

# RECURRENCE RELATION PRACTICE

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$$\begin{aligned} \textcircled{1} \quad T_n &= T(n-1) + n \\ T(n-1) &= T(n-2) + n-1 \\ T(n-2) &= T(n-3) + n-2 \end{aligned}$$

$$\begin{aligned} T_n &= T(n-2) + n-1 + n \\ T(n) &= T(n-3) + n-2 + n-1 + n \end{aligned}$$

⋮

$$T(n) = T(n-k) + kn - \frac{k(k-1)}{2} \quad \text{---} \textcircled{1}$$

for base case :-

$$n-k = 1$$

$$k = n-1$$

$$\text{subst in } \textcircled{1} \quad T(1) = 1$$

$$T(n) = T(1) + (n-1)n - \frac{(n-1)(n-2)}{2}$$

$$\therefore \text{Order} \Rightarrow n^2 [O(n^2)]$$

$$\begin{aligned} \textcircled{2} \quad T(n) &= T(n/2) + 1 \\ T(n) &= T(n/2) + 1 \end{aligned}$$

assuming  $n = 2^k$ , i.e.  $k = \log n$

$$T(2^k) = T\left(\frac{2^k}{2}\right) + 1$$

$$T(2^k) = T(2^{k-1}) + 1$$

$$= (T(2^{k-2}) + 1) + 1$$



$$= T(2^{k-2}) + 2$$

$$= [T(2^{k-3}) + 1] + 2$$

$$= T(2^{k-3}) + 3$$

$$T(2^k) = T(2^{k-k}) + k$$

$$= T(2^0) + k = T(1) + k$$

if  $T(1) = 1$  then,

$$T(2^k) = 1 + k$$

$$T(n) = \log n + 1$$

$$\therefore \Theta(\log n)$$

③

$$T(n) = 2T(n/2) + n \quad \text{--- (1)}$$

$$T(n/2) = 2T(n/2^2) + n/2$$

$$= 2[2T(n/2^2) + n/2] + n$$

$$= 2^2 T(n/2^2) + 2n \quad \text{--- (2)}$$

$$T(n/2^2) = 2T(n/2^3) + n/2^2$$

$$T(n/2) = 2^2 [2T(n/2^3) + \frac{n}{2^2}] + 2n$$

$$T(n) = 2^3 T(n/2^3) + 3n \quad \text{--- (3)}$$

Continuing,

$$T(n) = 2^k T\left(\frac{n}{2^k}\right) + kn$$

$$\text{let } \left\{ \begin{array}{l} T\left(\frac{n}{2^k}\right) = T(1) \\ n = 2^k \\ k = \log n \end{array} \right.$$

$$n = 2^k$$

$$k = \log n$$

$$T(n) = 2^k T(1) + kn$$

$$= n(1) + \underline{\underline{n \log n}}$$

$$O(n \log n)$$