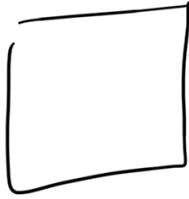


NOSQL - II



Inside a machine

① Unlike SQL, NoSQL is unstructured and has no fixed size (Update) → row

xyz → "abc: xyz"
 xyz → "_____"

key: value

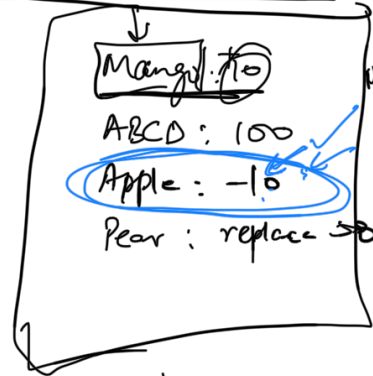
② SQL
 B+ trees write to log(N) read log(N)

Write-heavy ← write more optimise

① Immutable

update → create a new copy.

②



RAM
 mango: 10

ABCD: 100

Apple: -10

Pear: replace 50



Apple: 10

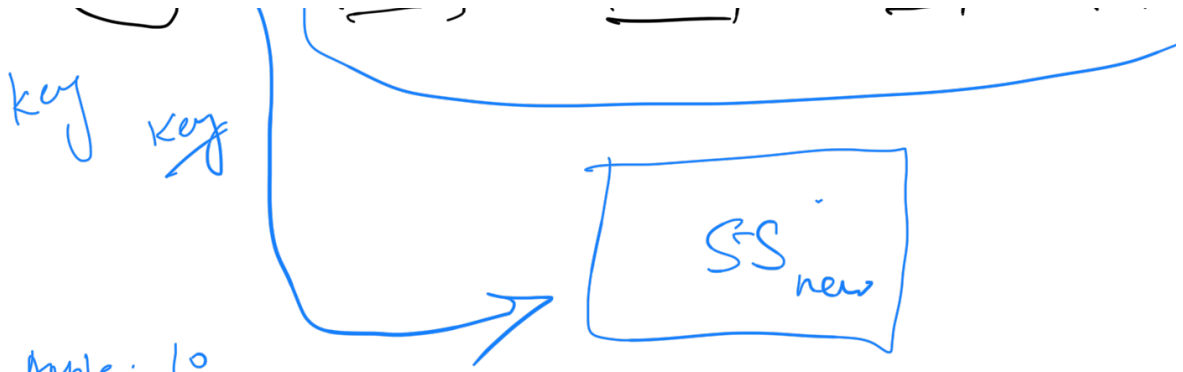
Banana: 20

Mango: 50

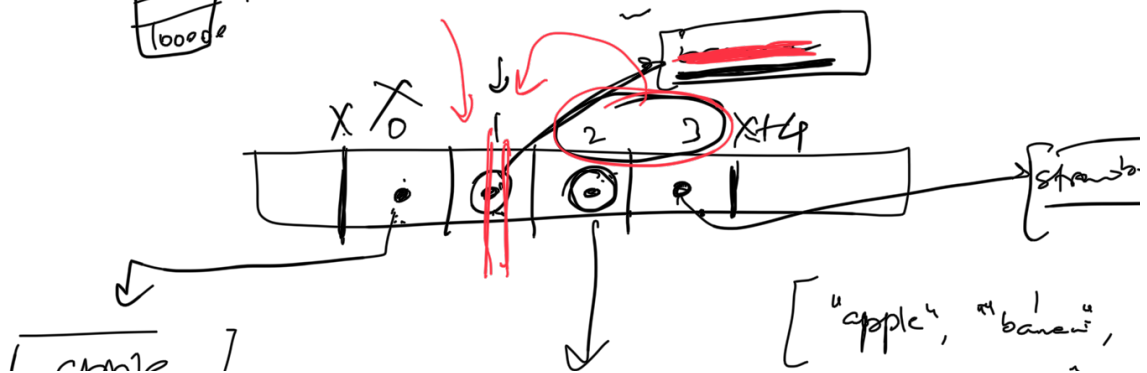
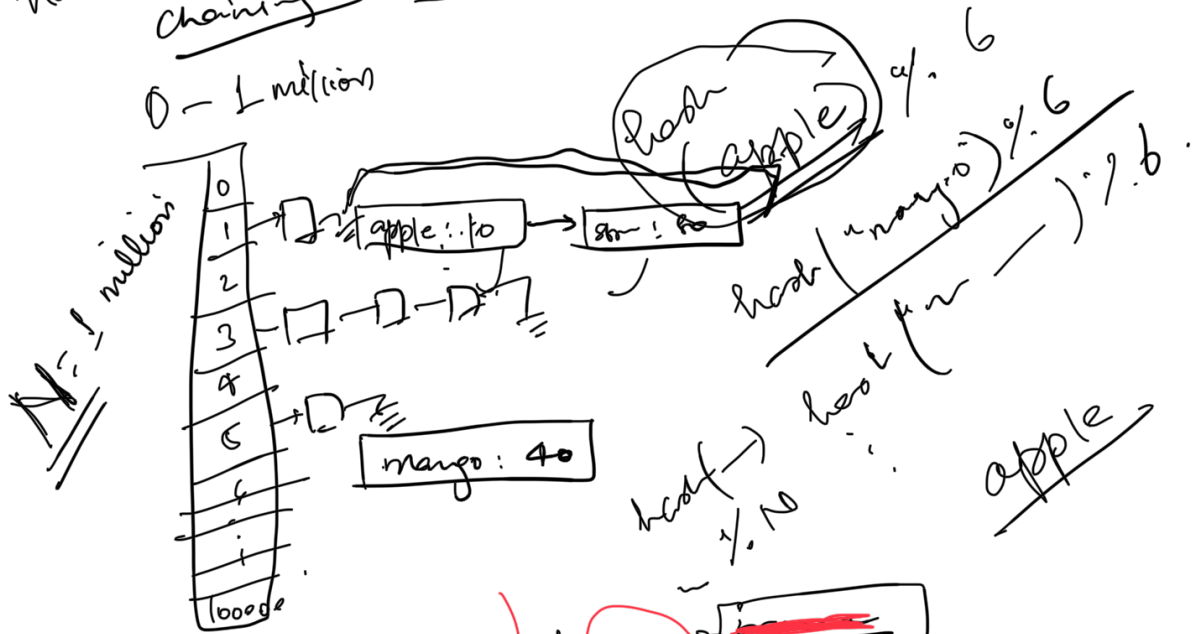
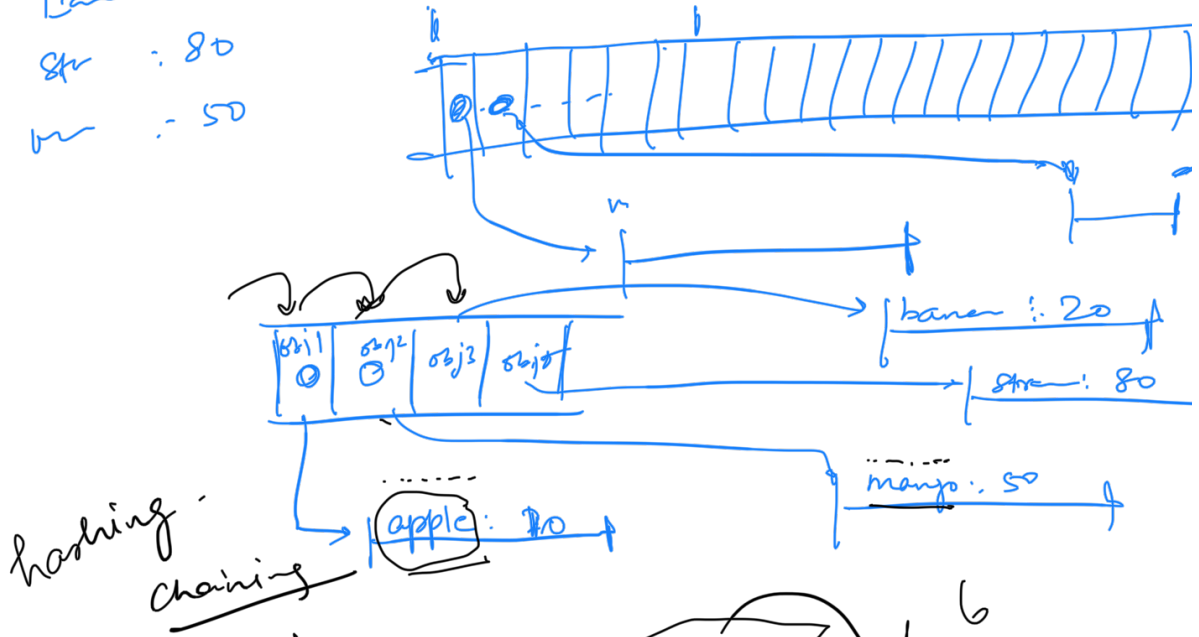
Pear: 30

Apple: 40

$$+5 -10 +50 = 45$$



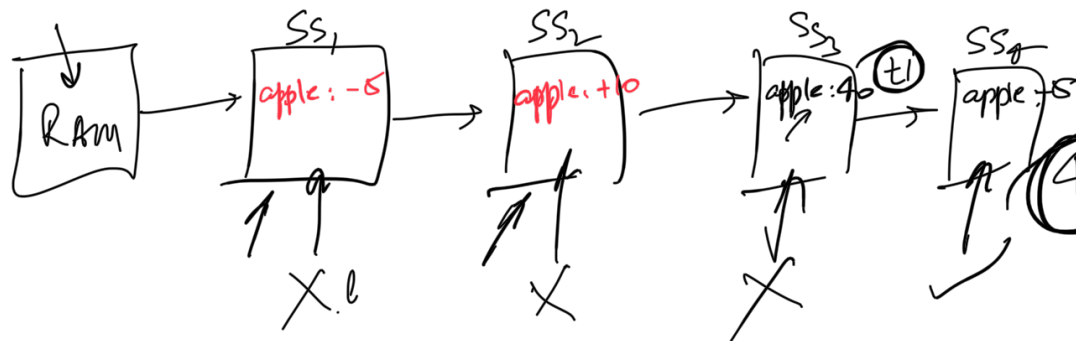
Apple: 10
Ban: 20
Str: 80
Mango: 50



apple

mango

"mango",
"strawberry"

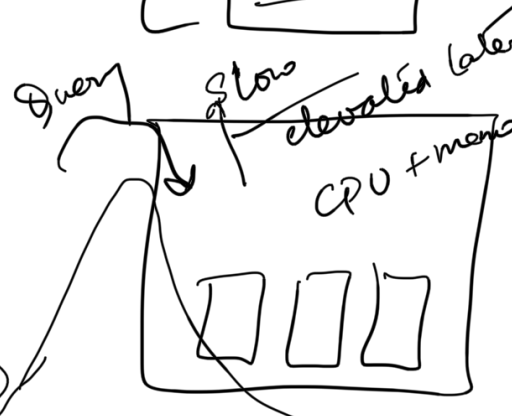
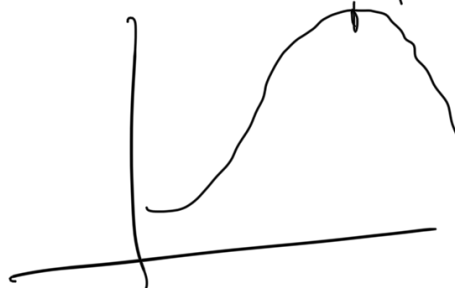


$\log(n) \times \# \text{ of sorted sets}$
2-3

Compaction



User configurable rpm

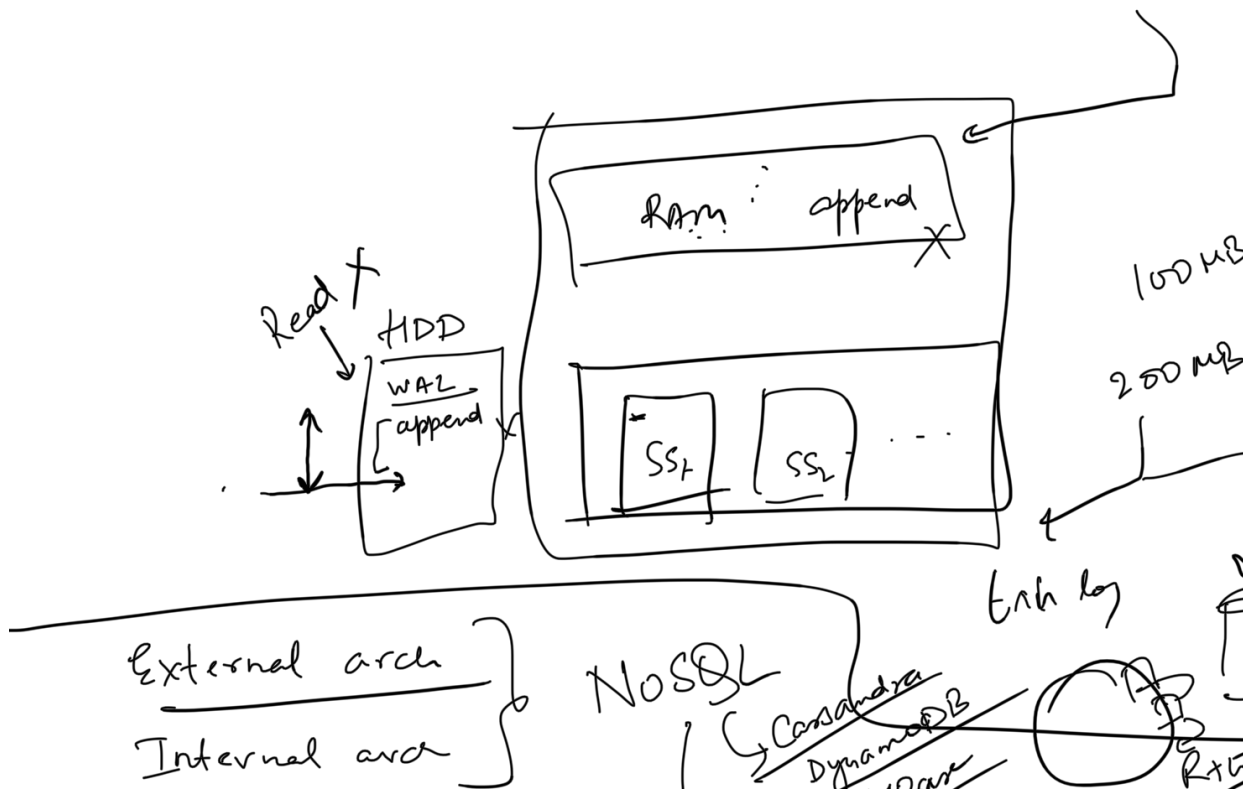


Transaction logs

50	SS1	apple : value = 50	t = 100
10	SS2	apple : value = 40	t = 200
50		apple : value : <u>incr 10</u>	t = 300
45		apple : value : <u>dec 5</u>	t = 400
<div style="border: 1px solid black; padding: 2px;">60</div>		apple : value : <u>60</u>	t = 500

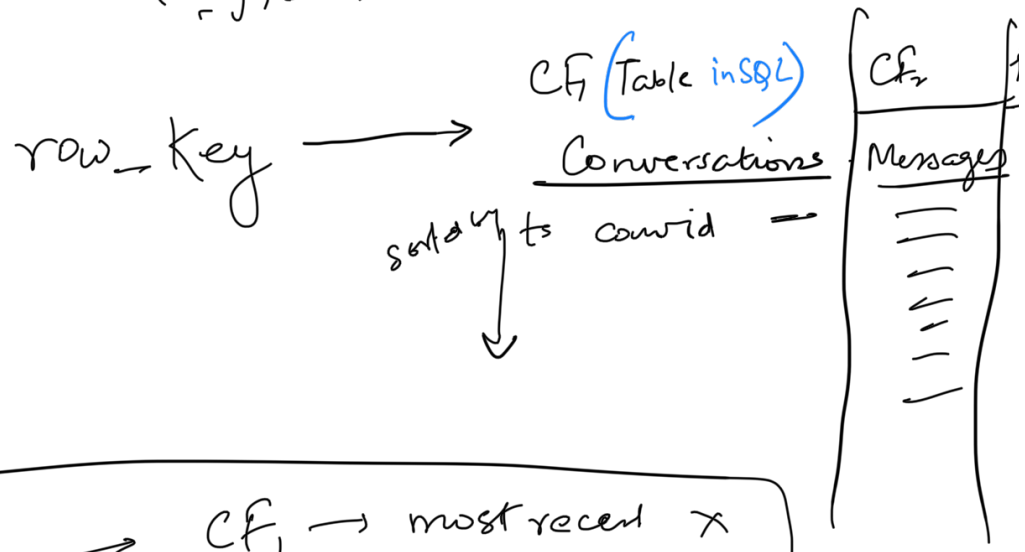
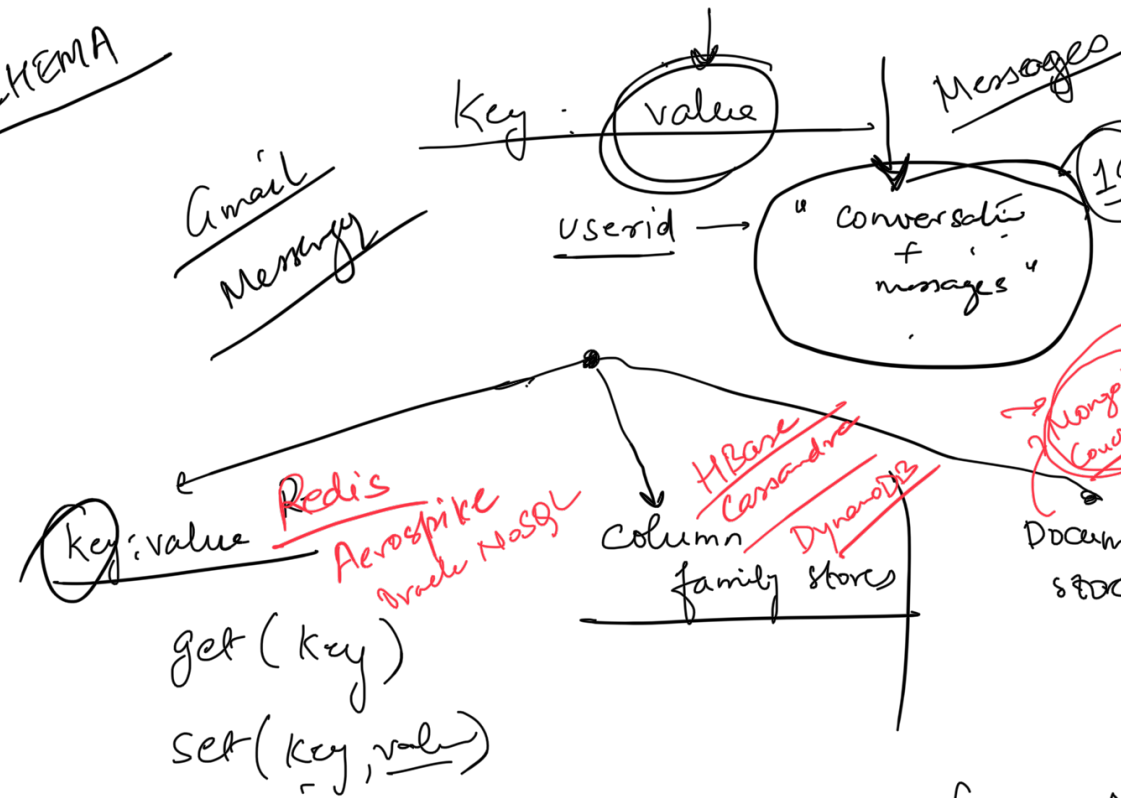
t_x , insert key \rightarrow value

write/update
 $\rightarrow O(1)$
 Read
 $\rightarrow O(\log(N))$
 compression
 tx, update key \rightarrow value 2
 tx, delete key
 update user, set score = score + 30
 where user-id = 10
 $\log(N)$ $\log(N)$



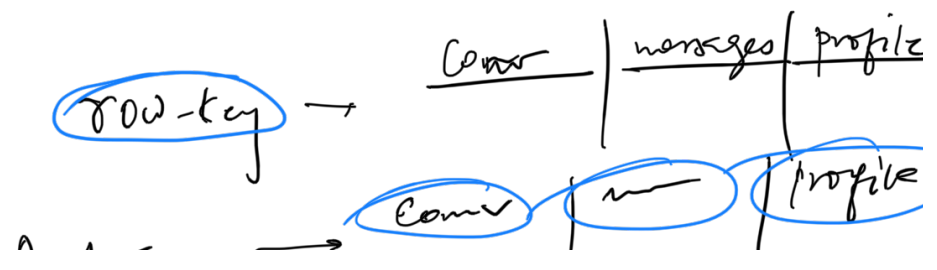
Sharding data / distributed data

SCHEMA



key → CF_i → most recent x entries

get(key, CF, x, offset)



Andhuman →

<u>Canv</u>	<u>rese</u>	<u>propM-</u>

Anghumen - con G
Anghu - versorge

```
{ "number": 20;  
  "class": "big" }
```



SON

Document Based I

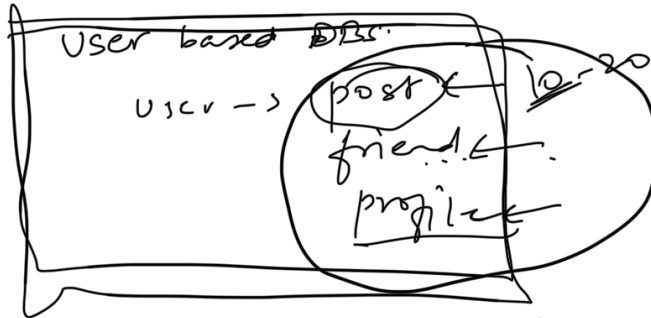
$\{ \text{a_type} : \text{car} \}$

mailbox → CP

pages

pagerank score

FB case



URL → score

(Key) → value

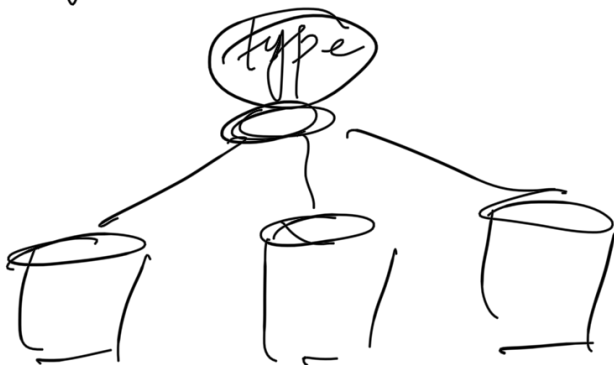
CF → Over
Do

CF store

$R=2$

{ type: car
model: Blue
price: 30L }

{ type: car
no.: 3012
price: 3L
number: 2
size: 6L }



Non-indexed



Google search typeahead

mic
michael j