

## Compiler Mid-Term Exam

1. What are the phases of a compiler?
2. Illustrate using the figure the difference between compiler and interpreter?
3. Consider the context-free grammar

$$S \rightarrow S S + \mid S S * \mid a$$

- a) Show how the **aa+a\*** string can be generated by this grammar (derivation).
- b) Construct a parse tree for this string.

4. Eliminate the left recursion of the following grammar:

$$S \rightarrow STS \mid ST \mid T$$

$$T \rightarrow Ta \mid Tb \mid U$$

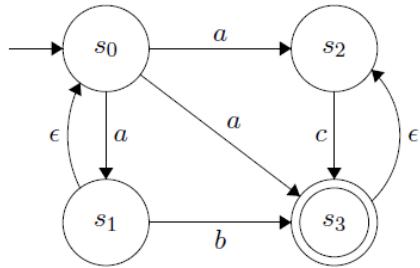
$$U \rightarrow T \mid c$$

Answer:

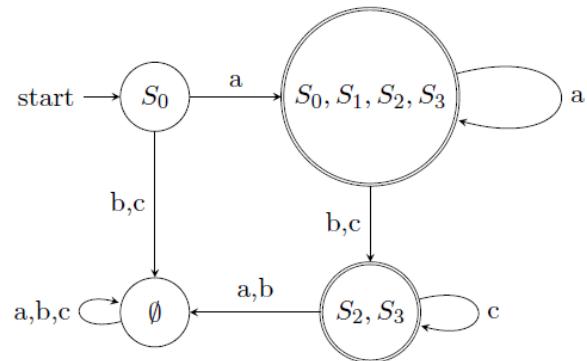
$$\begin{aligned} S &\rightarrow TS' \\ S' &\rightarrow TSS' \mid TS' \mid \epsilon \\ T &\rightarrow cT' \\ T' &\rightarrow aT' \mid bT' \mid \epsilon \end{aligned}$$

5. Convert the following NFA with  $\epsilon$ -transitions into equivalent DFA. **Note that** a DFA must have a transition defined for every state and symbol pair.

(a) Original NFA,  $\Sigma = \{a, b, c\}$ :



DFA:



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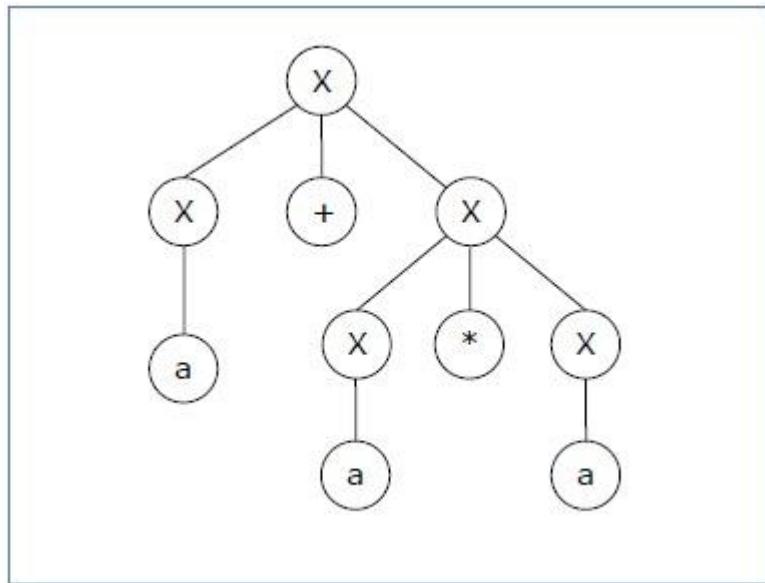
1. What is the meaning of grouping compiler phases into passes?
2. What are types of parsing?
3. Check whether the grammar G with production rules:  $X \rightarrow X+X \mid X^*X \mid X \mid a$  is ambiguous or not?

## Solution

Let's find out the derivation tree for the string "a+a\*a". It has two leftmost derivations.

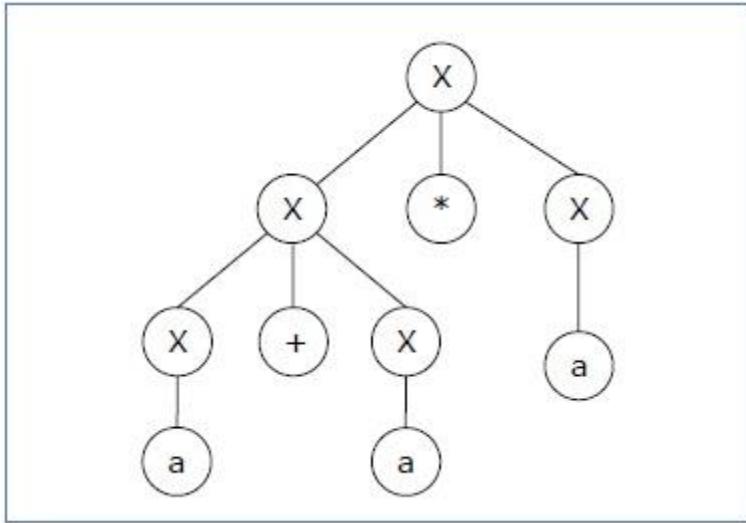
**Derivation 1** –  $X \rightarrow X+X \rightarrow a +X \rightarrow a+ X^*X \rightarrow a+a^*X \rightarrow a+a^*a$

**Parse tree 1** –



**Derivation 2** –  $X \rightarrow X^*X \rightarrow X+X^*X \rightarrow a+ X^*X \rightarrow a+a^*X \rightarrow a+a^*a$

**Parse tree 2** –



Since there are two parse trees for a single string "a+a\*a", the grammar **G** is ambiguous.

#### 4. Eliminate the left recursion of the following grammar:

$$E \rightarrow E+T \mid T$$

$$T \rightarrow T^*F \mid F$$

$$F \rightarrow (E) \mid id$$

Answer:

$$E \rightarrow TE'$$

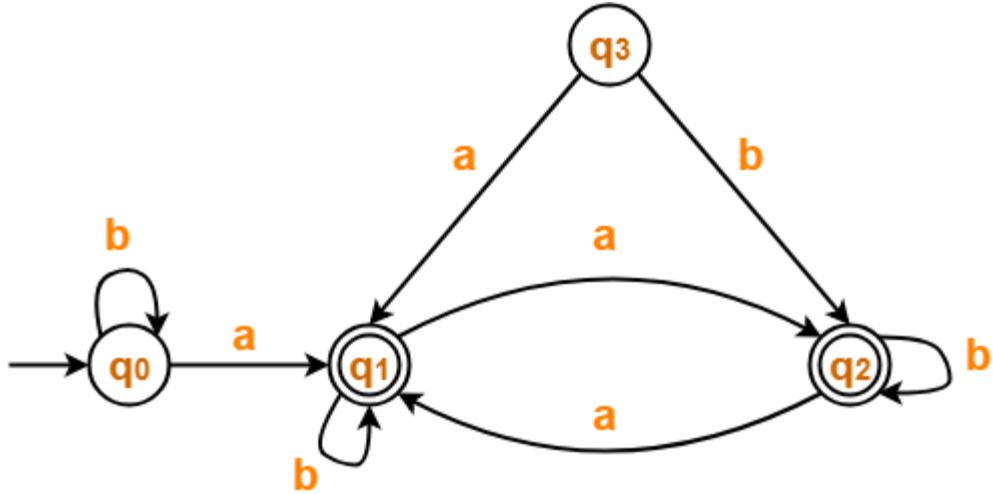
$$E' \rightarrow +TE' \mid \epsilon$$

Then eliminate for T as:

$$T \rightarrow FT'$$

$$T' \rightarrow *FT' \mid \epsilon$$

#### 5. Minimize the following DFA:



**Answer:**

Now using Equivalence Theorem, we have-

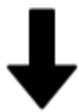
$$P_0 = \{ q_0 \} \{ q_1, q_2 \}$$

$$P_1 = \{ q_0 \} \{ q_1 \}$$

Since  $P_1 = P_0$ , so we stop.

From  $P_1$ , we infer that states  $q_1$  and  $q_2$  are equivalent and can be merged together.

So, Our minimal DFA is-



Minimal DFA