

Chapter

10

***Sorting and
Searching
Algorithms***

C++ Plus Data Structures

Third Edition

**C⁺⁺ *Plus* Data
Structures**

Nell Dale



Sorting means . . .

The values stored in an array have keys of a type for which the relational operators are defined. (We also assume unique keys.)

Sorting rearranges the elements into either ascending or descending order within the array. (We'll use ascending order.)



Straight Selection Sort

values [0]	36
[1]	24
[2]	10
[3]	6
[4]	12

Divides the array into two parts: already sorted, and not yet sorted.

On each pass, finds the smallest of the unsorted elements, and swaps it into its correct place, thereby increasing the number of sorted elements by one.



Selection Sort: Pass One

values [0]

36

[1]

24

[2]

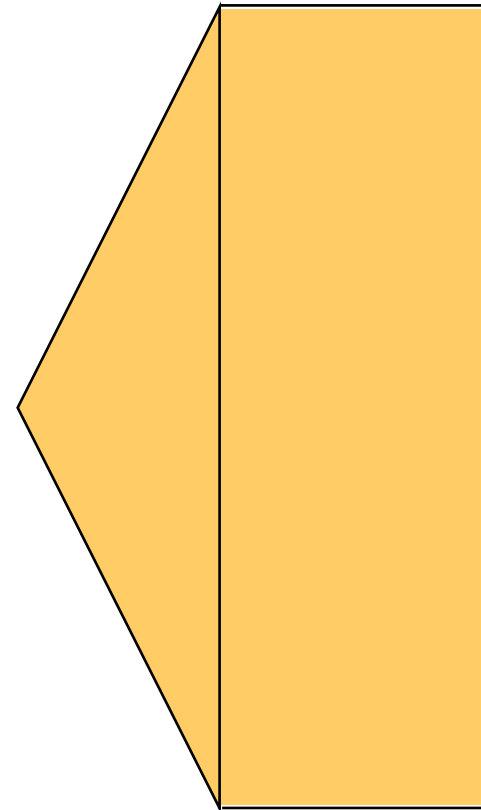
10

[3]

6

[4]

12

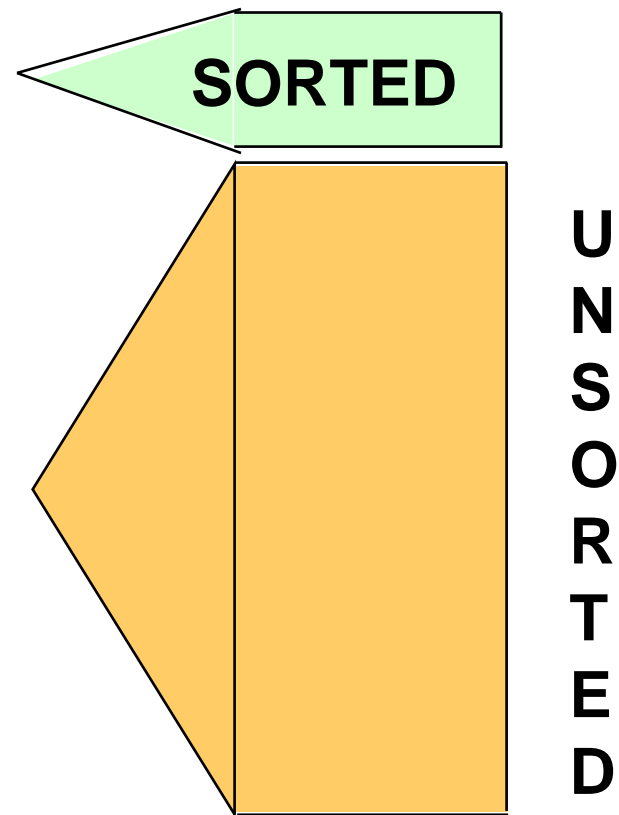


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Selection Sort: End Pass One

values [0]	6
[1]	24
[2]	10
[3]	36
[4]	12





Selection Sort: Pass Two

values [0]

6

[1]

24

[2]

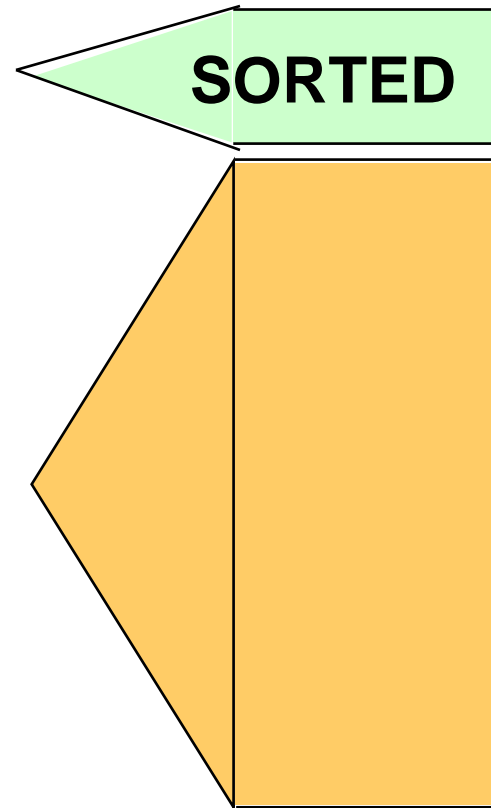
10

[3]

36

[4]

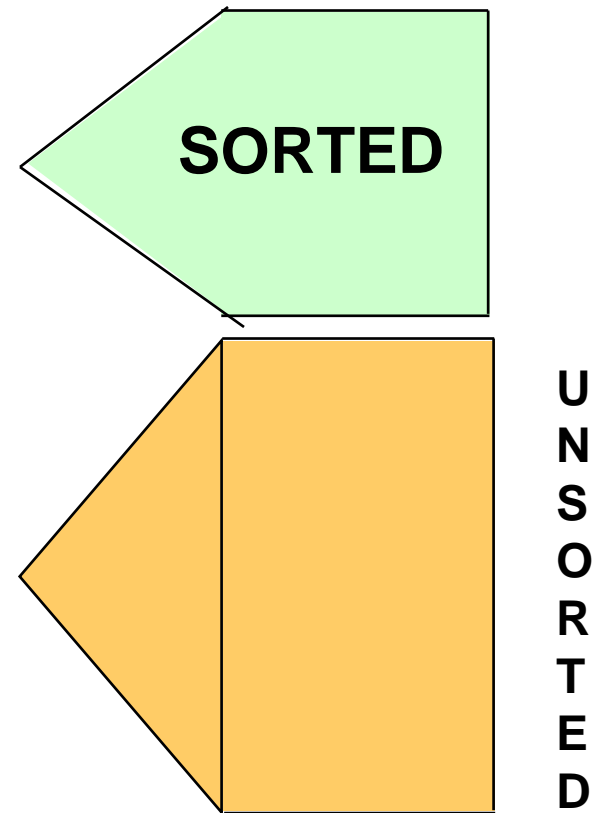
12





Selection Sort: End Pass Two

values [0]	6
[1]	10
[2]	24
[3]	36
[4]	12





Selection Sort: Pass Three

values [0]

6

[1]

10

[2]

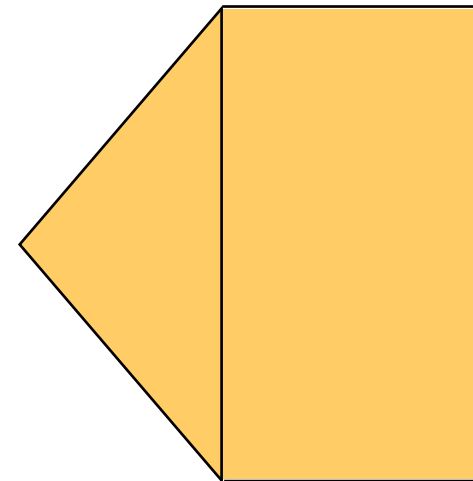
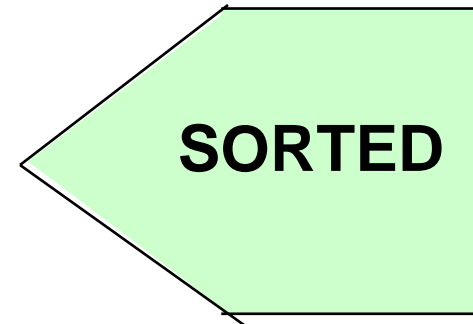
24

[3]

36

[4]

12



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Selection Sort: End Pass Three

values [0]

6

[1]

10

[2]

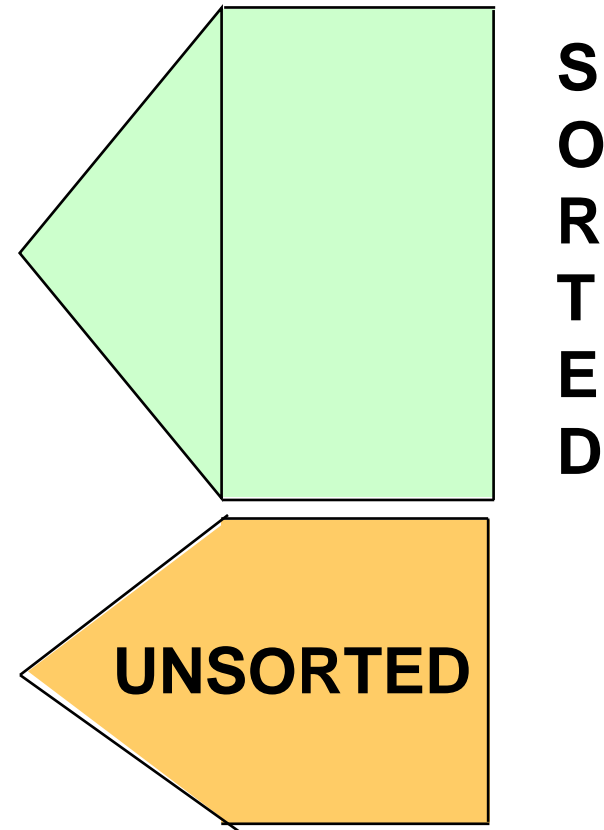
12

[3]

36

[4]

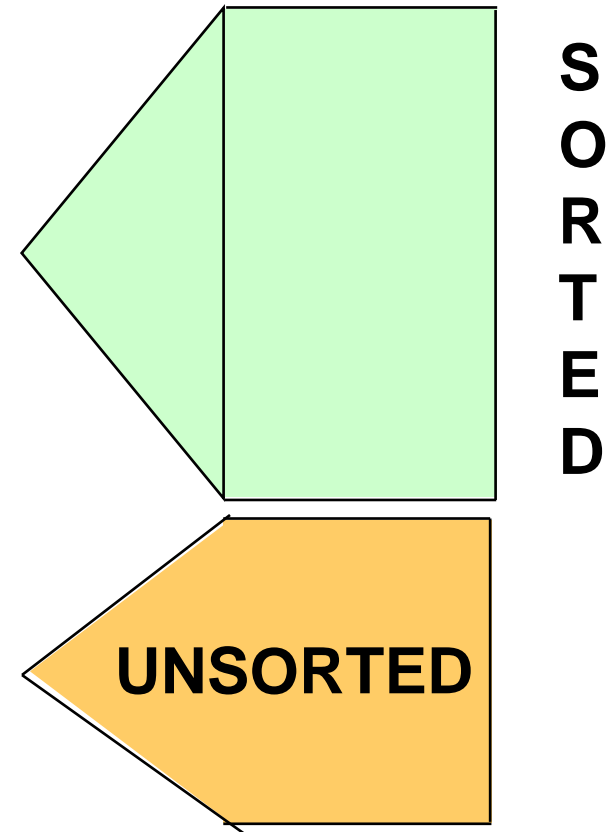
24





Selection Sort: Pass Four

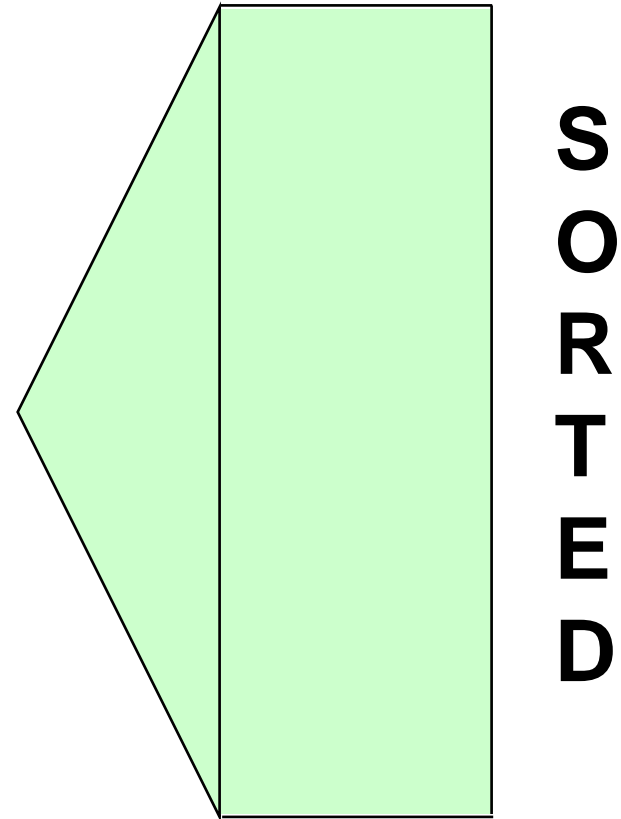
values	[0]	6
	[1]	10
	[2]	12
	[3]	36
	[4]	24





Selection Sort: End Pass Four

values	[0]	6
	[1]	10
	[2]	12
	[3]	24
	[4]	36





Selection Sort: How many comparisons?

values	[0]	6
	[1]	10
	[2]	12
	[3]	24
	[4]	36

4 compares for values[0]

3 compares for values[1]

2 compares for values[2]

1 compare for values[3]

$$= 4 + 3 + 2 + 1$$



For selection sort in general

The number of comparisons when the array contains N elements is

$$\text{Sum} = (N-1) + (N-2) + \dots + 2 + 1$$



Notice that . . .

$$\text{Sum} = (N-1) + (N-2) + \dots + 2 + 1$$

$$+ \text{Sum} = 1 + 2 + \dots + (N-2) + (N-1)$$

$$2 * \text{Sum} = N + N + \dots + N + N$$
A red bracket is drawn under the sequence of N terms, indicating that there are N terms of N being added.

$$2 * \text{Sum} = N * (N-1)$$

$$\text{Sum} = \frac{N * (N-1)}{2}$$



For selection sort in general

The number of comparisons when the array contains N elements is

$$\text{Sum} = (N-1) + (N-2) + \dots + 2 + 1$$

$$\text{Sum} = N * (N-1) / 2$$

$$\text{Sum} = .5 N^2 - .5 N$$

$$\text{Sum} = O(N^2)$$



```
template <class ItemType >
int  MinIndex(ItemType values [ ], int  start, int end)
//  Post: Function value = index of the smallest value
//  in values [start] . . values [end].
{
    int  indexOfMin = start ;

    for(int index = start + 1 ; index <= end ; index++)
        if  (values[ index] < values [indexOfMin])
            indexOfMin = index ;

    return  indexOfMin;
}
```




```
template <class ItemType >
void SelectionSort (ItemType values[ ],
    int numValues )

// Post: Sorts array values[0 . . numValues-1 ]
// into ascending order by key
{
    int endIndex = numValues - 1 ;

    for (int current = 0 ; current < endIndex;
        current++)

        Swap (values[current],
            values [MinIndex (values, current, endIndex) ] ) ;

}
```



Bubble Sort

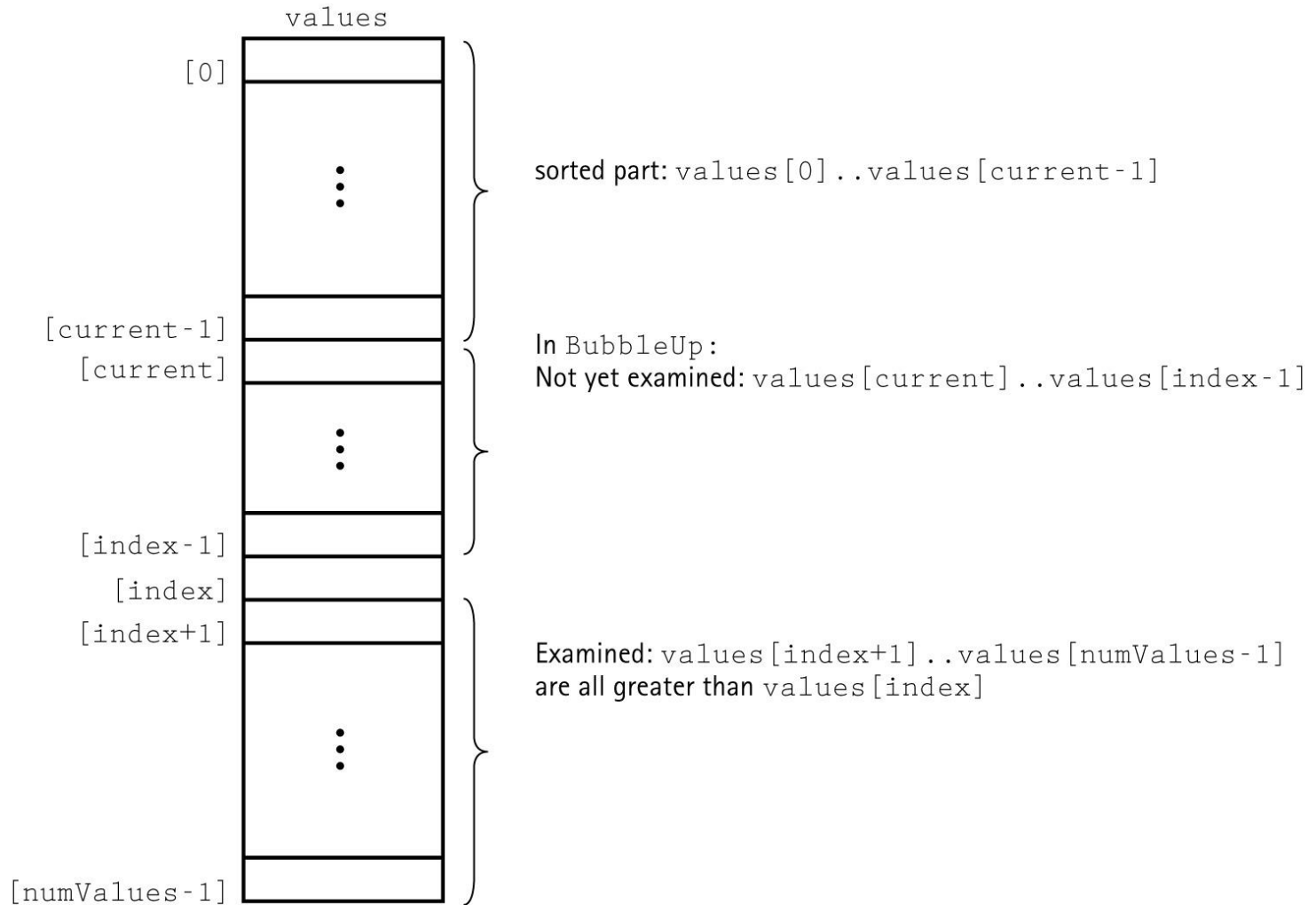
values [0]	36
[1]	24
[2]	10
[3]	6
[4]	12

Compares neighboring pairs of array elements, starting with the last array element, and swaps neighbors whenever they are not in correct order.

On each pass, this causes the smallest element to “bubble up” to its correct place in the array.



Snapshot of BubbleSort





Code for BubbleSort

```
template<class ItemType>
void BubbleSort(ItemType values[],
    int numValues)
{
    int current = 0;
    while (current < numValues - 1)
    {
        BubbleUp(values, current, numValues-1);
        current++;
    }
}
```



Code for BubbleUp

```
template<class ItemType>
void BubbleUp(ItemType values[],
    int startIndex, int endIndex)
// Post: Adjacent pairs that are out of
//       order have been switched between
//       values[startIndex]..values[endIndex]
//       beginning at values[endIndex].

{
    for (int index = endIndex;
        index > startIndex; index--)
        if (values[index] < values[index-1])
            Swap(values[index], values[index-1]);
}
```



Observations on BubbleSort

This algorithm is *always* $O(N^2)$.

There can be a large number of intermediate swaps.

Can this algorithm be improved?



Insertion Sort

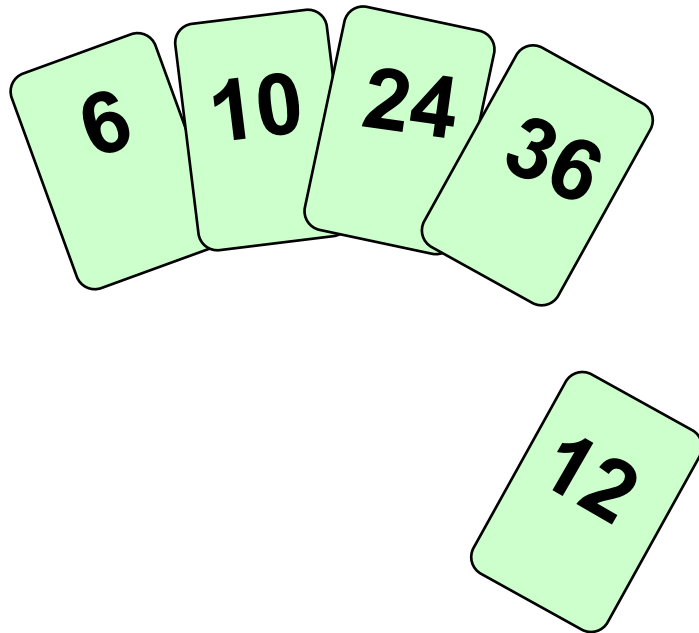
values [0]	36
[1]	24
[2]	10
[3]	6
[4]	12

One by one, each as yet unsorted array element is inserted into its proper place with respect to the already sorted elements.

On each pass, this causes the number of already sorted elements to increase by one.



Insertion Sort

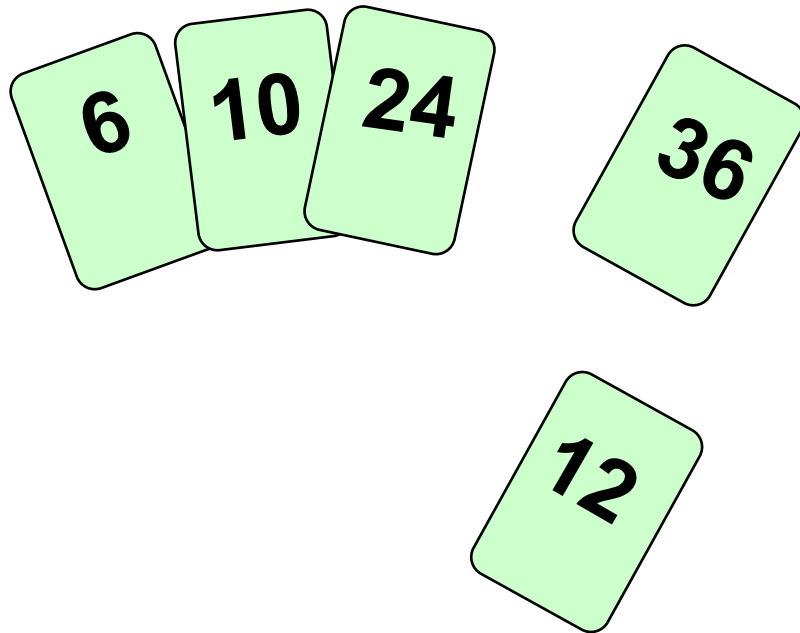


Works like someone who “inserts” one more card at a time into a hand of cards that are already sorted.

To insert 12, we need to make room for it by moving first 36 and then 24.



Insertion Sort

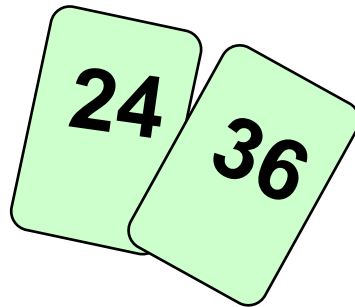
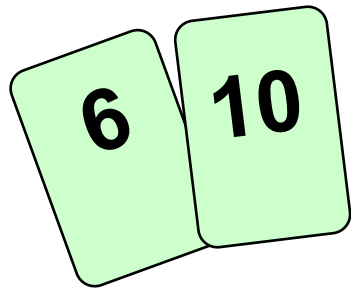


Works like someone who “inserts” one more card at a time into a hand of cards that are already sorted.

To insert 12, we need to make room for it by moving first 36 and then 24.



Insertion Sort

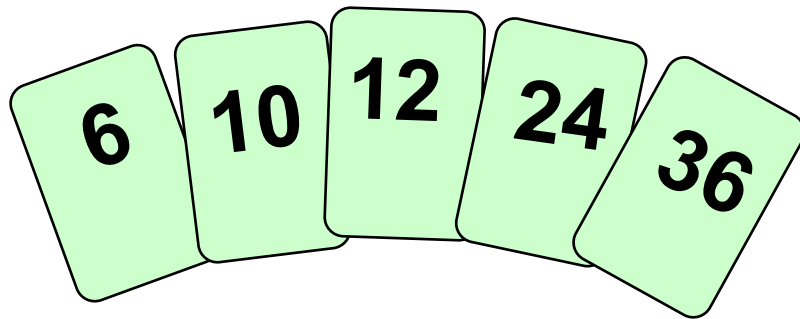


Works like someone who “inserts” one more card at a time into a hand of cards that are already sorted.

To insert 12, we need to make room for it by moving first 36 and then 24.



Insertion Sort

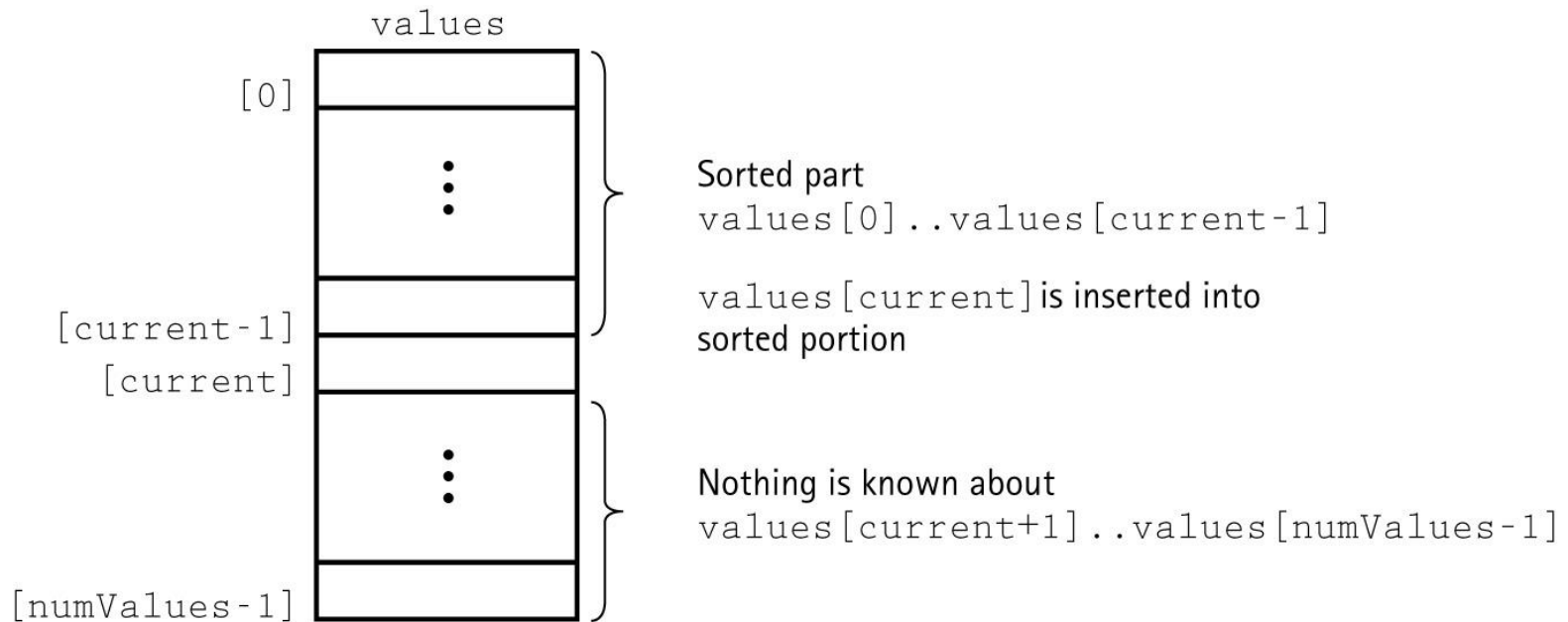


Works like someone who “inserts” one more card at a time into a hand of cards that are already sorted.

To insert 12, we need to make room for it by moving first 36 and then 24.



A Snapshot of the Insertion Sort Algorithm





```
template <class ItemType >
void InsertItem ( ItemType values [ ] , int start ,
int end )
// Post: Elements between values[start] and values
// [end] have been sorted into ascending order by key.
{
    bool finished = false ;
    int current = end ;
    bool moreToSearch = (current != start);

    while (moreToSearch && !finished )
    {
        if (values[current] < values[current - 1])
        {
            Swap(values[current], values[current - 1]);
            current--;
            moreToSearch = ( current != start );
        }
        else
            finished = true ;
    }
}
```



```
template <class  ItemType >
void InsertionSort ( ItemType  values [ ] ,
    int  numValues )

//  Post: Sorts array values[0 . . numValues-1 ] into
//  ascending order by key
{
    for (int count = 0 ; count < numValues; count++)

        InsertItem ( values , 0 , count ) ;
}
```



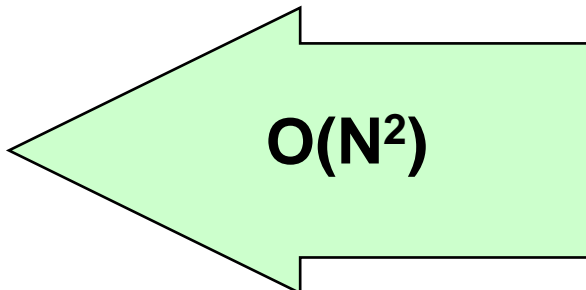
Sorting Algorithms and Average Case Number of Comparisons

Simple Sorts

Straight Selection Sort

Bubble Sort

Insertion Sort



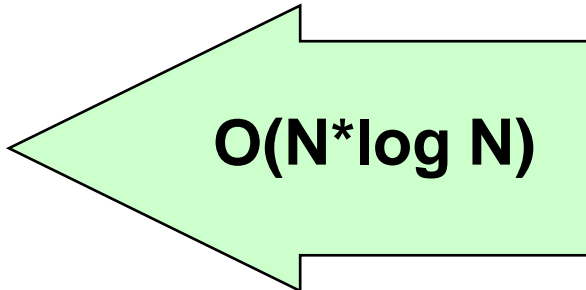
$O(N^2)$

More Complex Sorts

Quick Sort

Merge Sort

Heap Sort



$O(N \log N)$



Recall that . . .

A heap is a binary tree that satisfies these special **SHAPE** and **ORDER** properties:

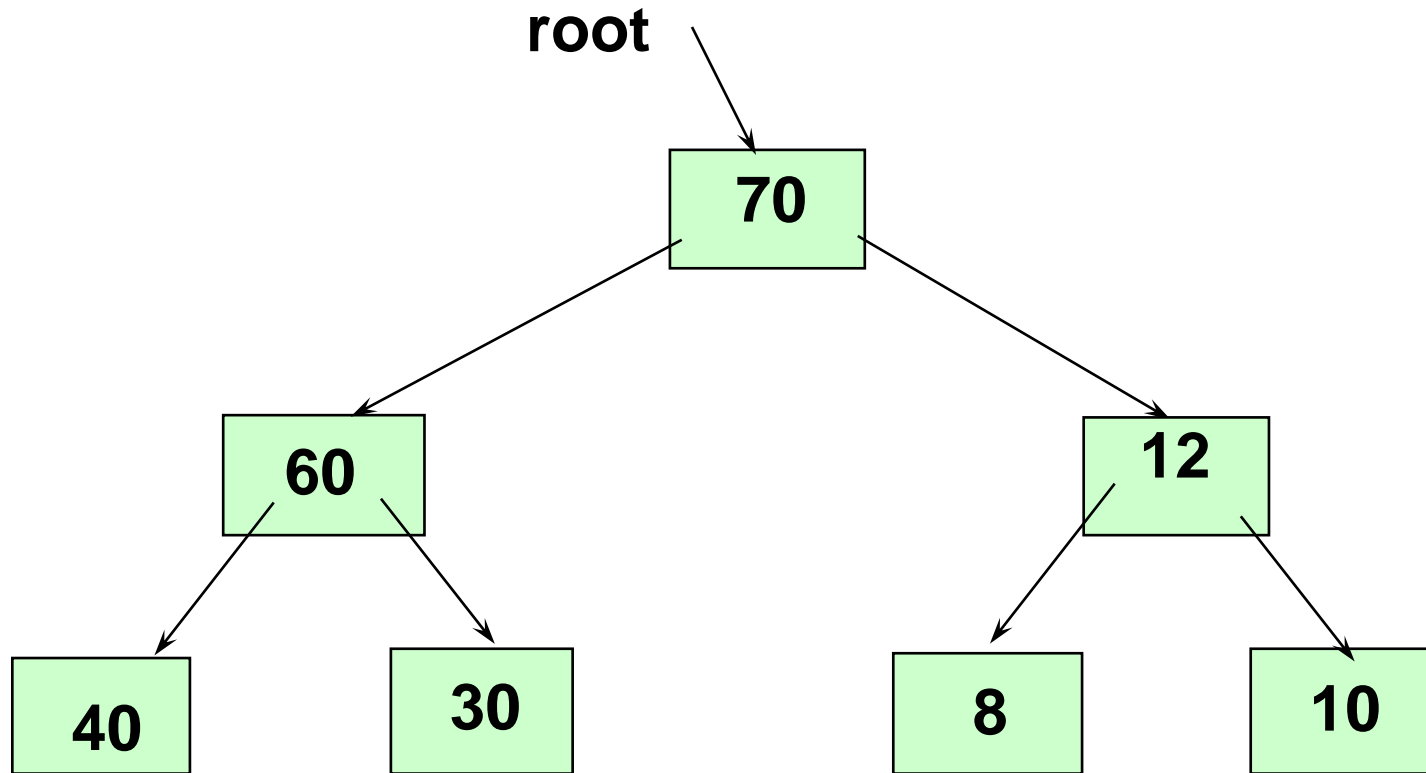
Its shape must be a complete binary tree.

For each node in the heap, the value stored in that node is greater than or equal to the value in each of its children.



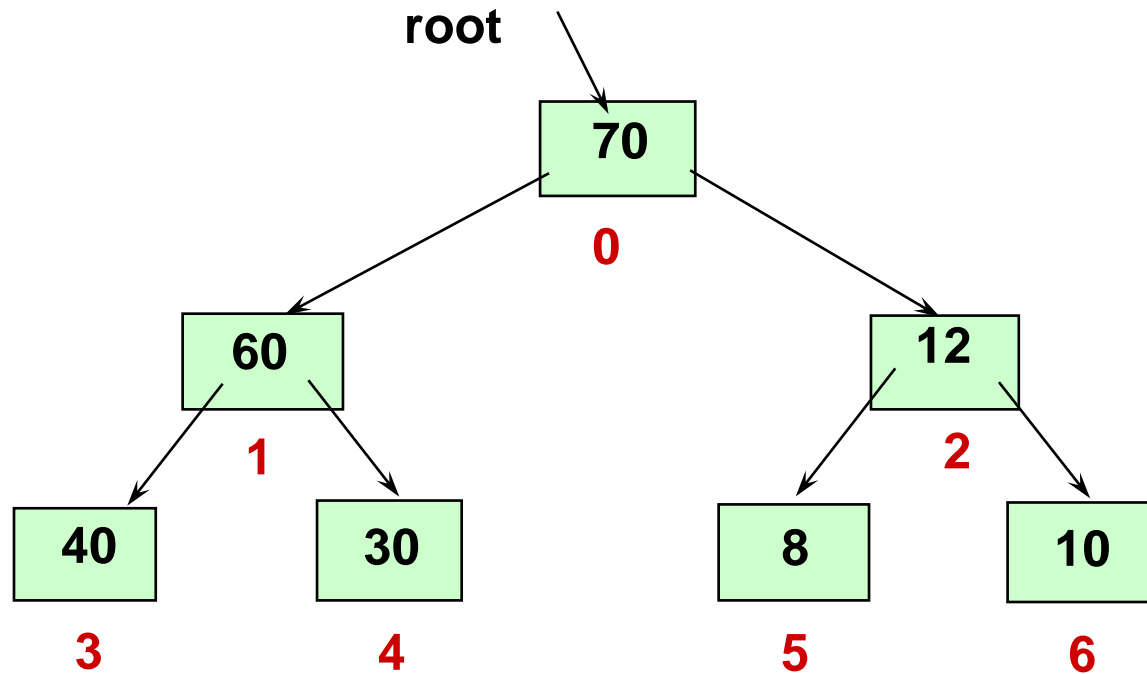
The largest element in a heap

is always found in the root node



The heap can be stored in an array

	values
[0]	70
[1]	60
[2]	12
[3]	40
[4]	30
[5]	8
[6]	10





Heap Sort Approach

First, make the unsorted array into a heap by satisfying the order property. Then repeat the steps below until there are no more unsorted elements.

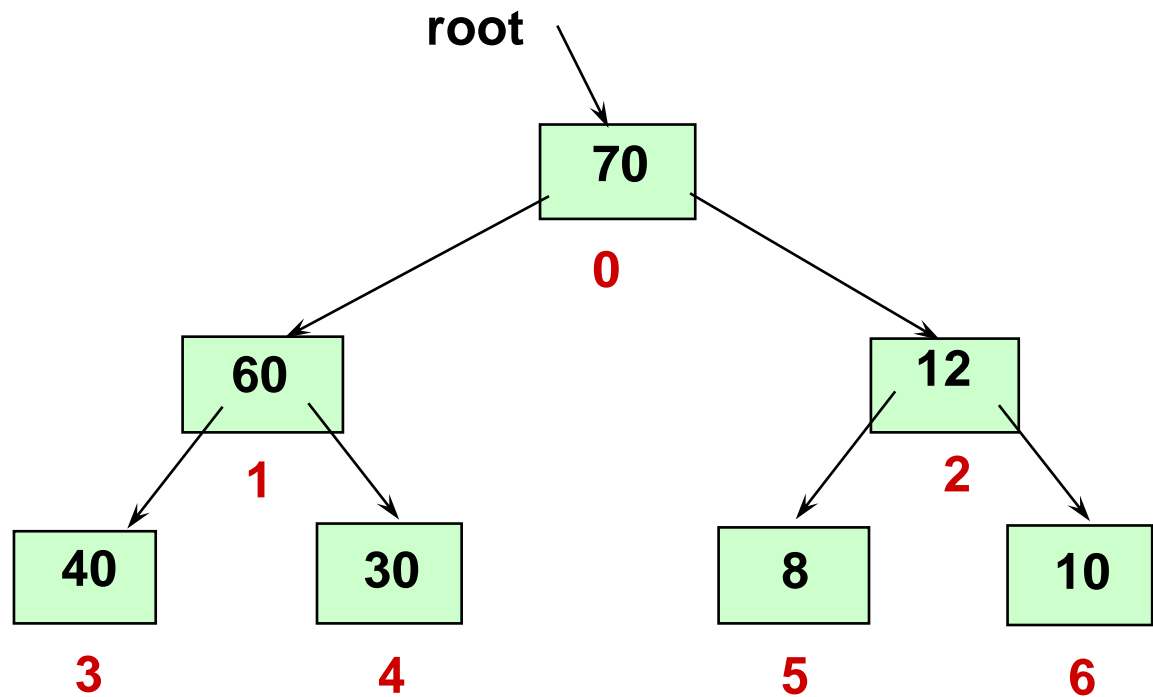
Take the root (maximum) element off the heap by swapping it into its correct place in the array at the end of the unsorted elements.

Reheap the remaining unsorted elements.
(This puts the next-largest element into the root position).



After creating the original heap

	values
[0]	70
[1]	60
[2]	12
[3]	40
[4]	30
[5]	8
[6]	10

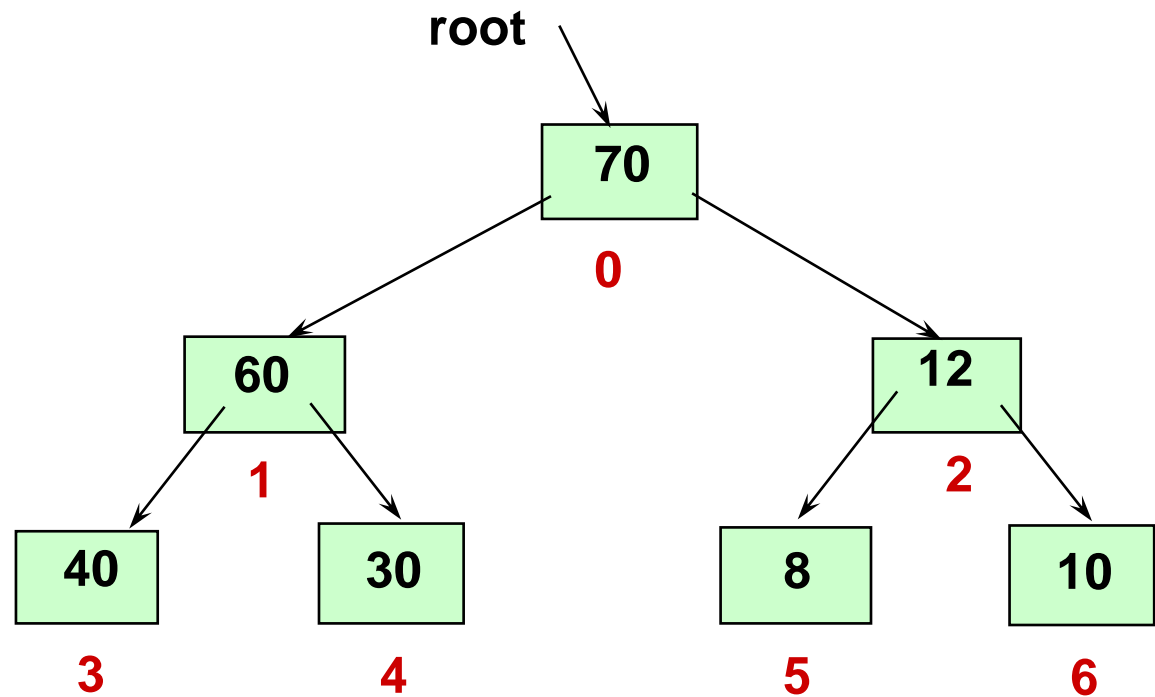




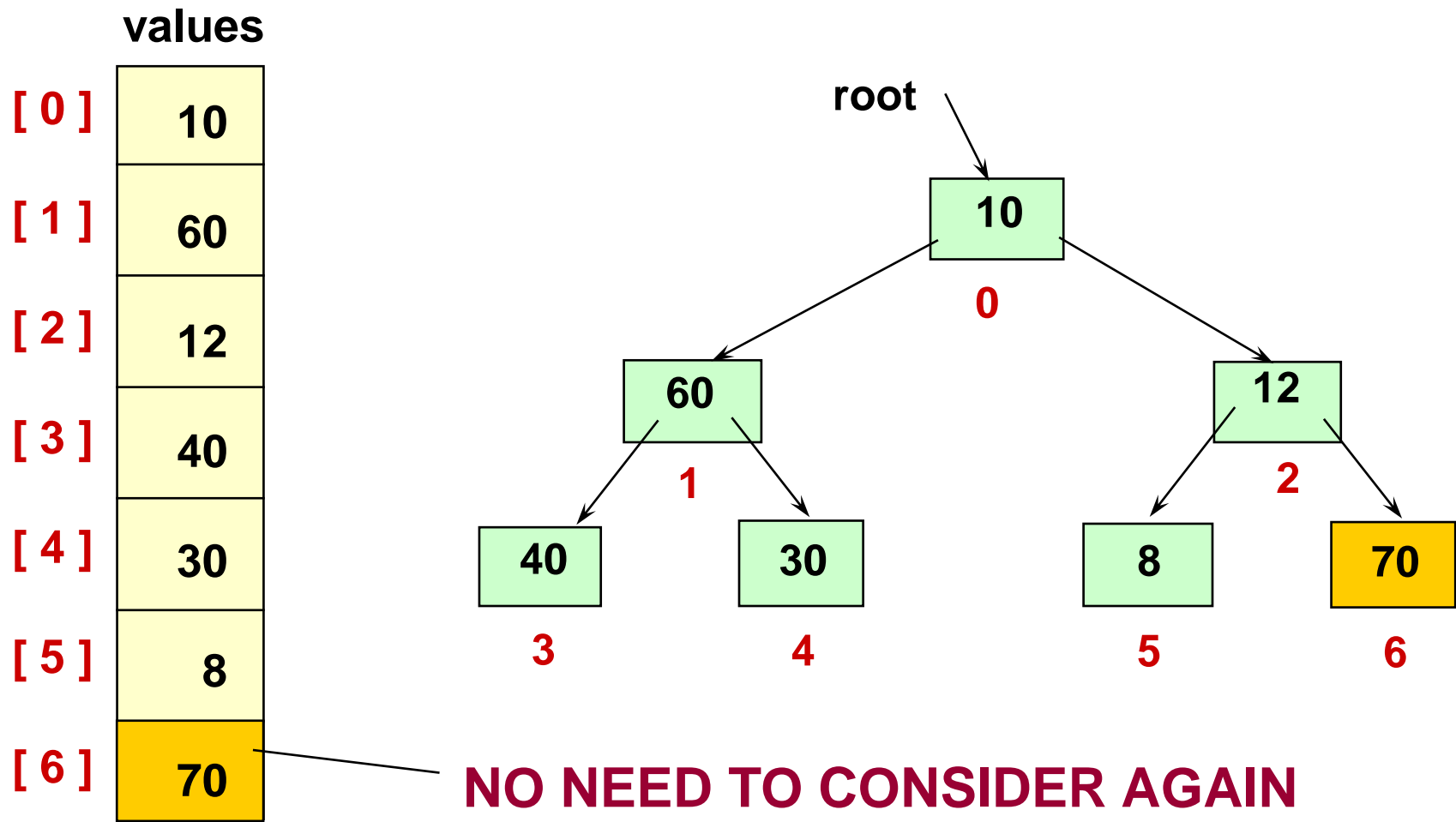
Swap root element into last place in unsorted array

values

[0]	70
[1]	60
[2]	12
[3]	40
[4]	30
[5]	8
[6]	10



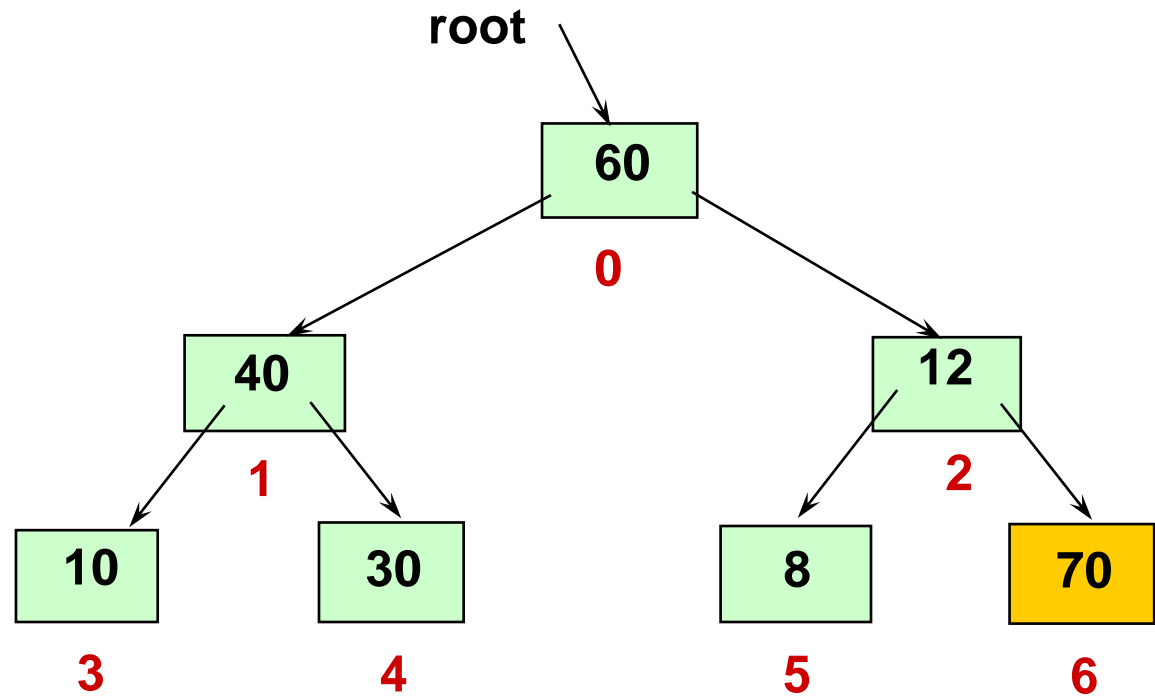
After swapping root element into its place





After reheaping remaining unsorted elements

	values
[0]	60
[1]	40
[2]	12
[3]	10
[4]	30
[5]	8
[6]	70

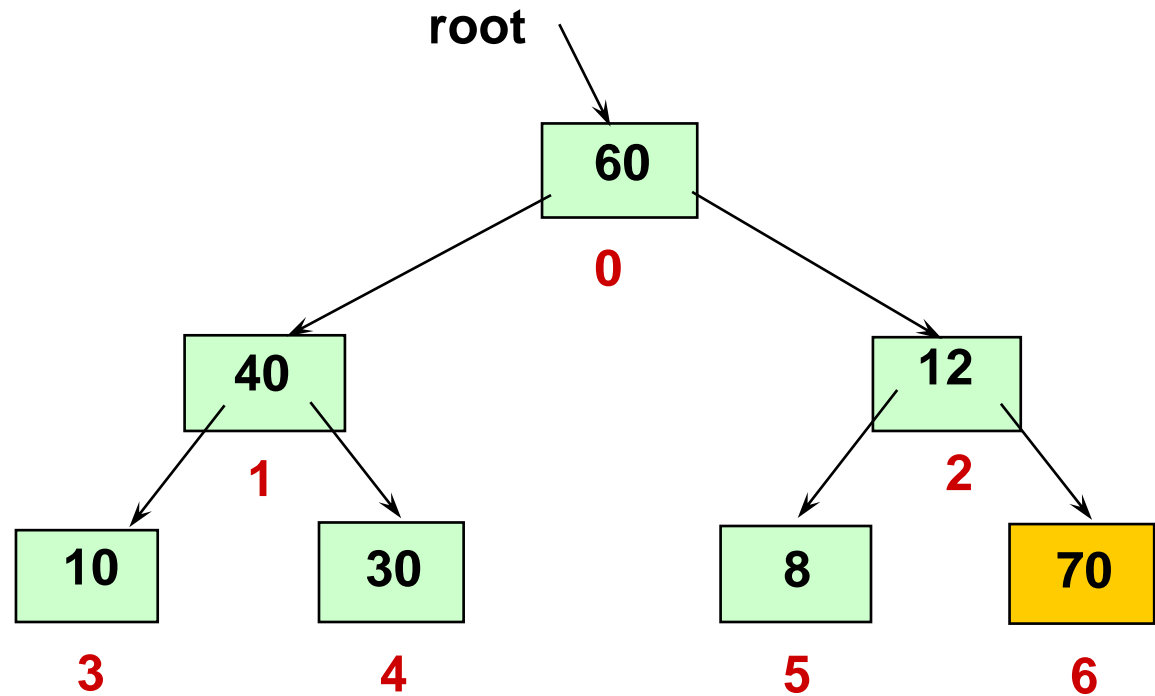




Swap root element into last place in unsorted array

values

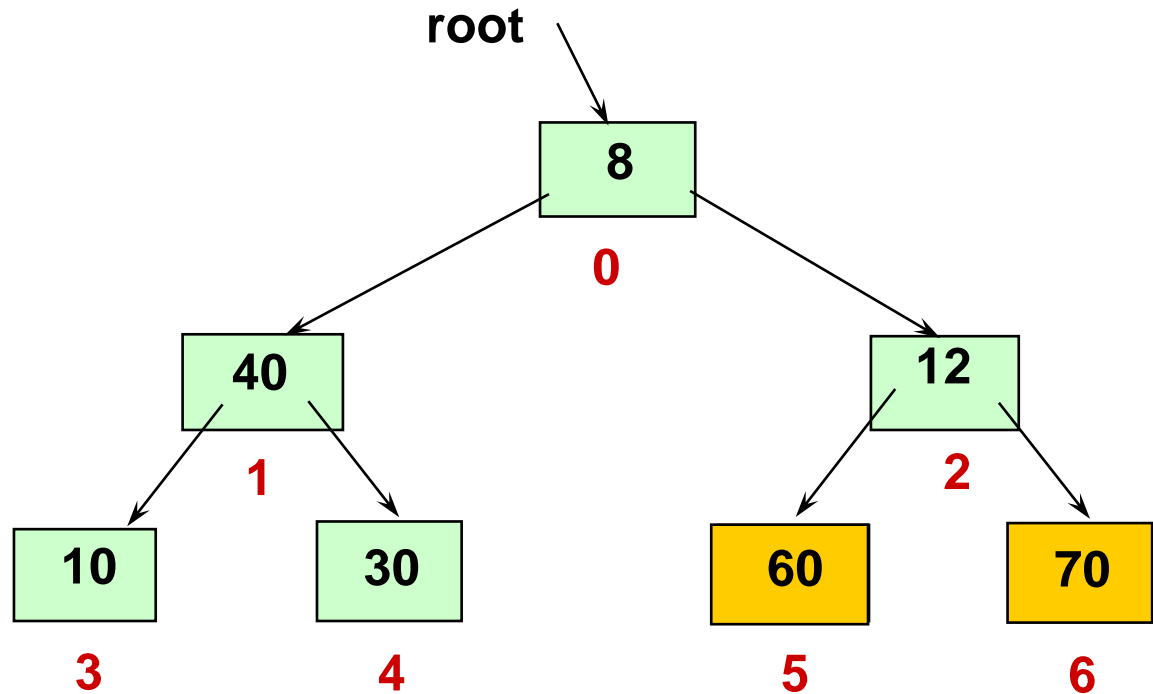
[0]	60
[1]	40
[2]	12
[3]	10
[4]	30
[5]	8
[6]	70





After swapping root element into its place

values	
[0]	8
[1]	40
[2]	12
[3]	10
[4]	30
[5]	60
[6]	70

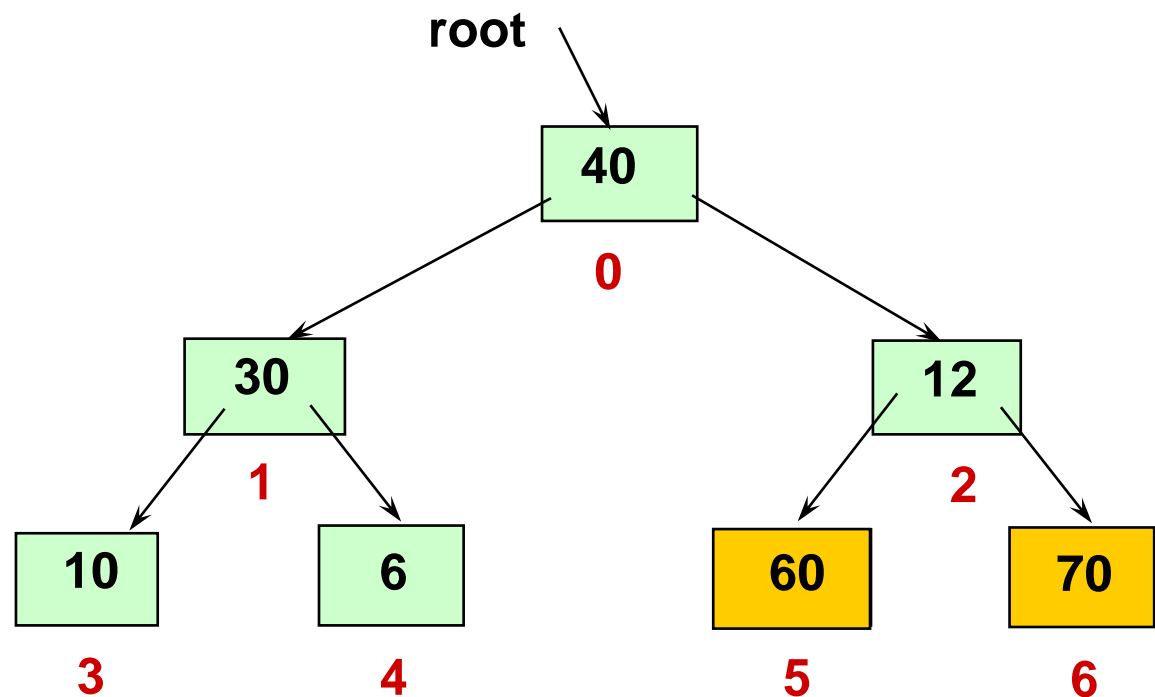


NO NEED TO CONSIDER AGAIN



After reheaping remaining unsorted elements

values	
[0]	40
[1]	30
[2]	12
[3]	10
[4]	6
[5]	60
[6]	70

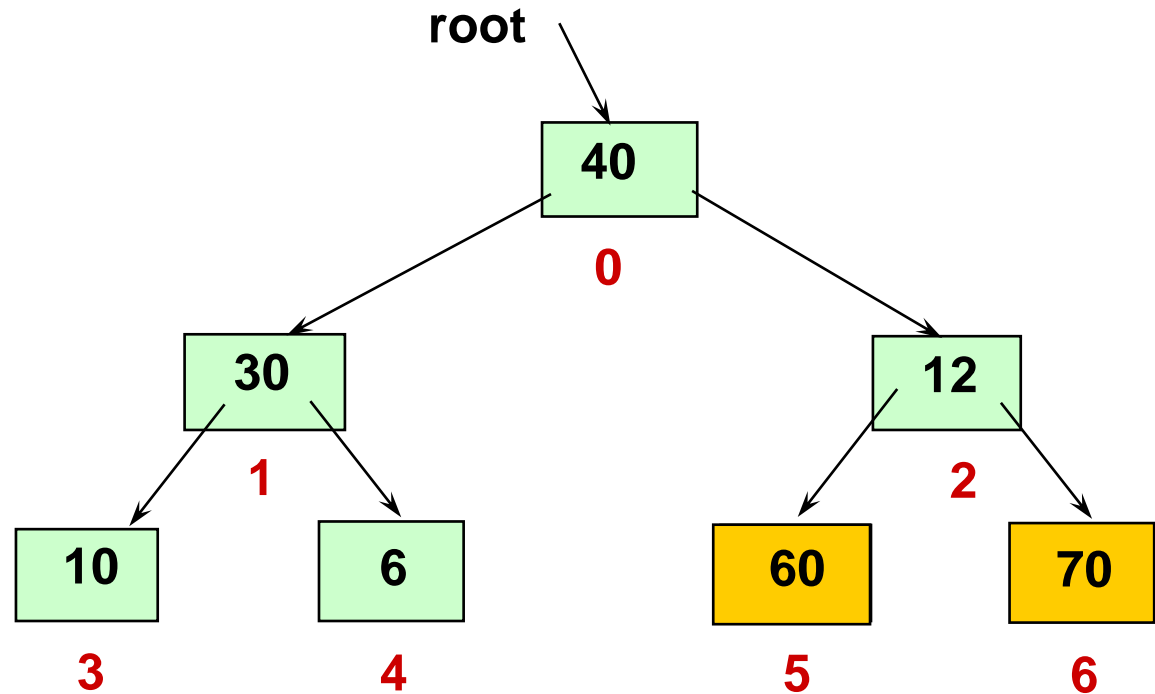




Swap root element into last place in unsorted array

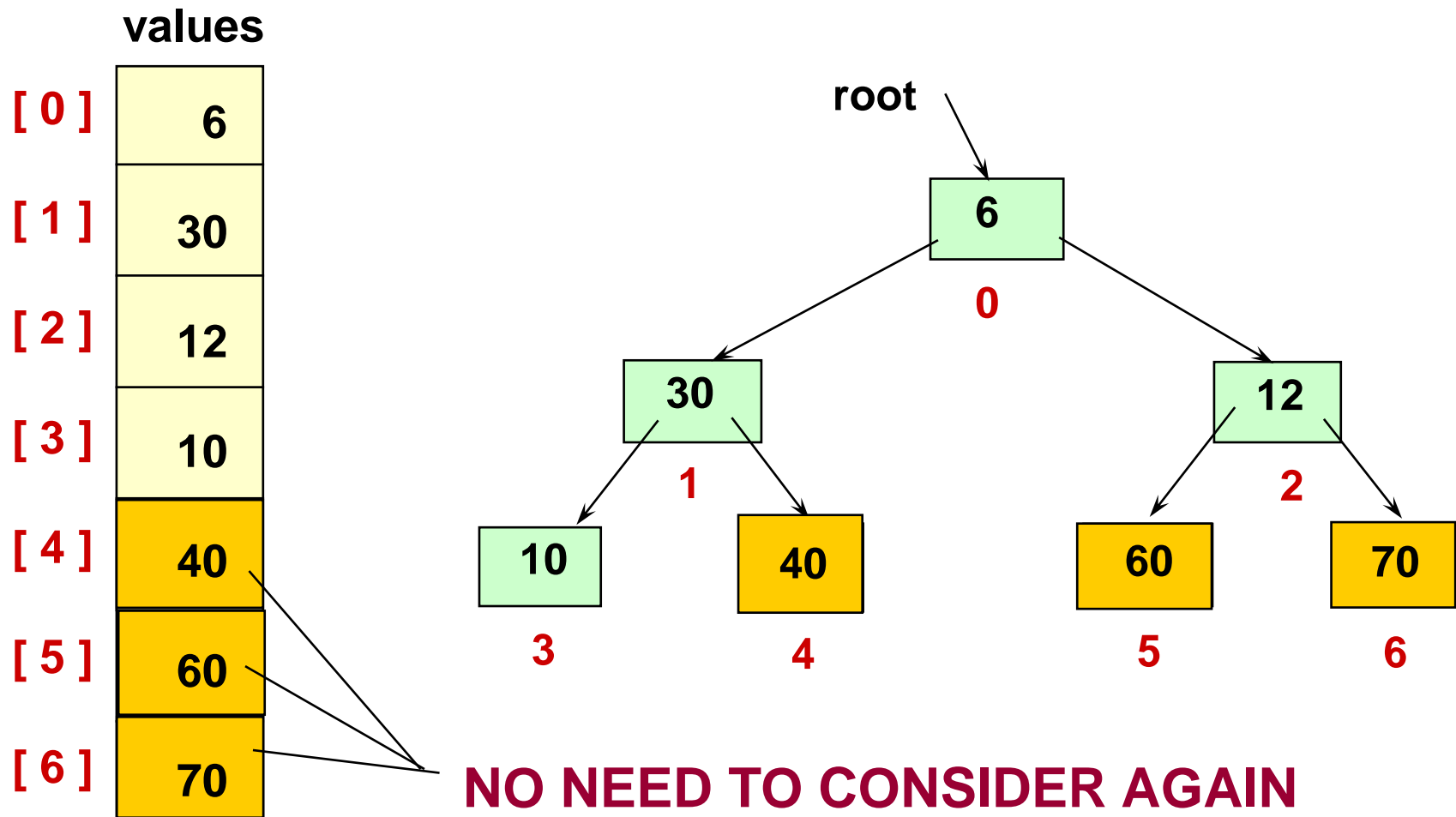
values

[0]	40
[1]	30
[2]	12
[3]	10
[4]	6
[5]	60
[6]	70





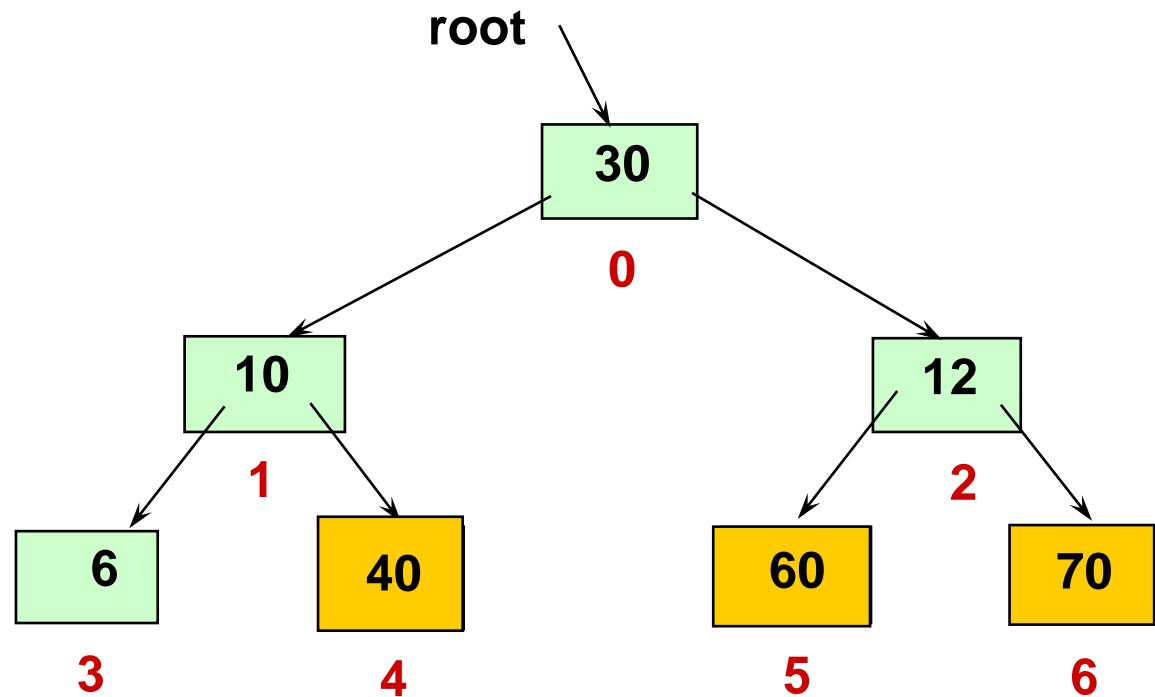
After swapping root element into its place





After reheaping remaining unsorted elements

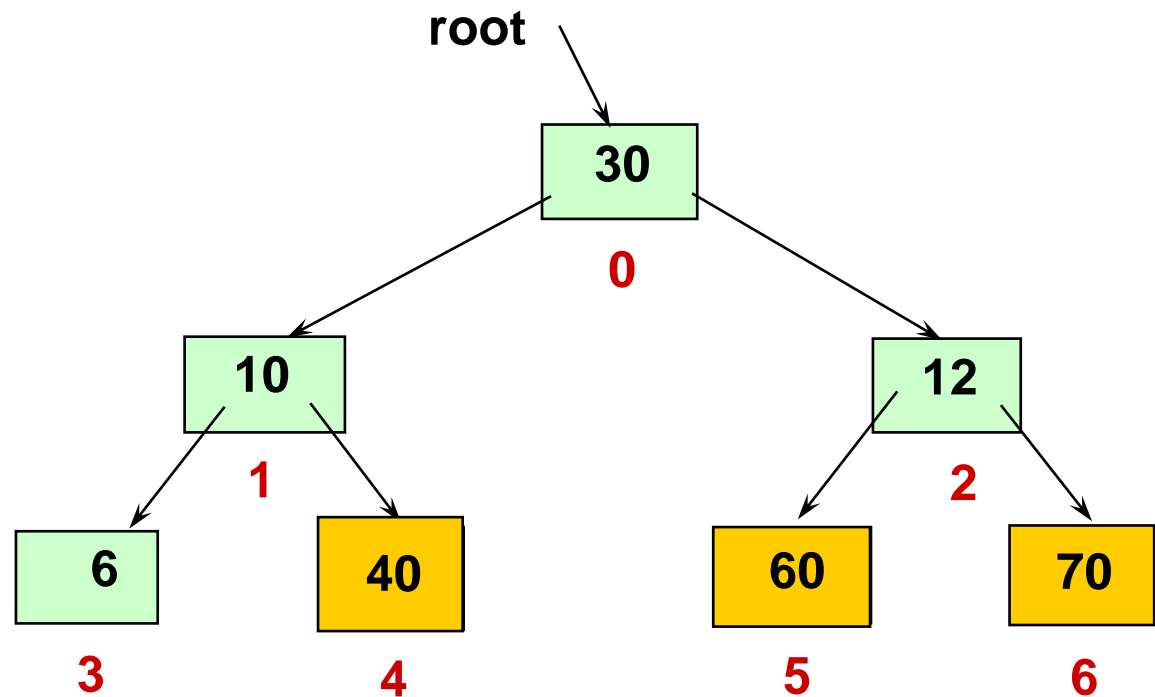
	values
[0]	30
[1]	10
[2]	12
[3]	6
[4]	40
[5]	60
[6]	70



Swap root element into last place in unsorted array

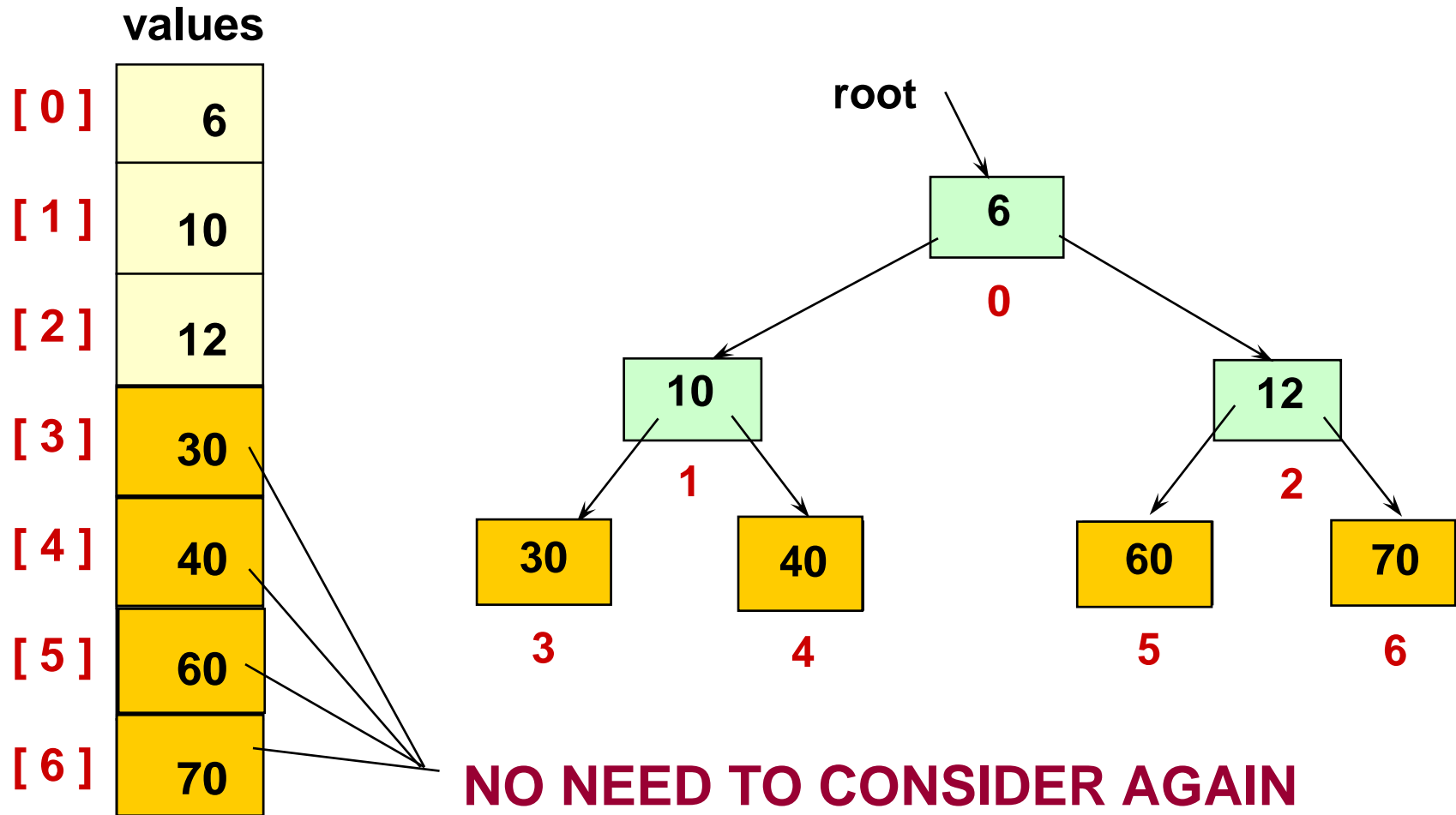
values

[0]	30
[1]	10
[2]	12
[3]	6
[4]	40
[5]	60
[6]	70





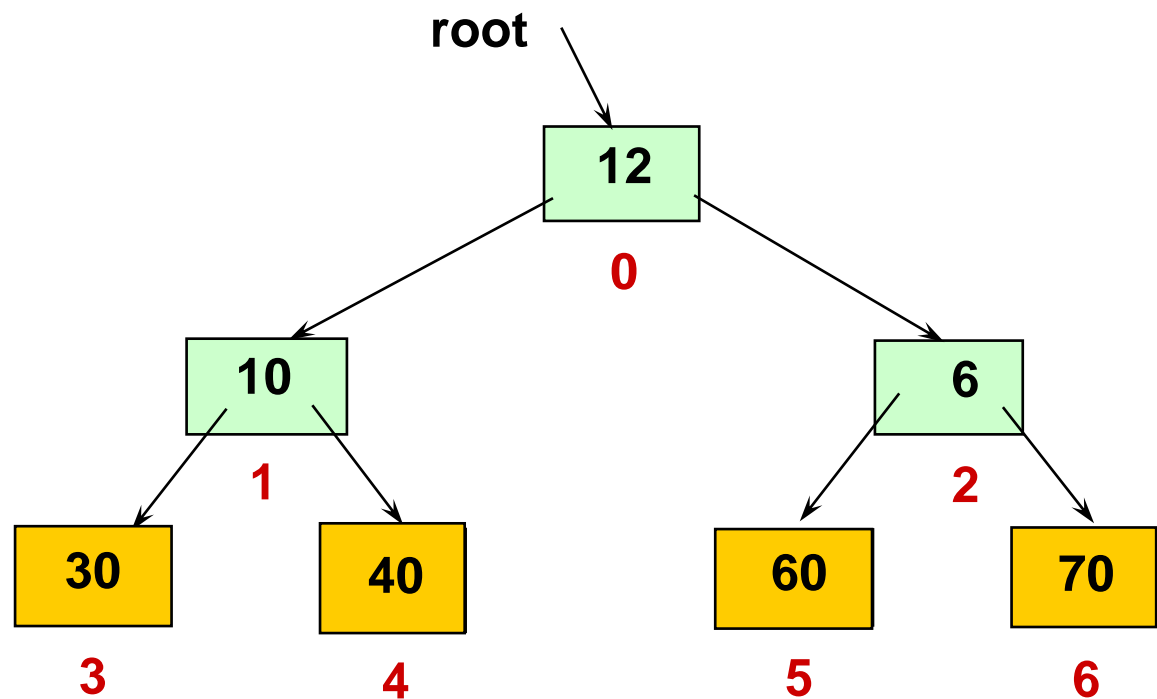
After swapping root element into its place





After reheaping remaining unsorted elements

values	
[0]	12
[1]	10
[2]	6
[3]	30
[4]	40
[5]	60
[6]	70

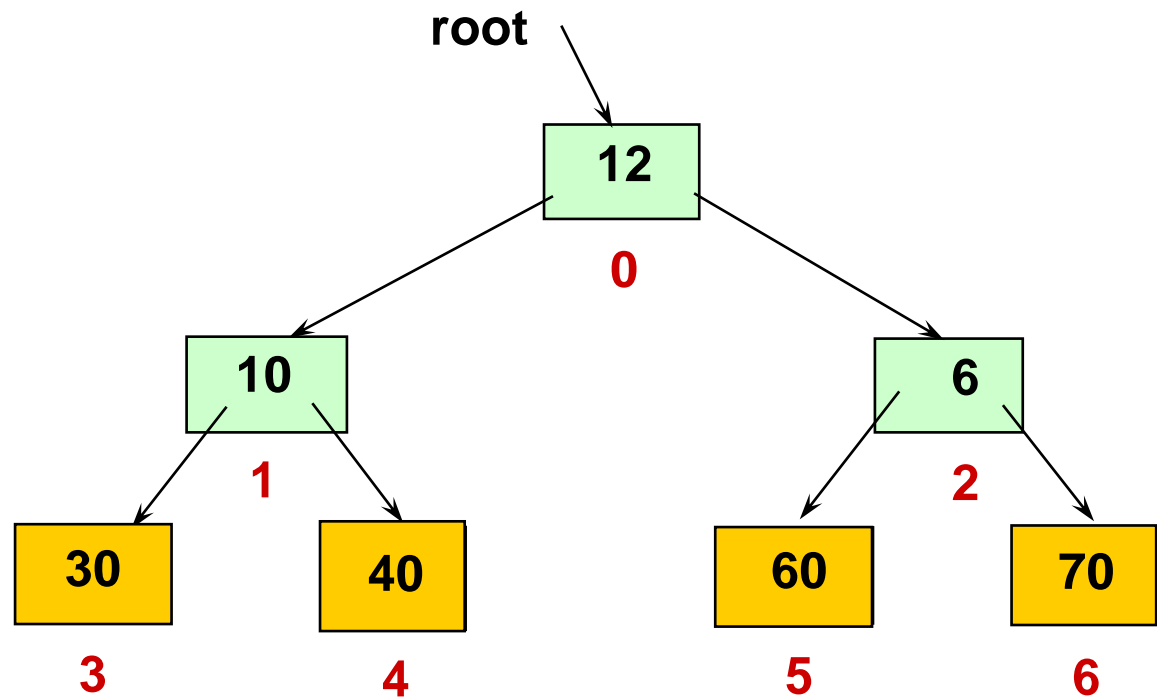




Swap root element into last place in unsorted array

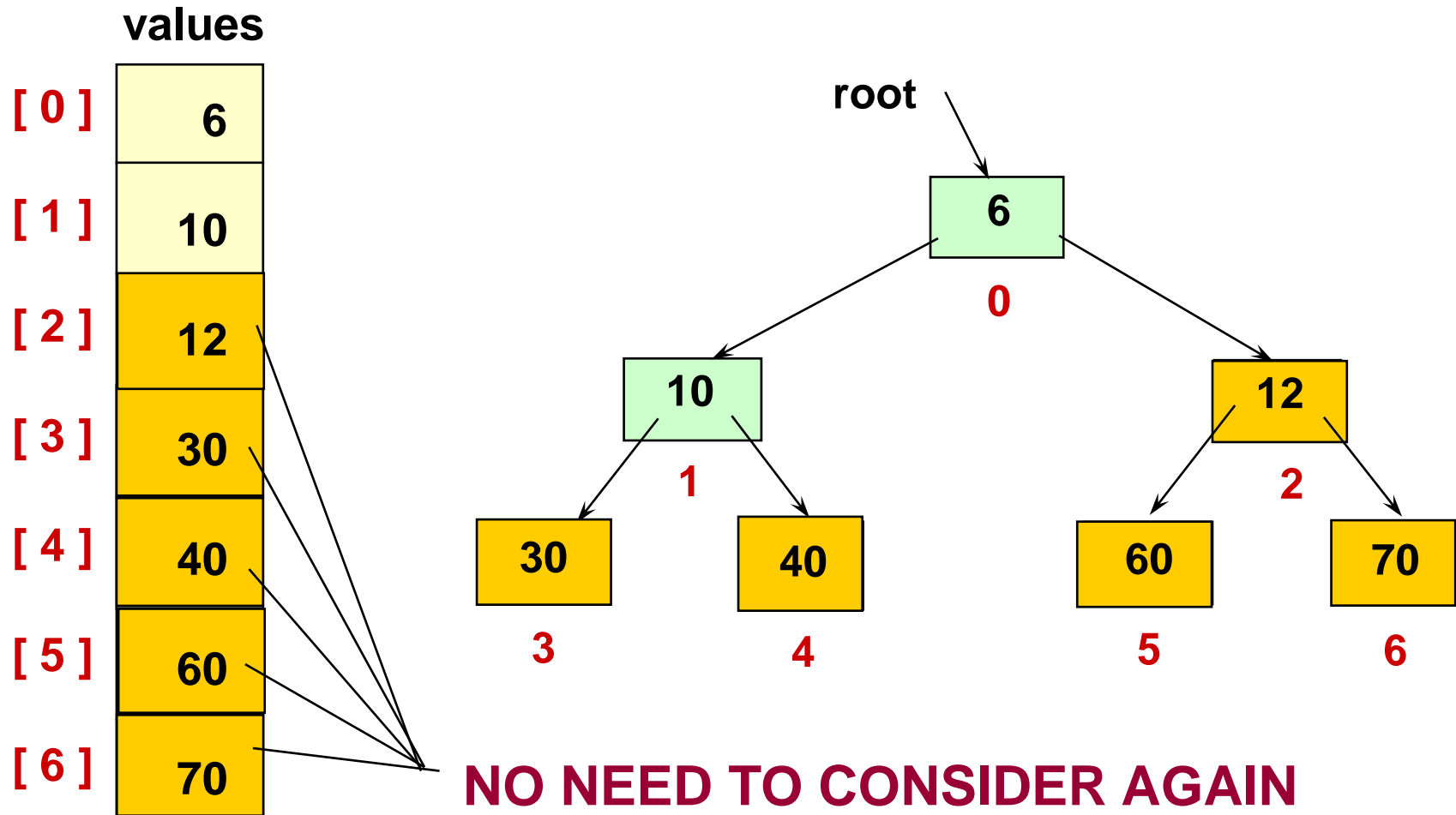
values

[0]	12
[1]	10
[2]	6
[3]	30
[4]	40
[5]	60
[6]	70



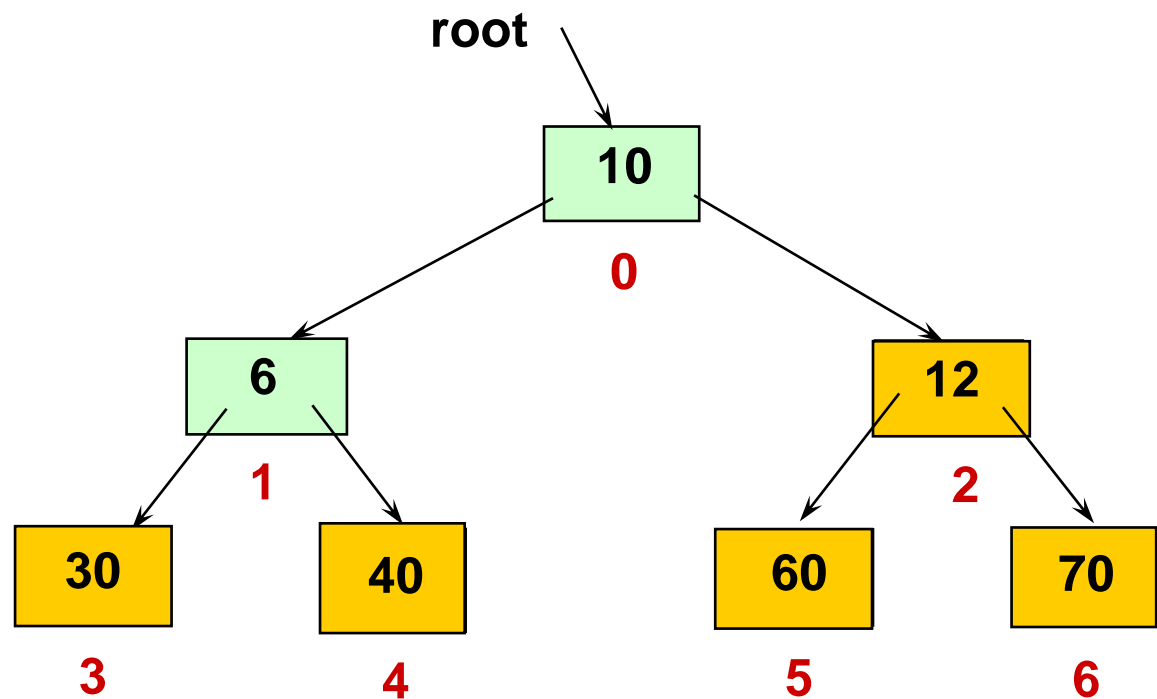


After swapping root element into its place



After reheaping remaining unsorted elements

	values
[0]	10
[1]	6
[2]	12
[3]	30
[4]	40
[5]	60
[6]	70

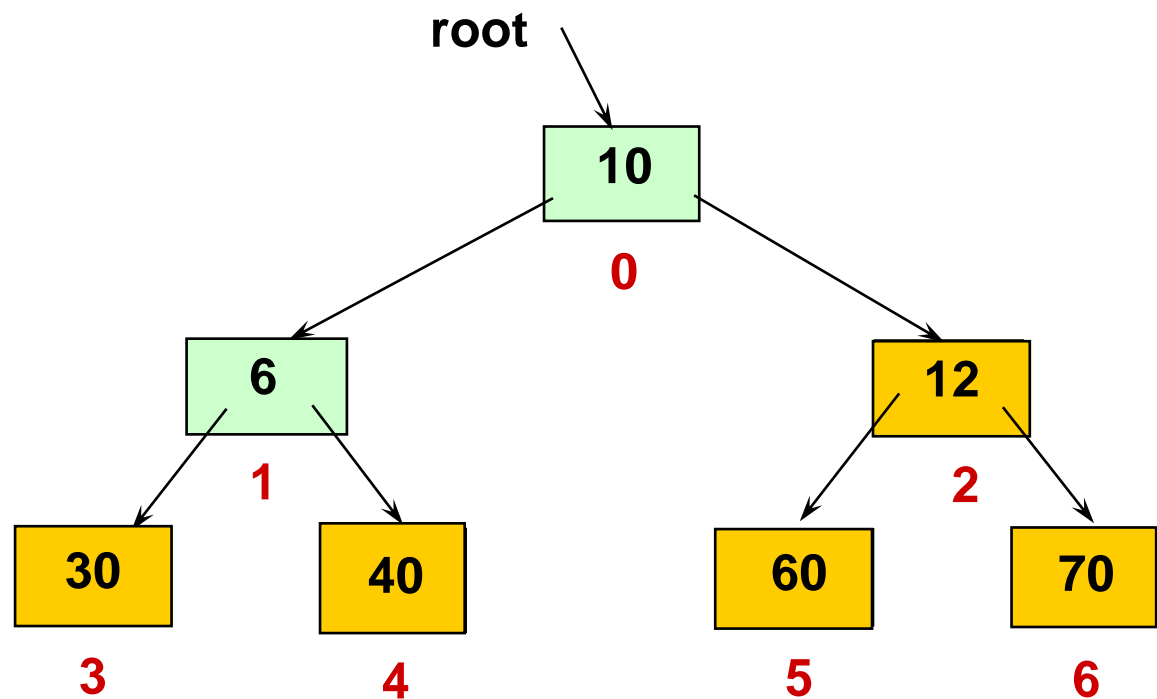




Swap root element into last place in unsorted array

values

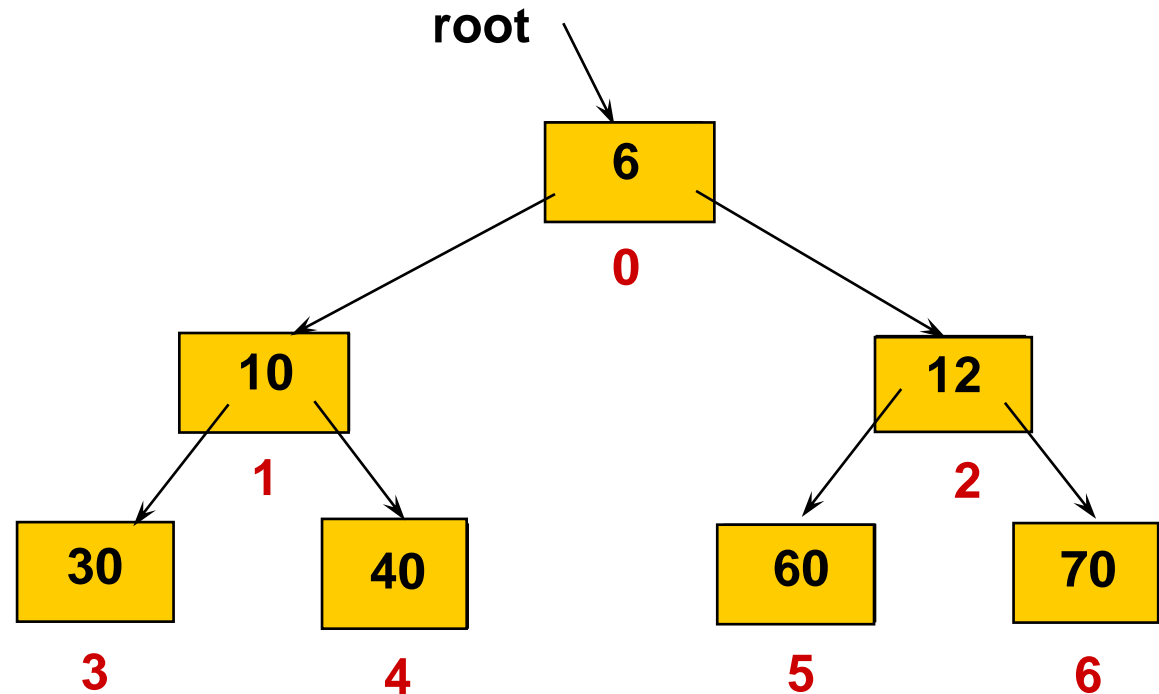
[0]	10
[1]	6
[2]	12
[3]	30
[4]	40
[5]	60
[6]	70





After swapping root element into its place

	values
[0]	6
[1]	10
[2]	12
[3]	30
[4]	40
[5]	60
[6]	70



ALL ELEMENTS ARE SORTED



```
template <class ItemType >
void HeapSort ( ItemType values [ ] , int
    numValues )
// Post: Sorts array values[ 0 . . numValues-1 ] into
// ascending order by key
{
    int index ;

    // Convert array values[0..numValues-1] into a heap
    for (index = numValues/2 - 1; index >= 0; index--)
        ReheapDown ( values , index , numValues - 1 ) ;

    // Sort the array.
    for (index = numValues - 1; index >= 1; index--)
    {
        Swap (values [0] , values[index]);
        ReheapDown (values , 0 , index - 1);
    }
}
```



ReheapDown

```
template< class  ItemType >
void ReheapDown ( ItemType  values [ ],  int  root,
                int  bottom )

//  Pre:  root is the index of a node that may violate the
//         heap order property
//  Post:  Heap order property is restored between root and
//         bottom

{
    int  maxChild ;
    int  rightChild ;
    int  leftChild ;

    leftChild  =  root * 2 + 1 ;
    rightChild =  root * 2 + 2 ;
```



```
if (leftChild <= bottom)           // ReheapDown continued
{
    if (leftChild == bottom)
        maxChild = leftChild;
    else
    {
        if (values[leftChild] <= values [rightChild])
            maxChild = rightChild ;
        else
            maxChild = leftChild ;
    }
    if (values[ root ] < values[maxChild])
    {
        Swap (values[root], values[maxChild]);
        ReheapDown ( maxChild, bottom  ;
    }
}
```


Building Heap From Unsorted Array

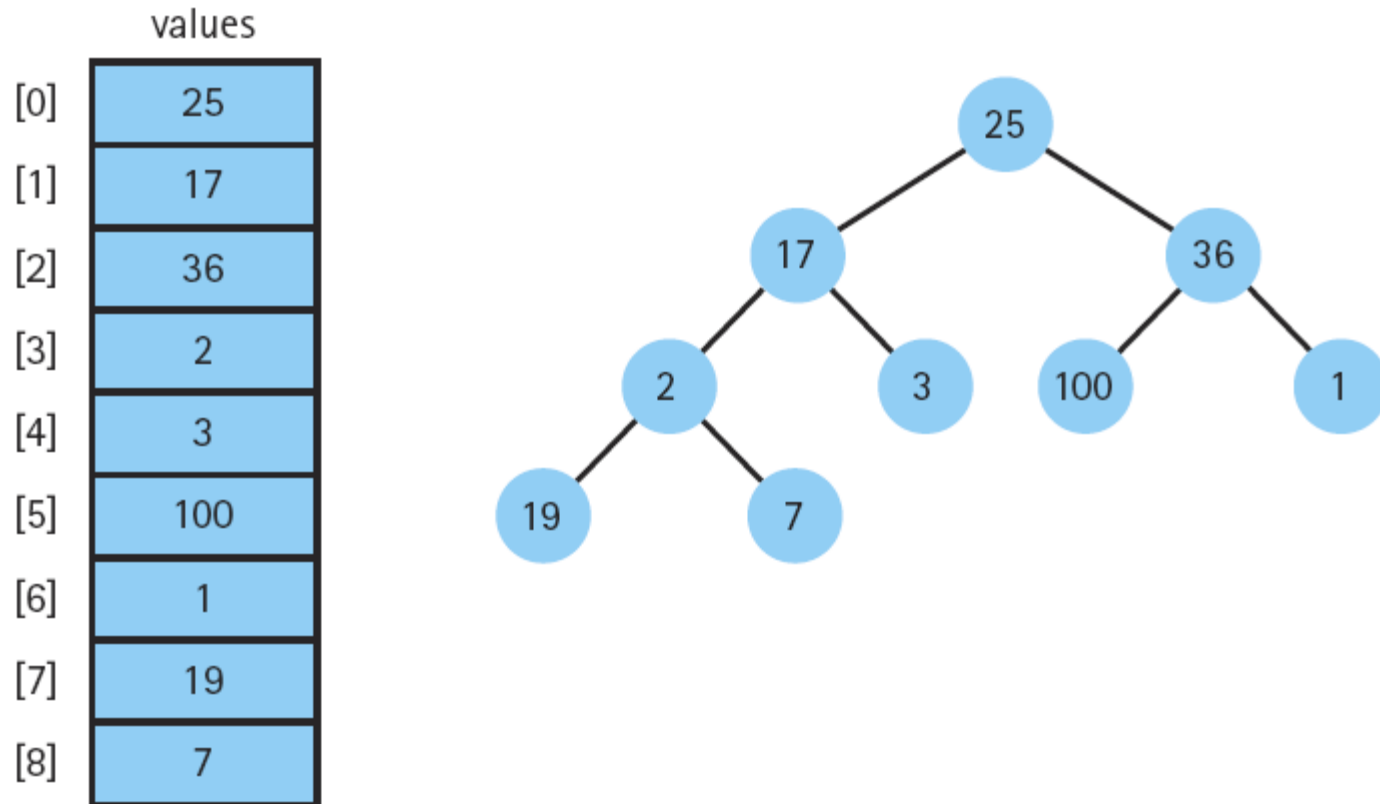
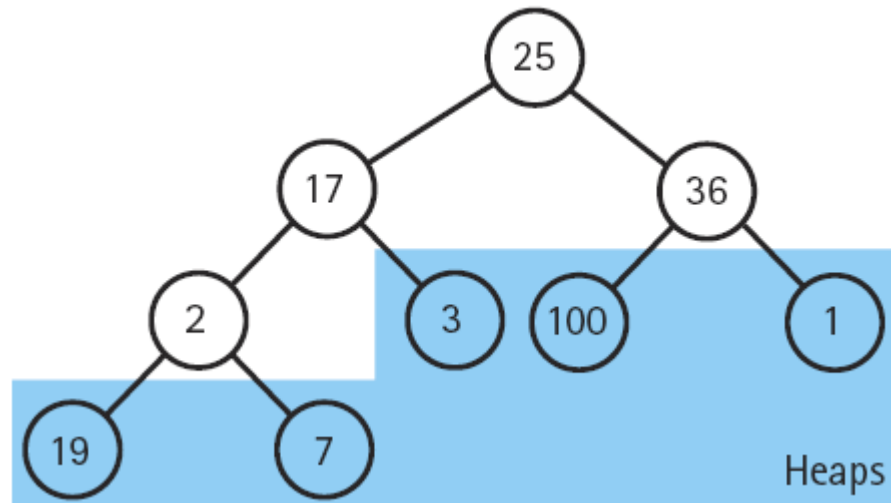


Figure 10.12 An unsorted array and its tree

Building Heap From Unsorted Array (cont'd)

Leaf nodes are already heaps

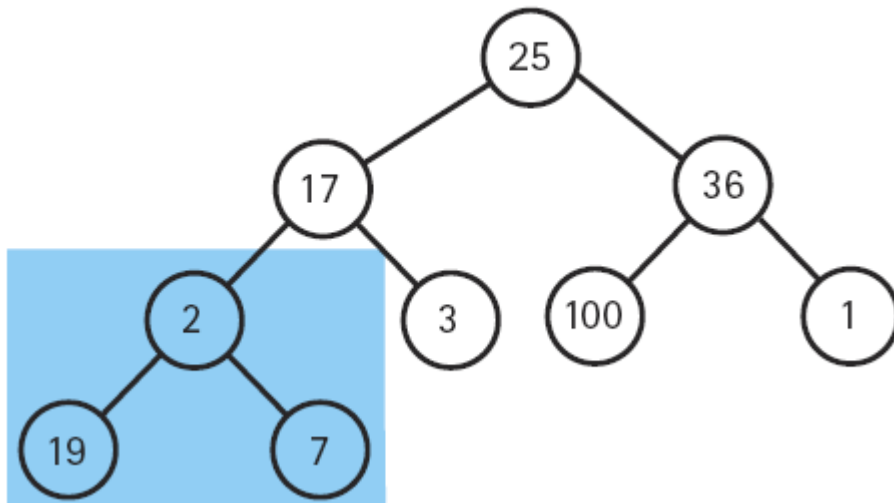
(a)



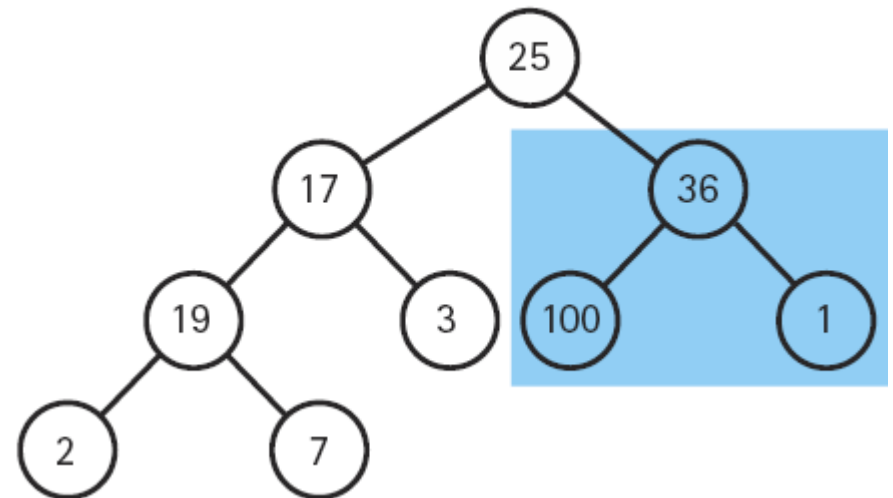
Building Heap From Unsorted Array (cont'd)

The subtrees rooted at first nonleaf nodes are almost heaps

(b)



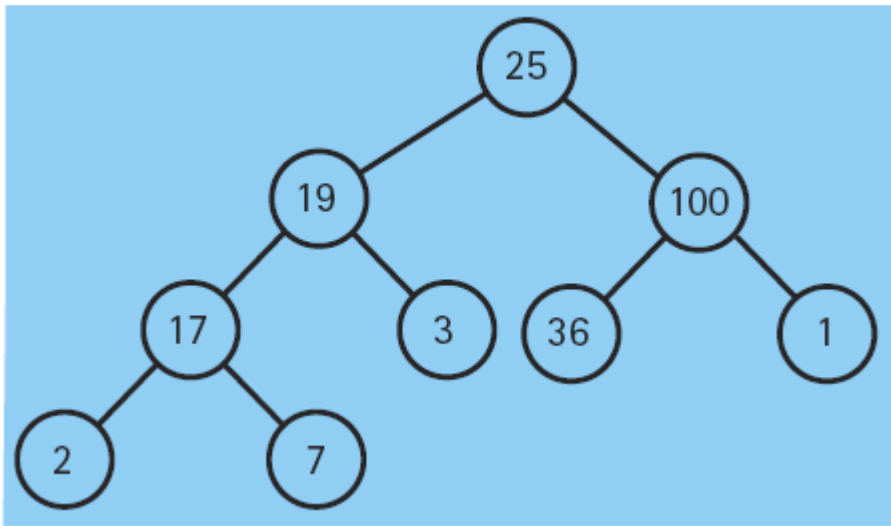
(c)



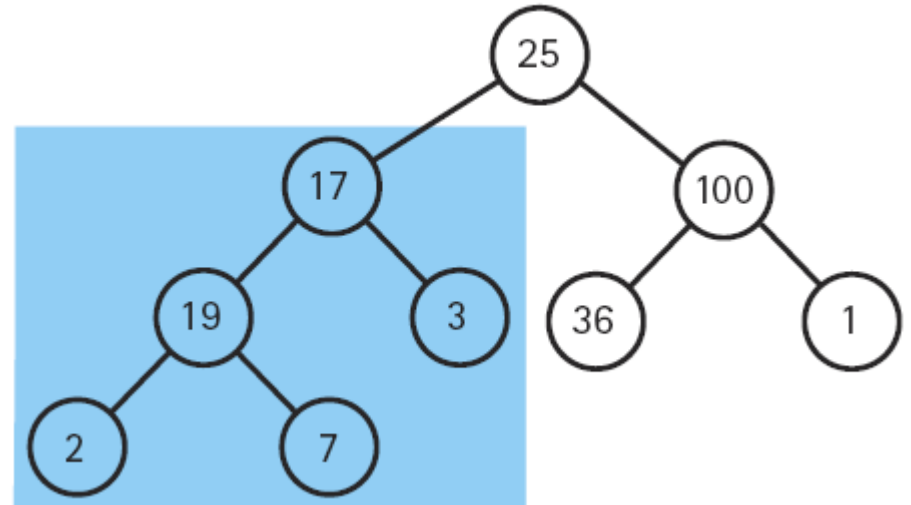
Building Heap From Unsorted Array (cont'd)

Move up a level in the tree and continue reheaping until we reach the root node

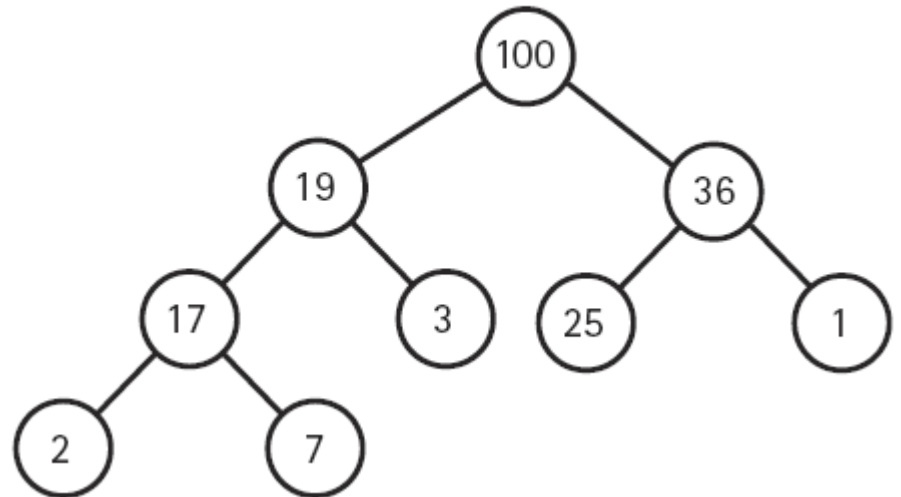
(e)



(d)



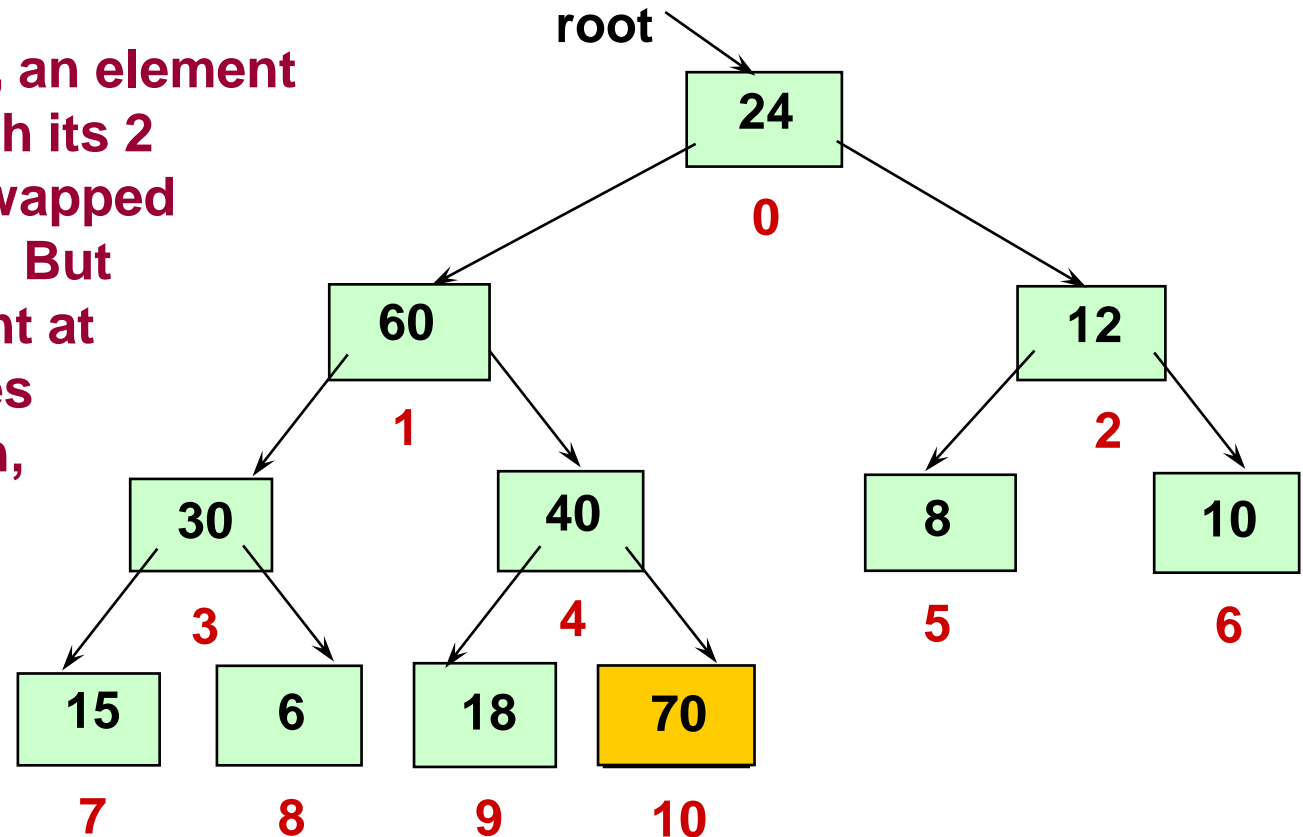
(f) Tree now represents a heap





Heap Sort: How many comparisons?

In reheap down, an element is compared with its 2 children (and swapped with the larger). But only one element at each level makes this comparison, and a complete binary tree with N nodes has only $O(\log_2 N)$ levels.





Heap Sort of N elements: How many comparisons?

$(N/2) * O(\log N)$ compares to create original heap

$(N-1) * O(\log N)$ compares for the sorting loop

$= O(N * \log N)$ compares total