# Star and snowflake schema

**DATABASE DESIGN** 



**Lis Sulmont**Curriculum Manager



#### Star schema

#### Dimensional modeling: star schema

#### **Fact tables**

- Holds records of a metric
- Changes regularly
- Connects to dimensions via foreign keys

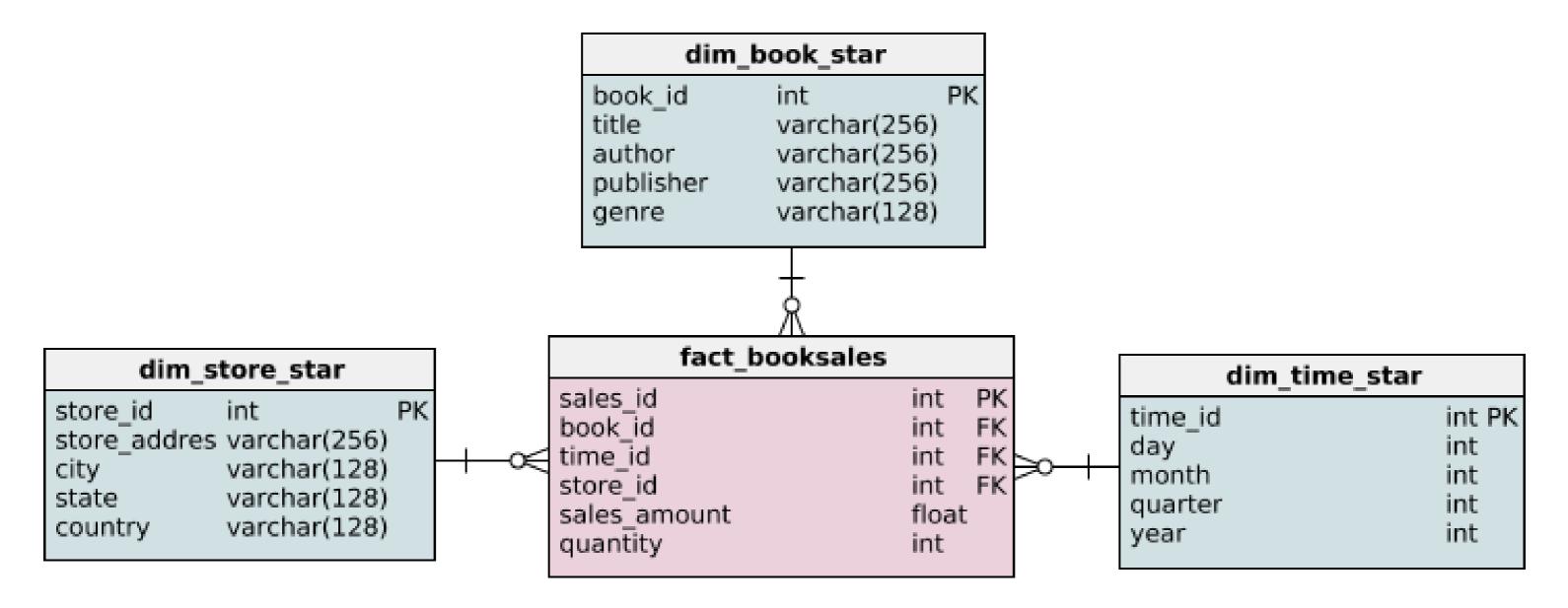
#### **Dimension tables**

- Holds descriptions of attributes
- Does not change as often

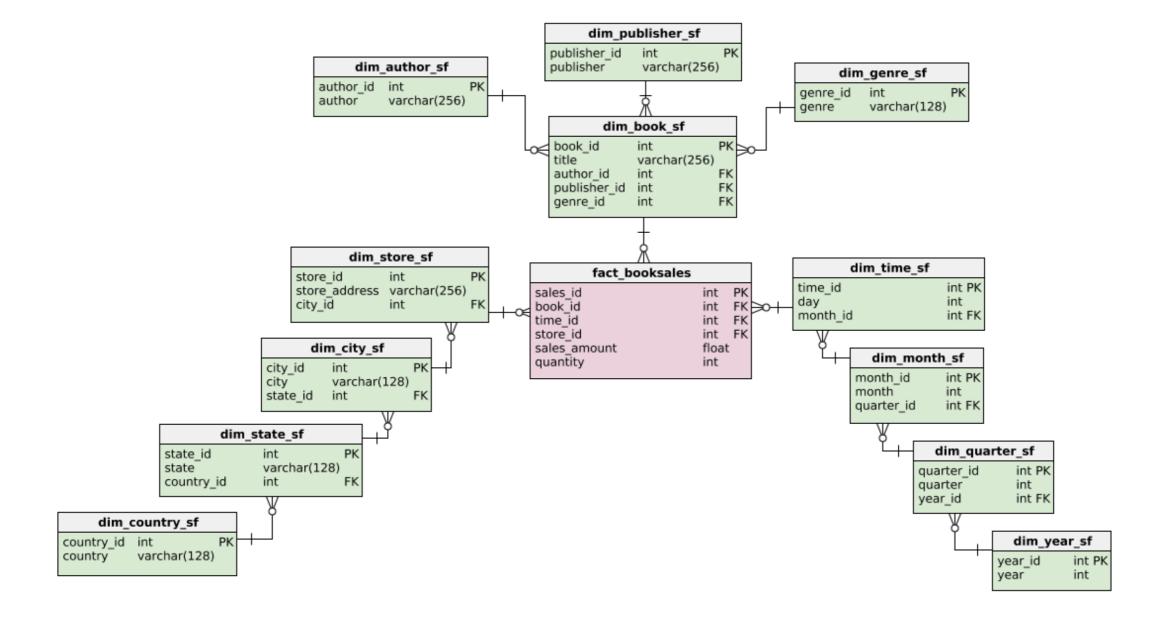
#### **Example:**

- Supply books to stores in USA and Canada
- Keep track of book sales

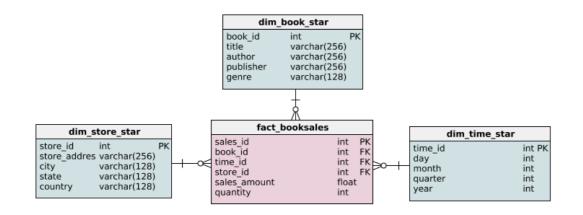
# Star schema example



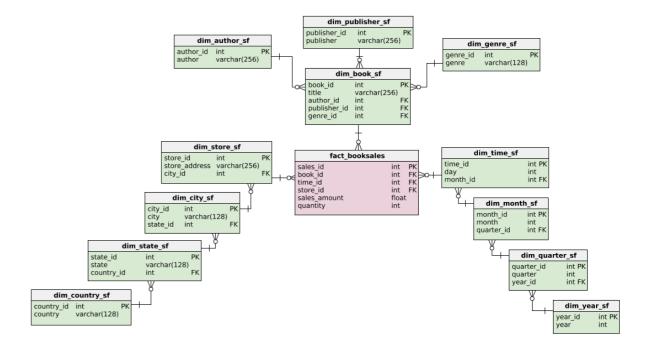
# Snowflake schema (an extension)



# Same fact table, different dimensions



Star schemas: one dimension



Snowflake schemas: more than one dimension

Because dimension tables are normalized

#### What is normalization?

- Database design technique
- Divides tables into smaller tables and connects them via relationships
- Goal: reduce redundancy and increase data integrity

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Identify repeating groups of data and create new tables for them

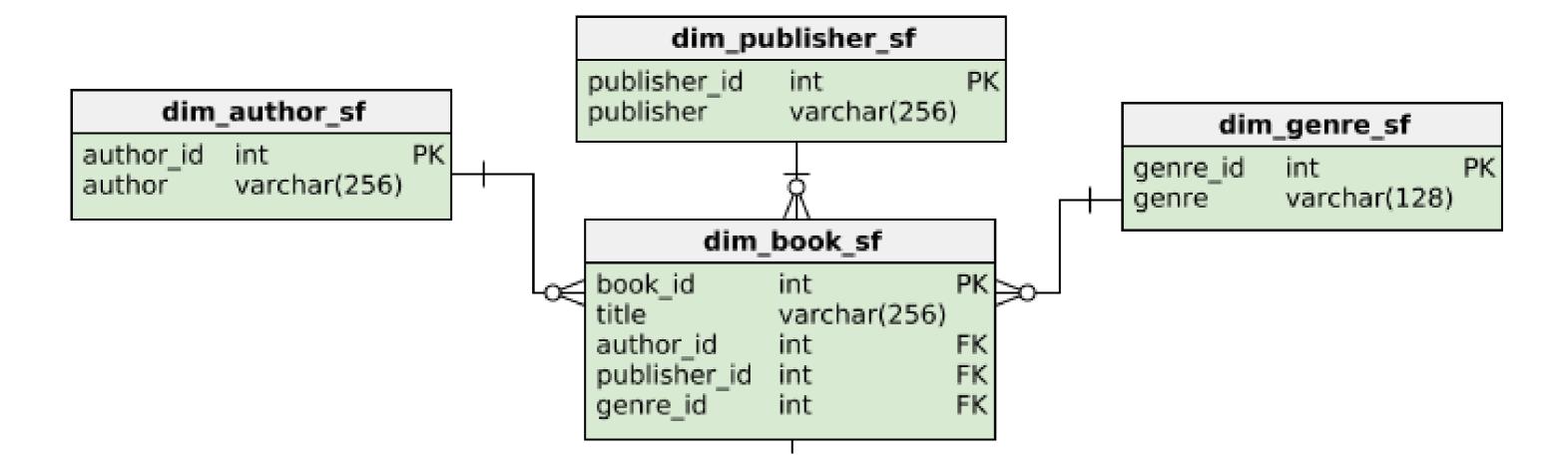
# Book dimension of the star schema

dim_	_book_star	
book_id	int	PK
title	varchar(256)	
author	varchar(256)	
publisher	varchar(256)	
genre	varchar(128)	

Most likely to have repeating values:

- Author
- Publisher
- Genre

# Book dimension of the snowflake schema

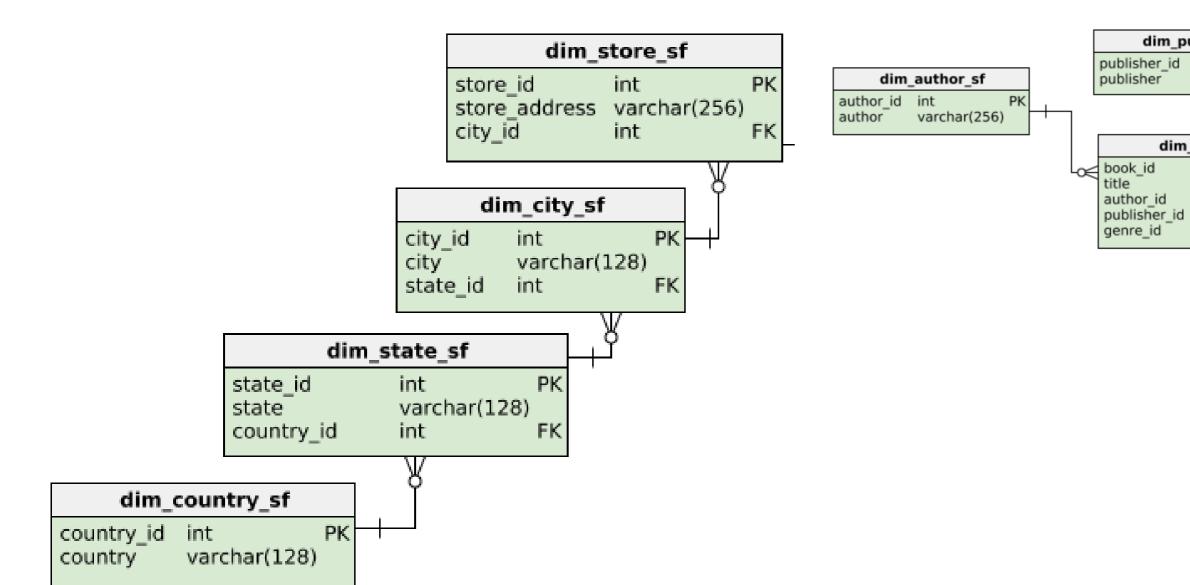


# Store dimension of the star schema

dim_st	tore_star	
store_id	int	PΚ
store_addres	varchar(256)	
city	varchar(128)	
	varchar(128)	
country	varchar(128)	

- City
- State
- Country

# Store dimension of the snowflake schema





dim genre sf

varchar(128)

genre\_id

dim\_publisher\_sf

dim book sf

int

int

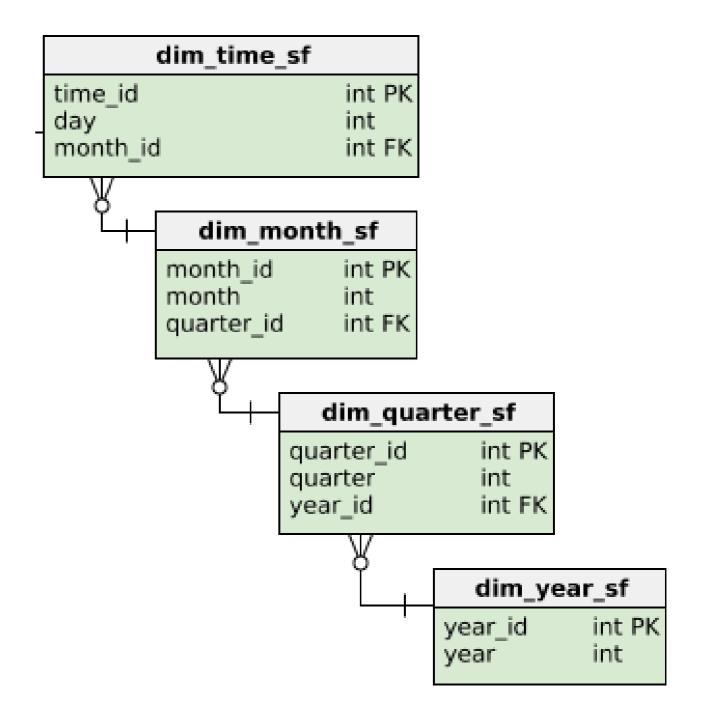
int

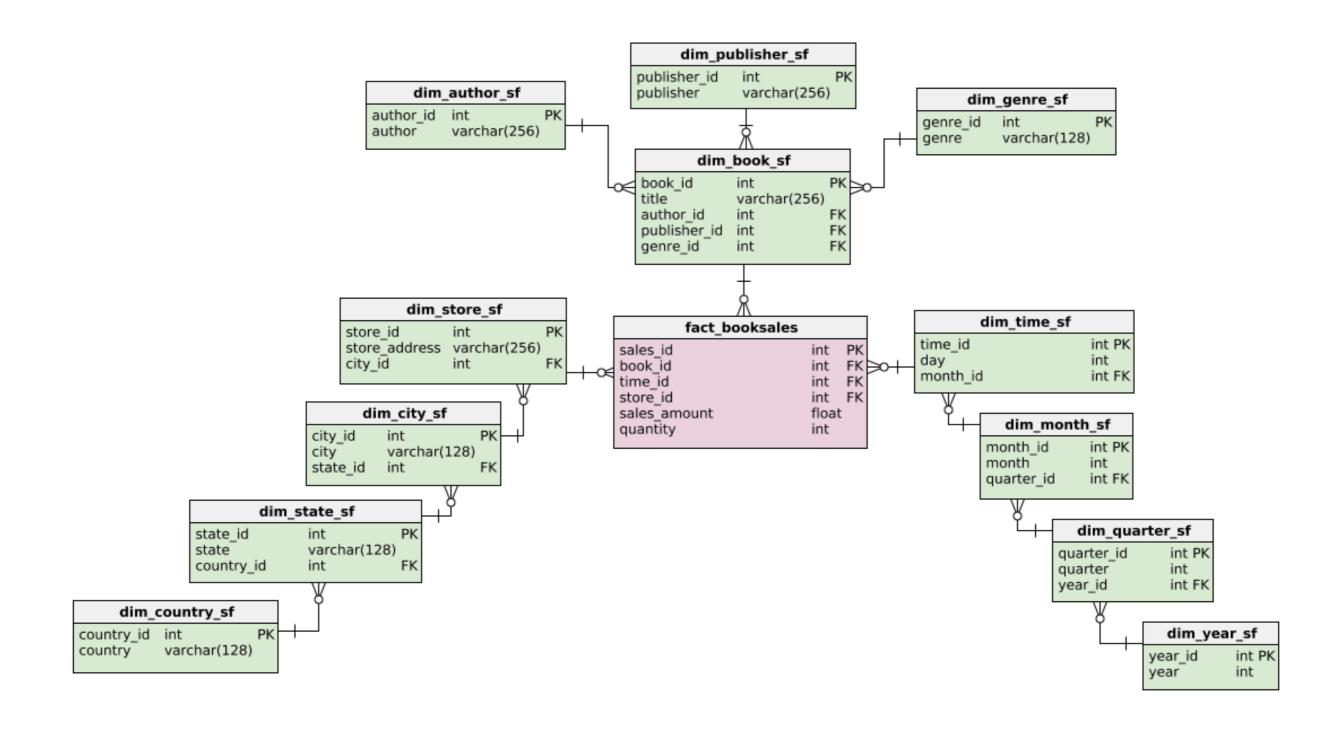
int

varchar(256)

varchar(256)

dim_time_star	
time_id	int PK
day	int
month	int
quarter	int
year	int





# Let's practice!

**DATABASE DESIGN** 



# Normalized and denormalized databases

**DATABASE DESIGN** 

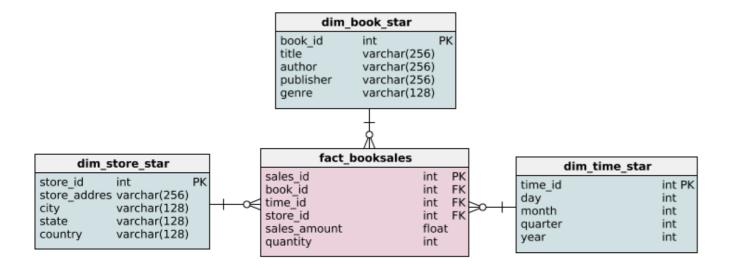


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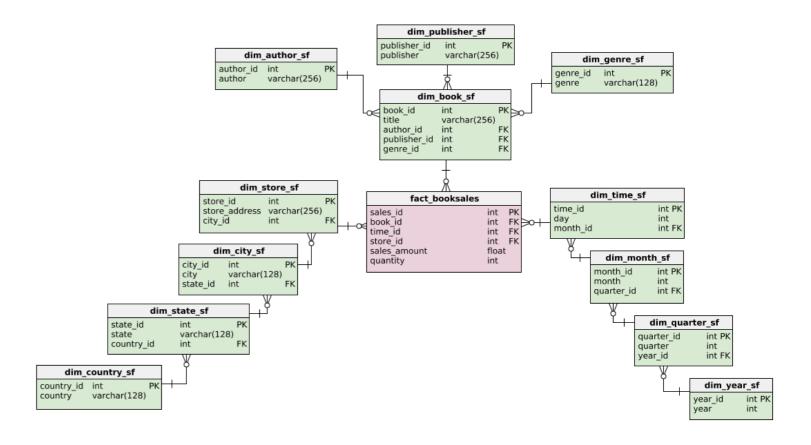


# Back to our book store example

Denormalized: star schema



#### Normalized: snowflake schema



#### **Denormalized Query**

Goal: get quantity of all Octavia E. Butler books sold in Vancouver in Q4 of 2018

```
SELECT SUM(quantity) FROM fact_booksales
    -- Join to get city
INNER JOIN dim_store_star on fact_booksales.store_id = dim_store_star.store_id
    -- Join to get author
INNER JOIN dim_book_star on fact_booksales.book_id = dim_book_star.book_id
    -- Join to get year and quarter
INNER JOIN dim_time_star on fact_booksales.time_id = dim_time_star.time_id
WHERE
dim_store_star.city = 'Vancouver' AND dim_book_star.author = 'Octavia E. Butler' AND
dim_time_star.year = 2018 AND dim_time_star.quarter = 4;
```

7600

#### Total of 3 joins



#### Normalized query

```
SEL FCT
  SUM(fact_booksales.quantity)
FROM
  fact booksales
  -- Join to get city
  INNER JOIN dim_store_sf ON fact_booksales.store_id = dim_store_sf.store_id
  INNER JOIN dim_city ON dim_store_sf.city_id = dim_city_sf.city_id
  -- Join to get author
  INNER JOIN dim_book_sf ON fact_booksales.book_id = dim_book_sf.book_id
  INNER JOIN dim_author_sf ON dim_book_sf.author_id = dim_author_sf.author_id
  -- Join to get year and quarter
  INNER JOIN dim_time_sf ON fact_booksales.time_id = dim_time_sf.time_id
  INNER JOIN dim_month_sf ON dim_time_sf.month_id = dim_month_sf.month_id
  INNER JOIN dim_quarter_sf ON dim_month_sf.quarter_id = dim_quarter_sf.quarter_id
  INNER JOIN dim_year_sf ON dim_quarter_sf.year_id = dim_year_sf.year_id
```

#### Normalized query (continued)

```
WHERE
   dim_city_sf.city = `Vancouver`
   AND
   dim_author_sf.author = `Octavia E. Butler`
   AND
   dim_year_sf.year = 2018 AND dim_quarter_sf.quarter = 4;
```

```
sum
7600
```

#### Total of 8 joins

So, why would we want to normalize a databases?

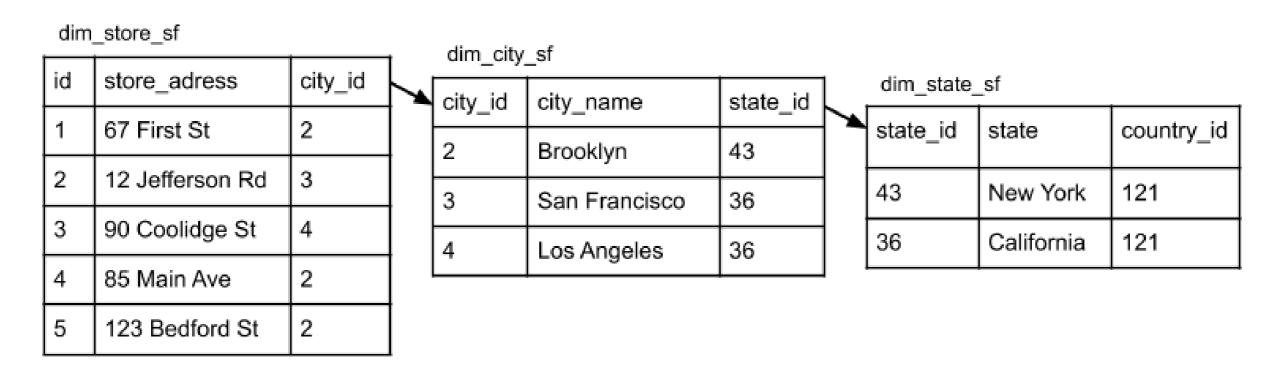
# Normalization saves space

dim\_store\_star

id	store_adress	city	state	country
1	67 First St	Brooklyn	New York	USA
2	12 Jefferson Rd	San Francisco	California	USA
3	90 Coolidge St	Los Angeles	California	USA
4	85 Main Ave	Brooklyn	New York	USA
5	123 Bedford St	Brooklyn	New York	USA

Denormalized databases enable data redundancy

# Normalization saves space



Normalization eliminates data redundancy

# Normalization ensures better data integrity

#### 1. Enforces data consistency

Must respect naming conventions because of referential integrity, e.g., 'California', not 'CA' or 'california'

#### 2. Safer updating, removing, and inserting

Less data redundancy = less records to alter

#### 3. Easier to redesign by extending

Smaller tables are easier to extend than larger tables

### Database normalization

#### Advantages

- Normalization eliminates data redundancy: save on storage
- Better data integrity: accurate and consistent data

#### Disadvantages

Complex queries require more CPU

#### Remember OLTP and OLAP?

**OLTP** 

**OLAP** 

e.g., Operational databases

e.g., Data warehouses

#### Typically highly normalized

- Write-intensive
- Prioritize quicker and safer insertion of data

#### Typically less normalized

- Read-intensive
- Prioritize quicker queries for analytics

# Let's practice!

**DATABASE DESIGN** 



# Normal forms

DATABASE DESIGN



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### Normalization

Identify repeating groups of data and create new tables for them

A more formal defintion:

The goals of normalization are to:

- Be able to characterize the level of redundancy in a relational schema
- Provide mechanisms for transforming schemas in order to remove redundancy

<sup>1</sup> https://opentextbc.ca/dbdesign01/chapter/chapter <sup>2</sup> 12 <sup>3</sup> normalization/

# Normal forms (NF)

Ordered from least to most normalized:

- First normal form (1NF)
- Second normal form (2NF)
- Third normal form (3NF)
- Elementary key normal form (EKNF)
- Boyce-Codd normal form (BCNF)

- Fourth normal form (4NF)
- Essential tuple normal form (ETNF)
- Fifth normal form (5NF)
- Domain-key Normal Form (DKNF)
- Sixth normal form (6NF)

<sup>&</sup>lt;sup>1</sup> https://en.wikipedia.org/wiki/Database\_normalization#Normal\_forms

#### 1NF rules

- Each record must be unique no duplicate rows
- Each cell must hold one value

#### **Initial data**

#### In 1NF form



### 2NF

- Must satisfy 1NF AND
  - If primary key is one column
    - then automatically satisfies 2NF
  - If there is a composite primary key
    - then each non-key column must be dependent on all the keys

#### **Initial data**

Student_id (PK)	Course_id (PK)	Instructor_id	Instructor	Progress
235	2001	560	Nick Carchedi	.55
455	2345	658	Ginger Grant	.10
767	6584	999	Chester Ismay	1.00

## In 2NF form

```
| Course_id (PK) | Instructor_id | Instructor
|------|
| 2001 | 560 | Nick Carchedi |
| 2345 | 658 | Ginger Grant |
| 6584 | 999 | Chester Ismay |
```



### 3NF

- Satisfies 2NF
- No transitive dependencies: non-key columns can't depend on other non-key columns

#### **Initial Data**

## In 3NF

```
| Course_id (PK) | Instructor | Tech |
|------|
| 2001 | Nick Carchedi | Python |
| 2345 | Ginger Grant | SQL |
| 6584 | Chester Ismay | R |
```



## Data anomalies

What is risked if we don't normalize enough?

- 1. Update anomaly
- 2. Insertion anomaly
- 3. Deletion anomaly

#### **Update anomaly**

Data inconsistency caused by data redundancy when updating

```
| Enrolled_in
Student_ID | Student_Email
                                                         Taught_by
                                                        | Nick Carchedi
230
             lisa@gmail.com
                              | Cleaning Data in R
367
             bob@hotmail.com |
                               Data Visualization in R |
                                                          Ronald Pearson
520
            | ken@yahoo.com
                              | Introduction to Python
                                                          Hugo Bowne-Anderson
            ken@yahoo.com
                              | Arima Modeling in R
                                                        | David Stoffer
520
```

To update student 520 's email:

- Need to update more than one record, otherwise, there will be inconsistency
- User updating needs to know about redundancy

#### Insertion anomaly

Unable to add a record due to missing attributes

```
Student_ID | Student_Email
                              | Enrolled_in
                                                         Taught_by
                                                        | Nick Carchedi
230
             lisa@gmail.com
                              | Cleaning Data in R
367
             bob@hotmail.com | Data Visualization in R |
                                                          Ronald Pearson
            | ken@yahoo.com
                                                         Hugo Bowne-Anderson
520
                              | Introduction to Python
            ken@yahoo.com
                              | Arima Modeling in R
                                                        | David Stoffer
520
```

Unable to insert a student who has signed up but not enrolled in any courses

#### **Deletion anomaly**

Deletion of record(s) causes unintentional loss of data

```
| Enrolled_in
Student_ID | Student_Email
                                                Taught_by
                                               | Nick Carchedi
230
          lisa@gmail.com
                         | Cleaning Data in R
367
          Ronald Pearson
          | ken@yahoo.com
                         | Introduction to Python | Hugo Bowne-Anderson
520
          | ken@yahoo.com
                         | Arima Modeling in R
                                               | David Stoffer
520
```

If we delete Student 230, what happens to the data on Cleaning Data in R?

#### Data anomalies

What is risked if we don't normalize enough?

- 1. Update anomaly
- 2. Insertion anomaly
- 3. Deletion anomaly

The more normalized the database, the less prone it will be to data anomalies

Don't forget the downsides of normalization from the last video

# Let's practice

**DATABASE DESIGN** 

