

Heaps

Bottom elements can be in any order
 One deleted top; heap reheapifies
 and max/min element appears on top.

Identification

k

largest / greatest
 smallest / closest

If comes with k
 build min Heap.

If mentioned
 guaranteed to
 be solved using
 heap.

Time Complexity - $O(n \log k)$

If comes with k
 build max Heap.

STL

priority - queue <int> - max Heap.

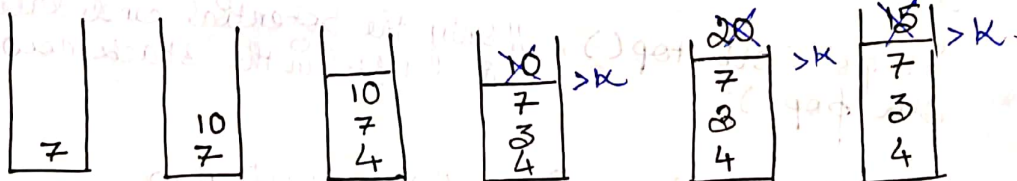
priority - queue <int, vector<int>, greater<int>>
 L min Heap

#1) Find k^{th} Smallest Element.

Smallest (Build Max Heap).

I/P = 7 10 4 3 20 15
 $k = 3$
 O/P = 7

Heap (Max)



Code

```
int kthSmallest (int* arr, int l, int r, int k) {
    priority_queue<int> maxH;
    for (int i = l; i <= r; i++) {
        maxH.push(arr[i]);
        if (maxH.size() > k)
            maxH.pop();
    }
    return maxH.top();
}
```

#2) Print the k largest elements — x —

Code

```
vector<int> kLargest(vector<int> arr, int n, int k) {
    priority_queue<int, vector<int>, greater<int>> minH;
    vector<int> res;
    for (int i = 0; i < n; i++) {
        minH.push(arr[i]);
        if (minH.size() > k)
            minH.pop();
    }
    while (!minH.empty()) {
        res.push_back(minH.top());
        minH.pop();
    }
    reverse(minH.begin(), minH.end());
    return res;
}
```

// On min Heap min on top.
Reversing to get greatest numbers first.

#3)

k Closest Numbers

Return k closest integers to x in the given array

arr = {1, 2, 3, 4, 6}

O/P = 1, 2, 3, 4.

k = 4; x = 3.

Code

```
vector<int> kClosest(vector<int> &arr, int k, int x) {
    priority_queue<pair<int, int>> maxH;
    for (int i = 0; i < arr.size(); i++) {
        maxH.push({abs(arr[i] - x), arr[i]});
        if (maxH.size() > k)
            maxH.pop();
    }
    while (!maxH.empty()) {
        res.push_back(maxH.top().second);
        maxH.pop();
    }
    sort(res.begin(), res.end());
    return res;
}
```


#4) Sort a k sorted array

Given an array of integers of size n ; where each element is at most k away from its target position. Print sorted array.

I/P -

$N=6$; $k=3$.

arr: 2 6 3 12 56 8



at most k .

O/P: 2 3 6 8 12 56

- ① Building minH; as sort in ascending order.
- ② Pop and push to O/P as soon as size $> k$.

Code

```
vector<int> nearlySorted (int *arr, int n, int k) {  
    priority_queue<int, vector<int>, greater<int>> >  
    minH;  
    vector<int> res;
```

```
    for (int i=0; i<n; i++) {  
        minH.push(arr[i]);
```

```
        if (minH.size() > k) {  
            res.push_back(minH.top());  
            minH.pop();  
        }
```

```
    }  
    while (!minH.empty()) {  
        res.push_back(minH.top());  
        minH.pop();  
    }
```

```
    return res;
```

```
}
```

#5) Top k frequent elements
 Given an integer array nums and an integer k;
 return the k most frequent elements (in any order).

I/P - nums : [1, 1, 1, 2, 2, 3] , k = 2.
 O/P - [1, 2] ↑ ↑
 times times

longest (20 min)

```

code
vector<int> topKFrequent(vector<int> &nums, int k) {
    unordered_map<int, int> hash;
    int size = nums.size();

    for (int i = 0; i < size; i++) {
        if (hash.find(nums[i]) == hash.end()) // Not found.
            hash[nums[i]] = 1;
        else
            hash[nums[i]]++;
    }

    vector<int> res;
    priority_queue<pair<int, int>, vector<pair<int, int>, greater<pair<int, int>>> minH;

    for (auto it : hash) {
        minH.push({it.second, it.first});
        if (minH.size() > k)
            minH.pop();
    }

    while (!minH.empty()) {
        res.push_back(minH.top().second);
        minH.pop();
    }
    reverse(res.begin(), res.end());

    return res;
}
    
```

#6) Frequency Sort Characters

I/P: s = 'tree'

O/P: "eetrt" / "etetr"

Code string frequency sort (string s) {

unordered_map<char, int> hash;

for (int i = 0; i < s.length(); i++) {

if (hash.find(s[i]) == hash.end())

hash[s[i]] = 1;

else

hash[s[i]]++;

}

priority_queue<pair<int, char>> maxH;

for (auto it : hash)

maxH.push({it.second, it.first});

string res;

while (!maxH.empty()) {

for (int i = 0; i < maxH.top().first; i++)

res += (maxH.top().second);

maxH.pop();

// push number of times it is present

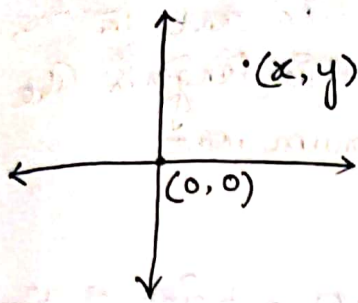
}

return res;

}

Using compare function coded in C++

#7) k Closest Points to the Origin



$$\text{distance} = \sqrt{x^2 + y^2}$$

Here we will maximize $(x^2 + y^2)$

Closest (MaxHeap).

I/P: points $[[1,3], [-2,2]]$; $k=1$

O/P: $[-2,2]$.

distance, (x,y)

Code

```
vector<vector<int>> kClosest(vector<vector<int>>& pts,
                             int k)
```

```
{
    priority_queue<pair<int, pair<int, int>>> maxH
    int size = pts.size(); // no. of column = 2.
```

```
    for(int i=0; i<size; i++)
```

```
    {
        maxH.push({pts[i][0]*pts[i][0] + pts[i][1]*pts[i][1],
                    {pts[i][0], pts[i][1]}});
```

```
        if (maxH.size() > k)
            maxH.pop();
    }
```

```
    vector<vector<int>> res;
```

```
    while(!maxH.empty())
```

```
    {
        vector<int> pt;
```

```
        pt.push_back(maxH.top().first);
```

```
        pt.push_back(maxH.top().second);
```

```
        res.push_back(pt);
```

```
        maxH.pop();
    }
```

```
    return res;
```

2

#8) Connect "n" ropes to minimise cost.

There are given N ropes of different lengths, we need to connect these ropes into one rope. The cost to connect two ropes is equal to sum of their lengths. The task is to connect ropes with minimum cost.

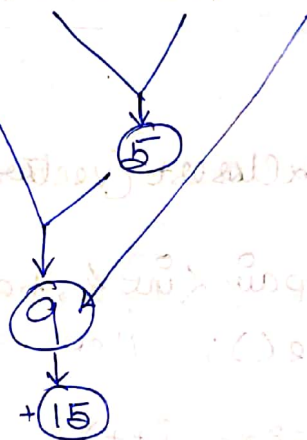
— x —

n = 4

arr = 4, 3, 2, 6.

** At a point of time, select the 2 minimum from the heap-min

4, 3, 2, 6.



$$5 + 9 + 15 = 29$$

Code

```

long long minCost(long long* arr, int n) {
    priority_queue<long long, vector<long long>,
        greater<long long>> minH;
    for (int i = 0; i < n; i++)
        minH.push(arr[i]);
    long long cost = 0;
    while (minH.size() >= 2) {
        long long a = minH.top(); minH.pop();
        long long b = minH.top(); minH.pop();
        cost = cost + a + b;
        minH.push(a + b);
    }
    return cost;
}
    
```


#) a) Reorganize Strings (leetcode) AMAZON
 Rearrange the characters in a string such that
 no 2 adjacent elements are same

I/P: "aab"
 O/P: "aba"

Idea is to add the most frequently
 occurring string; followed by the 2nd most
 occurring string.

string reorganizeString (string s)

unordered_map<char, int> hashmap;

for (int i = 0; i < s.length(); i++)

if (hashmap.find(s[i]) == hashmap.end())

hashmap[s[i]] = 1;

else

hashmap[s[i]]++;

}

priority_queue<pair<int, int>> maxH;

for (auto it : hashmap)

maxH.push({it.second, it.first});

string res = "";

while (!maxH.empty())

pair<int, int> a = maxH.top(); ~~second;~~
 maxH.pop();

pair<int, int> b = maxH.top(); ~~maxH.pop();~~
 maxH.pop();

res += a.second;

res += b.second;

int freq-most = ~~max~~ a.first - 1;

int freq-less = b.first - 1;

if (freq-most > 0) maxH.push({freq-most, a.second});

if (freq-less > 0) maxH.push({freq-less, b.second});

if (!maxH.empty()) // 1 element left

pair<int, int> top = maxH.top(); ~~maxH.pop();~~

~~res~~

maxH.pop();

if (top.first > 1) // If the last element has a frequency more than 1 then the result can't be achieved

else

res += top.second; (It has to be together)

}

return res;

}