

CS 754 - Advanced Image Processing

Assignment 2 - Report

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Question 1

1.1

If $\delta_{2s} = 1$, then for any $2s$ - sparse vector $h \in \mathbb{R}^n$, we have

$$0 \leq \|\Phi h\|_2^2 \leq 2 \|h\|_2^2 \quad (1)$$

If $2s$ columns of Φ are linearly dependent, then we have

$$\sum_i \alpha_i \phi_i = 0, \alpha_i \neq 0$$

Where $\{\phi_i\}$ are the $2s$ linearly dependent columns of Φ and $\{\alpha_i\}$ are some non-zero constants. This summation can be re-written as

$$\Phi v = 0, \quad \Phi = [\phi_1 | \dots | \phi_{2s}], v^T = [\alpha_1, \dots, \alpha_{2s}]$$

Now, we can extend v by adding rows of zeros so that the extended vector $v_1 \in \mathbb{R}^n$. Similarly, we can extend the columns of Φ by adding random columns (this works because we want the measurement matrix to be random in Compressed Sensing) so that $\Phi \in \mathbb{R}^{m \times n}$. Even after this extension, we still have

$$\Phi v_1 = 0$$

Clearly, we can see that v_1 is a $2s$ - sparse vector and also satisfies (1) trivially because $\|\Phi v_1\|_2 = 0$. Hence, the above justifies that when $\delta_{2s} = 1$, there exists a $2s$ - sparse vector h such that $\Phi h = 0$ and $2s$ columns of Φ are linearly dependent.

1.2

We have x^* and x which satisfy the convex optimization problem

$$\min_{\tilde{x} \in \mathbb{R}^n} \|\tilde{x}\|_1 \quad \text{subject to } \|y - \Phi \tilde{x}\|_2 \leq \epsilon \quad (2)$$

Consider

$$\|\Phi(x^* - x)\|_2 = \|(\Phi x^* - y) + (y - \Phi x)\|_2$$

Then, using the triangle inequality, we have

$$\|\Phi(x^* - x)\|_2 = \|(\Phi x^* - y) + (y - \Phi x)\|_2 \leq \|\Phi x^* - y\|_2 + \|y - \Phi x\|_2$$

But using (2) we can say that

$$\|\Phi x^* - y\|_2 + \|y - \Phi x\|_2 \leq \epsilon + \epsilon = 2\epsilon$$

Hence, we have

$$\|\Phi(x^* - x)\|_2 \leq \|\Phi x^* - y\|_2 + \|y - \Phi x\|_2 \leq 2\epsilon$$

1.3

We know that each h_{T_j} is s sparse. We also know that the ℓ_∞ norm of a vector denotes the maximum element of that vector. As h_{T_j} is s sparse, we have, for each non-zero entry α_i of h_{T_j}

$$\alpha_i^2 \leq \|h_{T_j}\|_\infty^2 \quad i = 1, 2, 3, \dots, s$$

Adding all the above s inequalities, we have

$$\begin{aligned} \|h_{T_j}\|_2^2 &\leq s \|h_{T_j}\|_\infty^2 \\ \implies \|h_{T_j}\|_2 &\leq s^{1/2} \|h_{T_j}\|_\infty \end{aligned} \quad (3)$$

We know that, by the construction of the vectors $\{h_{T_j}\}$, the largest element of h_{T_j} is smaller than the smallest element of $h_{T_{j-1}}$. In fact, the $\|h_{T_j}\|_\infty$ is smaller than every (non-zero) element of $h_{T_{j-1}}$. Re-writing this mathematically, we have

$$\|h_{T_j}\|_\infty \leq |\beta_i| \quad i = 1, 2, 3, \dots, s$$

Adding all the above s inequalities, we get

$$\begin{aligned} s \|h_{T_j}\|_\infty &\leq \|h_{T_{j-1}}\|_1 \\ \implies s^{1/2} \|h_{T_j}\|_\infty &\leq s^{-1/2} \|h_{T_{j-1}}\|_1 \end{aligned} \quad (4)$$

Combining the inequalities (3) and (4), we have

$$\|h_{T_j}\|_2 \leq s^{1/2} \|h_{T_j}\|_\infty \leq s^{-1/2} \|h_{T_{j-1}}\|_1$$

1.4

Using the relation between the 1st and the 3rd expression in the above inequality, we have

$$\|h_{T_j}\|_2 \leq s^{-1/2} \|h_{T_{j-1}}\|_1$$

Summing up all such above inequalities for $j \geq 2$, we have

$$\sum_{j \geq 2} \|h_{T_j}\|_2 \leq s^{-1/2} \left(\|h_{T_1}\|_1 + \|h_{T_2}\|_1 + \dots \right) \quad (5)$$

By construction of $\{h_{T_j}\}$, we can say that

$$\|h_{T_1}\|_1 + \|h_{T_2}\|_1 + \dots = \|h_{T_0^c}\|_1$$

Hence, (5) becomes

$$\sum_{j \geq 2} \|h_{T_j}\|_2 \leq s^{-1/2} \left(\|h_{T_1}\|_1 + \|h_{T_2}\|_1 + \dots \right) \leq s^{-1/2} \|h_{T_0^c}\|_1 \quad (6)$$

1.5

We know that

$$h_{(T_0 \cup T_1)^c} = \sum_{j \geq 2} h_{T_j}$$

Hence, we have

$$\|h_{(T_0 \cup T_1)^c}\|_2 = \left\| \sum_{j \geq 2} h_{T_j} \right\|_2$$

Using the extended version of triangle inequality for the above equation and using the inequality (6) above, we get

$$\|h_{(T_0 \cup T_1)^c}\|_2 = \left\| \sum_{j \geq 2} h_{T_j} \right\|_2 \leq \sum_{j \geq 2} \|h_{T_j}\|_2 \leq s^{-1/2} \|h_{T_0^c}\|_1 \quad (7)$$

1.6

We have the *reverse triangle inequality* for any two real numbers as follows:

$$|x - y| \geq \left| |x| - |y| \right|$$

Using this inequality for $|x_i + h_i|, i \in T_0$, we have

$$|x_i + h_i| \geq \left| |x_i| - |h_i| \right| \geq |x_i| - |h_i|$$

Summing over all $i \in T_0$, we have

$$\sum_{i \in T_0} |x_i + h_i| \geq \|x_{T_0}\|_1 - \|h_{T_0}\|_1 \quad (8)$$

Again using the reverse triangle inequality for $|h_i + x_i|, i \in T_0^c$, we have

$$|h_i + x_i| \geq \left| |h_i| - |x_i| \right| \geq |h_i| - |x_i|$$

Summing over all $i \in T_1$, we have

$$\sum_{i \in T_0^c} |x_i + h_i| \geq \|h_{T_0^c}\|_1 - \|x_{T_0^c}\|_1 \quad (9)$$

By adding the inequalities (8) and (9) we get

$$\sum_{i \in T_0} |x_i + h_i| + \sum_{i \in T_0^c} |x_i + h_i| \geq \|x_{T_0}\|_1 - \|h_{T_0}\|_1 + \|h_{T_0^c}\|_1 - \|x_{T_0^c}\|_1 \quad (10)$$

1.7

As per the convex optimization problem that we are solving, we have

$$\|x\|_1 \geq \|x + h\|_1$$

Using this in the inequality (10) we have

$$\|x\|_1 \geq \|x_{T_0}\|_1 - \|h_{T_0}\|_1 + \|h_{T_0^c}\|_1 - \|x_{T_0^c}\|_1 \quad (11)$$

By definition we have

$$\|x_{T_0^c}\|_1 = \|x - x_s\|_1$$

By using the reverse triangle inequality, we have

$$\|x_{T_0^c}\|_1 = \|x - x_s\|_1 \geq \|x\|_1 - \|x_s\|_1$$

Noting the fact that $\|x_s\| = \|x_{T_0}\|$ and using inequality (11) we have

$$\|x_{T_0^c}\|_1 \geq \|x_{T_0}\|_1 - \|h_{T_0}\|_1 + \|h_{T_0^c}\|_1 - \|x_{T_0^c}\|_1 - \|x_{T_0}\|_1$$

which is same as

$$\|h_{T_0^c}\|_1 \leq \|h_{T_0}\|_1 + 2 \|x_{T_0^c}\|_1 \quad (12)$$

1.8

Using inequalities (7) and (12) we have

$$\|h_{(T_0 \cup T_1)^c}\|_2 \leq s^{-1/2} \left(\|h_{T_0}\|_1 + 2 \|x_{T_0^c}\|_1 \right)$$

Using the known inequality $\|v\|_1 \leq \sqrt{k} \|v\|_2$ for a k sparse vector v , we have

$$\|h_{(T_0 \cup T_1)^c}\|_2 \leq s^{-1/2} \left(\|h_{T_0}\|_1 + 2 \|x_{T_0^c}\|_1 \right) \leq s^{-1/2} \left(\sqrt{s} \|h_{T_0}\|_2 + 2 \|x_{T_0^c}\|_1 \right)$$

Hence we have

$$\|h_{(T_0 \cup T_1)^c}\|_2 \leq \|h_{T_0}\|_2 + 2e_0, \quad e_0 \equiv s^{-1/2} \|x - x_s\|_1 \quad (13)$$

1.9

We know that by Cauchy-Schwartz inequality, we have

$$|\langle \Phi h_{T_0 \cup T_1}, \Phi h \rangle| \leq \|\Phi h_{T_0 \cup T_1}\|_2 \|\Phi h\|_2 \quad (14)$$

As $h = x^* - x$, we have by inequality justified in section 1.3,

$$\|\Phi h\|_2 = \|\Phi(x^* - x)\|_2 \leq 2\epsilon$$

also, as $h_{T_0 \cup T_1}$ is a $2s$ sparse vector and Φ follows RIP of order $2s$, we have

$$\|\Phi h_{T_0 \cup T_1}\|_2 \leq \sqrt{1 + \delta_{2s}} \|h_{T_0 \cup T_1}\|_2$$

Using the above two inequalities in (14), we have

$$|\langle \Phi h_{T_0 \cup T_1}, \Phi h \rangle| \leq \|\Phi h_{T_0 \cup T_1}\|_2 \|\Phi h\|_2 \leq 2\epsilon \sqrt{1 + \delta_{2s}} \|h_{T_0 \cup T_1}\|_2 \quad (15)$$

1.10

Lemma 2.1 from the paper states

$$|\langle \Phi x, \Phi x' \rangle| \leq \delta_{s+s'} \|x\|_2 \|x'\|_2$$

for all x, x' supported on disjoint subsets $T, T' \subset \{1, 2, \dots, n\}$ with $|T| \leq s, |T'| \leq s'$ Using this lemma for $x = h_{T_0}$ and $x' = h_{T_j}$, we have

$$|\langle \Phi h_{T_0}, \Phi h_{T_j} \rangle| \leq \delta_{2s} \|h_{T_0}\|_2 \|h_{T_j}\|_2 \quad (16)$$

1.11

We know that for two positive real numbers a and b , the QM-AM inequality is as follows:

$$\frac{a+b}{2} \leq \sqrt{\frac{a^2+b^2}{2}}$$

Taking $a = \|h_{T_0}\|_2$ and $b = \|h_{T_1}\|_2$ and substituting in the above inequality, we get

$$\|h_{T_0}\|_2 + \|h_{T_1}\|_2 \leq \sqrt{2} \sqrt{\|h_{T_0}\|_2^2 + \|h_{T_1}\|_2^2} = \sqrt{2} \|h_{T_0 \cup T_1}\|_2 \quad (17)$$

as T_0 and T_1 are disjoint.

1.12

By replacing T_0 with T_1 in inequality (16) and adding both, we get

$$|\langle \Phi h_{T_0}, \Phi h_{T_j} \rangle| + |\langle \Phi h_{T_1}, \Phi h_{T_j} \rangle| \leq \delta_{2s} (\|h_{T_0}\|_2 + \|h_{T_1}\|_2) \|h_{T_j}\|_2$$

Making use of triangle inequality for the LHS of above inequality, we get

$$|\langle \Phi h_{T_0} + \Phi h_{T_1}, \Phi h_{T_j} \rangle| \leq \delta_{2s} (\|h_{T_0}\|_2 + \|h_{T_1}\|_2) \|h_{T_j}\|_2$$

Replacing $\Phi h_{T_0} + \Phi h_{T_1} = \Phi h_{T_0 \cup T_1}$ and above inequalities for $j \geq 2$, we get

$$\sum_{j \geq 2} |\langle \Phi h_{T_0 \cup T_1}, \Phi h_{T_j} \rangle| \leq \delta_{2s} (\|h_{T_0}\|_2 + \|h_{T_1}\|_2) \left(\sum_{j \geq 2} \|h_{T_j}\|_2 \right)$$

Again making use of triangle inequality for the LHS of the above inequality, we get

$$\left| \left\langle \Phi h_{T_0 \cup T_1}, \sum_{j \geq 2} \Phi h_{T_j} \right\rangle \right| \leq \delta_{2s} (\|h_{T_0}\|_2 + \|h_{T_1}\|_2) \left(\sum_{j \geq 2} \|h_{T_j}\|_2 \right)$$

Finally, using (17), we finally have

$$\left| \left\langle \Phi h_{T_0 \cup T_1}, \sum_{j \geq 2} \Phi h_{T_j} \right\rangle \right| \leq \delta_{2s} \sqrt{2} \|h_{T_0 \cup T_1}\|_2 \left(\sum_{j \geq 2} \|h_{T_j}\|_2 \right) \quad (18)$$

Now, using the fact that $\Phi h = \Phi h_{T_0 \cup T_1} + \sum_{j \geq 2} \Phi h_{T_j}$ in the inequality (15) above, we get

$$\left| \left\langle \Phi h_{T_0 \cup T_1}, \Phi h_{T_0 \cup T_1} + \sum_{j \geq 2} \Phi h_{T_j} \right\rangle \right| \leq 2\epsilon \sqrt{1 + \delta_{2s}} \|h_{T_0 \cup T_1}\|_2$$

Using the linearity property of inner product, we have

$$\left| \|\Phi h_{T_0 \cup T_1}\|_2^2 + \left\langle \Phi h_{T_0 \cup T_1}, \sum_{j \geq 2} \Phi h_{T_j} \right\rangle \right| \leq 2\epsilon \sqrt{1 + \delta_{2s}} \|h_{T_0 \cup T_1}\|_2 \quad (19)$$

Now, adding the inequalities (18) and (19), we get

$$\begin{aligned} \left| \|\Phi h_{T_0 \cup T_1}\|_2^2 + \left\langle \Phi h_{T_0 \cup T_1}, \sum_{j \geq 2} \Phi h_{T_j} \right\rangle \right| + \left| \left\langle \Phi h_{T_0 \cup T_1}, \sum_{j \geq 2} \Phi h_{T_j} \right\rangle \right| &\leq 2\epsilon \sqrt{1 + \delta_{2s}} \|h_{T_0 \cup T_1}\|_2 + \\ &\quad \delta_{2s} \sqrt{2} \|h_{T_0 \cup T_1}\| \left(\sum_{j \geq 2} \|h_{T_j}\|_2 \right) \end{aligned}$$

We know from the triangle inequality that

$$\|\Phi h_{T_0 \cup T_1}\|_2^2 \leq \left| \|\Phi h_{T_0 \cup T_1}\|_2^2 + \left\langle \Phi h_{T_0 \cup T_1}, \sum_{j \geq 2} \Phi h_{T_j} \right\rangle \right| + \left| - \left\langle \Phi h_{T_0 \cup T_1}, \sum_{j \geq 2} \Phi h_{T_j} \right\rangle \right|$$

Using this, we finally have,

$$\|\Phi h_{T_0 \cup T_1}\|_2^2 \leq \|h_{T_0 \cup T_1}\|_2 \left(2\epsilon \sqrt{1 + \delta_{2s}} + \sqrt{2} \delta_{2s} \sum_{j \geq 2} \|h_{T_j}\|_2 \right) \quad (20)$$

From the restricted isometry property, we have

$$(1 - \delta_{2s}) \|h_{T_0 \cup T_1}\|_2^2 \leq \|\Phi h_{T_0 \cup T_1}\|_2^2$$

Hence, we have

$$(1 - \delta_{2s}) \|h_{T_0 \cup T_1}\|_2^2 \leq \|\Phi h_{T_0 \cup T_1}\|_2^2 \leq \|h_{T_0 \cup T_1}\|_2 \left(2\epsilon \sqrt{1 + \delta_{2s}} + \sqrt{2} \delta_{2s} \sum_{j \geq 2} \|h_{T_j}\|_2 \right) \quad (21)$$

1.13

By combining the inequality (6) with RHs of (21), we have

$$\|h_{T_0 \cup T_1}\|_2 \left(2\epsilon \sqrt{1 + \delta_{2s}} + \sqrt{2} \delta_{2s} \sum_{j \geq 2} \|h_{T_j}\|_2 \right) \leq \|h_{T_0 \cup T_1}\|_2 \left(2\epsilon \sqrt{1 + \delta_{2s}} + \sqrt{2} \delta_{2s} s^{-1/2} \|h_{T_0^c}\|_1 \right)$$

Using the above inequality with (21), we have

$$(1 - \delta_{2s}) \|h_{T_0 \cup T_1}\|_2^2 \leq \|h_{T_0 \cup T_1}\|_2 \left(2\epsilon \sqrt{1 + \delta_{2s}} + \sqrt{2} \delta_{2s} s^{-1/2} \|h_{T_0^c}\|_1 \right)$$

On simplifying the above inequality, we get

$$\|h_{T_0 \cup T_1}\|_2 \leq \alpha \epsilon + \rho s^{-1/2} \|h_{T_0^c}\|_1, \quad \alpha \equiv \frac{2\sqrt{1 + \delta_{2s}}}{1 - \delta_{2s}}, \rho \equiv \frac{\sqrt{2} \delta_{2s}}{1 - \delta_{2s}} \quad (22)$$

1.14

We now conclude from (22) and (12) that

$$\|h_{T_0 \cup T_1}\|_2 \leq \alpha\epsilon + \rho s^{-1/2} \|h_{T_0}\|_1 + 2\rho e_0 \leq \alpha\epsilon + \rho \|h_{T_0}\|_2 + 2\rho e_0$$

As T_0 and T_1 are disjoint and h_{T_0} and h_{T_1} are both s sparse, we can say that

$$\|h_{T_0}\|_2 \leq \|h_{T_0} + h_{T_1}\|_2 = \|h_{T_0 \cup T_1}\|_2$$

Using this in the above inequality, we get

$$\|h_{T_0 \cup T_1}\|_2 \leq \alpha\epsilon + \rho \|h_{T_0 \cup T_1}\|_2 + 2\rho e_0$$

On simplifying, we get

$$\|h_{T_0 \cup T_1}\|_2 \leq (1 - \rho)^{-1} (\alpha\epsilon + 2\rho e_0) \quad (23)$$

1.15

We can write h as

$$h = h_{T_0 \cup T_1} + h_{(T_0 \cup T_1)^c}$$

Hence, by triangle inequality, we have

$$\|h\|_2 \leq \|h_{T_0 \cup T_1}\|_2 + \|h_{(T_0 \cup T_1)^c}\|_2$$

But from (13), we know that

$$\|h_{(T_0 \cup T_1)^c}\|_2 \leq \|h_{T_0}\|_2 + 2e_0 \leq \|h_{T_0 \cup T_1}\|_2 + 2e_0$$

Combining the above two inequalities, we get

$$\|h\|_2 \leq 2 \|h_{T_0 \cup T_1}\|_2 + 2e_0 \quad (24)$$

Now combining (24) with (23) we get

$$\|h\|_2 \leq 2 \|h_{T_0 \cup T_1}\|_2 + 2e_0 \leq 2(1 - \rho)^{-1} (\alpha\epsilon + (1 + \rho)e_0) \quad (25)$$

1.16

We know that

$$\|h_{T_0}\|_1 \leq \rho \|h_{T_0^c}\|_1$$

Adding $\|h_{T_0^c}\|_1$ on both sides, we get

$$\|h\|_1 \leq (1 + \rho) \|h_{T_0^c}\|_1 \quad (26)$$

Using the result

$$\|h_{T_0^c}\|_1 \leq 2(1 - \rho)^{-1} \|x_{T_0^c}\|_1$$

in (26), we get

$$\|h\|_1 \leq 2(1 + \rho)(1 - \rho)^{-1} \|x_{T_0^c}\|_1 \quad (27)$$

Question 3

3.a

The oracular solution will be nothing but the solution of the linear equation

$$\mathbf{y} = \Phi_S \mathbf{x}_S$$

which is nothing but

$$\tilde{\mathbf{x}} = \Phi_S^\dagger \mathbf{y}$$

3.b

$$\begin{aligned} \|\tilde{\mathbf{x}} - \mathbf{x}\|_2 &= \|\Phi_S^\dagger \mathbf{y} - \mathbf{x}\|_2 \\ &= \|\Phi_S^\dagger (\Phi_S \mathbf{x} - \boldsymbol{\eta}) - \mathbf{x}\|_2 \\ &= \|\Phi_S^\dagger \boldsymbol{\eta}\|_2 \leq \|\Phi_S^\dagger\|_2 \|\boldsymbol{\eta}\|_2 \end{aligned}$$

3.c

We know that if λ_i are the non-zero singular values of Φ , then the singular values of Φ_S^\dagger will be λ_i^{-1} . Also, if the eigen values of Φ are denoted as e_i , then we know $\lambda_i = \sqrt{e_i}$.

As Φ satisfies restricted isometry property of order $2k$, we have

$$(1 - \delta_{2k}) \|\mathbf{x}\|_2^2 \leq \|\Phi \mathbf{x}\|_2^2 \leq (1 + \delta_{2k}) \|\mathbf{x}\|_2^2$$

Re-writing $\|A\|_2^2 = A^T A$, we can say that

$$1 - \delta_{2k} \leq e_i \leq 1 + \delta_{2k}$$

which is equivalent to

$$\frac{1}{\sqrt{1 + \delta_{2k}}} \leq \frac{1}{\lambda_i} \leq \frac{1}{\sqrt{1 - \delta_{2k}}}$$

Hence we can say that the largest singular value of Φ_S^\dagger lies between $\frac{1}{\sqrt{1 + \delta_{2k}}}$ and $\frac{1}{\sqrt{1 - \delta_{2k}}}$

3.d

From Noisy recovery theorem, we have

$$\|x^* - x\|_2 \leq C_0 k^{-1/2} \|x - \tilde{x}\|_1 + C_1 \epsilon$$

We know that for a $2k$ sparse vector $x - \tilde{x}$, we have

$$\|x - \tilde{x}\|_1 \leq \sqrt{2k} \|x - \tilde{x}\|_2$$

Hence, we have

$$\|x^* - x\|_2 \leq \sqrt{2} C_0 \|x - \tilde{x}\|_2 + C_1 \epsilon$$

But we have

$$\|x - \tilde{x}\|_2 \leq \frac{\epsilon}{\sqrt{1 - \delta_{2k}}}$$

Hence, we will have

$$\|x^* - x\|_2 \leq C_2 \frac{\epsilon}{\sqrt{1 - \delta_{2k}}}$$

where

$$C_2 = C_0\sqrt{2} + C_1\sqrt{1 - \delta_{2k}}$$

which is a constant factor.

Question 5

- Title: Boolean Compressed Sensing and Noisy Group Testing
- Link: <https://arxiv.org/pdf/0907.1061.pdf>
- Objective Function:

Here, we try to maximize the likelihood of decoder which in essence minimizes the error probability. For N items and a defective set of size K , the decoder goes through all $\binom{N}{K}$ possible sets of size K and chooses the set that is most likely. If we denote the tests' outcomes as Y^T , we choose ω^* for which

$$p(Y^T | \mathbf{X}_{S_{\omega^*}}) > p(Y^T | \mathbf{X}_{S_{\omega}}); \quad \forall \omega \neq \omega^*$$

Here

- N is the total number of items, k is the known number of defectives (or positive items), p denotes the probability that an item is part of a given test, and T is the total number of tests
- Codewords: For the j -th term, X_j^T is a binary vector $\in \{0, 1\}^T$, with the t -th entry $X_j(t) = 1$ if the j -th item is pooled in test t , and 0 otherwise. Following an information theoretic convention, we call it the j -th codeword. The observation vector Y^T is a binary vector of length T , with entries equal to 1 for the tests with positive outcome. Similarly $Y(t)$ denotes the t -th component of the vector Y^T .
- $\mathbf{X} \in \{0, 1\}^{N \times T}$ is the measurement matrix, or the codebook, which is a collection of N codewords defining the pool design, i.e., the assignment of items to tests.

$$\mathbf{X} = [X_1^T; X_2^T; \dots; X_N^T]$$

- Given a subset $S \subset \{1, 2, \dots, N\}$ with cardinality $|S|$, the matrix \mathbf{X}_S is an $|S| \times T$ matrix formed from the rows indexed by S . In other words, \mathbf{X}_S denotes the codewords (each of length T) corresponding to the items in S . Similarly, X_S denotes a vector, whose components are restricted to the set of components indexed by S . Thus, X_S is a column of the matrix \mathbf{X}_S . When indexing by test is needed, $X_S(t)$ is used to specifically denote the t -th column of the matrix \mathbf{X}_S , and $X_j(t)$ is the t -th component of the vector X_j^T .

- Index the different sets of items of size K as S_ω with index ω . Since there are N items in total, there are $\binom{N}{K}$ such sets, hence

$$\omega \in \mathcal{I} = \left\{ 1, 2, \dots, \binom{N}{K} \right\}$$

- Differences between Tapestry and the proposed method:
 1. Tapestry models the task from a linear algebra perspective whereas the proposed method models the task from an Information Theory perspective making use of channel coding concepts and probability and random processes.
 2. Tapestry solves a constrained convex optimization problem whereas the proposed method solves an un-constrained optimization problem.
 3. In Tapestry pooling, we model noise as a Poisson random variable whereas in the proposed method they model it as a Bernoulli random variable.