# THEORETICAL COMPUTER SCIENCE TUTORING (5)

Maurizio Fiusco

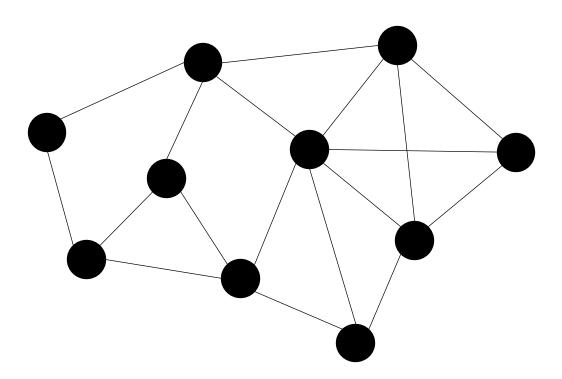


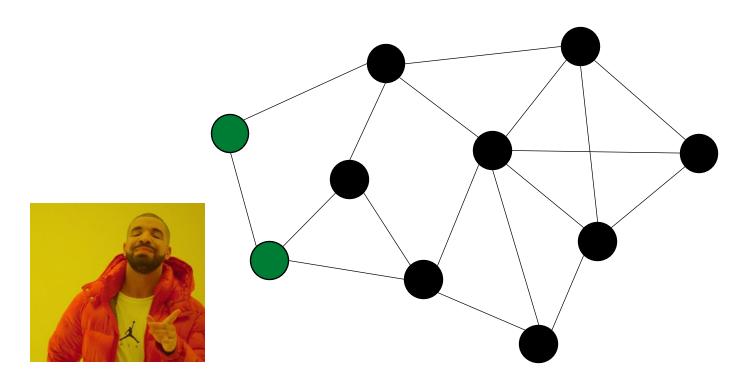
Given an undirected graph G = (V, E) and an integer h, we say that a subset V' of the nodes of G is  $\alpha$ -colorable with h colors if there exists a function  $c: V' \to \{1, ..., h\}$  such that, for every  $u, v \in V'$ , if c(u) = c(v) then  $(u, v) \in E$ .

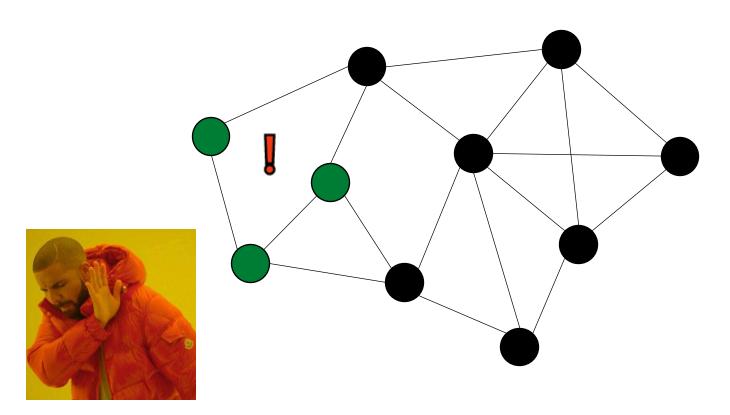
Consider the following decision problem: given an undirected graph G = (V, E) and an integer k, decide whether there does not exist a subset  $V' \subseteq V$  of cardinality at least k that is  $\alpha$ -colorable with only one color.

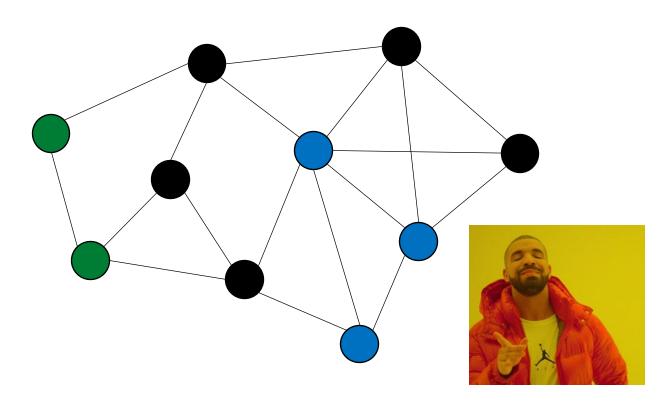
After formalizing the problem using the triple  $\langle I, S, \pi \rangle$ , answer the following questions (in the order deemed appropriate), providing justification for each response.

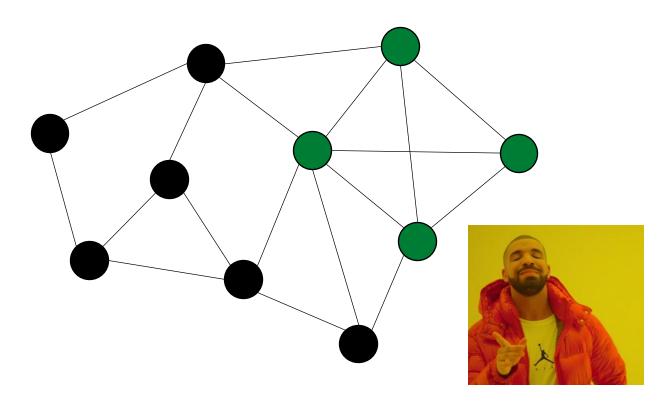
- Is the problem in P?
- Is the problem in NP?
- Is the problem in coNP?

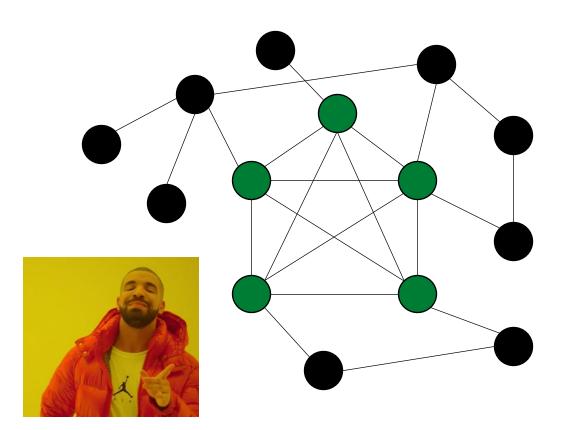








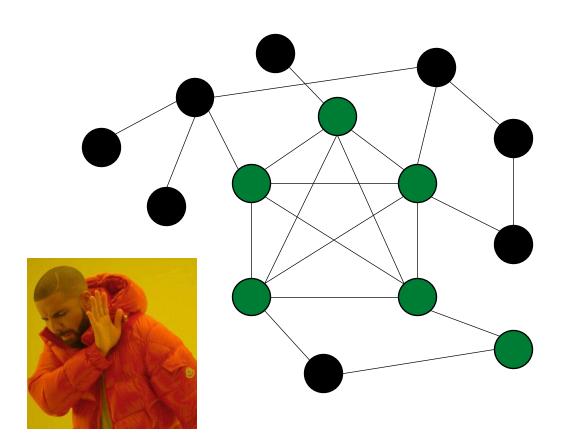


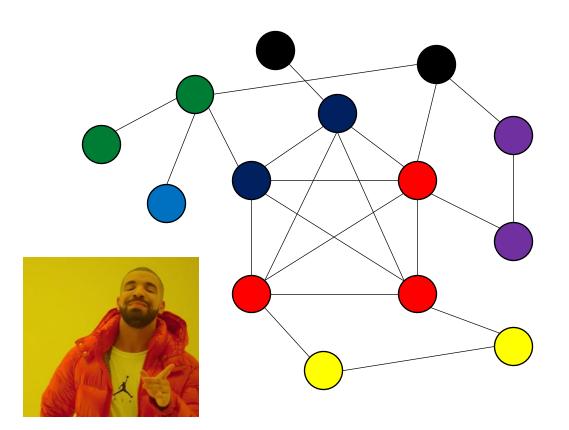


Given an undirected graph G = (V, E) and an integer h, we say that a subset V' of the nodes of G is  $\alpha$ -colorable with h colors if there exists a function  $c: V' \to \{1, ..., h\}$  such that, for every  $u, v \in V'$ , if c(u) = c(v) then  $(u, v) \in E$ .



I can only color the cliques



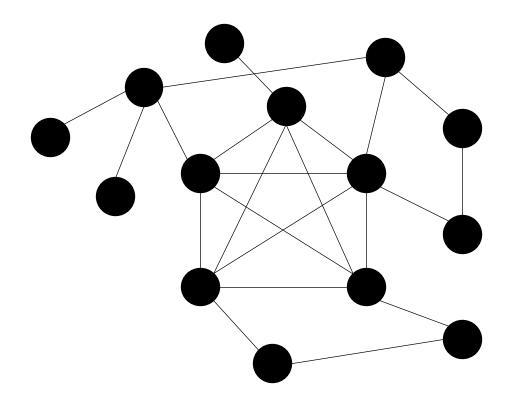


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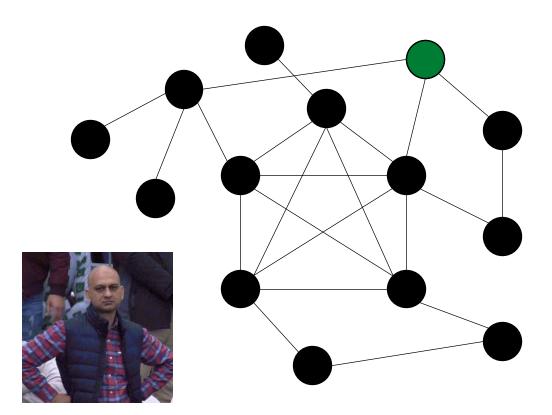
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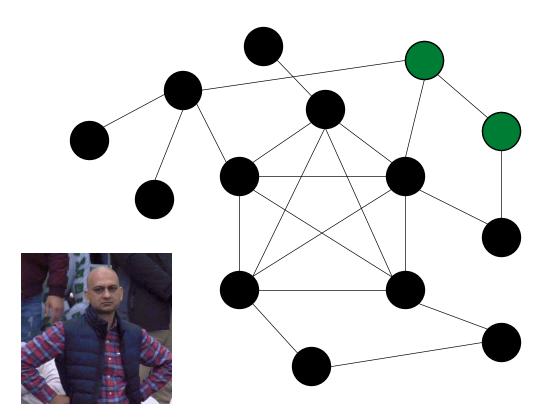
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- Is the problem in NP?
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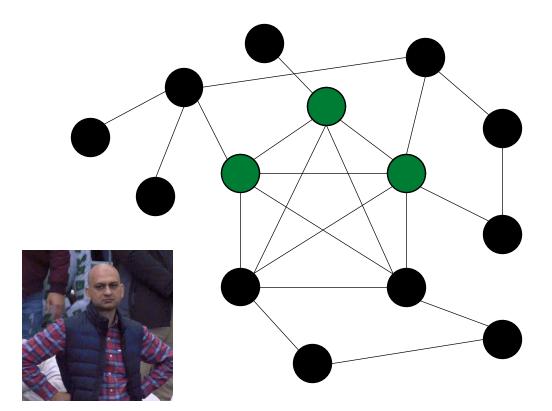
$$k = 1$$



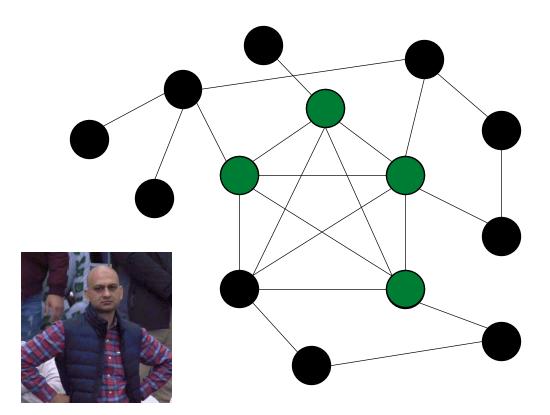
$$k = 2$$



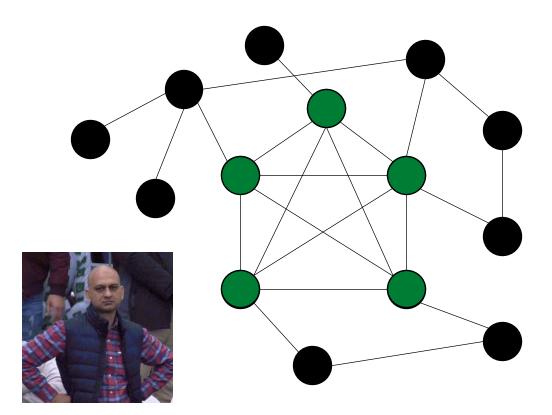
$$k = 3$$



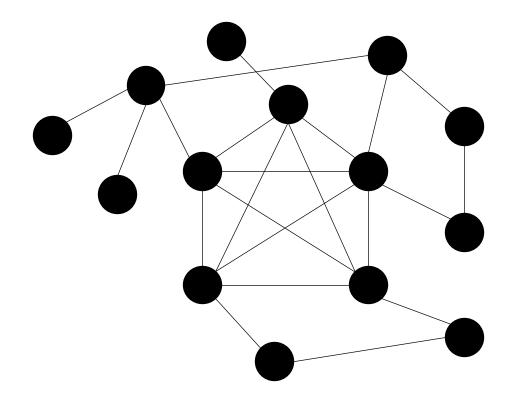
$$k = 4$$



$$k = 5$$



$$k \ge 6$$
 ?



- $I_{exam} = \{ \langle G = (V, E), k \rangle : G \text{ is an undirected graph } \land k \in \mathbb{N} \}$
- $S_{exam}(G,k) = \{V' \subseteq V\}$
- $\pi_{exam}(G, k, S_{exam}(G, k)) = \forall V' \in S_{exam}(G, k), c: V' \rightarrow \{1\}[|V'| < k \lor \exists u, v \in V' [c(u) = c(v) \land (u, v) \notin E]]$

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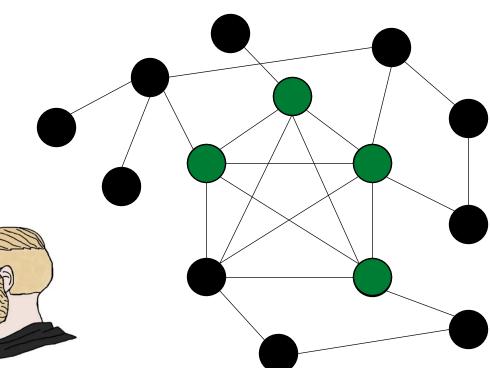
- $I_{\neg exam} = \{ \langle G = (V, E), k \rangle : G \text{ is an undirected graph } \land k \in \mathbb{N} \}$
- $S_{\neg exam}(G, k) = \{V' \subseteq V\}$
- $\pi_{\neg exam}(G, k, S_{\neg exam}(G, k)) = \exists V' \in S_{\neg exam}(G, k) [|V'| \ge k \land \forall u, v \in V' [(u, v) \in E]]$



- $I_{clique} = \{ \langle G = (V, E), k \rangle : G \text{ is an undirected graph } \land k \in \mathbb{N} \}$
- $S_{clique}(G, k) = \{V' \subseteq V\}$
- $\pi_{clique}\left(G, k, S_{clique}(G, k)\right) = \exists V' \in S_{clique}(G, k) \left[|V'| \ge k \land \forall u, v \in V' \left[(u, v) \in E\right]\right]$

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$$k = 4$$



Clique is NP-complete  $\Rightarrow \neg$ exam is NP-complete  $\Rightarrow$  exam is coNP-complete



To be more formal, you should formally specify the reduction between **Clique** and ¬exam, but the problems are practically identical

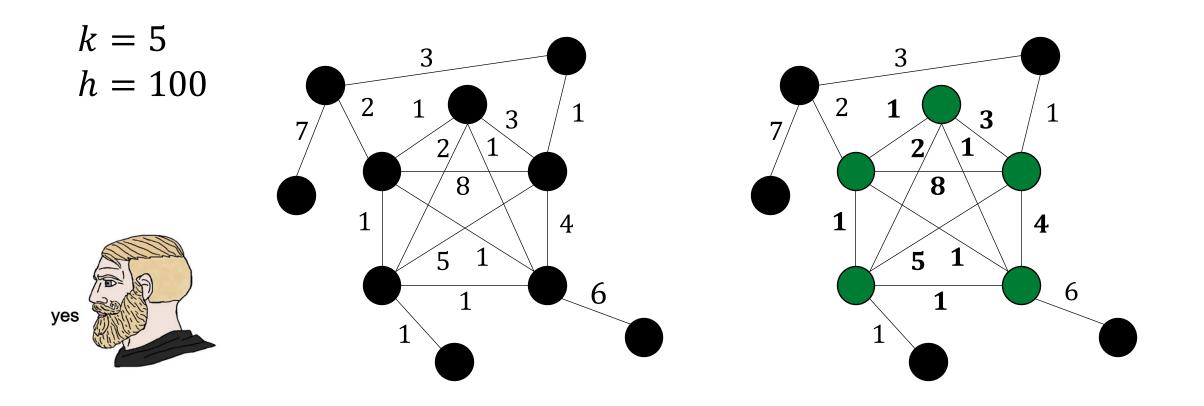
Consider the following decision problem: given an undirected and weighted graph G = (V, E, w) with  $w: E \to \mathbb{N}$  and two integers h and k, decide whether G has a clique of k nodes such that the sum of the weights of its edges is at most h.

After formalizing the problem using the triple  $\langle I, S, \pi \rangle$ , answer the following questions (in the order deemed appropriate), providing justification for each response.

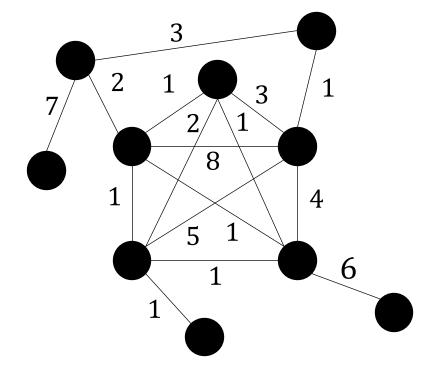
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**Note well**: Deciding whether a graph contains a clique of exactly k nodes is an NP-complete problem.

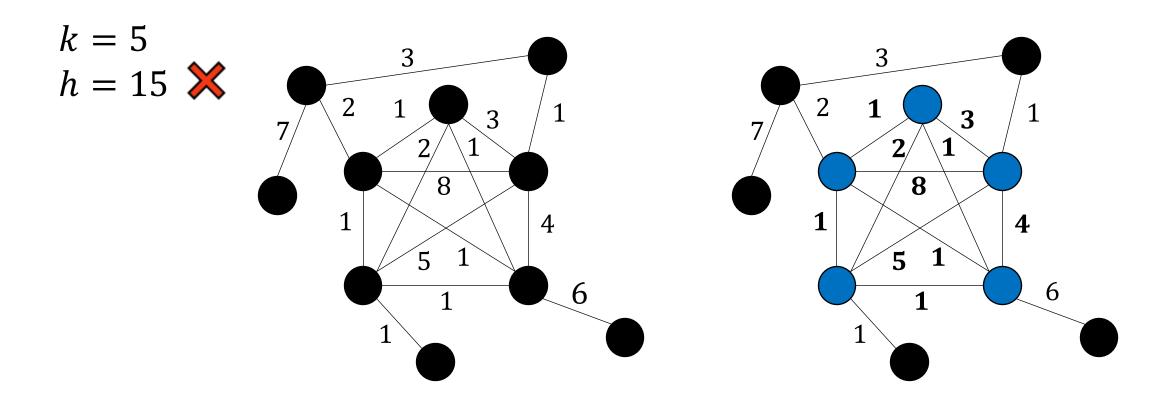
**Suggestion**: Remember the reduction that demonstrates that TSP is NP-complete.



$$k = 6$$
$$h = 100$$







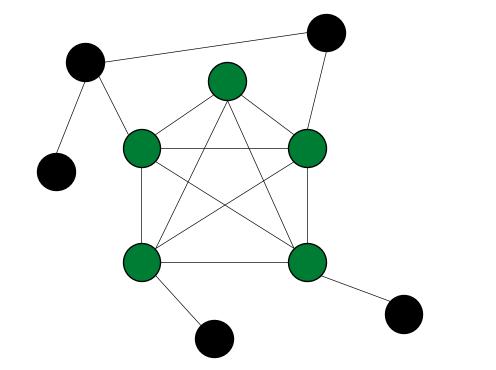
- $I_{ex} = \{ \langle G = (V, E, w), k, h \rangle : G \text{ is an undirected and weighted graph } \Lambda w : E \rightarrow \mathbb{N} \land k, h \in \mathbb{N} \}$
- $S_{ex}(G,k,h) = \{V' \subseteq V\}$
- $\pi_{ex}(G, k, S_{ex}(G, k)) = \exists V' \in S_{ex}(G, k) [|V'| = k \land \forall u, v \in V' [(u, v) \in E] \land \sum_{u,v \in V'} w(\{u,v\}) \leq h]$

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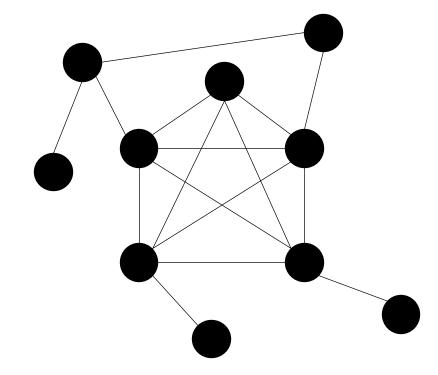
- $I_{nw} = \{ \langle G = (V, E, w), k \rangle : G \text{ is an undirected graph } \land w : E \rightarrow \mathbb{N} \land k \in \mathbb{N} \}$
- $S_{nw}(G,k) = \{V' \subseteq V\}$
- $\pi_{nw}(G, k, S_{nw}(G, k)) = \exists V' \in S_{nw}(G, k) [|V'| = k \land \forall u, v \in V' [(u, v) \in E]]$  $\succeq \text{ (it's Clique)}$

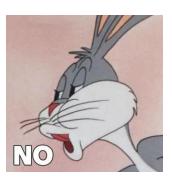
$$k = 5$$



- $I_{nw} = \{ \langle G = (V, E, w), k \rangle : G \text{ is an undirected graph } \land w : E \rightarrow \mathbb{N} \land k \in \mathbb{N} \}$
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$$k = 6$$

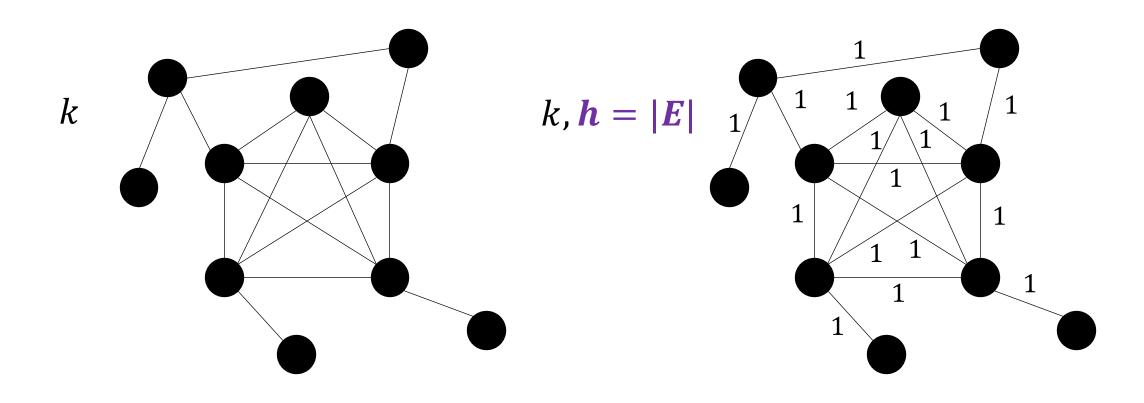




- $I_{ex} = \{ \langle G = (V, E, w), k, h \rangle : G \text{ is an undirected and weighted graph } \Lambda w : E \rightarrow \mathbb{N} \land k, h \in \mathbb{N} \}$
- $S_{ex}(G,k,h) = \{V' \subseteq V\}$
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#### Let's reduce from Clique

- $I_{Clique} = \{ \langle G = (V, E, w), k \rangle : G \text{ is an undirected graph } \land w : E \rightarrow \mathbb{N} \land k \in \mathbb{N} \}$
- $S_{Clique}(G, k) = \{V' \subseteq V\}$
- $\pi_{Clique}\left(G, k, S_{Clique}(G, k)\right) = \exists V' \in S_{Clique}(G, k) \left[|V'| = k \land \forall u, v \in V' \left[(u, v) \in E\right]\right]$



Instance of Clique

Instance of our problem

#### Instance of Clique

 $I_{Clique} = \{ \langle G = (V, E), k \rangle : G \text{ is an undirected graph } \land k \in \mathbb{N} \}$ 

#### Instance of our problem

$$I_{ex} = \{ \langle G_{ex} = (V, E, w), k, h = |E| \rangle$$
  
:  $G_{ex}$  is an undirected and weighted graph  $\land w : E \rightarrow \mathbb{N} \land k \in \mathbb{N} \}$ 

If G has a **clique** of k vertices,  $G_{ex}$  has a **clique** of k vertices, such that the **sum** of the weights of its edges is at most k,  $I_{ex}$  is a "yes" instance

If G doesn't have a **clique** of k vertices,  $G_{ex}$  doesn't have a **clique** of k vertices,  $I_{ex}$  is a "no" instance