FINITE-VALUED STREAMING STRING TRANSDUCERS

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ismaeljecker.github.io

Finite-valued regular relations form an expressive set of binary relations with good algorithmic properties

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Theorem: We can decide in polynomial space whether a given SST defines a **finite-valued regular relation**

Theorem: Every **finite-valued regular relation** can be decomposed into a finite union of **regular functions**

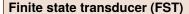
Transducers are abstract machines that recognise **relations**

$$\mathsf{R} \subseteq \Sigma^* \times \Gamma^*$$

Transducers are abstract machines that recognise **relations**

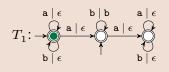
Regular relations

Transducers are abstract machines that recognise **relations**



Input: a b b a a a b b b a b a a a

Output:



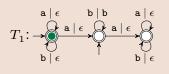
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Regular relations

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Finite state transducer (FST)

Input: a b b a a a b b b a b a a a

Output:



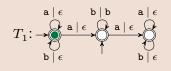
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Finite state transducer (FST)

Input: abbaaaabbbabaaa

Output:



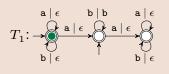
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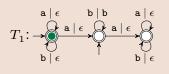
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Input: abbaaaabbaaaa

Output:



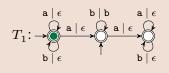
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Finite state transducer (FST)

Input: abbaa a bbbabaaa

Output:



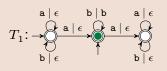
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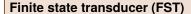
Input: abbaaa bbbabaaa

Output:



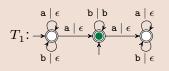
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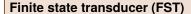
Input: abbaaabbbabaaa

Output: b



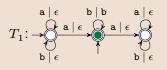
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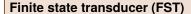
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Output: b b



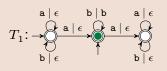
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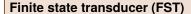
Input: abbaaabbbabaaa

Output: bbb



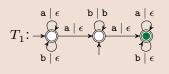
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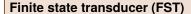
Input: a b b a a a b b b a b a a a

Output: bbb



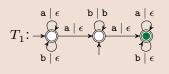
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Input: abbaaabbbabaa

Output: bbb



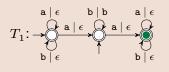
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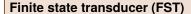
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Output: bbb



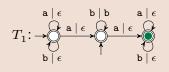
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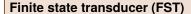
Input: abbaaabbbabaa

Output: b b b



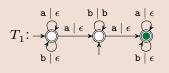
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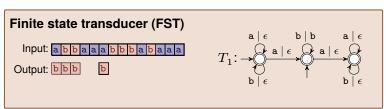
Input: abbaaabbbabaaa

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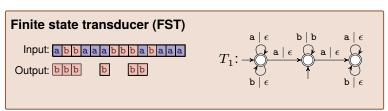
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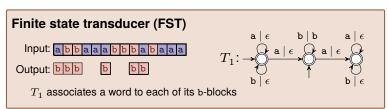
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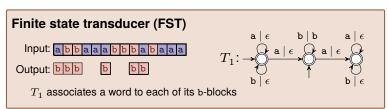
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Streaming string transducer (SST)

Input: a b b a a a b b b

 R_1 : R_2 :

$$a \mid R_1 := R_1 a$$

$$T_2 : \longrightarrow R_1 R_2$$

$$b \mid R_2 := R_2 b$$

Regular relations

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Streaming string transducer (SST)

Input: a b b a a a b b b

 $R_1: \boxed{a}$ R₂:

$$a \mid R_1 := R_1 a$$

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R₂: b

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 $R_1: \boxed{a}$

R₂: b b

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$$T_2 : \longrightarrow R_1 R_2 R_3 A$$
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$$R_2 R_3 = R_3 b$$

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Streaming string transducer (SST)

Input: a b b a a a b b b

R₁: a a a a R₂: bbbbb

$$\begin{array}{c|c} a \mid R_1 \coloneqq R_1 a \\ \hline T_2 \colon \longrightarrow \begin{array}{c} R_1 R_2 \\ R_2 R_3 \end{array}$$

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Streaming string transducer (SST)

Input: a b b a a a b b b a b a a a

 R_1 : aaaaaaaa R_2 : bbbbbb

$$A \mid R_1 := R_1 A$$

$$T_2 : \longrightarrow R_1 R_2 R_3$$
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Rational relations
Regular relations

recognised by finite state transducers recognised by streaming string transducers

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Streaming string transducer (SST)

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Input: a b b a a a b b b a

Output: a a a a a a

 T_2 groups the a's and the b's of its input

$$A \mid R_1 := R_1 A$$

$$T_2 : \longrightarrow R_1 R_2$$
 $B_2 \mid R_2 \mid R_2 \mid R_3 \mid R_4 \mid R_5 \mid R_5$

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Regular relations

Rational relations
Regular relations

recognised by finite state transducers recognised by streaming string transducers

Equivalence of rational relations is undecidable

Regular relations

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Equivalence of **regular functions** is **decidable**

→ each input is mapped to at most 1 output

Rational relations
Regular relations

recognised by finite state transducers recognised by streaming string transducers

functions

Equivalence of rational relations is undecidable

Equivalence of **regular functions** is **decidable**

→ each input is mapped to at most 1 output

Equivalence of finite-valued regular relations is decidable

 $\rightarrow \exists \mathbf{k} \in \mathbb{N} \text{ s.t. each input is mapped to at most } \mathbf{k} \text{ outputs}$

Rational relations
Regular relations

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Theorem: We can decide in polynomial space whether a given

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Theorem: Every **finite-valued regular relation** can be decomposed

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Raised as open problems by [2011. Alur, Deshmukh]

Known to hold for FST [1989. Weber], [1993. Weber]

and for SST with a single register [2017. Gallot et al.]

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⇒ The equivalence problem for finite-valued SST is in Elementary

⇒ Finite-valued 2-way FST are as expressive as finite-valued SST

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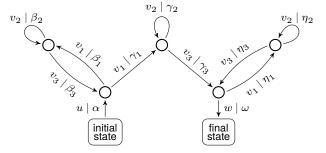
recognised by finite state transducers recognised by streaming string transducers

DECIDING FINITE VALUEDNESS

Theorem: We can decide in polynomial space whether a given

SST defines a finite-valued regular relation

Forbidden pattern: (inspired by [2008. De Souza])



- The relation recognised by the pattern is not 1-valued
- The substitutions produced on the loops have idempotent structure

Theorem: Every finite-valued regular relation can be decomposed

into a finite union of regular functions

Proof: (inspired by [2008. Sakarovitch, de Souza] and relying on [2023. FJLW])

Theorem: Every **finite-valued regular relation** can be decomposed into a finite union of **regular functions**

Proof: (inspired by [2008. Sakarovitch, de Souza] and relying on [2023. FJLW]) that associates at most \mathbf{k} outputs to each input

 $\mathcal{T}_1,\,...,\,\mathcal{T}_k$ —that all associate at most 1 output to each input

Theorem: Every **finite-valued regular relation** can be decomposed into a finite union of **regular functions**

Proof: (inspired by [2008. Sakarovitch, de Souza] and relying on [2023. FJLW])

**That associates at most k outputs to each input

 \mathcal{T}' that has at most **k** runs on each input

 $\mathcal{T}_1',\,...,\,\mathcal{T}_k'$ that all have at most 1 run on each input

 \mathcal{T}_1 , ..., \mathcal{T}_k —that all associate at most 1 output to each input

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That associates at most k outputs to each input

that has at most $\mathbf k$ runs on each input $\mathcal T'_i$ copies the ith lexicographically smallest run of $\mathcal T'$ $\mathcal T'_1, \, ..., \, \mathcal T'_k$ that all have at most $\mathbf 1$ run on each input

 \mathcal{T}_1 , ..., \mathcal{T}_k —that all associate at most 1 output to each input

Theorem: Every **finite-valued regular relation** can be decomposed into a finite union of **regular functions**

Proof: (inspired by [2008. Sakarovitch, de Souza] and relying on [2023. FJLW])

 ${\mathcal T}$ that associates at most ${\mathbf k}$ outputs to each input

 \mathcal{T}' that has at most **k** runs on each input

 $\mathcal{T}'_1, \ldots, \mathcal{T}'_{\nu}$ that all have at most 1 run on each input

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Theorem: Every **finite-valued regular relation** can be decomposed into a finite union of **regular functions**

Proof: (inspired by [2008. Sakarovitch, de Souza] and relying on [2023. FJLW])

that associates at most \mathbf{k} outputs to each input \mathcal{T}' only keeps the runs of \mathcal{T} that are far from each other that has at most \mathbf{k} runs on each input

 $\mathcal{T}_1',\,...,\,\mathcal{T}_k'$ that all have at most 1 run on each input

 ${\mathcal T}_1,...,{\mathcal T}_{\mathsf k}$ that all associate at most 1 output to each input

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