

Algorithms Required for Heapsort

1. **Heapify:** This function ensures that a subtree rooted at index i satisfies the heap property. It compares the node with its children and swaps it with the largest child if necessary, then recursively heapifies the affected subtree.
2. **Build Heap:** This is done by calling **heapify** on all non-leaf nodes, starting from the last non-leaf node down to the root. This rearranges the array into a max-heap.
3. **Sort:** After building the heap, the largest element (root of the heap) is swapped with the last element of the array. The size of the heap is reduced, and **heapify** is called on the root to maintain the heap property. This process is repeated until the heap is empty.

Analysis of the Heapsort Algorithm

1. **Time Complexity:**
 - **Building the Heap:** The **heapify** function is called $n/2$ times, and each call takes $O(\log n)$ time in the worst case. Thus, building the heap takes $O(n)$.
 - **Sorting:** The sorting process involves n calls to **heapify**, each taking $O(\log n)$ time. Therefore, the sorting step takes $O(n \log n)$.
 - Overall, the time complexity of Heapsort is $O(n \log n)$.