

1. Draw PDAs for the following languages:

(a)  $\{w \mid w \text{ has balanced brackets, } w \in \{(\,,\,)\}^*\}$  [5]

(b)  $\{w \mid w = w^R, |w| \geq 1, w \in \{0, 1\}^*\}$  [5]

2. Design the production rules for a context-free grammar that generates the following languages over  $\{0, 1\}$

(a)  $\{w \mid w = w^R, |w| \geq 1, w \in \{0, 1\}^*\}$  [5]

(b)  $\{0^n 1^m \mid n, m \geq 1, n \leq m\}$  [5]

3. Show that  $\{w \mid w \in \{0,1\}^*, w = w^R\}$  is not a regular language using the Pumping Lemma. [5]

4. Consider the following production rules for a context-free grammar:

$$\begin{aligned} E &\rightarrow E \wedge T \mid E \vee T \mid T \\ T &\rightarrow T \implies F \mid T \iff F \mid F \\ F &\rightarrow (E) \mid \neg F \mid \forall F \mid \exists F \mid V \\ V &\rightarrow p \mid q \mid r \mid 1 \mid 0 \end{aligned}$$

- (a) What are the non-terminals of the above grammar? [2]
- (b) What are the terminals of the above grammar? [3]
- (c) Using  $E$  as the start symbol, show the parse tree of  $\forall(0 \vee 1)$ . [3]
- (d) Using  $E$  as the start symbol, show the left-most derivation of  $\exists p$ . [2]