

Network Information Security

Lecture six: mode of operation

Rachel yuan
2019.10

1

Try to compare DES and AES

2

| Factors | AES | DES |
|-------------------------|-----|-----|
| Cipher type | | |
| Key length | | |
| Block size | | |
| Possible keys | | |
| Security | | |
| S-Box | | |
| Template | | |
| Reversibility | | |
| Number of Rounds | | |
| Implementation | | |
| Developed | | |
| Confusion in each round | | |
| Diffusion in each round | | |

3

How to encrypt long messages using block cipher?

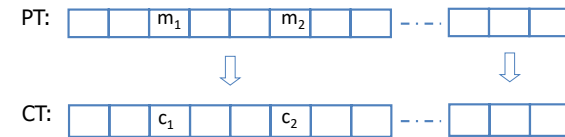
4

A **mode of operation** describes how to repeatedly apply a cipher's single-block operation to securely transform amounts of data larger than a block

5

A straightforward method

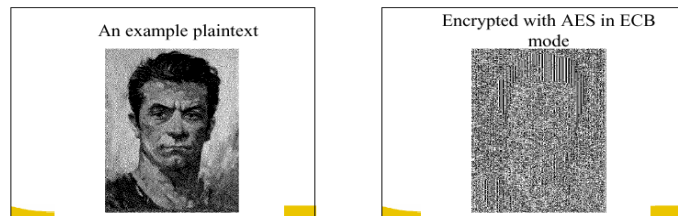
- Electronic Code Book (ECB):



- Problem:
 - if $m_1 = m_2$ then $c_1 = c_2$

6

In pictures

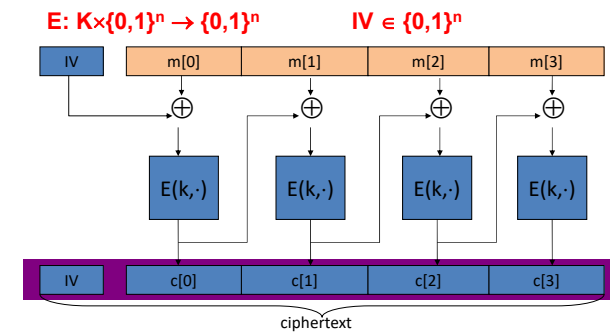


(courtesy B. Preneel)

7

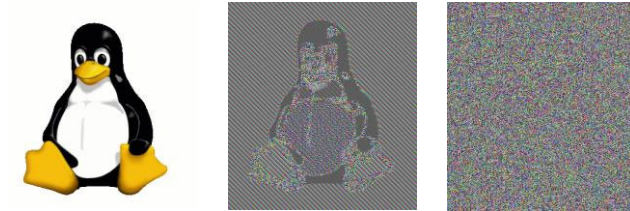
Construction 1: CBC with random IV

Let (E, D) be a block cipher. $E_{CBC}(k, m)$: choose **random** $IV \in X$ and do:



CBC: cipher block chaining

8



CBC properties

- Parallel processing possible?
- Do **errors** propagate?

9

10

$$c_i = E(m_i \oplus c_{i-1}, k), \text{ for } i = 1, 2, \dots$$

$$c_0 = E(m_0 \oplus IV, k)$$

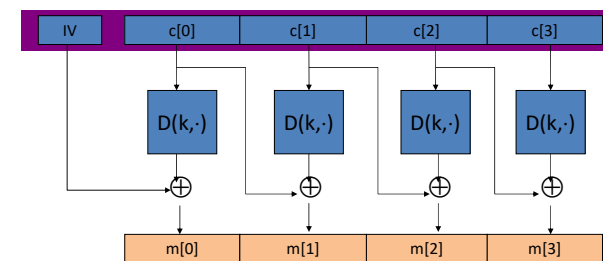
$$m_i = D(c_i, k) \oplus c_{i-1}, \text{ for } i = 1, 2, \dots$$

$$m_0 = D(c_0, k) \oplus IV$$

11

Decryption circuit

In symbols: $c[0] = E(k, IV \oplus m[0]) \Rightarrow m[0] = D(k, c[0]) \oplus IV$

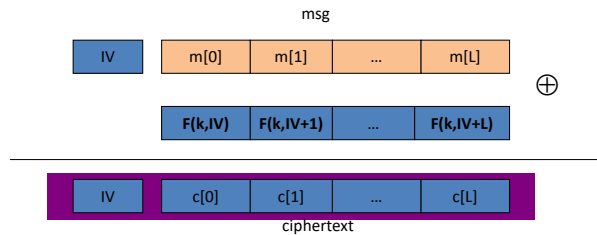


12

Construction 2: rand ctr-mode

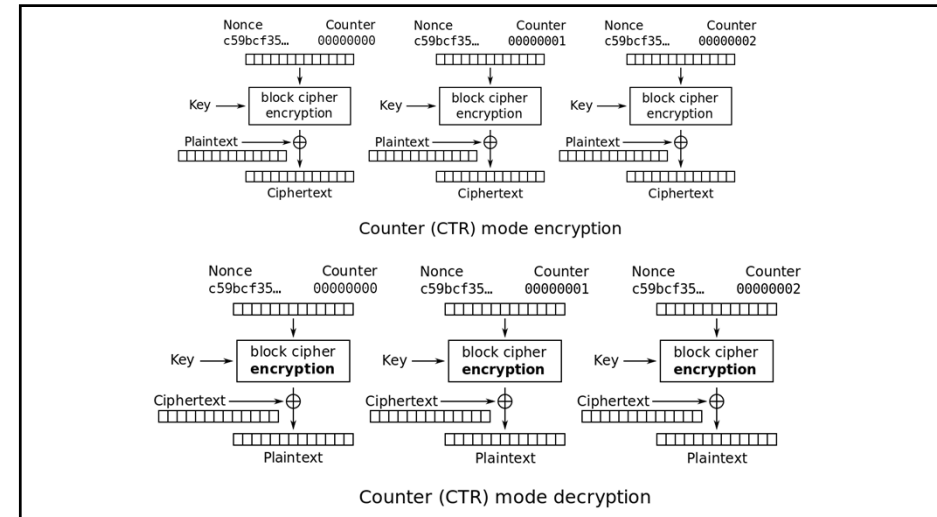
Let $F: K \times \{0,1\}^n \rightarrow \{0,1\}^n$ be a secure block cipher.

$E(k,m)$: choose a random $IV \in \{0,1\}^n$ and do:



note: parallelizable (unlike CBC)

13



14

$$c_i = m_i \oplus E(IV + i, k), \text{ for } i = 0, 1, 2, \dots$$

$$m_i = c_i \oplus E(IV + i, k), \text{ for } i = 0, 1, 2, \dots$$

15

Along with CBC, CTR mode is one of two block cipher modes recommended by Niels Ferguson and Bruce Schneier.

16