Expectation is not everything....

Which game would you prefer?

1. With probability $\frac{1}{2}$ win \$1, with probability $\frac{1}{2}$ pay \$1.
2. With probability $\frac{1}{2}$ win \$100,000, with probability $\frac{1}{2}$ pay \$100,000.
3. With probability $\frac{1}{1{,}000{,}000}$ win \$1,000,000, with probability $\frac{1}{2}$ pay \$5, else \$0.

(1) $E(X) = 1 \cdot \frac{1}{2} + (-1) \cdot \frac{1}{2} = \emptyset$

(2) $E(X) = 10^5 \cdot \frac{1}{2} + (-10^5) \cdot \frac{1}{2} = 0$

(3) $E(X) = 10^6 \cdot \frac{1}{10^6} + (-5) \cdot \frac{1}{2} + 0 \cdot \left(1 - \frac{1}{2} - \frac{1}{10^6}\right) = 1 - \frac{5}{2}$

Which Job Would You Prefer?

- A job that pays \$1000 a week. $\longrightarrow E(X) = 1000$
- A job that pays \$1 a week plus a bonus of \$1,000,000 with probability $\frac{1}{1000}$. $\longrightarrow E(X) = 1 + 10^3$

Bounding Deviation from Expectation

Theorem

[Markov Inequality] *For any non-negative random variable*

$$Pr(X \geq a) \leq \frac{E[X]}{a}.$$

Proof.

$$E[X] = \sum iPr(X = i) \geq a \sum_{i \geq a} Pr(X = i) = aPr(X \geq a).$$

□

Example: What is the probability of getting more than $\frac{3N}{4}$ heads in N coin flips? $\leq \frac{N/2}{3N/4} \leq \frac{2}{3}$.

Variance

Definition

The **variance** of a random variable X is

$$Var[X] = \mathbf{E}[(X - \mathbf{E}[X])^2] = \underbrace{\mathbf{E}[X^2]}_{\text{hard to compute}} - (\mathbf{E}[X])^2.$$

Easy to compute

Definition

The **standard deviation** of a random variable X is

$$\sigma(X) = \sqrt{Var[X]}.$$

Example: Let X be a 0-1 random variable with $Pr(X = 0) = Pr(X = 1) = 1/2$.

$$\left[E[X] = 1/2.\right]^{*}$$

$$Var[X] = \frac{1}{2}(1 - \frac{1}{2})^2 + \frac{1}{2}(0 - \frac{1}{2})^2 = \frac{1}{4}.$$

$$= E\left(\left(X - \overset{*}{E}(X)\right)^2\right)$$

Chebyshev's Inequality

Theorem 1

For **any** *random variable*

$$\boxed{Pr(|X - E[X]| \geq a) \leq \frac{Var[X]}{a^2}.}$$

Proof.

$$Pr(|X - E[X]| \geq a) = Pr((X - E[X])^2 \geq a^2)$$

By Markov inequality

$$Pr((X - E[X])^2 \geq a^2) \leq \frac{E[(X - E[X])^2]}{a^2}$$

$$= \frac{Var[X]}{a^2}$$

GENERAL USEFUL CONCENTRATION BOUNDS

Theorem 2

For **any** *random variable*

$\sigma[X] = \sqrt{VAR[X]}$

$$Pr(|X - E[X]| \geq a\overbrace{\sigma[X]}) \leq \frac{1}{a^2}.$$

Theorem 3

For **any** *random variable*

$$Pr(|X - E[X]| \geq \epsilon E[X]) \leq \frac{Var[X]}{\epsilon^2 (E[X])^2}.$$

Theorem 4

If X and Y are independent random variable

$$E[XY] = E[X] \cdot E[Y],$$

Proof.

$$E[XY] = \sum_i \sum_j i \cdot jPr((X = i) \cap (Y = j)) =$$

$$\sum_i \sum_j ijPr(X = i) \cdot Pr(Y = j) =$$

$$(\sum_i iPr(X = i))(\sum_j jPr(Y = j)).$$

$\square$

Theorem 5

If X *and* Y *are independent random variable*

$$Var[X+Y] = Var[X] + Var[Y].$$

Proof.

$$Var[X+Y] = E[(X+Y-E[X]-E[Y])^2] =$$

$$E[(X-E[X])^2 + (Y-E[Y])^2 + 2(X-E[X])(Y-E[Y])] =$$

$$Var[X] + Var[Y] + 2E[X-E[X]]E[Y-E[Y]]$$

Since the random variables $X - E[X]$ and $Y - E[Y]$ are independent.
But $E[X-E[X]] = E[X] - E[X] = 0.$ □

Back to Coin Flips

Assume again that we flip N coins. Let X be the number of heads.
$X_i = 1$ if the i-th flip was a head else $X_i = 0$.
$E[X_i] = 1/2$. $Var[X_i] = 1/4$.

$$Pr(X \geq 3N/4) \leq Pr(|X - E[X]| \geq N/4) =$$

THM 3

$$Pr(|X - E[X]| \geq E[X]/2) \leq \frac{Var[X]}{(E[X])^2(1/4)} =$$

THM 5

$$\frac{N/4}{(N^2/4)(1/4)} = \boxed{4/N.}$$

!!

A significantly better bound than $3/4$.

Bernoulli Trial

Let X be a 0-1 random variable such that

$$Pr(X = 1) = p, \qquad Pr(X = 0) = 1 - p.$$

$$E[X] = 1 \cdot p + 0 \cdot (1 - p) = p.$$

$$Var[X] = p(1 - p)^2 + (1 - p)(0 - p)^2 = p(1 - p)[1 - p + p] =$$

$$p(1 - p).$$

A Binomial Random variable

Consider a sequence of n independent Bernoulli trials $X_1, ..., X_n$. Let

$$X = \sum_{i=1}^{n} X_i.$$

X has a **Binomial** distribution $X \sim B(n, p)$.

$$Pr(X = k) = \binom{n}{k} p^k (1-p)^{n-k}.$$

$$E[X] = np.$$

$$Var[X] = np(1-p).$$

Algorithm for Computing the Median

The **median** of a set X of n distinct elements is the $\lceil \frac{n}{2} \rceil$ largest element in the set.

If $n = 2k + 1$, the median element is the $k + 1$-th element in the sorted order.

Easily computed through sorting in $O(n \log n)$ time. There exists a complicated $O(n)$ deterministic algorithm.

↳ not suitable for most

Randomized Median Algorithm

S

Input: A set of $n = 2k + 1$ elements from a totally ordered universe.

Output: The $k + 1$-th largest element in the set.

$|R| = n^{3/4}$

1. Pick a (multi)-set R of $s = n^{3/4}$ elements in S, chosen independently and uniformly at random with replacement. Sort the set R. *
2. Let d be the $(\frac{1}{2}n^{3/4} - \sqrt{n})$th smallest element in the sorted set R.
3. Let u be the $(\frac{1}{2}n^{3/4} + \sqrt{n})$th smallest element in the sorted set R.
4. By comparing every element in S to d and u compute the set $C = \{x \in S \ : \ d \leq x \leq u\}$, and the numbers $\ell_d = |\{x \in S \ : \ x < d\}|$ and $\ell_u = |\{x \in S \ : \ x > u\}|$.

   linear time → | ← linear time
5. If $\ell_d > n/2$ or $\ell_u > n/2$ then FAIL.
6. If $|C| \leq 4n^{3/4}$ then sort the set C, otherwise FAIL.
7. Output the $(\lfloor \frac{n}{2} \rfloor - \ell_d + 1)$st element in the sorted order of C.

   → C "starts" from ℓ_d

CONSTRUCTION of R from S' $\quad |S|=n$

- $R=\emptyset$
- For $i=1$ to $n^{3/4}$ do

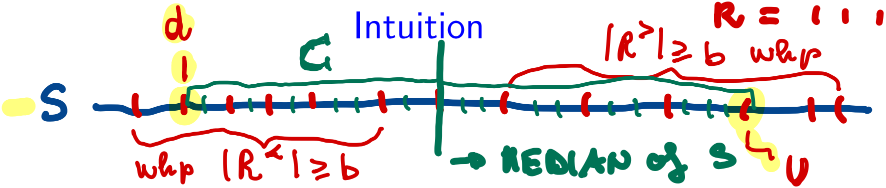
  pick $v \in_U S'$ and do $\left[\begin{array}{l} R=R\cup\{v\} \\ \text{Keep } v \text{ in } S \end{array}\right]$
- STOP

FACT. $\forall v\in S:\ \Pr[v\in R] = 1-\left(1-\frac{1}{n}\right)^{n^{3/4}} \cdot (1)$

for suff. large n, $(1) \sim 1-e^{-\frac{n^{3/4}}{n}} \sim \frac{1}{n^{1/4}}$

$\Pr[v\in R] \sim \frac{1}{n^{1/4}} \quad \forall v \in S'$

Intuition



- We can sort sets of size $< n/\log n$ in linear time.
- The sample of R elements are spaced "more or less" evenly among the elements of X.
- W.h.p. more than $\underbrace{\frac{1}{2}n^{3/4} - \sqrt{n}}_{b}$ samples are smaller than the median. $\Rightarrow |R^<| \geq b$
- W.h.p. more than $\underbrace{\frac{1}{2}n^{3/4} - \sqrt{n}}_{b}$ samples are larger than the median. $|R^>| \geq b$
- W.h.p. the median is in the set C, and $|C| < n/logn$.

→ Put all elements x of S s.t. $d \leq x \leq u$ into C

What might turn WRONG? BAD EVENT

$Y_1 = \sum_{i=1}^{n^{3/4}} Y_1^i \qquad Y_1^i = 1$ iff $x_i <$ median i-th sample

Let Y_1 be the number of samples below the median. in R
Let Y_2 be the number of samples above the median. in R
The algorithm fails to compute the median in $O(n)$ time iff at least one of the following three events occurs: BAD EVENTS

1. $E_1: \ Y_1 < \frac{1}{2}n^{3/4} - \sqrt{n}$. → R is too shifted on the Right
   d is too large
2. $E_2: \ Y_2 < \frac{1}{2}n^{3/4} - \sqrt{n}$. → R is ... symmetric
3. $E_3: \ |C| > n/\log n$. → R is not well distributed

What is the probability that the three random variables Y_1, Y_2 and $|C|$ are all within the required ranges?.

The sample space in execution of this algorithm is the set of all possible choices of $n^{3/4}$ elements from n, with repetitions. (The sample space has $n^{n^{3/4}}$ points.)
Each point in the sample space defines values for Y_1, Y_2 and $|C|$.
Computing the probabilities directly is too complicated, instead we use bounds on deviation from the expectation.

$Y_1^i = \begin{cases} 1 & i\text{-th} < \text{median} \\ 0 \end{cases}$

$Y_1 \equiv \sum_{i=1}^{n^{3/4}} Y_1^i$

Y_1 be the number of samples below the median.

What is the probability that $Y_1 < \frac{1}{2}n^{3/4} - \sqrt{n}$

Viewing Y_1 as the sum of $n^{3/4}$ independent 0-1 random variable, each with expectation $1/2$ and variance $1/4$ we prove (not counting the median itself):

$$E[Y_1] = \frac{1}{2}n^{3/4}.$$

$$Var[Y_1] = \frac{1}{4}n^{3/4}.$$

$VAR[Y_1] =$

$= \sum VAR[Y_1^i] =$

$n^{3/4} \cdot \frac{1}{4}$

THM 1

Applying $\overbrace{\text{Chebyshev Inequality}}$ we get:

$$Pr(E_1 : Y_1 < \frac{1}{2}n^{3/4} - \sqrt{n}) \leq Pr(\overbrace{|Y_1 - E[Y_1]|} > \overbrace{\sqrt{n}}^{a}) \leq$$

$$\frac{Var[Y_1]}{\underbrace{n}_{a^2}} = \frac{n^{3/4}/4}{n} = \frac{1}{4}n^{-1/4}. \qquad \text{whp}$$

Similarly

$$Pr(E_2 : Y_2 < \frac{1}{2}n^{3/4} - \sqrt{n}) \leq \frac{1}{4}n^{-1/4}.$$

$$Pr(E_1 \cup E_2) \leq \frac{2}{4}n^{-1/4}.$$

$C \equiv \{ x \in S \mid d \leq x \leq u \}$

$d \equiv (\frac{1}{2} n^{3/4} - \sqrt{n})$ elm in R

$U \equiv \ldots + \sqrt{n} \ldots$

Recall: $E_3 : \ |C| > n / \log n$.

Lemma

$$\Pr(E_3) \leq \frac{1}{2} n^{-1/4}.$$

Define the following two events:

1. $\mathcal{E}_{3,1}$: at least $2n^{3/4}$ elements of C are greater than the median;
2. $\mathcal{E}_{3,2}$: at least $2n^{3/4}$ elements of C are smaller than the median.

If $|C| > 4n^{3/4}$, then at least one of the above two events occurs.

↳ this is deterministically true by def. of MEDIAN.

u ≡ is the $\left(\frac{1}{2}n^{3/4}+\sqrt{n}\right)$-th element in SORTED R

* by def of C ⇒ u is the largest element of C

We bound $\mathcal{E}_{3,1}$: at least $2n^{3/4}$ elements of C are greater than the median;

At least $2n^{3/4}$ elements of C above the median $\Rightarrow$ *

u is at least the $\frac{1}{2}n + 2n^{3/4}$ largest in S $\Rightarrow$

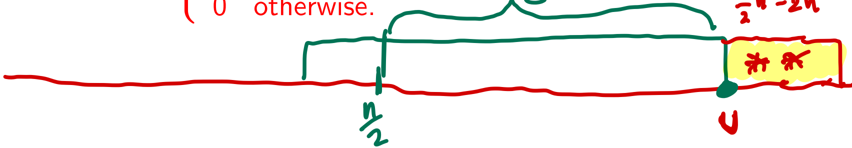
R had at least $\frac{1}{2}n^{3/4} - \sqrt{n}$ samples among the $\frac{1}{2}n - 2n^{3/4}$ largest **

elements in S.

Let X be the number of samples among the $\frac{1}{2}n - 2n^{3/4}$ largest **

elements in S. Let $X = \sum_{i=1}^{n^{3/4}} X_i$ where

$$X_i = \begin{cases} 1 & \text{the } i\text{-th sample in } \frac{1}{2}n - 2n^{3/4} \\ & \text{largest elements in } S \\ 0 & \text{otherwise.} \end{cases}$$



$$E[X_i] = E[(X_i)^2] = \frac{1}{2} - 2n^{-1/4}$$

$$\frac{\frac{1}{2}n - 2n^{3/4}}{n} = P_n$$

$$Var[X_i] = E[(X_i)^2] - (E[X_i])^2 \leq \frac{1}{4}.$$

$$E[X] = \frac{1}{2}n^{3/4} - 2\sqrt{n}$$

$$Var[X] \leq \frac{1}{4}n^{3/4}$$

Applying Chebyshev's Inequality yields

$$\begin{aligned}
\Pr(\mathcal{E}_{3,1}) &= \Pr(X \geq \frac{1}{2}n^{3/4} - \sqrt{n}) \\
&\leq \Pr(|X - E[X]| \geq \sqrt{n}) \\
&\leq \frac{Var[X]}{n} = \frac{\frac{n^{\frac{3}{4}}}{4}}{n} = \frac{1}{4}n^{-\frac{1}{4}}.
\end{aligned}$$

Similarly,

$$\Pr(\mathcal{E}_{3,2}) \leq \frac{1}{4} n^{-\frac{1}{4}},$$

and

$$\Pr(E_3) \leq \Pr(\mathcal{E}_{3,1}) + \Pr(\mathcal{E}_{3,2}) \leq \frac{1}{2} n^{-\frac{1}{4}}.$$

The probability that the algorithm succeeds is

$$\geq 1 - (Pr(E_1) + Pr(E_2) + Pr(E_3)) \geq 1 - \frac{1}{n^{1/4}}.$$

The Geometric Distribution

- How many times we need to perform a trial with probability p for success till we get the first success?
- How many times do we need to roll a dice until we get the first 6?

Definition

A geometric random variable X with parameter p is given by the following probability distribution on $n = 1, 2, \ldots$.

$$\Pr(X = n) = (1 - p)^{n-1} p.$$

Memoryless Distribution

Lemma

For a geometric random variable with parameter p and $n > 0$,

$$\Pr(X = n + k \mid X > k) \quad = \quad \Pr(X = n).$$

Proof.

$$\begin{aligned}
\Pr(X = n + k \quad \mid \quad X > k) &= \frac{\Pr((X = n + k) \cap (X > k))}{\Pr(X > k)} \\
&= \frac{\Pr(X = n + k)}{\Pr(X > k)} = \frac{(1 - p)^{n+k-1} p}{\sum_{i=k}^{\infty} (1 - p)^i p} \\
&= \frac{(1 - p)^{n+k-1} p}{(1 - p)^k} = (1 - p)^{n-1} p = \Pr(X = n).
\end{aligned}$$

$\square$

Expectation

- Let X be a geometric random variable with parameter p.
- Let $Y = 1$ if the first trail is a success, $Y = 0$ otherwise.
- $$\begin{aligned} \mathbf{E}[X] &= \Pr(Y=0)\mathbf{E}[X \mid Y=0] + \Pr(Y=1)\mathbf{E}[X \mid Y=1] \\ &= (1-p)\mathbf{E}[X \mid Y=0] + p\mathbf{E}[X \mid Y=1]. \end{aligned}$$
- If $Y = 0$ let Z be the number of trials after the first one.
- $\mathbf{E}[X] = (1-p)\mathbf{E}[Z+1] + p \cdot 1 = (1-p)\mathbf{E}[Z] + 1$
- But $\mathbf{E}[Z] = \mathbf{E}[X]$, giving $\mathbf{E}[X] = 1/p$.

Example: Coupon Collector's Problem

- We place balls independently and uniformly at random in n boxes.
- Let X be the number of balls placed until all boxes are not empty.
- What is $\mathbf{E}[X]$?

- Let X_i = number of balls placed when there were exactly $i-1$ non-empty boxes.
- $X = \sum_{i=1}^{n} X_i$.
- X_i is a geometric random variable with parameter $p_i = 1 - \frac{i-1}{n}$.
-

$$\mathbf{E}[X_i] = \frac{1}{p_i} = \frac{n}{n-i+1}.$$

$$\begin{aligned}
\mathbf{E}[X] &= E\left[\sum_{i=1}^{n} X_i\right] = \sum_{i=1}^{n} \mathbf{E}[X_i] \\
&= \sum_{i=1}^{n} \frac{n}{n-i+1} = n\sum_{i=1}^{n} \frac{1}{i} = n \ln n + \Theta(n).
\end{aligned}$$

Variance of a Geometric Random Variable

- We use
$$Var[X] = \mathbf{E}[(X - \mathbf{E}[X])^2] = \mathbf{E}[X^2] - (\mathbf{E}[X])^2.$$

- To compute $\mathbf{E}[X^2]$, let $Y = 1$ if the first trail is a success, $Y = 0$ otherwise.
-
$$\begin{aligned}\mathbf{E}[X^2] &= \Pr(Y=0)\mathbf{E}[X^2 \mid Y=0] + \Pr(Y=1)\mathbf{E}[X^2 \mid Y=1] \\ &= (1-p)\mathbf{E}[X^2 \mid Y=0] + p\mathbf{E}[X^2 \mid Y=1].\end{aligned}$$

- If $Y = 0$ let Z be the number of trials after the first one.
-
$$\begin{aligned}\mathbf{E}[X^2] &= (1-p)\mathbf{E}[(Z+1)^2] + p \cdot 1 \\ &= (1-p)\mathbf{E}[Z^2] + 2(1-p)\mathbf{E}[Z] + 1,\end{aligned}$$

- $\mathbf{E}[Z] = 1/p$ and $\mathbf{E}[Z^2] = \mathbf{E}[X^2]$.
- $$\begin{aligned} \mathbf{E}[X^2] &= (1-p)\mathbf{E}[(Z+1)^2] + p \cdot 1 \\ &= (1-p)\mathbf{E}[Z^2] + 2(1-p)\mathbf{E}[Z] + 1, \end{aligned}$$
- $$\mathbf{E}[X^2] = (1-p)\mathbf{E}[X^2] + 2(1-p)/p + 1 = (1-p)\mathbf{E}[X^2] + (2-p)/p,$$
- $\mathbf{E}[X^2] = (2-p)/p^2$.

$$
\begin{aligned}
Var[X] &= \mathbf{E}[X^2] - \mathbf{E}[X]^2 \\
&= \frac{2-p}{p^2} - \frac{1}{p^2} \\
&= \frac{1-p}{p^2}.
\end{aligned}
$$

Back to the Coupon Collector's Problem

- We place balls independently and uniformly at random in n boxes.
- Let X be the number of balls placed until all boxes are not empty.
- $E[X] = nH_n = n \ln n + \Theta(n)$
- What is $\Pr(X \geq 2\mathbf{E}[X])$?
- Applying Markov's inequality

$$\Pr(X \geq 2nH_n) \leq \frac{1}{2}.$$

- Can we do better?

- Let X_i = number of balls placed when there were exactly $i-1$ non-empty boxes.
- $X = \sum_{i=1}^{n} X_i$.
- X_i is a geometric random variable with parameter $p_i = 1 - \frac{i-1}{n}$.
- $Var[X_i] \le \frac{1}{p^2} \le (\frac{n}{n-i+1})^2$.
- $$Var[X] = \sum_{i=1}^{n} Var[X_i] \le \sum_{i=1}^{n} \left(\frac{n}{n-i+1}\right)^2 = n^2 \sum_{i=1}^{n} \left(\frac{1}{i}\right)^2 \le \frac{\pi^2 n^2}{6}.$$
- By Chebyshev's inequality

$$\Pr(|X - nH_n| \ge nH_n) \le \frac{n^2\pi^2/6}{(nH_n)^2} = \frac{\pi^2}{6(H_n)^2} = O\left(\frac{1}{\ln^2 n}\right).$$

Direct Bound

- The probability of not obtaining the i-th coupon after $n \ln n + cn$ steps:

$$\left(1 - \frac{1}{n}\right)^{n(\ln n + c)} < \mathrm{e}^{-(\ln n + c)} = \frac{1}{\mathrm{e}^c n}.$$

- By a union bound, the probability that some coupon has not been collected after $n \ln n + cn$ step is e^{-c}.
- The probability that all coupons are not collected after $2n \ln n$ steps is at most $1/n$.

The Advantage of Multiple Samples

Theorem

For any constant a,

$$Var[aX] = a^2 Var[x].$$

Proof.

$$\begin{aligned} Var[aX] &= E[(aX - E[aX])^2] = E[a^2(X - E[X])^2] \\ &= a^2 E[(X - E[X])^2] = a^2 Var[X]. \end{aligned}$$

$\square$

Theorem

Let $X_1, \ldots, X_n$ *be* n *independent, identically distributed random variable. Let* $\bar{X} = \frac{1}{n}\sum_{i=1}^{n} X_i$.

$$Var[\bar{X}] = Var[\frac{1}{n}\sum_{i=1}^{n} X_i] = \frac{1}{n^2} Var[\sum_{i=1}^{n} X_i] = \frac{1}{n} Var[X_i].$$

The (Weak) Law of Large Numbers

Theorem

Let $X_1, \dots, X_n$ be independent, identically distributed, random variables. Let $\bar{X}_n = \frac{1}{n}\sum_{i=1}^{n} X_i$. For any constant $\epsilon > 0$,

$$\lim_{n\to\infty} \Pr(|\bar{X}_n - \mathbf{E}[X]| \le \epsilon) = 1.$$

Proof.

$Var[\bar{X}_n] = \frac{1}{n} Var[X_i]$. Applying Chebyshev's bound

$$\Pr(|\bar{X}_n - \mathbf{E}[X]| > \epsilon) \le \frac{Var[X_i]}{n\epsilon^2}.$$

□

[Can be proven even when $Var[X_i]$ is not bounded.]