

Compiler Mid-Term Exam

1. What are the phases of a compiler?
2. Illustrate using the figure the difference between compiler and interpreter?
3. Consider the context-free grammar

$$S \rightarrow S S + \mid S S * \mid a$$

- a) Show how the $aa+a*$ string can be generated by this grammar (derivation).
- b) Construct a parse tree for this string.

4. Eliminate the left recursion of the following grammar:

$$S \rightarrow STS \mid ST \mid T$$

$$T \rightarrow Ta \mid Tb \mid U$$

$$U \rightarrow T \mid c$$

Answer:

$$S \rightarrow TS'$$

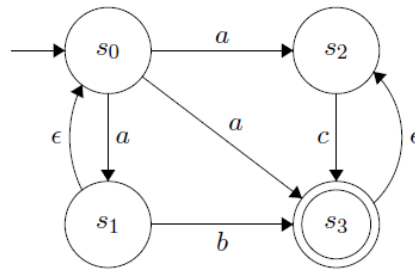
$$S' \rightarrow TSS' \mid TS' \mid \epsilon$$

$$T \rightarrow cT'$$

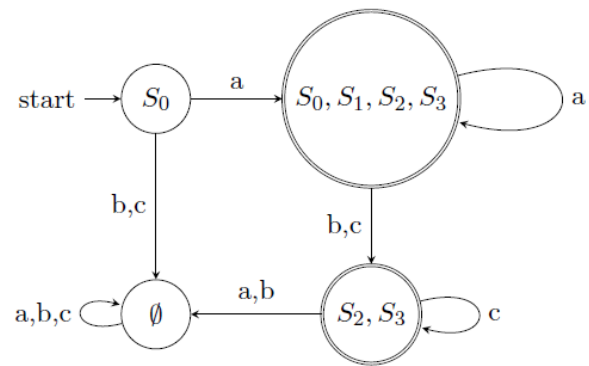
$$T' \rightarrow aT' \mid bT' \mid \epsilon$$

5. Convert the following NFA with ϵ -transitions into equivalent DFA. **Note that** a DFA must have a transition defined for every state and symbol pair.

(a) Original NFA, $\Sigma = \{a, b, c\}$:



DFA:



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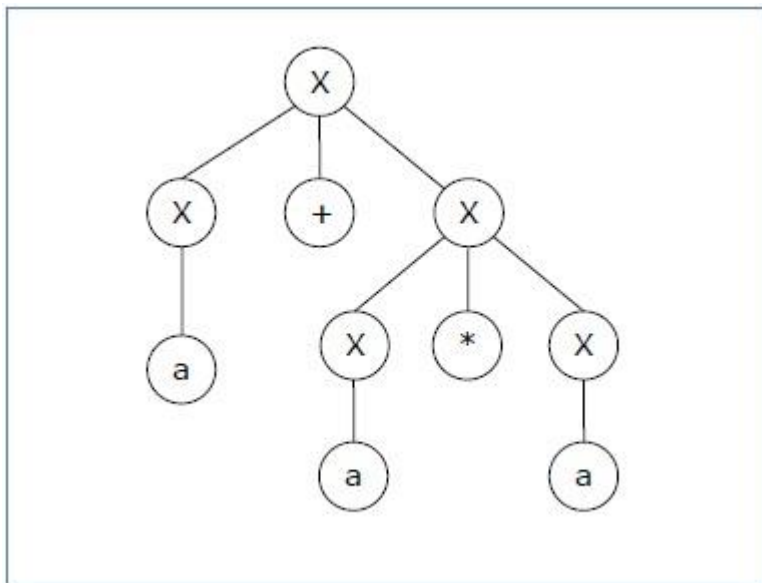
1. What is the meaning of grouping compiler phases into passes?
2. What are types of parsing?
3. Check whether the grammar G with production rules: $X \rightarrow X+X \mid X^*X \mid X \mid a$ is ambiguous or not?

Solution

Let's find out the derivation tree for the string "a+a*a". It has two leftmost derivations.

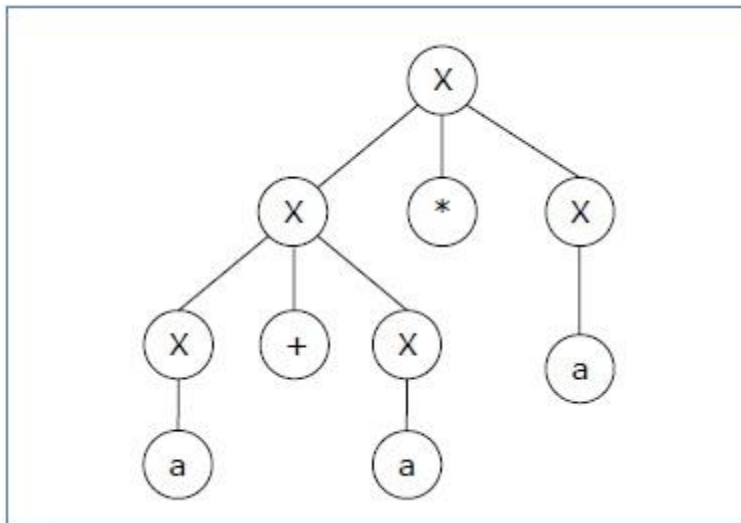
Derivation 1 – $X \rightarrow X+X \rightarrow a+X \rightarrow a+X^*X \rightarrow a+a^*X \rightarrow a+a^*a$

Parse tree 1 –



Derivation 2 – $X \rightarrow X^*X \rightarrow X+X^*X \rightarrow a+X^*X \rightarrow a+a^*X \rightarrow a+a^*a$

Parse tree 2 –



Since there are two parse trees for a single string "a+a*a", the grammar **G** is ambiguous.

4. Eliminate the left recursion of the following grammar:

$E \rightarrow E+T \mid T$

$T \rightarrow T*F \mid F$

$F \rightarrow (E) \mid id$

Answer:

$E \rightarrow TE'$

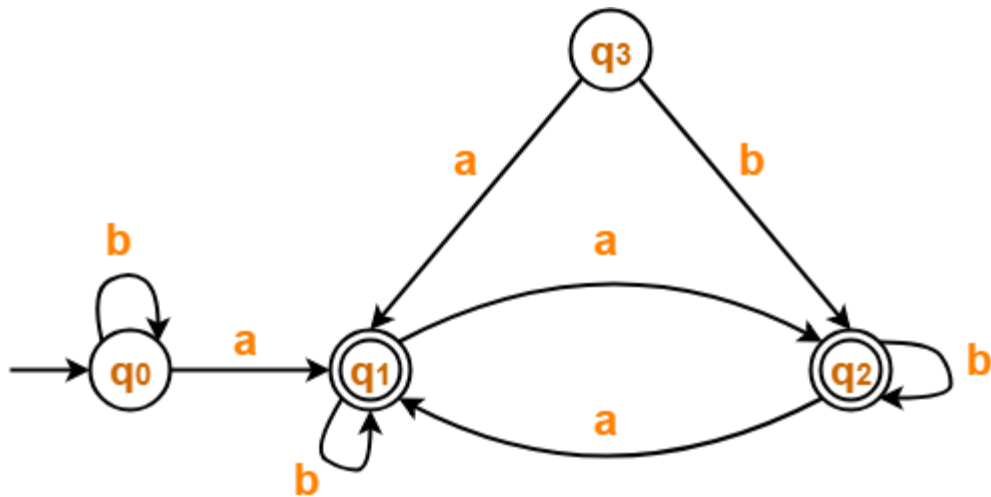
$E' \rightarrow +TE' \mid \epsilon$

Then eliminate for T as:

$T \rightarrow FT'$

$T' \rightarrow *FT' \mid \epsilon$

5. Minimize the following DFA:



Answer:

Now using Equivalence Theorem, we have-

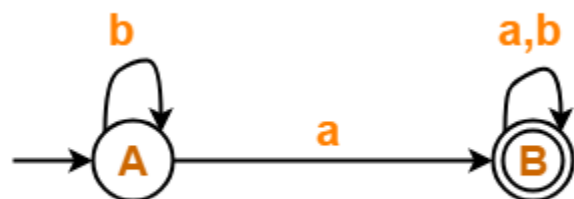
$$P_0 = \{ q_0 \} \{ q_1, q_2 \}$$

$$P_1 = \{ q_0 \} \{ q_1, q_2 \}$$

Since $P_1 = P_0$, so we stop.

From P_1 , we infer that states q_1 and q_2 are equivalent and can be merged together.

So, Our minimal DFA is-



Minimal DFA