

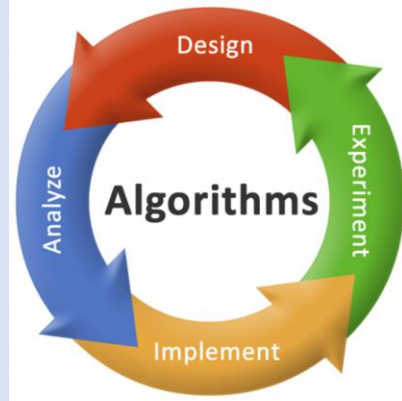
Dynamic Programming

Introduction

COP 3503

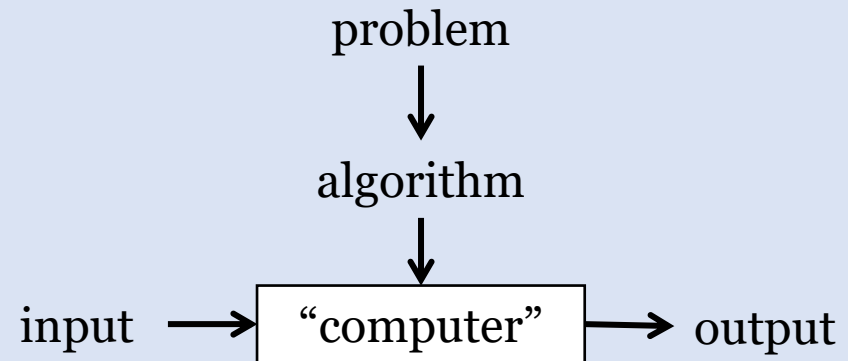
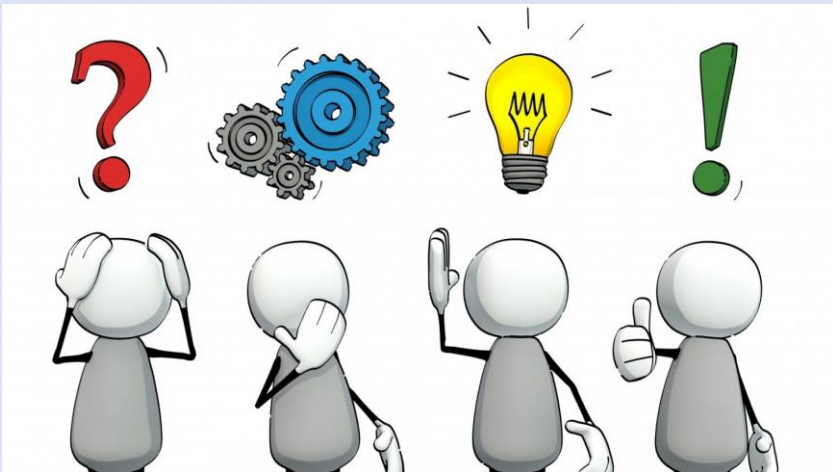
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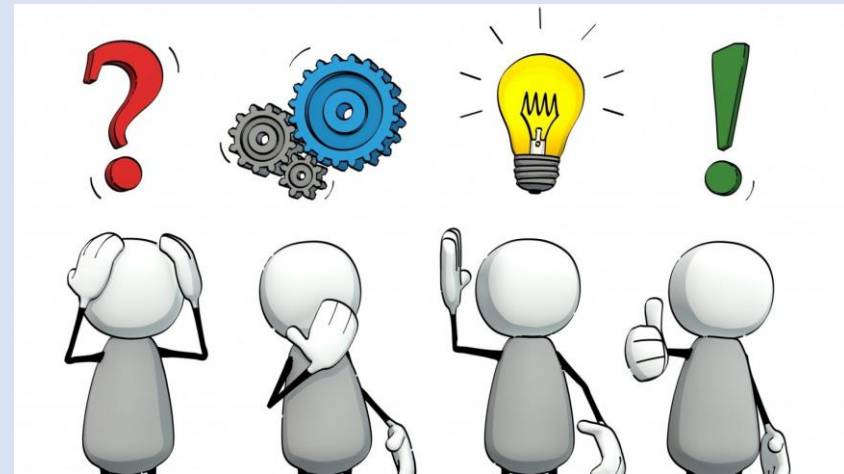
What is an Algorithm? (review)

- A well-defined computational procedure which takes a value (or even set of values) as input and produces a value (or set of values) as output.
- An algorithm is said to be **correct** if, for every input instance, it halts with the correct output.



The Output Produced by our Algorithms

- Something to consider with our problems we are solving as programmers and computer scientists.
- Does there exist a group of solutions to a problem?



Greedy Algorithms (Review)

- Our objective is to produce the best output to a solution.
- Greedy algorithms incorporate the concept of making the best the decision at the current moment (without looking at the big picture overall).
- Greedy algorithms make a greedy choice
 - This results in looking at only one subproblem.
- Does a greedy algorithm produce the optimal solution always?
 - **NO!**

Dynamic Programming Introduction

- Dynamic Programming is a technique. Not an algorithm.
 - Like Divide and Conquer and Backtracking
- Dynamic Programming is applied to optimization problems.
 - Finding the Maximum
 - Finding the Minimum
- Dynamic Programming is applicable when the subproblems are not independent. The subproblems share subsubproblems.
 - Dynamic Programming solves the subproblem and stores the result to be used later.
- This allows for optimal solutions always.

Fibonacci Series

- The series
 - 0, 1, 1, 2, 3, 5, 8, 13, 21, 34, ...
- Many Approaches to solving this problem.
- Recursive
- Dynamic Programming

Recursive Approach

Algorithm 1 Fibonacci (n)

```
1: if  $n \leq 1$  then  
2:   return  $n$   
3: end if  
4: return Fibonacci ( $n - 1$ ) + Fibonacci ( $n - 2$ )
```

Running Time Complexity: $O(2^n)$

Memoization

- Relieve the potential inefficiency of recursion by using the basic idea of dynamic programming.
- With Fibonacci, we were using recursion without storing the result.
 - This can be very bad in terms of large recursion calls.
- The idea is we can store a previous result that will be used in a later subproblem. (Dynamic Programming)
- Memoization can help reduce running time complexity

DP Approach with Memoization

Algorithm 1 memoized_fibonacci (n)

```
1: for  $i = 1$  to  $n$  do  
2:    $\text{results}[i] = -1$   
3: end for  
4: return memoized_fibonacci_rekurs(results,  $n$ )
```

Algorithm 2 memoized_fibonacci_rekurs (results, n)

```
1: if  $\text{results}[n] \neq -1$  then  
2:   return  $\text{results}[n]$   
3: end if  
4: if  $n == 1$  then  
5:    $\text{val} = 1$   
6: else if  $n == 2$  then  
7:    $\text{val} = 1$   
8: else  
9:    $\text{val} = \text{memoized\_fibonacci\_rekurs}(\text{results}, n - 2)$   
10:   $\text{val} = \text{memoized\_fibonacci\_rekurs}(\text{results}, n - 1)$   
11: end if  
12:  $\text{results}[n] = \text{val}$   
13: return  $\text{val}$ 
```

Remember making change?



Remember the Greedy Algorithm

MakeChangeGreedy(n)

$$q = \left\lfloor \frac{n}{25} \right\rfloor$$

$$n_q = n \bmod 25$$

$$d = \left\lfloor \frac{n_q}{10} \right\rfloor$$

$$n_d = n_q \bmod 10$$

$$k = \left\lfloor \frac{n_d}{5} \right\rfloor$$

$$n_k = n_d \bmod 5$$

$$p = n_k$$

Let's Derive the Dynamic
Programming Solution

Dynamic Programming Problems We Will Observe

- 0-1 Knapsack
- Longest Common Subsequence (LCS)
- Sequence Alignment
- Matrix Chain Multiplication