

1 $F = \{ \pm m / \beta^t \cdot \beta^e \mid \beta^{t-1} \leq m \leq \beta^t, m \in \mathbb{Z}, e \in \mathbb{Z} \}$.

As written in the textbook, we can restrict the range of m in order to make 1-1 correspondence between floating pt & IEEE representation.

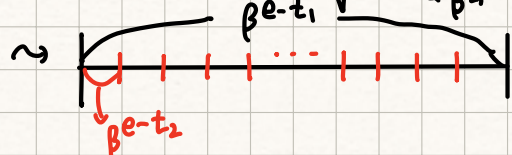
So, the gap size between each floating pts is $1/\beta^t \cdot \beta^e = \beta^{e-t}$.

Let F_1, F_2 be F with single precision t_1 & double precision t_2 , respectively.

Single: $t_1 = 24$ Double: $t_2 = 53$

Let any IEEE single precision $\# \frac{m}{\beta^{t_1}} \times \beta^e$ be given. ($\beta^{t_1-1} \leq m \leq \beta^{t_1}$) (WLOG, let's consider positive $\#$)

Now we are considering the $[\frac{m}{\beta^{t_1}} \times \beta^e, \frac{m+1}{\beta^{t_1}} \times \beta^e]$. Note that $\frac{m}{\beta^{t_1}} \times \beta^e = \frac{m \beta^{t_2-t_1}}{\beta^{t_2}} \times \beta^e \in F_2$ as $\beta^{t_2-1} \leq m \beta^{t_2-t_1} \leq \beta^{t_2}$



So, $\#$ of double precision numbers including the endpts is $\beta^{e-t_1} / \beta^{e-t_2} + 1 = \beta^{t_2-t_1} + 1$.

\therefore Between a pair of adjacent single precision numbers (so, excluding endpts), there are $\beta^{t_2-t_1} - 1$ double precision $\#$ s.

2 (a) For $n = 1, 2, \dots, \beta^t$, it is representable as $0.d_1 \dots d_t \times \beta^e$ is our elements of F .

So, we can not represent the binary number more than t digits: $\underbrace{10 \dots 0}_t 1 = \beta^t + 1$.

Note that β^t is representable.

(b) substitute t as 24 & 53

(c) construct $2^t-4, 2^t-3, 2^t-2, 2^t-1, 2^t, 2^t+1, 2^t+2, 2^t+3, 2^t+4$

$$\begin{aligned} \underbrace{10 \dots 0}_t 1 &\sim \underbrace{10 \dots 0}_t 00 \\ \underbrace{10 \dots 1}_t 1 &\sim \underbrace{10 \dots 1}_t 0 \end{aligned}$$

Note that 2^t+2 representable, so gap between n must be 2 as $fl(2^t+1) = 2^t$.
& calculate the gap. If num representable, gap must be 1.