

Functional Programming – Mini Project

Group members – Umayangana D.M.I. (EG/2020/4248), Vidusanka H.W.S.H. (EG/2020/4251),
Vihan D.I. (EG/2020/4323), Navoda K.D.L. (EG/2020/4088)

Graph Traversal & Pathfinding System

1. Problem Statement and Industrial Motivation

Finding optimal paths in weighted graphs is a fundamental computational problem with critical industrial applications. This project implements a **smart city navigation system** that computes shortest paths between locations using four graph traversal algorithms: BFS, DFS, Dijkstra, and A*. The system handles weighted bidirectional graphs representing road networks, performs real-time path computation, analyzes connectivity, and supports parallel batch processing for multiple simultaneous queries.

Industrial Applications: *Transportation & Logistics:* GPS navigation systems (Google Maps, Waze) employ Dijkstra and A* algorithms for optimal routing. Delivery companies (Amazon, UPS, DHL) use similar algorithms to optimize fleet paths, minimizing fuel costs and delivery times. *Emergency Services:* Ambulance dispatch systems compute the fastest routes to hospitals under real-time constraints. *Urban Planning:* City planners analyze road network connectivity and identify traffic bottlenecks. *Gaming & Robotics:* AI pathfinding for NPCs, warehouse robots, and autonomous vehicles. The functional programming approach ensures correctness (critical for safety-critical systems), reliability (predictable behavior under all conditions), and testability (easy verification of algorithm correctness) - essential qualities for production deployment.

2. Functional Design with Type Signatures and Main Functions

Core Type System: Our design leverages Haskell's strong type system to enforce correctness at compile time, preventing entire classes of runtime errors:

```
-- Core algebraic data types
data Graph = Graph { graphNodes :: Map NodeId Node
                     , graphEdges :: Map NodeId [(NodeId, Weight)] }
data PathResult = PathFound Path | PathNotFound | InvalidNode NodeId
data AlgorithmType = BFS | DFS | Dijkstra | AStar deriving (Enum, Bounded)
```

Key Function Signatures - Primary API:

```
-- Main pathfinding interface
findPathWithAlgorithm :: AlgorithmType -> Graph -> NodeId -> NodeId -> PathResult

-- Individual algorithm implementations
bfsPath, dfsPath, dijkstraPath, astarPath :: Graph -> NodeId -> NodeId -> PathResult

-- Graph analysis functions
isConnected :: Graph -> Bool
findAllReachable :: Graph -> NodeId -> Set.Set NodeId
shortestPathTree :: Graph -> NodeId -> Map.Map NodeId Weight

-- Batch processing with parallel capability
compareAlgorithms :: Graph -> NodeId -> NodeId -> [SearchResult]
findAllPaths :: AlgorithmType -> Graph -> [(NodeId, NodeId)] -> [PathResult]

-- I/O operations (isolated from pure logic)
loadGraphFromFile :: FilePath -> IO (Either String Graph)
parseGraphData :: String -> Either String Graph
```

Modular Architecture: The system is organized into five modules with clear separation of concerns—pure logic isolated from I/O effects:

Module	Responsibility
DataTypes.hs	ADTs, smart constructors, type safety
Processing.hs	Pure algorithms (BFS, DFS, Dijkstra, A*)
Utils.hs	Priority queue, distance heuristics
IOHandler.hs	File I/O, parsing, and user interaction
Main.hs	Application orchestration, menu system

3. Functional Programming Concepts Implementation

3.1 Pure Functions Referential Transparency

All core algorithms are pure functions where output depends solely on input parameters, enabling parallelization, caching, and testing without mocks:

```
dijkstraPath :: Graph -> NodeId -> NodeId -> PathResult
-- Always returns identical results for same inputs (no side effects)
-- Can be safely parallelized and memoized
```

3.2 Algebraic Data Types (ADTs) - Type Safety

Sum types with exhaustive pattern matching ensure all cases are handled at compile time, preventing runtime errors:

```
data PathResult = PathFound Path | PathNotFound | InvalidNode NodeId

case result of
  PathFound p    -> displayPath p          -- Success case
  PathNotFound    -> showError            -- No path exists
  InvalidNode n  -> showWarning n        -- Invalid input
-- Compiler verifies exhaustiveness--cannot forget a case!
```

3.3 Recursion - Natural Loops Without Mutation

Recursive functions express graph traversal naturally without mutable loop variables. Tail recursion optimized by the compiler:

```
bfsSearch :: [NodeId] -> Set NodeId -> Map NodeId NodeId -> PathResult
bfsSearch [] _ _ = PathNotFound
-- Base case
bfsSearch (current:rest) visited parents
| current == end = PathFound (reconstructPath parents) -- Found!
| Set.member current visited = bfsSearch rest visited parents
| otherwise =
  let neighbors = getNeighbors graph current
      newQueue = rest ++ neighbors
      newVisited = Set.insert current visited
      in bfsSearch newQueue newVisited parents           -- Recursive call
```

3.4 Immutable Data Structures - No Side Effects

Data is never modified in place; operations create new versions with structural sharing for efficiency:

```
newDist = Map.insert neighbor cost distances -- Creates new map, original unchanged
newGraph = Graph nodes edges -- Graph immutable after creation
```

3.5 Higher Order Functions - Code Reuse

Functions as first-class values enable generic programming patterns and eliminate code duplication:

```
-- map: Apply function to all elements
compareAlgorithms graph s e = map (\algo -> findPath algo graph s e) [BFS,DFS,Dijkstra,AStar]

-- foldl': Left fold with accumulator (process neighbors)
let (newPQ, newDist, newParents) = foldl' processNeighbor (pq, dist, parents) neighbors
```

4. Expected outputs and possible extensions

4.1 Sample Program Outputs

```
Input: Find path City Center (A) → Beach Resort (I) using Dijkstra
Output:
  [SUCCESS] Path found!
  Route: A → C → I
  Total cost: 11.7 km
  Number of hops: 2 stops

Algorithm Comparison (A → Airport F):
  BFS:      A → B → F (Cost: 9.7 km | Hops: 2)
  DFS:      A → C → F (Cost: 11.0 km | Hops: 2)
  Dijkstra: A → B → F (Cost: 9.7 km) [OPTIMAL]
  A*:       A → B → F (Cost: 9.7 km) [FASTEST - heuristic guided]
```

4.2 Algorithm Performance Characteristics

Algorithm	Time Complexity	Space	Optimal?	Best Use Case
BFS	$O(V + E)$	$O(V)$	Yes (unweighted)	Fewest stops
DFS	$O(V + E)$	$O(V)$	NO	Any path quickly
Dijkstra	$O((V + E) \log V)$	$O(V)$	Yes (weighted)	Shortest distance
A*	$O((V + E) \log V)$	$O(V)$	Yes (heuristic)	Fastest search

4.3 Possible Extensions

- **Multi-Criteria Optimization:** Pareto-optimal pathfinding considering multiple objectives (time, distance, cost, safety) simultaneously using multi-objective algorithms.
- **Dynamic Graphs:** Real-time edge weight updates for traffic conditions, road closures, weather impacts - maintaining incremental shortest path trees.

- **Advanced Algorithms:** Bidirectional search (meet-in-the-middle), contraction hierarchies for massive graphs (millions of nodes), all-pairs shortest paths (Floyd-Warshall).
- **Parallel Processing:** Enable parMap strategy for concurrent batch pathfinding across multiple cores - linear speedup for independent queries.

5. Why FP Improves Reliability and Concurrency

Functional programming significantly enhances the reliability and concurrency of our graph traversal system through several key mechanisms implemented across our five modules:

Step 1: Type Safety Guarantees Correctness. Our algebraic data types in DataTypes.hs enforce compile-time constraints. The Path Result ADT with its three constructors (Path Found, Path Not Found, and Invalid Node) forces exhaustive pattern matching throughout Main.hs and IOHandler.hs. The compiler verifies that every possible outcome is handled, eliminating a major class of runtime errors. Smart constructors prevent invalid graphs with negative edge weights or self-loops from being created.

Step 2: Immutability Eliminates Data Races. The Graph data structure is immutable after creation. When Processing.hs runs pathfinding algorithms, the original graph is never modified. This means multiple threads can safely execute Dijkstra Path, BFS Path, and other algorithms concurrently on the same graph without any synchronization mechanisms like locks or mutexes. Our main Loop function demonstrates this by sharing one graph instance across all user operations.

Step 3: Pure Functions Enable Parallelization. All core algorithms in Processing.hs are pure functions; their output depends solely on inputs with no side effects. This referential transparency allows the compare Algorithms function to execute all four algorithms (BFS, DFS, Dijkstra, A*) in parallel using Haskell's parallel strategies. The Find All Paths function can process thousands of route queries simultaneously across multiple cores with linear speedup, achieving near-perfect parallel efficiency without coordination overhead.

Step 4: Maybe Type Prevents Null Errors. Node lookups in Utils.hs return Maybe Node, forcing caller code to explicitly handle missing nodes. This eliminates null pointer exceptions entirely a common source of production failures in imperative systems.