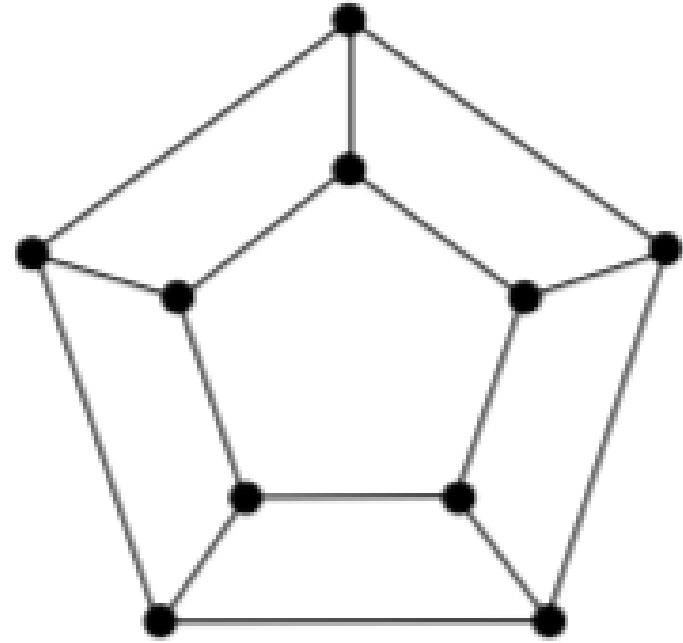
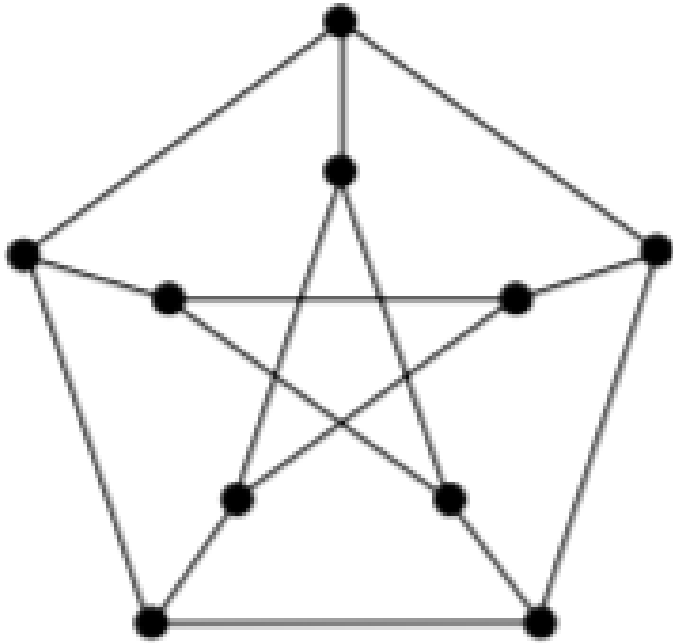


Ans:- Hint: Check Cycles

Not isomorphic, the right-hand side graph contains a triangle, but the left-hand side graph does not.

# Graph Isomorphism

Ques:- Are the pairs of graphs are isomorphic?

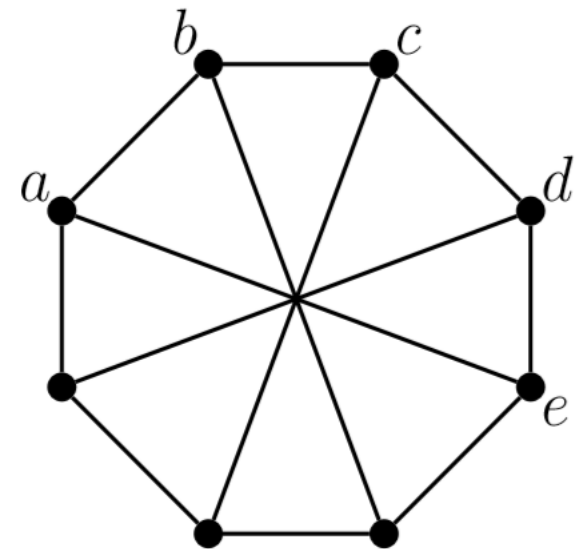
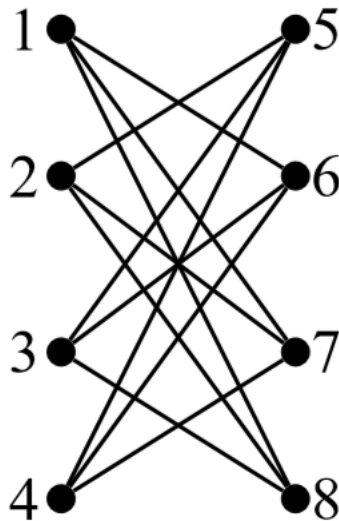
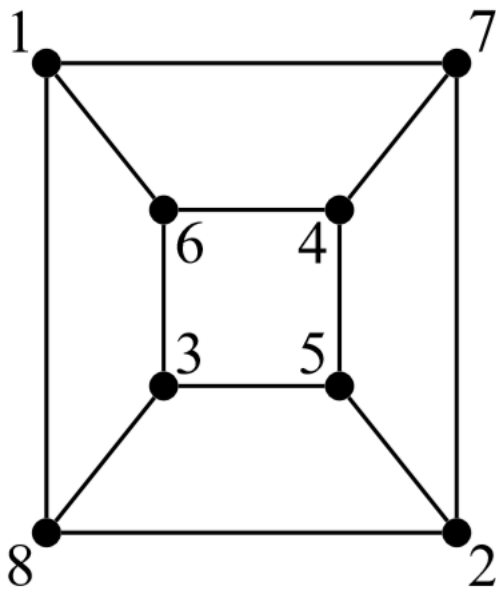


Ans:- Hint: Check Cycles  
Not isomorphic.

First does not contain a cycle of length 4, while the right graph contains a cycle of length 4.

# Graph Isomorphism

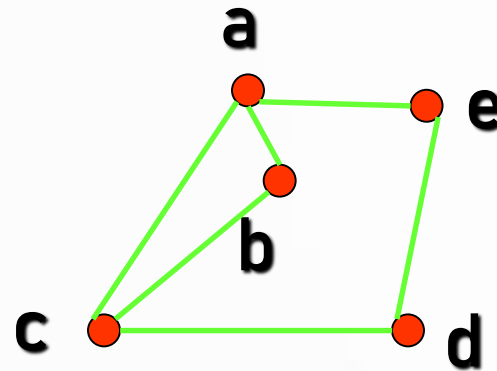
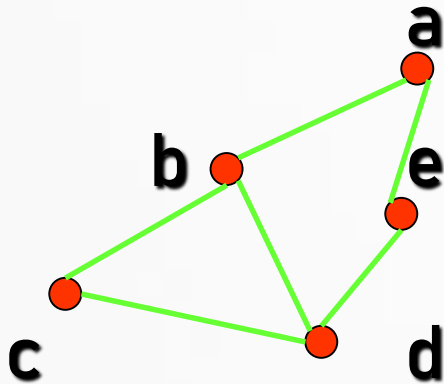
**Ques:-** Determine which pairs of graphs below are isomorphic. Justify your answer!



**Ans:-** The first two graphs are isomorphic.  
On the other hand, the third graph contains an odd cycle on 5 vertices  $a, b, c, d, e$ , thus, this graph is not isomorphic to the first two.

# Graph Isomorphism

Ques:- Are the following two graphs isomorphic?



Sol:- Yes, they are isomorphic, because they can be arranged to look identical. You can see this if in the right graph you move vertex b to the left of the edge  $\{a, c\}$ . Then the isomorphism  $f$  from the left to the right graph is:  $f(a) = e$ ,  $f(b) = a$ ,  $f(c) = b$ ,  $f(d) = c$ ,  $f(e) = d$ .



That's all for now...