CSE 3318

Week of 08/07/2023

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Linked List vs Binary Tree

Linked List Node

Binary Tree Node

```
typedef struct node
typedef struct node
   int node number;
                                   int node number;
   struct node *next ptr;
                                   struct node *left ptr;
                                   struct node *right ptr;
NODE;
                               NODE;
NODE *LinkedListHead;
                                     *root;
                               NODE
```

Binary Tree vs Linked List

```
NewNode = malloc(sizeof(NODE));
          NewNode->node number = NodeNumber;
Linked List
          NewNode->next ptr = NULL;
          NewNode = malloc(sizeof(NODE));
          NewNode->node number = NodeNumber;
Binary Tree
          NewNode->left ptr = NULL;
          NewNode->right ptr = NULL;
```

Binary Tree vs Linked List

Add a node to the end of a linked list

```
NewNode = malloc(sizeof(NODE));
NewNode->node_number = NodeNumber;
```

Set the pointer of the last node to the new node

```
TempPtr->next_ptr = NewNode;
```

Add a node to a binary tree

```
NewNode = malloc(sizeof(NODE));
NewNode->node_number = NodeNumber;
NewNode->left_ptr = NULL;
NewNode->right ptr = NULL;
```

Set the parent node's left or right ptr to the address of the new child

```
/* Allocates memory for a new node with the given data and sets the left and
   right pointers to NULL */
NODE *CreateNewNode(int NodeNumber)
                                               typedef struct node
      // Allocate memory and assign pointers
      NODE *node = malloc(sizeof(NODE));
                                                  int node number;
                                                  struct node *left ptr;
      node->left ptr = NULL;
                                                  struct node *right ptr;
      node->right ptr = NULL;
                                               NODE;
      // Assign data to this node
      node->node number = NodeNumber;
      printf("Node Number %d %p\n", NodeNumber, node);
      return (node);
```

Node Number 1 0x16a67010

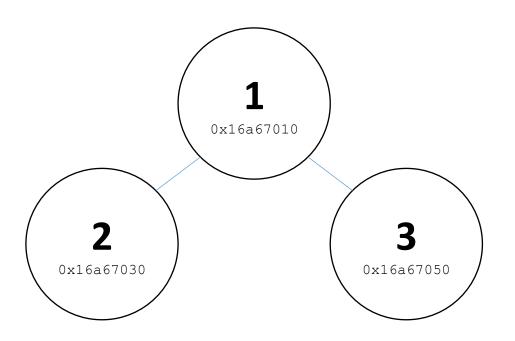
Reminder!!

When you typedef a structure that contains a pointer to the structure, you must use typedef struct node and not just typedef struct

```
/* declare root of tree */
NODE *root;
                                                0x16a67010
/* create root and label with "1" */
root = CreateNewNode(1);
                                           NULL
// Print pointer values
printf("\nleft ptr(1) %p\tright ptr(1) %p\n",
        root->left ptr, root->right ptr);
                           right ptr(1) (nil)
 left ptr(1) (nil)
```

NULL

```
root->left_ptr = CreateNewNode(2);
root->right_ptr = CreateNewNode(3);
```



```
printf("\nleft_ptr(2) %p\tright_ptr(3) %p\n",root->left_ptr, root->right_ptr);
```

```
Node Number 2 0x16a67030
Node Number 3 0x16a67050

left_ptr(2) 0x16a67030 right_ptr(3) 0x16a67050
```

BinaryTreeDemo.c

Binary Tree vs Binary Search Tree

Binary Tree

Each node can have a maximum of two child nodes and there is no order to how the nodes are organized in the tree.

Binary Search Tree

Each node can have a maximum of two child nodes and there is a relative order to how the nodes are organized in the tree.

Binary Search Tree

A binary search tree (with no duplicate node values) has the characteristic that the values in any left subtree are less than the value in its parent node, and the values in any right subtree are greater than the value in its parent node.

The shape of the binary search tree that corresponds to a set of data can vary, depending on the order in which the values are inserted into the tree.

Binary Search Tree

What makes a binary tree a binary search tree?

- All nodes in the left subtree are less than the root
- All nodes in the right subtree are greater than the root
- Each subtree is itself a binary search tree
- No duplicates allowed*

Binary Search Tree (BST)

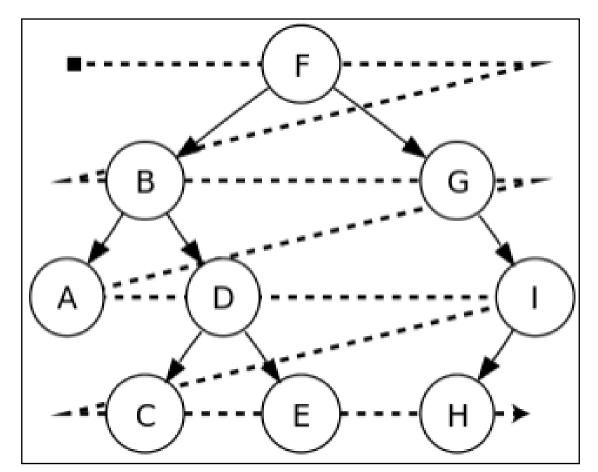
• Linear data structures (linked list, queues and stacks) are traversed in a linear order. Tree structures are traversed in multiple ways - from any given node, there is more than one possible next node in the traversal path

• Tree structures may be traversed in

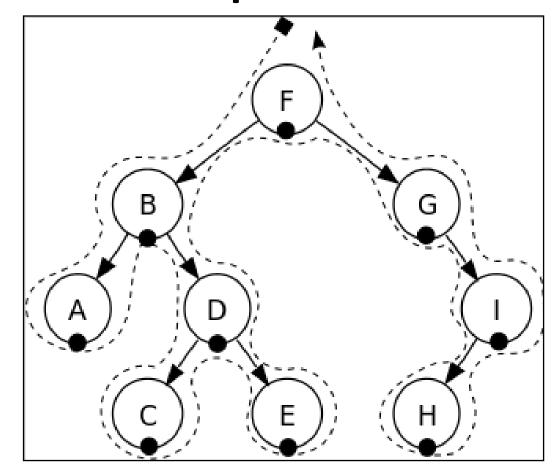
- Breadth-first Order
- Depth-first Order

BST Breadth-first vs Depth-first Traversal

Breadth-first



Depth-first



BST Depth-first Traversals

- Inorder Traversal
 - Gives us the nodes in increasing order
- Preorder Traversal
 - Parent nodes are visited before any of its child nodes
 - Used to create a copy of the tree
 - File systems use it to track your movement through directories
- Postorder Traversal
 - Used to delete the tree
 - File systems use it to delete folders and the files under them

BST Depth-first Traversals

Depth First Tree Traversals

Preorder

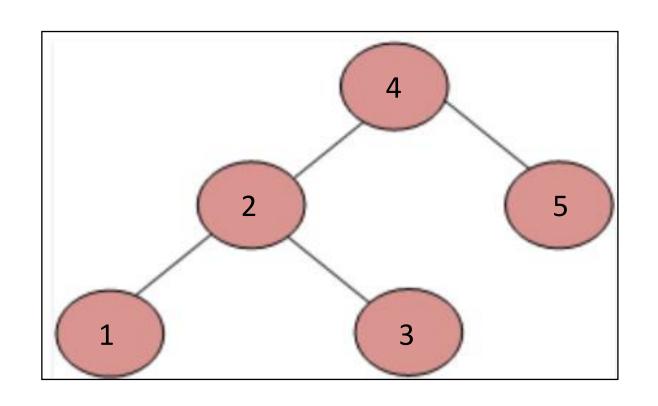
Root, Left, Right 4 2 1 3 5

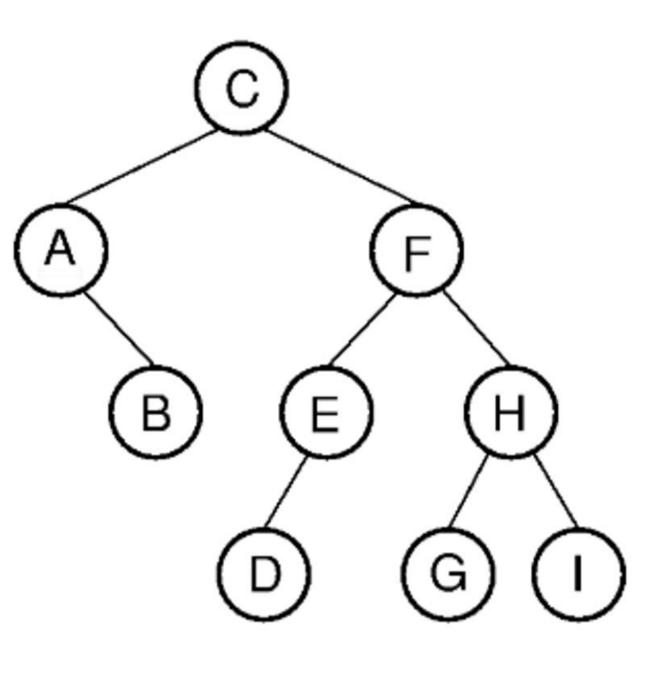
Postorder

Left, Right, Root 1 3 2 5 4

Inorder

Left, Root, Right 12345





Depth First Tree Traversals

Preorder

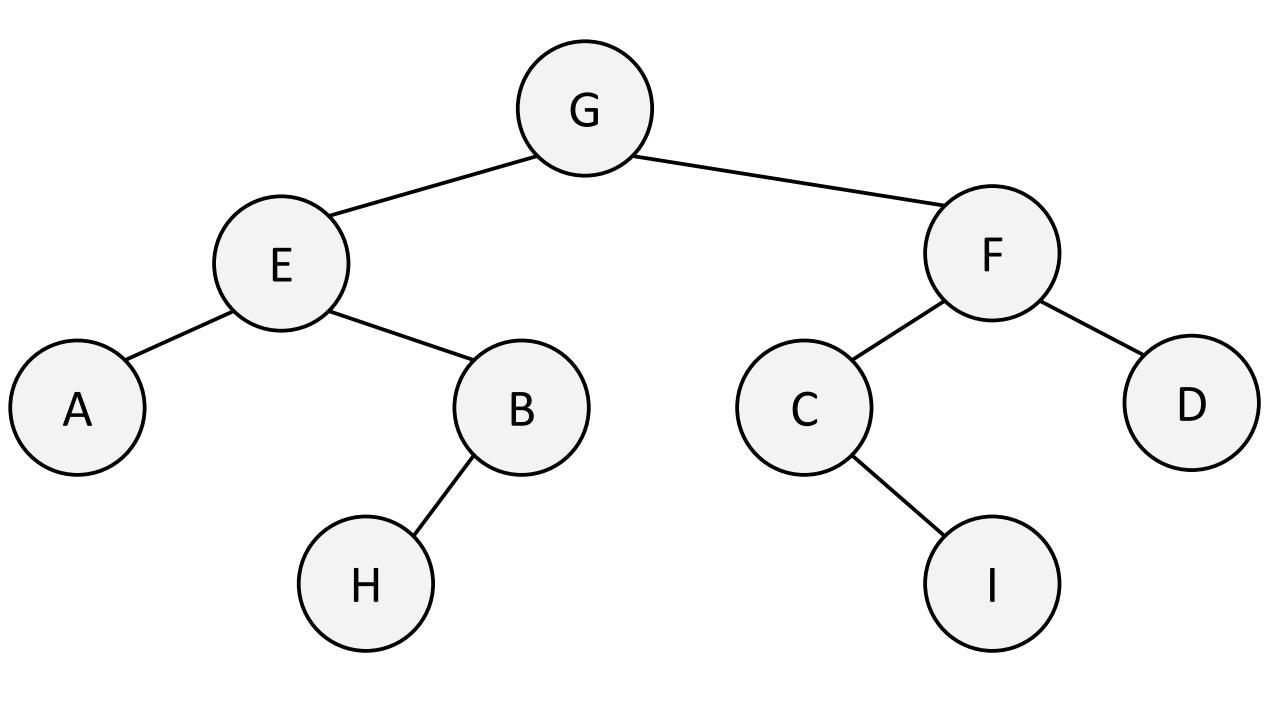
Root, Left, Right CABFEDHGI

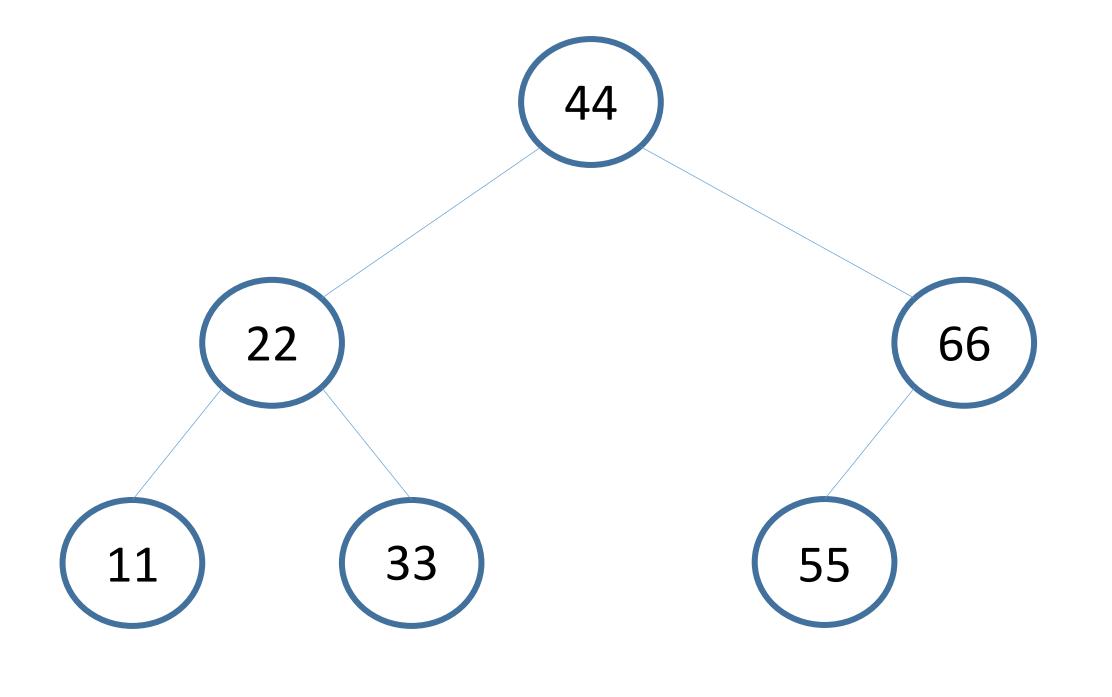
Postorder

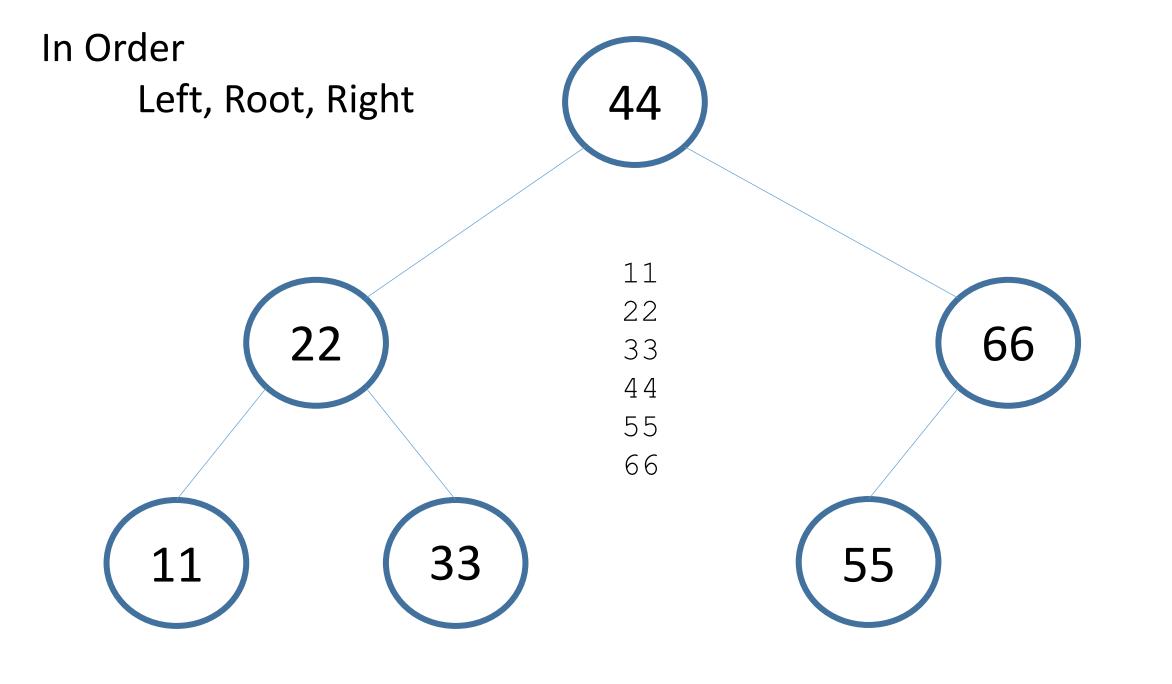
Left, Right, Root B A D E G I H F C

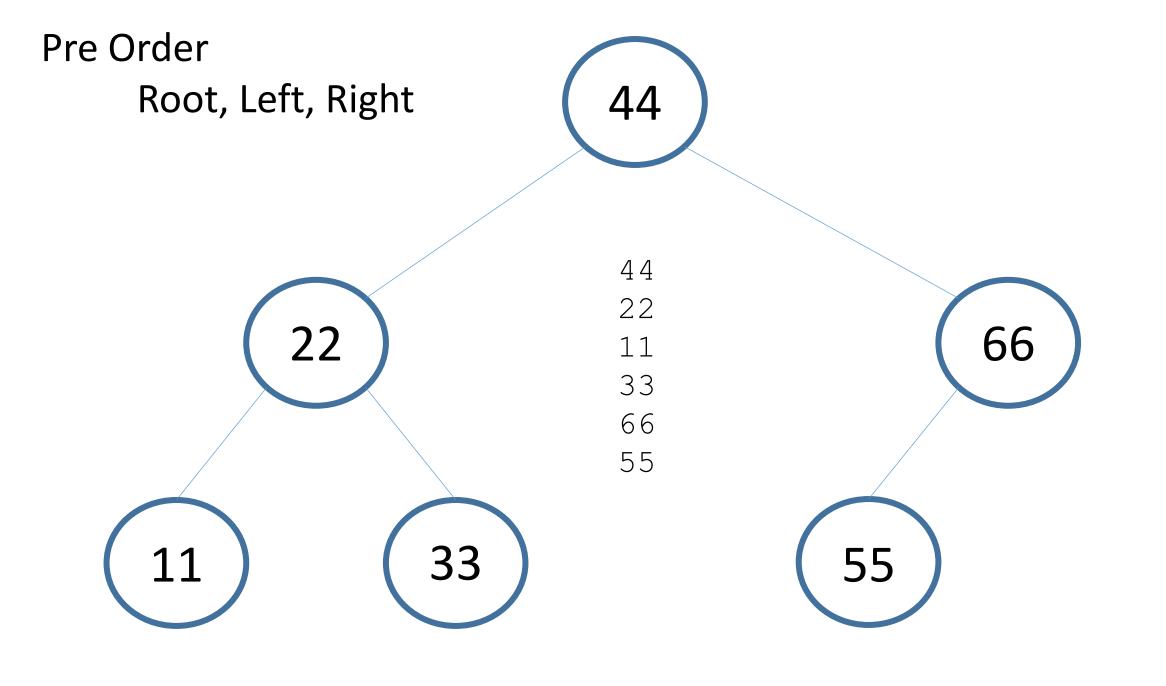
Inorder

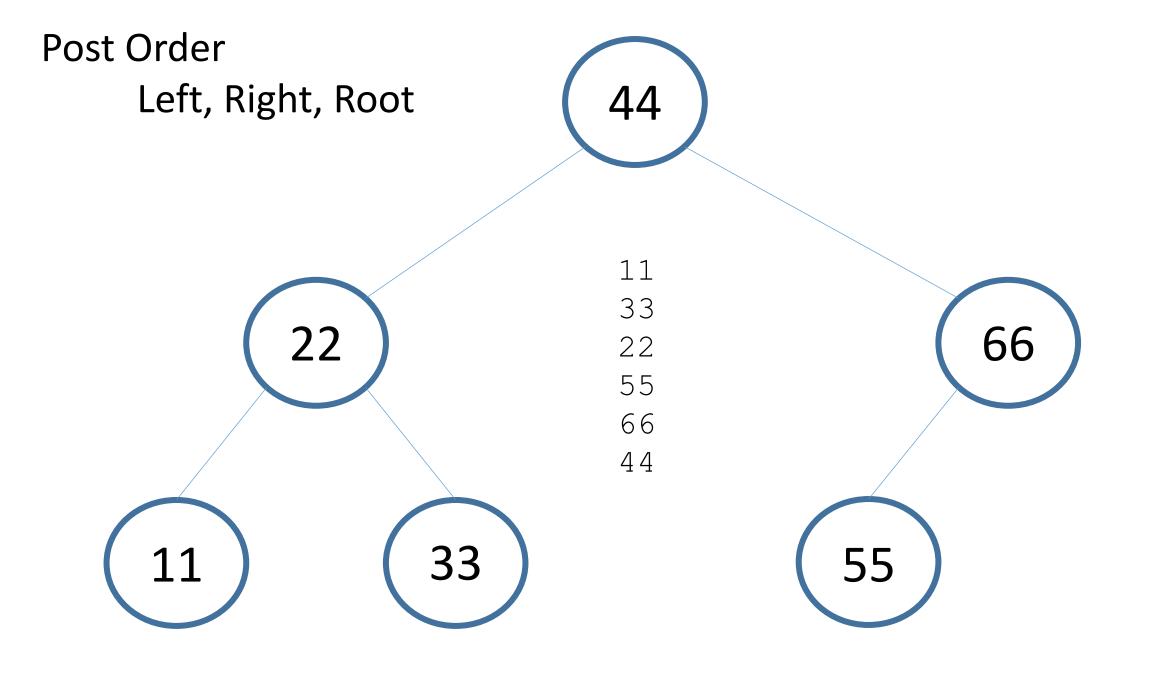
Left, Root, Right ABCDEFGHI

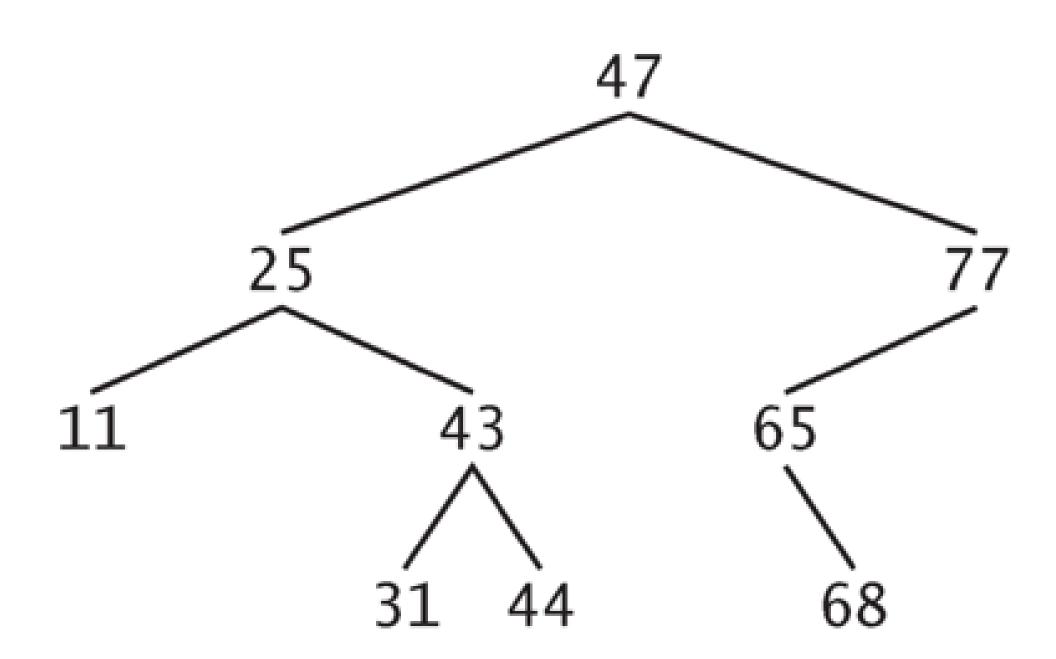












How many nodes in the tree? 9

Enter data for node 1: 47

Enter data for node 2 : 25

Enter data for node 3: 77

Enter data for node 4: 11

Enter data for node 5: 43

Enter data for node 6: 65

Enter data for node 7: 31

Enter data for node 8: 44

Enter data for node 9 : 68

BST Traversal in Inorder

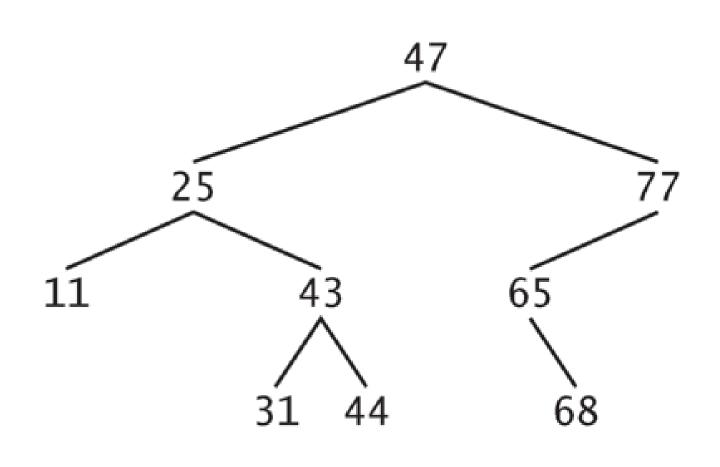
Node4-11	Node2-25	Node7-31
Node5-43	Node8-44	Node1-47
Node6-65	Node9-68	Node3-77

BST Traversal in Preorder

Nodel-47	Node2-25	Node4-11
Node5-43	Node7-31	Node8-44
Node3-77	Node6-65	Node9-68

BST Traversal in Postorder

Node4-11	Node7-31	Node8-44	
Node5-43	Node2-25	Node9-68	
Node6-65	Node3-77	Node1-47	



```
typedef struct node
    int node data;
    struct node *right;
    struct node *left;
NODE;
NODE *root = NULL;
AddBSTNode (&root, node data);
```

```
void AddBSTNode(NODE **current node, int add data)
    if (*current node == NULL)
        *current node = malloc(sizeof(NODE));
        (*current node) -> left = (*current node) -> right = NULL;
        (*current node) -> node data = add data;
    else
        if (add data < (*current node) ->node data )
            AddBSTNode(&(*current node)->left, add data);
        else if (add data > (*current node) ->node data )
            AddBSTNode(&(*current node)->right, add data);
        else
            printf(" Duplicate Element !! Not Allowed !!!");
```

```
void Inorder(NODE *tree node)
                                              if(tree node != NULL)
            Inorder(root);
                                                 Inorder(tree node->left);
            Preorder (root);
                                                 printf("Node%d",
                                                         tree node->node data);
            Postorder (root);
                                                 Inorder(tree node->right);
void Preorder(NODE *tree node)
                                         void Postorder(NODE *tree node)
   if(tree node != NULL)
                                            if(tree node != NULL)
                                                Postorder(tree node->left);
      printf("Node%d",
                                                Postorder(tree node->right);
              tree node->node data);
      Preorder(tree node->left);
                                                printf("Node%d",
      Preorder(tree node->right);
                                                         tree node->node data);
```

Priority Queue

Making a priority queue

Use a heap structure

Insert new elements into the heap

Remove the top element to get the highest priority element

Change priorities by removing the element and then re-inserting it

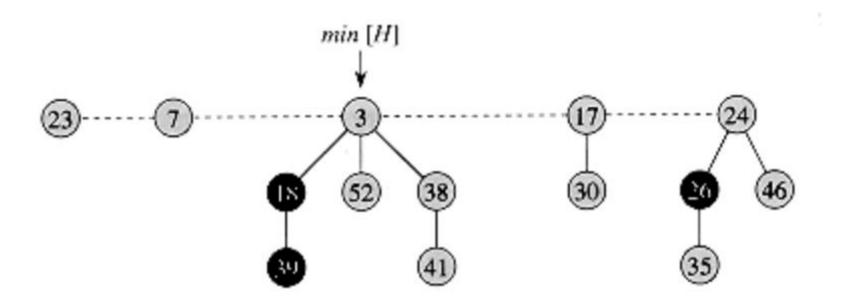
Fibonacci Heap

A fibonacci heap is a data structure that consists of a collection of trees which follow min heap or max heap property.

In a fibonacci heap, a node can have more than two children or no children at all.

Fibonacci heaps are linked lists of heap ordered trees.

A pointer to the minimum element is maintained



Fibonacci Heap

Fibonacci Heap

The root nodes of a Fibonacci heap are connected with a circular doubly linked list.

Each layer of the Fibonacci heap are connected with circular doubly linked list.

Parent and child nodes are connected with circular doubly linked list.

Fibonacci Heap

A node can be removed from a circular doubly linked list in a Fibonacci heap in O(1) time.

Two of these types of heaps can be concatenated in O(1) time

An item is added to a Fibonacci Heap by creating a new heap and connecting to the circular doubly linked list of roots.

The pointer to the minimum element is updated if needed when a new item is added.

Nodes are "marked" when they lose a child node and are moved up to the root level when they lose another child node.

Question: What is an algorithm? What is the need for an algorithm?

An algorithm is a well-defined computational procedure that takes some values or the set of values, as an input and produces a set of values or some values, as an output.

Algorithms help us measure and analyze the complexity time and space of the problems.

We can compare the performance of the algorithms with respect to other techniques.

Algorithms provide the basic idea of the problem and an approach to solve it. Some reasons to use algorithms are...

- improves the efficiency of an existing technique.
- gives a strong description of requirements and goal of the problems to the designer.
- provides a reasonable understanding of the flow of the program.
- measures the performance of the methods in different cases (Best cases, worst cases, average cases).
- identifies the resources (input/output, memory) cycles required by the algorithm.
- reduces the cost of design

Question: What is the Complexity of Algorithm?

The complexity of the algorithm is a way to classify how efficient an algorithm is compared to alternative ones. Its focus is on how execution time increases with the data set to be processed. The computational complexity of the algorithm is important in computing.

It is very suitable to classify algorithm based on the relative amount of time or relative amount of space they required and specify the growth of time/ space requirement as a function of input size.

Time complexity

Time complexity is the running time of a program as a function of the size of the input.

Space complexity

Space complexity analyzes the algorithm, based on how much space an algorithm needs to complete its task. Space complexity analysis was critical in the early days of computing (when storage space on the computer was limited) – not as important today.

Worst-case: f(n)

It is defined by the maximum number of steps taken on any instance of size n.

Best-case: f(n)

It is defined by the minimum number of steps taken on any instance of size n.

Average-case: f(n)

It is defined by the average number of steps taken on any instance of size n.

Write an algorithm to reverse a string. For example, if my string is "French" then my result will be "honerf".

Step1: Create two variables i and j

Step2: set i equal to 0 and set j equal to the length of the string - 1

Step3: swap string [i] with string[j]

Step4: increment i by 1 and decrement j by 1

Step5: repeat steps 3 and 4 while i < j

Write an algorithm to insert a node in a sorted linked list.

What are the 2 possibilities?

Linked list is empty

Linked list is not empty

Linked list is empty

Create a new node

If linked list is empty (linked list head is NULL), then set linked list head to the new node.

Linked list is not empty

Create a new node

Traverse the linked list until you find where to insert the new node.

Save the previous node's next pointer. Set previous node's pointer to the address of the new node. Set the new node's next pointer to the saved address that was saved from previous node's next pointer.

Add new node to linked list

Create a new node

Traverse the linked list to the end.

Set the next pointer of the end node to the new node's address.

What are the Asymptotic Notations?

Asymptotic analysis is used to measure the efficiency of an algorithm that doesn't depend on machine-specific constants and prevents the algorithm from comparing the time taking algorithm.

Asymptotic notation is a mathematical tool that is used to represent the time complexity of algorithms for asymptotic analysis.

What are the Asymptotic Notations?

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Asymptotic notation is a mathematical tool that is used to represent the time complexity of algorithms for asymptotic analysis.

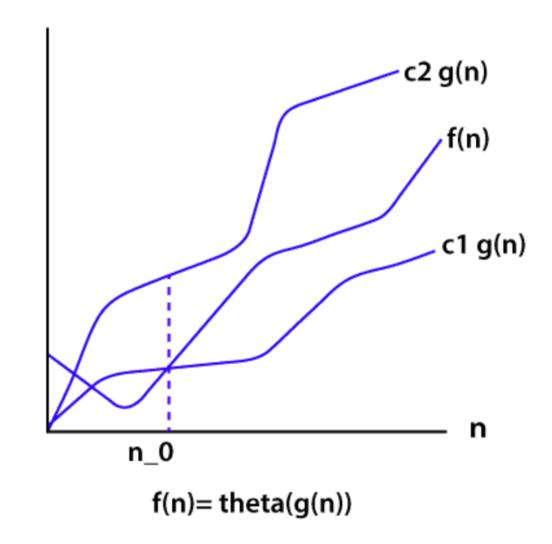
θ Notation

 θ Notation defines the exact asymptotic behavior.

To define a behavior, it bounds functions from above and below.

A convenient way to get Theta notation of an expression is to drop low order terms and ignore leading constants.

θ Notation



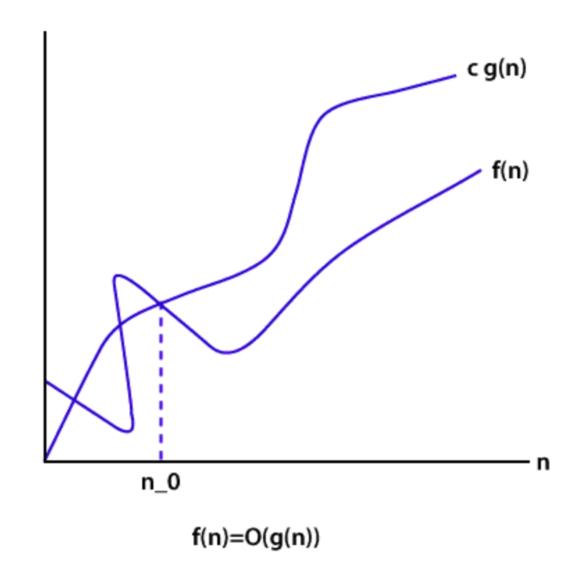
Big O Notation

The Big O notation bounds a function from above, it defines an upper bound of an algorithm.

Let's consider the case of insertion sort; it takes linear time in the best case and quadratic time in the worst case.

The time complexity of insertion sort is O(n²). It is useful when we only have upper bound on time complexity of an algorithm.

Big O Notation

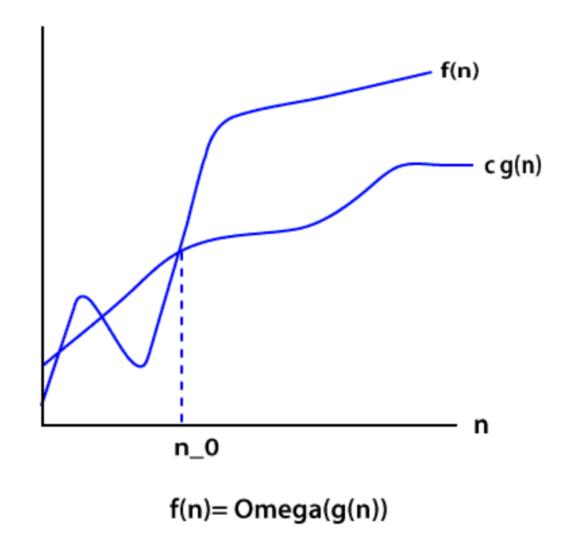


Ω Notation

Just like Big O notation provides an asymptotic upper bound, the Ω **Notation** provides an asymptotic lower bound on a function.

It is useful when we have lower bound on time complexity of an algorithm.

Ω Notation



Explain the Bubble Sort algorithm

Bubble sort is the simplest sorting algorithm among all sorting algorithm. It repeatedly works by swapping the adjacent elements if they are in the wrong order.

How would you use bubble sort to sort this array?

{7,2,5,3,8}

Pass1:

- (72538) -> (27538) swap 7 and 2.
- (2**75**38) -> (25738) swap 7 and 5.
- (25**73**8) -> (25378) swap 7 and 3.
- (25378) -> (25378) algorithm does not swap 7 and 8 because 7<8.

Pass2:

- (25378) -> (25378) algorithm does not swap 2 and 5 because 2<5.
- (2**53**78) -> (23578) swap 3 and 5.
- (23**57**8) -> (23578) algorithm does not swap 5 and 7 because 5<7.
- (23578) -> (23578) algorithm does not swap 7 and 8 because 7<8.

```
void BubbleSort(int arr[], int n)
 for (int i = 0; i < n-1; i++)
  for (int j = 0; j < n-i-1; j++)
    if (arr[j] > arr[j+1])
     swap(&arr[j], &arr[j+1]);
```

$$n = 6$$

i	j
0	0,1,2,3,4
1	0,1,2,3
2	0,1,2
3	0,1
4	0

```
void BubbleSort(int arr[], int n)
 for (int i = 0; i < n-1; i++)
   for (int j = 0; j < n-i-1; j++)
```

What is the time complexity of Bubble Sort?

How to swap two integers without swapping the temporary variable?

Suppose we have two integers i and j, the value of i=7 and j=8 then how will you swap them without using a third variable?

Looks good but...

The integer will overflow if the addition is more than the maximum value of int as defined by INT_MAX and if subtraction is less than minimum value, INT_MIN .

The integer will overflow if the addition is more than the maximum value of int as defined by INT_MAX and if subtraction is less than minimum value, INT_MIN.

```
/* Minimum and maximum values a `signed int' can hold. */
# define INT_MIN (-INT_MAX - 1)

i = i + j # define INT_MAX 2147483647

j = i - j Swap 2147483647 and 1

i = i - j

i = 2147483647 + 1 = -2147483648

j = -2147483648 - 1 = 2147483647

i = -2147483648 - 2147483649 = 1
```

This works in C but would not work in a language like Java that does not handle integer overflow.

Want to really impress?

Use the XOR method...

```
i = i ^ j;
j = i ^ j;
i = i ^ j;
```

```
Set i = 7 and j = 8
i = i ^ j
i = 7 ^ 8
i = 00000111 ^ 00001000 = 00001111
i = 15
j = i ^ j
j = 15 ^ 8
j = 00001111 ^ 00001000 = 00000111
\dot{1} = 7
i = i ^ j
i = 15 ^ 7
i = 00001111 ^ 00000111 = 00001000
i = 8
```

What are Divide and Conquer algorithms?

Divide and Conquer is not an algorithm; it's a pattern for an algorithm.

It is an algorithm that breaks up a large input into smaller pieces and solves the problem for each of the small pieces.

Then, all of the piecewise solutions are merged into a global solution.

This strategy is called divide and conquer.

Divide and conquer uses the following steps

Divide: This step divides the original problem into a set of subproblems.

Conquer: This step solves every subproblem individually.

Combine: This step puts together the solutions of the subproblems to get the solution to the whole problem.

Give some examples of Divide and Conquer algorithms

Merge Sort

Quick Sort

Binary Search

Explain the BFS algorithm?

BFS (Breadth First Search) is a graph traversal algorithm.

It starts traversing the graph from the root node and explores all the neighboring nodes.

It selects the nearest node and visits all the unexplored nodes.

The algorithm follows the same procedure for each of the closest nodes until it reaches the goal state.

What is Dijkstra's shortest path algorithm?

Dijkstra's algorithm is an algorithm for finding the shortest path from a starting node to the target node in a weighted graph. The algorithm makes a tree of shortest paths from the starting vertex and source vertex to all other nodes in the graph.

Suppose you want to go from home to school using the shortest possible way. You know some roads are heavily congested and using those routes will take more time (meaning these edges have a larger weight).

Dijkstra's algorithm can be used to find the shortest path.

What are Greedy algorithms?

A greedy algorithm is an algorithmic strategy which is made for the best optimal choice at each sub stage with the goal of this, eventually leading to a globally optimum solution. This means that the algorithm chooses the best solution at the moment without regard for consequences.

In other words, an algorithm that always takes the best immediate, or local, solution while finding an answer.

Greedy algorithms find the overall, ideal solution for some idealistic problems, but may discover less-than-ideal solutions for some instances of other problems.

List some Greedy algorithms?

Prim

Kruskal

Dijkstra

What is a linear search?

Linear search is used on a group of items. It relies on the technique of traversing a list from start to end by visiting properties of all the elements that are found on the way.

Step1: Traverse the array using **for loop**.

Step2: In every iteration, compare the target value with the current value of the array

Step3: If the values match, return the current index of the array

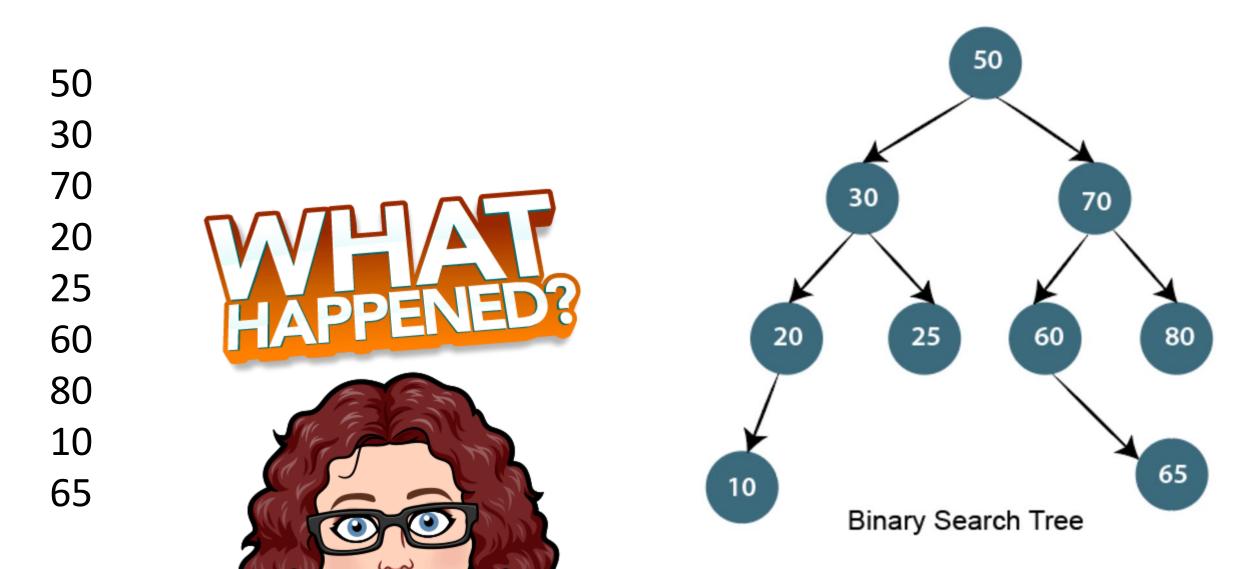
Step4: If the values do not match, shift on to the next array element.

Step5: If no match is found, return -1

What is a Binary Search Tree?

The binary search tree is a special type of data structure which has the following properties.

- Nodes which are less than root will be in the left subtree.
- Nodes which are greater than root will be right subtree.
- A binary search tree should not have duplicate nodes.
- Both sides subtree (left and right) also should be a binary search tree.



Write an algorithm to insert a node in the Binary search tree

Compare the node to be inserted with the root node and traverse left (if smaller) or right (if greater) according to the value of the node to be inserted.

Algorithm

Set root node as the current node

If the node to be inserted < root

If it has left child, then traverse left

If it does not have left child, insert node here

If the node to be inserted > root

If it has the right child, traverse right

If it does not have the right child, insert node here.

Count the leaf nodes of the binary tree

Algorithm

Traverse the tree (inorder, preorder or post order)

If a node is a leaf (both right and left pointers are NULL), then increment leaf counter

What is the difference between the Singly Linked List and Doubly Linked List?

This is a traditional interview question on the data structure. The major difference between the singly linked list and the doubly linked list is the ability to traverse.

You cannot traverse back in a singly linked list because in it a node only points towards the next node and there is no pointer to the previous node.

On the other hand, the doubly linked list allows you to navigate in both directions in any linked list because it maintains two pointers towards the next and previous node.

What is a hash algorithm and how is it used?

You will want to get comfortable answering this question because hash algorithms are popular now due to their use in cryptography.

A hash algorithm refers to a hash function, which takes a string and converts it to a fixed length regardless of how long it was to begin with.

You can use it for a wide range of applications, from cryptocurrency to passwords and a range of other validation tools.

What are the three laws that govern a recursive algorithm?

These kinds of algorithm interview questions may come as follow-ups for the "What is a recursive algorithm?" question.

A recursive algorithm needs to follow these laws:

It has to have a base case.

It has to call itself.

It needs to change its state and shift towards the base case.

What Are the Different Types of Data Structures?

A collection of data values stored sequentially Last-in-first-out (LIFO) data structures where the element placed last is accessed first.	Arrays Stacks
A first-in-first-out data structure.	Queues
A collection of data values stored in a linear order and connected to each other	Linked lists
A data structure in which data values are placed in nodes connected by edges	Graphs
Similar to a linked list, but with data values linked in a hierarchical fashion	Trees
A binary tree data structure wherein parent data values can be compared child data values	to Heaps
A table where each value is assigned a key and then stored, making access individual values easy.	sing Hash table
<i>1</i>	