### **Analysis of the Algorithm**

Kruskal's algorithm is a greedy algorithm used to find the Minimum Spanning Tree (MST) of a connected, undirected graph with weighted edges. The algorithm follows these steps:

### 1. Initialization:

- Create a graph with a specified number of vertices.
- Store edges in a list, where each edge is represented as a tuple of two vertices and a weight.

## 2. Sorting Edges:

• Sort all edges in non-decreasing order based on their weights. This is crucial because the algorithm processes the smallest edges first to ensure that the MST remains minimal.

#### 3. Union-Find Structure:

- Use a union-find (disjoint-set) data structure to keep track of which vertices are in which
  components. This helps in efficiently checking whether adding an edge would create a
  cycle.
- The **find** function implements path compression to flatten the structure of the tree whenever **find** is called, which speeds up future operations.
- The **union** function uses union by rank to attach smaller trees under larger trees, keeping the overall structure balanced.

## 4. Building the MST:

- Iterate through the sorted edges and for each edge, check if the vertices of the edge belong to different components using the **find** function.
- If they do, add the edge to the MST and unite the two components using the **union** function.
- Stop when the number of edges in the MST equals **V 1**, where **V** is the number of vertices.

# 5. **Output**:

Print the edges included in the MST and the total cost of the MST.

### **Complexity Analysis**

### Time Complexity:

- Sorting the edges takes (O(E \log E)), where (E) is the number of edges.
- Each union and find operation takes nearly constant time, (O(\alpha(V))), where (\alpha) is the inverse Ackermann function, which grows very slowly.
- Overall, the time complexity is (O(E \log E)).

- Space Complexity:
  - The space complexity is (O(V + E))