

# Post Midsem

## The two point Connectivity function (on $\mathbb{Z}^d$ )

For  $u, v \in \mathbb{Z}^d$ , define,  $T_d(u, v) = P_p(u \leftrightarrow v)$

$$= P_p(0 \leftrightarrow v-u) \quad \boxed{\text{[by Transl invariance]}}$$

Notation: We say  $a_n \approx b_n$  as  $n \rightarrow \infty$  if  $\frac{\log a_n}{\log b_n} \rightarrow 1$   
as  $n \rightarrow \infty$

We have already discussed that

$$C_1(n, d) e^{-n\phi(p)} \leq P_p(0 \leftrightarrow \delta B_n) \leq C_2(n, d) e^{-n\phi(p)}$$

We showed that,

$$P_p(0 \leftrightarrow \delta B_n) \approx e^{-n\phi(p)}$$

Shown using  
methods of  
reliability theory.

Thm: Take  $0 < p \leq 1$  and  $\phi(p)$  as before. Let

$e_n = (n, 0, 0 \dots 0)$ , then we have

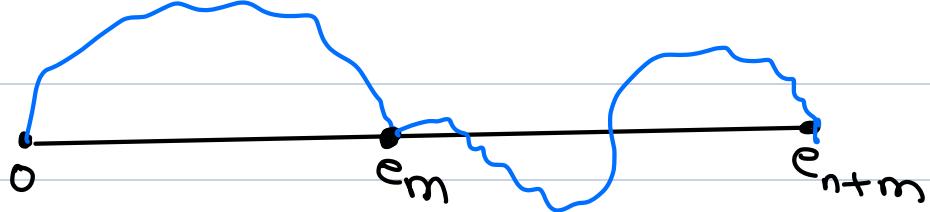
$$\lim_{n \rightarrow \infty} \left\{ -\frac{1}{n} \log (\tau_p(0, e_n)) \right\} = \phi(p)$$

And,  $\exists$  a positive constant  $\xi > 0$  (<sup>not dependent on p</sup>) s.t.

$$\xi p n^{d(1-d)} e^{-n\phi(p)} \leq \tau_p(0, e_n) \leq e^{-n\phi(p)}$$

And hence,  $\tau_p(0, e_n) \approx e^{-n\phi(p)}$

Proof :  $\{0 \longleftrightarrow e_{n+m}\} \supseteq \{0 \longleftrightarrow e_m\} \cap \{e_m \longleftrightarrow e_{m+n}\}$



∴ Taking  $t_k = -\log \tau_p(0, e_k)$ , then we get

$$t_{m+n} \leq t_m + t_n$$

Then by Fekete's Subadditive lemma,

$$\lim_{n \rightarrow \infty} \frac{t_n}{n} = \inf_{n \geq 1} \frac{t_n}{n}$$

Def<sup>n</sup> : Let  $\gamma(p) := \lim_{n \rightarrow \infty} -\frac{1}{n} \log (\tau_p(0, e_n))$

$$= \inf_{n \geq 1} \left\{ -\frac{\log(T_p(o, e_n))}{n} \right\}$$

Note:  $\log(T_p(o, e_n)) \leq -n \varphi(p) \quad \forall n \geq 1$

Now we show,  $\varphi(p) = \phi(p)$

Obviously,  $\{o \leftrightarrow e_n\} \subseteq \{o \leftrightarrow \delta B_n\} \Rightarrow \varphi(p) \geq \phi(p)$ .

$$\left[ \therefore \frac{-\log(P_p(o \leftrightarrow \delta B_n))}{n} \leq \frac{-\log P_p(o \leftrightarrow e_n)}{n} \right]$$

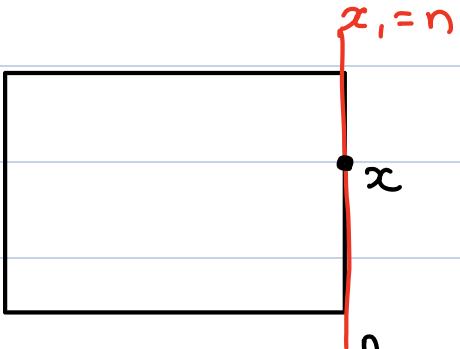
So, we need to show  $\varphi(p) \leq \phi(p)$ .

A simple union bound tells us that,

$$P_p(o \leftrightarrow x) \geq \frac{1}{\# \delta B_n} P_p(o \leftrightarrow \delta B_n) \quad -(1)$$

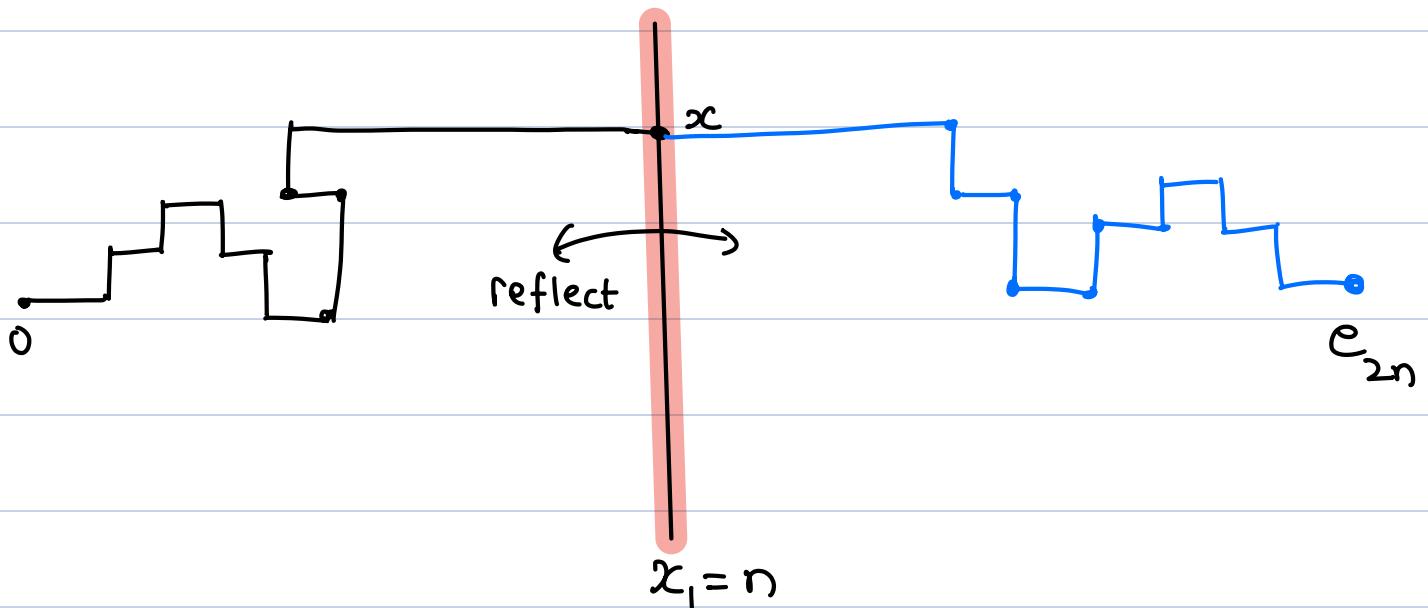
for some  $x \in \delta B_n$

By translation invariance, assume that the first coord of  $x$  is  $x_1 = n$ .



Now consider the reflected path, along the line

$x_1 = n$ . Then,



Clearly,  $\{0 \leftrightarrow x\} \cap \{x \leftrightarrow e_{2n}\} \subseteq \{0 \leftrightarrow e_{2n}\}$

Thus, by FKG

$$\mathbb{P}_p(0 \leftrightarrow x) \mathbb{P}_p(x \leftrightarrow e_{2n}) \leq \mathbb{P}_p(0 \leftrightarrow e_{2n})$$

Since  $\mathbb{P}_p(0 \leftrightarrow x) = \mathbb{P}_p(x \leftrightarrow e_{2n})$  (See the figure above)

$$(T_p(0, x))^2 \leq \mathbb{P}_p(0 \leftrightarrow e_{2n})$$

$$\therefore T_p(0, e_{2n}) \geq \frac{1}{(2d(2n)^{d-1})^2} (\mathbb{P}_p(0 \leftrightarrow \delta B_n))^2$$

(by (1))

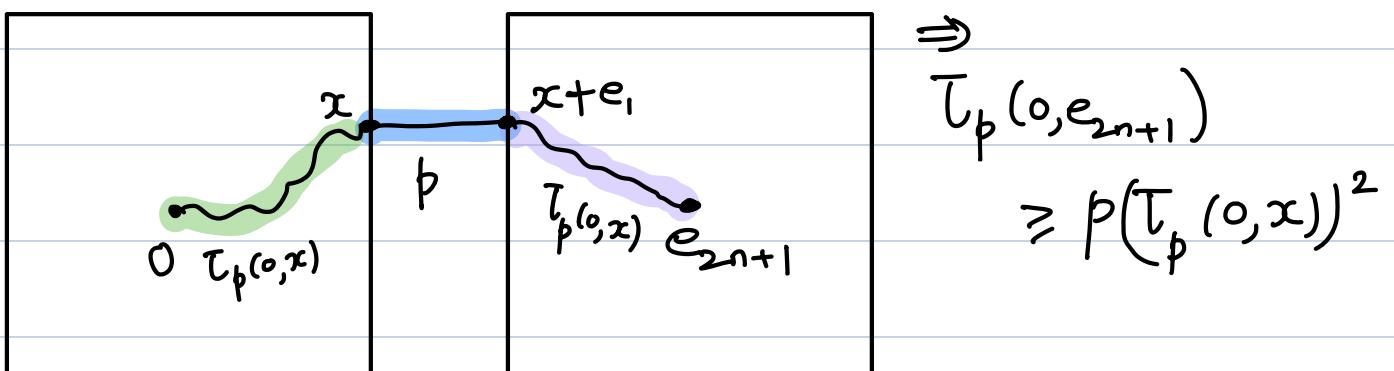
$$T_p(0, e_{2n}) \geq C n^{4(1-d)} \exp(-2n\phi(p))$$

$$\therefore -\frac{1}{2n} T_p(0, e_{2n}) \leq \log(\text{poly}) + \phi(p)$$

Now let's move to the odd case

Firstly,

$$\begin{aligned} \{0 \longleftrightarrow x\} \cap \{x \longleftrightarrow x+e_i\} \cap \{x+e_i \longleftrightarrow e_{2n+1}\} \\ \subseteq \{0 \longleftrightarrow e_{2n+1}\} \end{aligned}$$



And hence by the same logic  $-\frac{1}{2n+1} \log(T_p(0, e_{2n+1})) \leq o(n) + \phi(p)$

And hence  $\lim_{n \rightarrow \infty} -\frac{1}{n} \log(T_p(0, e_n)) \leq \phi(p)$

Thus,  $I(p) = \phi(p)$

This also shows the theorem



General Connectivity in  $\mathbb{Z}^d$ :

We are interested in  $T_p(0, x)$ ,  $x \in \mathbb{Z}^d$ .

Prop<sup>n</sup> (\*): Let  $p \in (0, 1)$  and  $\phi(p)$  as before. Then  $\exists \lambda > 0$  independent of  $p$  s.t.

$$\lambda p^d |x|^{4d(1-d)} e^{-|x|\phi(p)} \leq T_p(0, x) \leq \exp(-\|x\|\phi(p))$$

Where,  $|x| = \sum_{i=1}^d |x_i| = d_{\mathbb{Z}^d}(0, x)$  (graph dist)

$\|x\| = \max \{|x_i| : i = 1, \dots, d\}$  = length of the smallest box centered at 0 containing  $x$

Proof:

## ① Upper bound

Let  $\|x\|=n$ , then by the same argument as above

$$T_p(0, x) \leq \sqrt{T_p(0, e_{2n})} \leq \exp(-n\phi(p))$$

## ② Lower bound

Say in  $\mathbb{Z}^3$ , we wish to have a path from 0 to  $(x_1, -x_2, -x_3)$

This has the same prob as having paths from 0 to  $(x_1, x_2, -x_3)$  or to  $(x_1, x_2, x_3)$

[**reflection principle**]

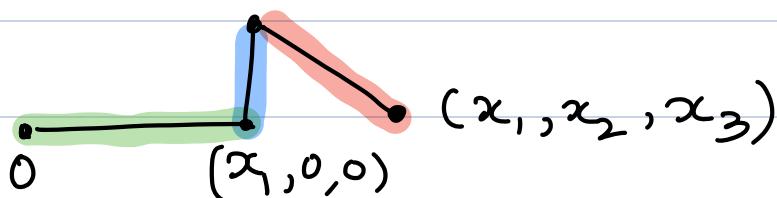
For  $x \in \mathbb{Z}^d$  with  $|x|=n$  and  $\tilde{x} = (|x_1|, \dots, |x_d|)$  then

$|\tilde{x}|=n$ , also from the above observation

$$T_p(0, x) = T_p(0, \tilde{x}). \text{ So WLOG, let } x_i \geq 0 \forall i$$

Let  $x^{(k)} = \{|x_1|, \dots, |x_k|, 0, \dots, 0\}$  for  $1 \leq k \leq d$

Then,  $\{0 \leftrightarrow \tilde{x}\} \supseteq \{0 \leftrightarrow x^{(1)}\} \cap \{x^{(1)} \leftrightarrow x^{(2)}\} \cap \dots \cap \{x^{(d-1)} \cap x^{(d)}\}$



$\therefore$  By FKG inequality,

$$T_p(0, x) \geq \prod_{i=1}^d T_p(x^{(i-1)}, x^{(i)})$$

$$= \prod_{i=1}^d T_p(0, (0, \dots |x_i|, 0 \dots 0))$$

$$= \prod_{i=1}^d T_p(0, e_{|x_i|}) \left( \approx e^{-|x_i| \phi(p)} \right)$$

$$\geq \prod_{i=1}^d c p^{|x_i|^{1-d}} e^{-|x_i| \phi(p)}$$

(By previous result)

$$\therefore T_p(0, x) \geq c p^d (|x_1| \dots |x_d|)^{1-d} e^{-|x| \phi(p)}$$

Since  $1-d < 0$ , by AM-GM inequality,

$$T_p(0, x) \geq \lambda p^d |x|^{4d(1-d)} e^{-|x| \phi(p)}$$



Why "Correlation length"?

At  $p = p_c$ , we know that  $\phi(p) = 0$ . At criticality, the physicists have a guess:

$$\tau_{p_c}(0, e_n) \sim n^{2-d-\gamma} \quad \text{for some } \gamma > 0$$

- At criticality  $\leftarrow$  polynomial decay
- Below Criticality  $\leftarrow$  exponential decay

When can we distinguish between a polynomial and a exponential decay?

Ans : When  $n\phi(p)$  is large

$$\xi(p) = \frac{1}{\phi(p)} \quad \text{— Correlation length}$$

$$P_p(0 \longleftrightarrow \delta B_n) \approx e^{-\frac{n}{\xi(p)}} \quad \curvearrowright \quad \begin{array}{l} \text{The exponential term} \\ \text{becomes significant only} \\ \text{when } n \gg \xi(p) \end{array}$$

$$\tau_p(0 \longleftrightarrow e_n) \approx e^{-\frac{n}{\xi(p)}} \quad \text{— Gives a notion of lengthscale at which}$$

natural length scale

where for distances smaller than  $\xi(p)$ , sites are likely to be correlated, for sites at a dist  $> \xi(p)$  they are effectively uncorrelated.

Recall,

$$\Theta(p) = \mathbb{P}_p(\# C(o) = \infty)$$

$$\chi(p) = \mathbb{E}_p(\# C(o)), \quad p_c := \inf\{p : \chi(p) = \infty\}$$

Exc: Show  $p_c = p_T$  (was open for > 40 years)

Prop<sup>n</sup>: For  $p < p_c$ ,  $T_p(0, x) \leq \left(1 - \frac{1}{\chi(p)}\right)^{|x|}$

$$x \in \mathbb{Z}^d$$

Proof:

Let  $S_n := \{v \in \mathbb{Z}^d : |v| \leq n\}$  — ball in graph metric

$$\partial S_n = \{v \in \mathbb{Z}^d : |v| = n\}$$

Let  $M_n = \#(C \cap \partial S_n) = \text{no. of vertices on the boundary conn to the origin}$

$$\therefore \mathbb{E}_p(M_n) = \mathbb{E}_p\left(\sum_{v \in \partial S_n} \mathbb{P}_{S_0 \leftrightarrow v}\right)$$

$$\sum_{n=0}^{\infty} \mathbb{E}_p(M_n) = \sum_{n=0}^{\infty} \sum_{v \in \partial S_n} T_p(0 \leftrightarrow v)$$

$$= \sum_{w \in \mathbb{Z}^d} T_p(0, w) = \mathbb{E}_p(\# C(o)) = \chi(p)$$

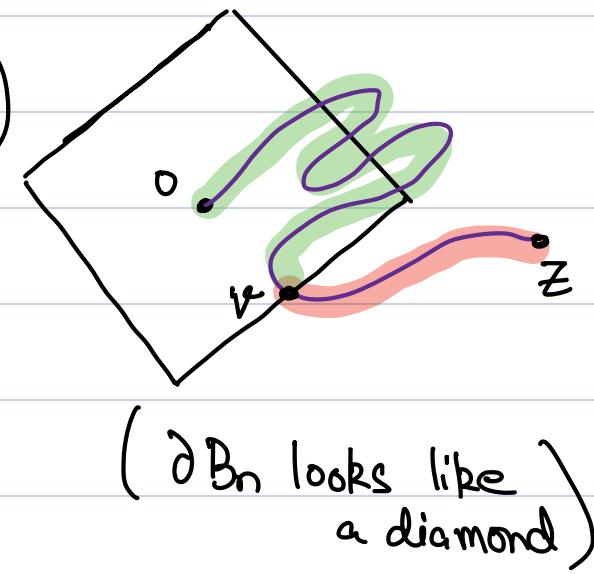
Let  $z \in S_m^c = \{v \in \mathbb{Z}^d : |v| > m\}$ . For a path from  $0 \leftrightarrow z$ , let  $v$  be the last vertex in  $S_m$

$$\therefore T_p(0, z) = P_p \left( \bigcup_{v \in \partial S_m} \{0 \leftrightarrow v\} \cap \{v \leftrightarrow z\} \right)$$

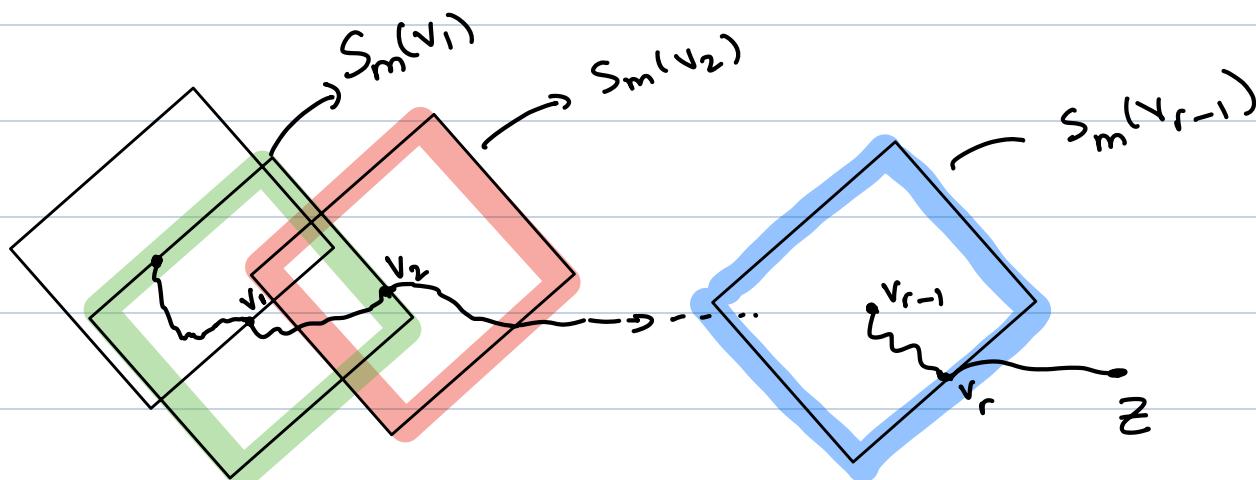
$$\stackrel{\leq}{\substack{(BK + \text{union bd})}} \sum_{v \in \partial S_m} P_p(0 \leftrightarrow v) P_p(v \leftrightarrow z)$$

$$= \sum_{v \in \partial S_m} T_p(0, v) T_p(v, z)$$

$$\leq \sum_{v \in \partial S_m} T_p(0, v) = F_p(M_m) \quad — (2)$$



Suppose  $|z| = \gamma m + \delta$ ,  $0 \leq \delta \leq m$ . Then



$$T_p(0, z) \leq \sum_{\{v_1, \dots, v_r\}} P_p(0 \leftrightarrow v_1) \cdots P_p(v_{r-1} \leftrightarrow v_r) P_p(v_r \leftrightarrow z)$$

( BK multiple times )

$$\leq \sum_{\{v_1 \dots v_r\}} P_p(v_0 \leftrightarrow v_1) \dots P_p(v_{r-1} \leftrightarrow v_r) \\ = (\mathbb{E}_p(M_m))^r - (3)$$

For  $u \in \mathbb{Z}^d$ ,  $R \geq 1$ ,

$$T_p(0, Ru) \geq P_p(ju \leftrightarrow (j+1)u \text{ for } 0 \leq j \leq R-1)$$



$$\geq (T_p(0, u))^R \quad (\text{By FKG and translation inv})$$

$$\begin{aligned} \therefore T_p(0, u) &\leq (T_p(0, Ru))^{\frac{1}{R}} \\ &\leq (\mathbb{E}_p(M_m))^{\left\lfloor \frac{|Ru|}{m} \right\rfloor \frac{1}{R}} \quad (\text{By prev argument}) \\ &\longrightarrow (\mathbb{E}_p(M_m))^{\left\lfloor \frac{|u|}{m} \right\rfloor} - (4) \quad \text{as } R \rightarrow \infty \end{aligned}$$

Finally for  $0 < p < p_c = p_T$ ,  $\chi(p) < \infty$

We'll show  $\exists m \text{ s.t. } \mathbb{E}_p(M_m) \leq \left(1 - \frac{1}{\chi(p)}\right)^m$

Assuming the existence of  $m$  as in (‡), using that  $m$  in (4) we get,

$$T_p(0, u) \leq \left(1 - \frac{1}{\chi(p)}\right)^{|u|} \leq \left(1 - \frac{1}{\chi(p)}\right)^{|u|}$$

(‡) is pretty easy to show since, if  $\forall m \quad \mathbb{F}_p(M_m) > \left(1 - \frac{1}{\chi(p)}\right)^m$ , then using (2)

$$\chi(p) = \sum_{m \geq 1} \mathbb{F}_p(M_m) > \sum_m \left(1 - \frac{1}{\chi(p)}\right)^m = \chi(p)$$

$\Rightarrow \Leftarrow$

So (‡) holds and hence our prob<sup>n</sup>

~~OK~~

So what did we achieve,

$$\phi(p) = \lim_n \left( -\frac{1}{n} \log T_p(0, e_n) \right)$$

(first thing we showed)

$$\begin{aligned} &\geq \lim_n \left( -\frac{1}{n} \log \left[ \left(1 - \frac{1}{\chi(p)}\right)^n \right] \right) \\ &\text{(Prob}^n \text{ above)} \end{aligned}$$

$$\geq \lim_{n \rightarrow \infty} \left( -\frac{1}{n} \log \left( e^{-n/\chi(p)} \right) \right)$$

$$[1-x \geq e^{-x} \forall x]$$

$$= \lim_{n \rightarrow \infty} \left( -\frac{1}{n} \cdot -\frac{n}{\chi(p)} \right) = \frac{1}{\chi(p)}$$

And hence  $\phi(p) \geq \frac{1}{\chi(p)} \Rightarrow \boxed{\xi(p) \geq \chi(p)}$