



# EMTH403

## Mathematical Foundation for Computer Science

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# Lecture Outcomes



After this lecture, you will be able to

- understand what is a biconditional Statements
- learn what is equivalence of formulas
- understand what is duality law.

# BICONDITIONALS

Let  $p$  and  $q$  be propositions. The biconditional statement  $p \Leftrightarrow q$  is the proposition  
“ $p$  if and only if  $q$ .”

The biconditional statement  $p \Leftrightarrow q$  is true when  $p$  and  $q$  have the same truth values, and is false otherwise.

**TABLE 6** The Truth Table for the Biconditional  $p \Leftrightarrow q$ .

$p$	$q$	$p \Leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

# BICONDITIONALS

Biconditional statements are also called **bi-implications**.

**TABLE 6** The Truth Table for the Biconditional  $p \leftrightarrow q$ .

$p$	$q$	$p \leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

# BICONDITIONALS – Example 1

Let  $p$  be the statement “You can take the flight,” and let  $q$  be the statement “You buy a ticket.”

Then  $p \Leftrightarrow q$  is the statements “You can take the flight if and only if you buy a ticket.”

**TABLE 6** The Truth Table for the Biconditional  $p \Leftrightarrow q$ .

$p$	$q$	$p \Leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

# BICONDITIONALS – Example 1

This statement is true if  $p$  and  $q$  are either both true or both false, that is, if you buy a ticket and can take the flight or if you do not buy a ticket and you cannot take the flight.

**TABLE 6** The Truth Table for the Biconditional  $p \leftrightarrow q$ .

$p$	$q$	$p \leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

# BICONDITIONALS – Example 1

It is false when  $p$  and  $q$  have opposite truth values, that is, when you do not buy a ticket, but you can take the flight (such as when you get a free trip) and when you buy a ticket but you cannot take the light (such as when the airline bumps you).

**TABLE 6** The Truth Table for the Biconditional  $p \leftrightarrow q$ .

$p$	$q$	$p \leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

# BICONDITIONALS – Example 2

**Ques:-**Let p and q be the propositions “Swimming at the New Jersey shore is allowed” and “Sharks have been spotted near the shore,” respectively. Express

$p \Leftrightarrow \neg q$  compound proposition as an English sentence.

**Ans:-**Swimming at the New Jersey shore is allowed if and only if sharks have not been spotted near the shore.

# Truth Tables of Compound Propositions

Construct the truth table of the compound proposition

$$(p \vee \neg q) \rightarrow (p \wedge q).$$

The Truth Table of  $(p \vee \neg q) \rightarrow (p \wedge q)$

is

**TABLE 1** The Truth Table for the Negation of a Proposition.

$p$	$\neg p$
T	F
F	T

$p$	$q$	$\neg q$	$p \vee \neg q$	$p \wedge q$	$(p \vee \neg q) \rightarrow (p \wedge q)$
T	T	F	T	T	T
T	F	T	T	F	F
F	T	F	F	F	T
F	F	T	T	F	F

# Truth Tables of Compound Propositions

Construct the truth table of the compound proposition

$$(p \vee \neg q) \rightarrow (p \wedge q).$$

The Truth Table of  $(p \vee \neg q) \rightarrow (p \wedge q)$

is

**TABLE 3** The Truth Table for the Disjunction of Two Propositions.

$p$	$q$	$p \vee q$
T	T	T
T	F	T
F	T	T
F	F	F

$p$	$q$	$\neg q$	$p \vee \neg q$	$p \wedge q$	$(p \vee \neg q) \rightarrow (p \wedge q)$
T	T	F	T	T	T
T	F	T	T	F	F
F	T	F	F	F	T
F	F	T	T	F	F

# Truth Tables of Compound Propositions

Construct the truth table of the compound proposition

$$(p \vee \neg q) \rightarrow (p \wedge q).$$

The Truth Table of  $(p \vee \neg q) \rightarrow (p \wedge q)$

is

**TABLE 2** The Truth Table for the Conjunction of Two Propositions.

$p$	$q$	$p \wedge q$
T	T	T
T	F	F
F	T	F
F	F	F

$p$	$q$	$\neg q$	$p \vee \neg q$	$p \wedge q$	$(p \vee \neg q) \rightarrow (p \wedge q)$
T	T	F	T	T	T
T	F	T	T	F	F
F	T	F	F	F	T
F	F	T	T	F	F

# Truth Tables of Compound Propositions

Construct the truth table of the compound proposition

$$(p \vee \neg q) \rightarrow (p \wedge q).$$

The Truth Table of  $(p \vee \neg q) \rightarrow (p \wedge q)$

is

**TABLE 5** The Truth Table for the Conditional Statement  $p \rightarrow q$ .

$p$	$q$	$p \rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

$p$	$q$	$\neg q$	$p \vee \neg q$	$p \wedge q$	$(p \vee \neg q) \rightarrow (p \wedge q)$
T	T	F	T	T	T
T	F	T	T	F	F
F	T	F	F	F	T
F	F	T	T	F	F

# Tautology

**DEFINITION:-** A compound proposition that is always true, no matter what the truth values of the propositional variables that occur in it, is called a tautology.

Example of a **Tautology**

$p$	$\neg p$	$p \vee \neg p$
T	F	T
F	T	T

# Contradiction

**DEFINITION:-** A compound proposition that is always false is called a contradiction.

Example of a **Contradiction**.

$p$	$\neg p$	$p \wedge \neg p$
T	F	F
F	T	F

# Contingency

**DEFINITION:-** A compound proposition that is neither a tautology nor a contradiction is called a **contingency**. As we can see in the truth table for conditional statement  $p \rightarrow q$ .

$p$	$q$	$p \rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

# Equivalence of formulas

The compound propositions  $p$  and  $q$  are called logically equivalent if  $p \Leftrightarrow q$  is a tautology.

**TABLE 6** The Truth Table for the Biconditional  $p \leftrightarrow q$ .

$p$	$q$	$p \leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

The notation  $p \equiv q$  denotes that  $p$  and  $q$  are logically equivalent.

# Equivalence of formulas

**Remark 1:** The symbol  $\equiv$  is not a logical connective, and  $p \equiv q$  is not a compound proposition but rather is the statement that  $p \Leftrightarrow q$  is a tautology.

**TABLE 6** The Truth Table for the Biconditional  $p \leftrightarrow q$ .

$p$	$q$	$p \leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

# Equivalence of formulas

**Remark 2:** The symbol  $\Leftrightarrow$  is sometimes used instead of  $\equiv$  to denote logical equivalence.

One way to determine whether two compound propositions are equivalent is to use a truth table.

In particular, the compound propositions  $p$  and  $q$  are equivalent if and only if the columns giving their truth values agree.

# Equivalence of formulas

Following Example illustrates this method to establish an extremely important and useful logical equivalence, namely, that of  $\neg(p \vee q)$  with  $\neg p \wedge \neg q$ .

This logical equivalence is one of the two De Morgan laws, shown in the following Table, named after the English mathematician Augustus De Morgan, of the mid-nineteenth century.

# Equivalence of formulas – Example 1

De Morgan's Laws.

Ques:-

Show

that

$$1. \neg(p \wedge q) \equiv \neg p \vee \neg q$$

$\neg(p \vee q)$  and  $\neg p \wedge \neg q$  are

$$2. \neg(p \vee q) \equiv \neg p \wedge \neg q$$

logically equivalent.

The truth table for  $\neg(p \vee q)$  and  $\neg p \wedge \neg q$  is:-

$p$	$q$	$p \vee q$	$\neg(p \vee q)$	$\neg p$	$\neg q$	$\neg p \wedge \neg q$
T	T	T	F	F	F	F
T	F	T	F	F	T	F
F	T	T	F	T	F	F
F	F	F	T	T	T	T

# Equivalence of formulas – Example 2

Show that  $p \rightarrow q$  and  $\neg p \vee q$  are logically equivalent.

Truth Tables for  $\neg p \vee q$  and  $p \rightarrow q$  is

$p$	$q$	$\neg p$	$\neg p \vee q$	$p \rightarrow q$
T	T	F	T	T
T	F	F	F	F
F	T	T	T	T
F	F	T	T	T

# Duality law

The dual of a compound proposition that contains only the logical operators  $\vee$ ,  $\wedge$ , and  $\neg$  is the compound proposition obtained by replacing each  $\vee$  by  $\wedge$ , each  $\wedge$  by  $\vee$ , each  $T$  by  $F$ , and each  $F$  by  $T$ .

The dual of  $s$  is denoted by  $s^*$ .

# Duality law – Example 1

Ques:- Find the dual of the compound proposition

$$p \wedge \neg q \wedge \neg r.$$

Ans:- $\neg p \vee \neg q \vee \neg r$ (replacing each  $\vee$  by  $\wedge$ , each  $\wedge$  by  $\vee$ )

# Duality law – Example 2

Ques:- Find the dual of the compound proposition (p

$\wedge q \wedge r) \vee s.$

Ans:– $(p \vee q \vee r) \wedge s$ (replacing each  $\vee$  by  $\wedge$ , each  $\wedge$  by  $\vee$ )

# Duality law – Example 3

Ques:- Find the dual of the compound proposition (p  
v F)  $\wedge$  (q v T).

Ans:-(p  $\wedge$  T) v(q  $\wedge$  F)(replacing each v by  $\wedge$ , each  $\wedge$   
by v, each T by F, and each F by T. )

That's all for now...