树状数组

一组数据,要求能够区间修改和区间询问.同时要完成两种操作,只要维护两个树状数组就好了.一个差分数组,另一个是对差分数组进行求和的数组.

$$Sum(n) = \sum_{i=1}^{n} a_i = a_n + a_{n-1} + \ldots + a_1$$

$$= (d_n + \ldots + d_1) + (d_{n-1} + \ldots + d_1) + \ldots + d_1$$

$$= (n - (n-1)) * d_n + (n - (n-2)) * d_{n-1} + \ldots + (n-0) * d_1$$

$$= n * \sum_{i=1}^{n} d_i - \sum_{i=1}^{n} d_i * (i-1)$$

```
ll lowbit(ll x) { return x&(-x);}
void update(ll i,ll val)
    11 x=i;
    while(i<=N){</pre>
        a[i]+=val;
        b[i] += (val*(x-1));
        i+=lowbit(i);
    }
}
11 getsum(11 i)
    11 res=0,x=i;
    while(i>0){
        res+=(a[i]*x-b[i]);
        i-=lowbit(i);
    return res;
}
int main()
    11 Q, num1, num2=0;
    scanf ("%11d %11d",&N,&Q);
    for (int i=1; i <= N; i++){
        scanf ("%11d",&num1);
        update(i,num1-num2);
        num2=num1;
    }e
    for (int i=1; i \leftarrow Q; i++)
    {
        char op=getchar();
        11 x,y,k;
        if (op=='C') {
             scanf ("%11d %11d %11d",&x,&y,&k);
             update(x,k);
             update(y+1,-k);
        else if (op=='Q') {
```

```
scanf("%11d %11d",&x,&y);
    printf("%11d\n",getsum(y)-getsum(x-1));
}
    else i--;
}
return 0;
}
```

ST表

```
int Min(int a,int b)
{
    return s[a] \le s[b]?a:b;
}
void getST()
    len=strlen(s);
    memset(ST,0,sizeof(ST));
    for (int j=0; j<=11; j++){
        for (int i=0;i+bin[j]-1<len;i++){
            if (j==0) ST[i][j]=i;
            else ST[i][j]=Min(ST[i][j-1],ST[i+bin[j-1]][j-1]);
        }
    }
}
int STsearch(int x,int y)
{
    if (x==y) return x;
    int t=Log[y-x+1];
    return Min(ST[x][t],ST[y-bin[t]+1][t]);
}
```

线段树标记

```
void build(int o,int 1,int r)
{
    lazy1[o]=1;
   if (l==r) { sum[o]=num[l]%p; return ;}
   build(lo,1,M);
   build(ro,M+1,r);
    sum[o]=(sum[lo]+sum[ro])%p;
}
void pushdown(int o,int 1,int r)
                           lazy1[lo]%=p;
    lazy1[lo]*=lazy1[o];
    lazy1[ro]*=lazy1[o];
                          lazy1[ro]%=p;
   lazy2[lo]=lazy2[lo]*lazy1[o]+lazy2[o]; lazy2[lo]%=p;
   lazy2[ro]=lazy2[ro]*lazy1[o]+lazy2[o]; lazy2[ro]%=p;
    sum[lo]=((sum[lo]*lazy1[o])+lazy2[o]*len(l,M))%p;
    sum[ro]=((sum[ro]*lazy1[o])+lazy2[o]*len(M+1,r))%p;
    lazy1[o]=1;
    lazy2[o]=0;
```

```
11 getsum(int o,int 1,int r)
    11 res=0;
    if (x \le 1 \& r \le y) return sum[o]\%p;
    pushdown(o,1,r);
    if (M >= X) {
        res+=getsum(lo,1,M);
        res%=p;
    }
    if ( M < y) {
        res = getsum(ro, M+1, r);
        res%=p;
    }
    return res%p;
}
void addval(int o,int l,int r)
    if (x<=1 && r<=y) {
        sum[o]=(sum[o]+k*len(l,r))%p;
        lazy2[o]+=k; lazy2[o]%=p;
        return ;
    }
    pushdown(o,1,r);
    if (M \ge X) addval(lo, l, M);
    if ( M < y ) addval(ro, M+1, r);
    sum[o]=sum[lo]+sum[ro];
    sum[o]%=p;
}
void multival(int o,int l,int r)
{
    if (x<=1 \&\& r<=y) {
        sum[o]=(sum[o]*k)%p;
        lazy1[o]*=k; lazy1[o]%=p;
        lazy2[o]*=k; lazy2[o]%=p;
        return ;
    }
    pushdown(o,1,r);
    if (M \ge x) multival(lo, l, M);
    if (M < y) multival(ro, M+1, r);
    sum[o]=sum[lo]+sum[ro];
    sum[o]%=p;
}
```

扫描线

```
else
         tree[o].sum1 = tree[lo].sum1 + tree[ro].sum1;
void pushdown(ll o, ll L, ll R)
    if (tree[o].lazy2) {
        11 mid = (L + R) >> 1;
         do_odd(lo, L, mid);
         do_odd(ro, mid + 1, R);
         tree[lo].lazy2 \wedge= 1;
         tree[ro].lazy2 \wedge= 1;
         tree[o].lazy2 = 0;
    }
}
void addval(ll o, ll L, ll R)
    if (num[R + 1] \leftarrow q] \mid | num[L] \rightarrow qr)
        return ;
    if (q) \leftarrow num[L] \&\& num[R + 1] \leftarrow qr) {
         tree[o].lazy += k;
         do_odd(o, L, R);
         tree[o].lazy2 \wedge= 1;
         pushup(o, L, R);
         return ;
    pushdown(o, L, R);
    11 M = (L + R) / 2;
    addval(lo, L, M);
    addval(ro, M + 1, R);
    pushup(o, L, R);
    tree[o].sum2 = tree[lo].sum2 + tree[ro].sum2;
}
```

树链剖分

```
void dfs1(int x, int f, int d)
{
    fa[x] = f;
    siz[x] = 1;
    dep[x] = d;
    int maxson = -1;
    for (int i = first[x]; i; i = inext[i]) {
        if (to[i] == f)
            continue;
        dfs1(to[i], x, d + 1);
        siz[x] += siz[to[i]];
        if (siz[to[i]] > maxson)
            maxson = siz[to[i]], son[x] = to[i];
    }
void dfs2(int x, int topf)
{
    id[x] = ++cur;
    top[x] = topf;
    w2[cur] = w1[x];
    if (!son[x])
        return ;
```

```
dfs2(son[x], topf);
    for (int i = first[x]; i; i = inext[i]) {
        if (to[i] != fa[x] && to[i] != son[x])
            dfs2(to[i], to[i]);
    }
}
int lss(int beg, int end)
    int ans = 0;
    while (top[beg] != top[end]) {
        if (dep[top[beg]] < dep[top[end]])</pre>
            swap(beg, end);
        ans +=sum(1, 1, n, id[top[beg]], id[beg]);
        beg = fa[top[beg]];
    if (dep[beg] > dep[end])
        swap(beg, end);
    ans += sum(1, 1, n, id[beg], id[end]);
    return ans;
}
int imax(int beg, int end)
    int ans = -inf;
    while (top[beg] != top[end]) {
        if (dep[top[beg]] < dep[top[end]])</pre>
            swap(beg, end);
        ans = max(ans, query(1, 1, n, id[top[beg]], id[beg]));
        beg = fa[top[beg]];
    }
    if (dep[beg] > dep[end])
        swap(beg, end);
    ans = max(ans, query(1, 1, n, id[beg], id[end]));
    return ans ;
}
```

可持久化并查集

```
inline int find(int root, int x)
{
    node now = Tree[root];
    if (now.1 == now.r)
        return now.fa;
    int mid = (now.1 + now.r) >> 1;
    if (x \ll mid)
        return find(now.ls,x);
    else
        return find(now.rs,x);
}
inline int getfa(int root, int x)
    int fa = find(root, x);
    if (fa == x)
        return x;
    else
        return getfa(root, fa);
}
```

```
inline int update(const int pre, int loc, int val)
{
   int now = tot++;
   copy(now, pre);
   if (Tree[now].l == Tree[now].r ) {
        Tree[now].fa = val;
        return now;
   }
   int mid = (Tree[now].l + Tree[now].r) >> 1;
   if (loc <= mid)
        Tree[now].ls = update(Tree[now].ls,loc,val);
   else
        Tree[now].rs = update(Tree[now].rs,loc,val);
   return now;
}</pre>
```

SBT

```
void rotate(int &rt, int k)
    int nrt = son[rt][!k]; //new root
    son[rt][!k] = son[nrt][k];
    son[nrt][k] = rt;
    siz[nrt] = siz[rt];
    siz[rt] = siz[son[rt][0]] + siz[son[rt][1]] + 1;
    rt = nrt;
}
void maintain(int &rt, bool k)
    if (siz[son[son[rt][k]][k]] > siz[son[rt][!k]])
        rotate(son[rt][k], !k);
    else if (siz[son[son[rt][k]][!k]] > siz[son[rt][!k]])
        rotate(son[rt][k], k), rotate(rt, !k);
    else
        return ;
    maintain(son[rt][0], 0);
    maintain(son[rt][1], 1);
    maintain(rt, 0);
   maintain(rt, 1);
}
void insert(int &rt, int val)
{
    if (rt == 0) {
        rt = ++tails;
        datas[rt] = val;
        son[rt][0] = son[rt][1] = 0;
        siz[rt] = 1;
        return ;
    ++siz[rt];
    if (val <= datas[rt])</pre>
        insert(son[rt][0], val);
    else
```

```
insert(son[rt][1], val);
    maintain(rt, datas[rt] > val);
}
void del(int &rt, int val)
    if (datas[rt] != val) {
        del(son[rt][datas[rt] < val], val);</pre>
        siz[rt] = siz[son[rt][0]] + siz[son[rt][1]] + 1;
        return ;
    }
    --siz[rt];
    if (son[rt][0] == 0)
        rt = son[rt][1];
    else if (son[rt][1] == 0)
        rt = son[rt][0];
    else {
        int p = son[rt][1];
        while (son[p][0] != 0)
            p = son[p][0];
        datas[rt] = datas[p];
        del(son[rt][1], datas[p]);
    }
}
int pre (int val)
    int cur = sbt, ans = 0;
    while (cur != 0) {
        if (datas[cur] < val)</pre>
            ans = cur, cur = son[cur][1];
        else
            cur = son[cur][0];
    }
    return datas[ans];
}
int nxt (int val)
    int cur = sbt, ans = 0;
    while (cur != 0) {
        if (datas[cur] > val)
            ans = cur, cur = son[cur][0];
        else
            cur = son[cur][1];
    return datas[ans];
}
int xth(int x)
    int cur = sbt;
    while (siz[son[cur][0]] + 1 != x) {
        if (siz[son[cur][0]] + 1 < x)
            x = (siz[son[cur][0]] + 1), cur = son[cur][1];
        else
            cur = son[cur][0];
    }
```

```
return datas[cur];
}

int ranks(int val)
{
    int cur = sbt, ans = 0;
    while (cur != 0) {
        if (datas[cur] < val)
            ans += (siz[son[cur][0]] + 1), cur = son[cur][1];
        else
            cur = son[cur][0];
    }
    return ans + 1;
}</pre>
```

莫队 回滚莫队

```
for (int k = 0; k \leftarrow blockn; ++k)
{
    int l = ranger[k] + 1, r = ranger[k];
    memset(cnt, 0, sizeof(cnt));
    11 temp = 0, now = 0;
    for (; block[q[i].1] == k; ++i) {
        int q1 = q[i].1, qr = q[i].r;
        if (block[q[i].1] == block[q[i].r]) {
            temp = 0;
            for (int j = ql; j \leftarrow qr; ++j)
                cnt2[id[j]] = 0;
            for (int j = q1; j \leftarrow qr; ++j) {
                if (!cnt2[id[j]])
                     cnt2[id[j]] = j;
                else
                    temp = max(temp, 1]] * abs(j - cnt2[id[j]]));
            ans[q[i].id] = temp;
            continue;
        }
        while (r < qr) {
            if (!cnt[id[++r]][0])
                cnt[id[r]][1] = cnt[id[r]][0] = r;
            else {
                cnt[id[r]][1] = r;
                now = max(now, 111 * abs(r - cnt[id[r]][0]));
            }
        }
        temp = now;
        while (q1 < 1) {
            if (!cnt[id[--1]][1])
                cnt[id[1]][1] = 1;
            else
                temp = max(temp, 111 * abs(1 - cnt[id[1]][1]));
        }
        ans[q[i].id] = temp;
        while (1 < ranger[k] + 1) {
            if (cnt[id[1]][1] < ranger[k] + 1)
                cnt[id[1]][1] = 0;
            ++1;
```

```
}
}
}
```

字符串

KMP

```
void getnextarr(char* str)
{
    int n=strlen(str);
    nexts[0]=-1, nexts[1]=0;
    for (int i=2;i<=n;i++)
    {
        int k=i-1;
        while (1){
            if (nexts[k]==-1 || str[i-1]==str[nexts[k]])
                nexts[i]=nexts[k]+1;
                break;
            }
            else k=nexts[k];
        }
    }
    return ;
}
void kmp()
    int len1=strlen(str1);
    int len2=strlen(str2);
    int i=0, j=0;
    getnextarr(str2);
    while (i<len1)
    {
        if (str1[i]==str2[j]){
                if (j!=len2-1) i++,j++;
                ans [cnt++]=i-j+1;
                j=nexts[j];
            }
        }
        else if (nexts[j]==-1) i++;
        else j=nexts[j];
    }
    for (int t=0;t<cnt;t++) printf("%d\n",ans[t]);</pre>
    for (int t=0;t<len2;t++) {}
         printf("%d ",nexts[t+1]);
    return ;
}
```

AC自动机

```
struct node{
    int f;
    int book;
    int next[26];
}tree[1000051];
struct imark{
    int num,str;
}ans[1000051];
bool cmp(imark a,imark b){
    if (a.num==b.num) return a.str<b.str;</pre>
    else return a.num>b.num;
}
char a[155][1000];
void clean(int x){
    memset(tree[x].next,0,sizeof(tree[x].next));
    tree[x].f=0;
    tree[x].book=0;
}
char s[1000506];
int ta=0;
void build(int x){
    int cur=0,l=strlen(a[x]);
    for (int i=0; i<1; ++i){
        if (!tree[cur].next[a[x][i]-'a']) {
            tree[cur].next[a[x][i]-'a']=++ta;
            clean(ta);
        }
        cur=tree[cur].next[a[x][i]-'a'];
    tree[cur].book=x;
}
```

马拉车

```
string Mannacher(string s)
   //插入"#"
   string t="$#";
   for(int i=0;i<s.size();++i)</pre>
       t+=s[i];
       t+="#";
   }
   vector<int> p(t.size(),0);
   //mx表示某个回文串延伸在最右端半径的下标,id表示这个回文子串最中间位置下标
   //resLen表示对应在s中的最大子回文串的半径, resCenter表示最大子回文串的中间位置
   int mx=0,id=0,resLen=0,resCenter=0;
    //建立p数组
   for(int i=1;i<t.size();++i)</pre>
       p[i]=mx>i?min(p[2*id-i],mx-i):1;
       //遇到三种特殊的情况,需要利用中心扩展法
       while(t[i+p[i]]==t[i-p[i]])++p[i];
```

```
//半径下标i+p[i]超过边界mx,需要更新
if(mx<i+p[i]){
    mx=i+p[i];
    id=i;
}
//更新最大回文子串的信息,半径及中间位置
if(resLen<p[i]){
    resLen=p[i];
    resCenter=i;
}
}
//最长回文子串长度为半径-1,起始位置为中间位置减去半径再除以2
return s.substr((resCenter-resLen)/2,resLen-1);
}
```