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Assignment Name: CSE201 - Data Structure

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- 1. What is data structure? What are the type of data structures? What are the notation of performance of an algorithm?
- Ans: Data structure can be defined as the group of data elements which provides an efficient way of storing and organizing data in the computer so that it can be used efficiently. Example: Annay, linked list, Stack, queue etc.

2 types of Data structure

- A data structure is called linear if all elements of its are arranged in the linear order. In linear structure, the elements are stoned non-hierarchical way where each element has the successor and predecessors except the 1st and last element. example: Array, linked list, stack, queues etc.
- (ii) Non linear Dota Structure: This data structure does not form a Sequence i.e. each item on elements are connected with two or more other item in a non linear arrangement.

 The data elements are not arranged in sequential structure.

 Example: Graph, Tree etc.

2nd Pant: to measure the Penformance of an algorithm, we generally used three asymptotic notation.

These are:

- (1) Big On Notation (0)
- 1 Omega (12) Notation
- (1) Theta (0) Notation

(1) Big Oh Notation (0):

Def: O(g(m)) = off(m): there exist positive constants c and no such that o≤f(m) ≤ c g(m) for all n≥no?

11 Omega (12) Notation:

Def: -2(g(n)) = f(n): there exist positive constants cand no such that $0 \le cg(n) \le f(n)$ for all $n \ge n_0$?

(11) Theta (8) Notation:

Def: O(g(n)) = f(n): there exist positive constants c_1, c_2 and no such that $o \le c_1 g(n) \le f(n) \le c_2 g(n)$ for all $n \ge n_0$?

2. What is Bubble sort? Write an algorithm for sorting an armay of N elements using Bubble sort.

Sort this armay using Bubble sort - 15, 18, 4, 5, 2

Ans! Bubble sort! Bubble sort is a sorting algorithm that compares two adjacent elements and swaps them until they are not in the intended order.

Bubble Sort Algorithmo

- 1. Repeat step 2 to 3 for i=1 to N-1
- 2. Set 0=1
- 3. Repeat while j = N-i

(a) If Data[J] > Data[J+1]

the interchange Data[J+1] and Data[J+1]

(b) set J=J+1

[End of inner loop]

[End of owder Loop]

4. Exit.

2nd Part: Given Armay 15, 18, 4, 5, 2

total 5 elements are here. So, the size of armay in 5.

let, Arm[5] = 215, 18, 4, 5, 23

Now Discuss the Procedure of Bubble sort:

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(15) (18) 4 5 2

compare them. if 15>18 them swap else not.

15 (19) (4) 5 2

compare them. since 18>4. swap them.

15 4 (18) (5) 2

18>5 then swap them

15 4 5 (18) (2)

18>2. swap them

15 4 5 2 18 [and finaly 18 is stored permonarily at n-1 th position.

Again penform the overall operation among the 1st index to 3thindex and then 2,1 and so on.

(5) 4) 5 2 18 4 (5) 5) 2 18 4 5 (5) 2 18 (9) 5) 2 15 18 4 (5) 2 15 18 4 2 5 15 18 2 4 5 15 18 and this is the sorted annay.

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3. What are the type of linked list? Explain with diagram. White an Algorithm to delete element (node) from "pos" position in the list.

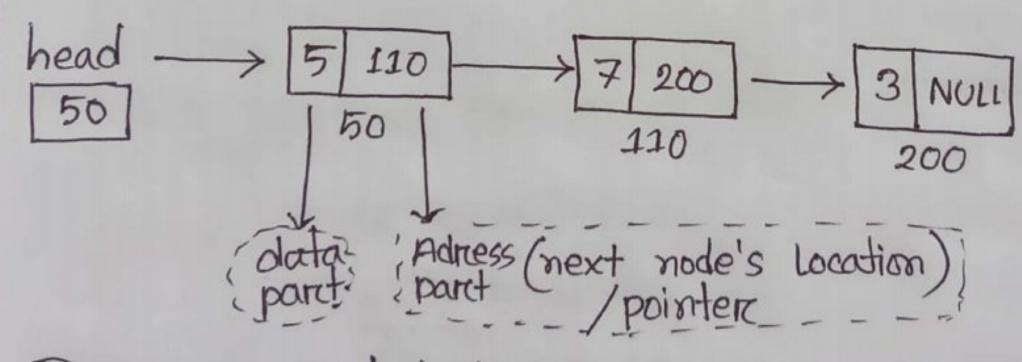
Ans: linked list is a linear data structure, in which the elements are not stoned at contiguous memory location. The elements are linked using pointer.

Types of linked list:

Singly linked list:

It is commonly used in program.

The singly linked list contain two paret one is data paret and another is adress paret which contain the adress of next node. The adress of next node is also known as pointer.

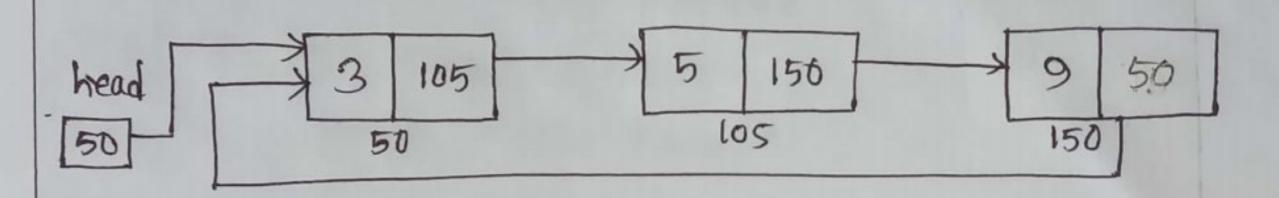


Doubly Linked list:

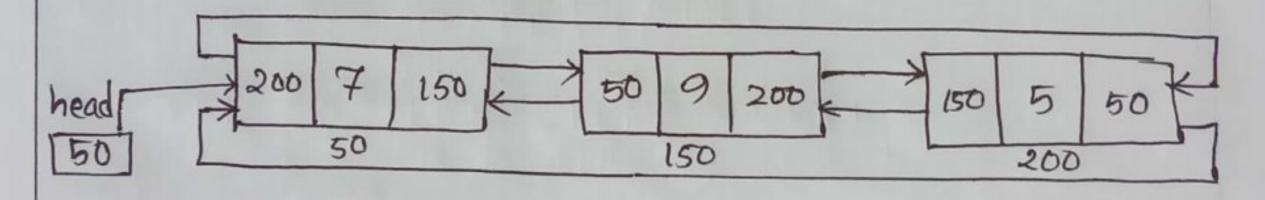
The doubly linked list contains too pointer. It has three part where one data part and two adress part which points to it's previous nodes and next nodes.

Circular linked list of It's a variation of singly linked list.

The only difference is that the last node does connect to the first node. So the link part of the last node holds the first node's adness.



Doubly Circular linked list: It's the combination of doubly linked list and doubly linked list.



and Part: An Algorithm to delete element from "pos" position:

STEP-1: IF HEAD = NULL
Write underflow
Go to step 11
END of If

STEP-2: SET TEMP = HEAD

STEP-3: SET I=0

STEP-4: REPEAT STEP 5 TO 8

UNTIL IZ POS

STEP-5: TEMP1 = TEMP

STEP-6: TEMP = TEMP -> NEXT

STEP-7: JF TEMP = NULL

Write "NOT Found"

GO to STEP 11

END of IF

STEP-8: J=]+1

END OF LOOP

STEP-9: TEMP1 -> NEXT = TEMP -> NEXT

STEP-10: FREE TEMP

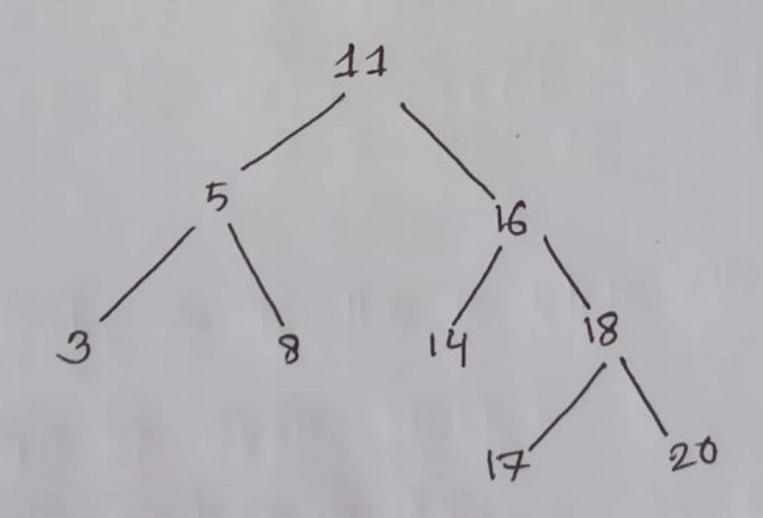
STEP-11 : EXIT

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in order = $\sqrt{3}$, 5, 8, 11, 14, 16, 17, 18, 20

Ans: If T is any Binary free, then T is called complete Binary tree if is Each node of T has atmost 2 children.

(ii) All nodes at the last level as fare left as possible.



5. convent the infix equation to post-fix using stack - (A+B^D)/(E-F)*G

Ans: Let Given equation, P= (A+BND)/(E-F)+G

Character Scamed	Stack	Postfix Expression
((
A	(A
+	(+	A
B	(+	AB
^	(+1	AB
D	(+1	ABD
)		ABDA+
		ABDA+
(10	ABDA+
E	10	ABDA+ E
_	16-	ABDA+E
F	1(-	ABDA+EF
)	1	ABDA+EF-
+	+	ABDN+EF-/
G	+	ABDN+EF-/G
		ABDN+ EF-/G1+

6. What is the difference between sequential (lineari) Search and Binary Search? Show the step of searching element '5' Using Binary Search Technique in this array-2,5,8,15,20.

Ans:

Discription: Linear Search is a technique where the searching the element sequentially until the element is found. But binary search is a method that finds the middle element in the list recursively until the middle element is matched with the searched element.

Sorted Data: In linear Search, the elements can be any order. but the precondition of binary search is that the array must be sorted.

Approach: linear search is a sequential approach but binary search is the devide and conquer approach.

Size: linear Search is preferable for small size and binary search fore large size.

Implementation: linear search can be implemented on any linear data structure such as array, linked list, stack, etc. On the other hand, the implementation of binarry search is limited as it can be implemented only on those data structures that have two.

element 2 searched element; find the right

echnique. Let, Annay = 2, 5, 8, 15, 202

Applying divide and conquer, we will find 5. Let F=5

so, 1sly, divide the annay and every time find the middle element and compare with the searched element. If

mid-element = searched element; "Successfully find"

return index.

else find from the right

2 5 8 15 20 here, mid = (0+4)/2 = 2

Annay [mid]>F so, searched element is within

Left em elements.

0 1 2 5 here, mid = (0+1)/2 = 0

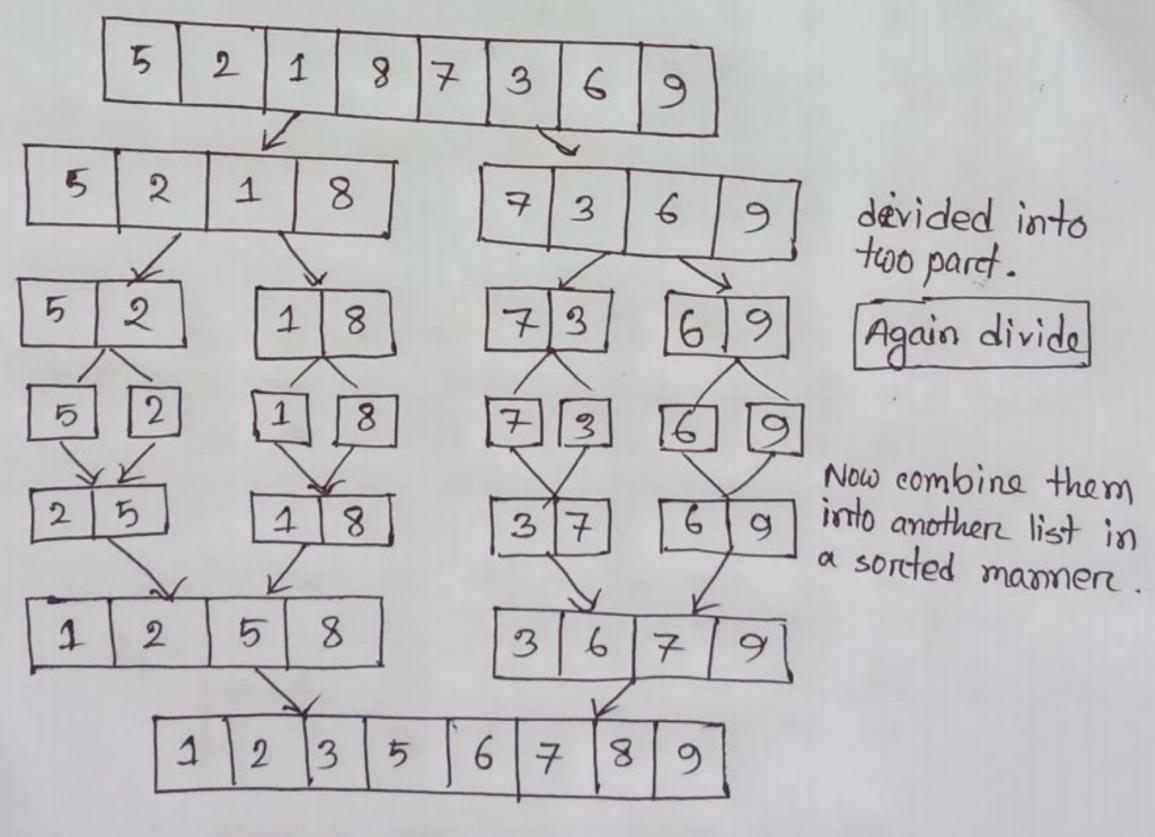
Annay [mid] < F so, searched element is within right elements.

 $\frac{1}{5}$ here, mid = $(1+1)/_2 = 1$

Annay [mid] = = F: the index of 5 is 1. 7. Explain merge sort using an example.

Ans: Menge sort is a sorting technique based on devide and conquen technique. In menge sort, first devide the armay into equals halves and then combine them in a sorted manner.

Let's A=35,2,1,8,7,36,9 is an array which is unsorted.



And this is the final sorted element of Arcray A.

- 8. Consider the weighted graph of in Fig. 1. Suppose the nodes are stored in an array Data as follows:

 DATA: X,Y,S,T
 - (a) Find the weight matrix of G

(b) Find the modrix of of shoretest paths using Warshall Algo.

Ans: (a) the weighted matrix of $G_1 = \frac{3}{5} \begin{bmatrix} 0 & 7 & 0 & 0 \\ 3 & 0 & 2 & 0 \\ 0 & 5 & 0 & 0 \end{bmatrix}$

$$91 = \begin{bmatrix} 0 & 7 & 0 & 0 \\ 3 & 10 & 2 & 0 \\ 0 & 0 & 0 & 5 \\ 6 & 1 & 4 & 0 \end{bmatrix}$$

$$92 = \begin{bmatrix} 10 & 7 & 9 & \infty \\ 3 & 10 & 2 & \infty \\ \infty & \infty & \infty & 5 \\ 4 & 1 & 3 & \infty \end{bmatrix}$$

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{ 91[2,2] = MIN[9/2,2), 9.(2,1) + 9.(1,2)] = MIN(x, 3+7)=10 (92[1,1] = MIN[Q,(1,1),Q(1,2)+Q,(2,1)] = MIN(x,7+3)=10 92[4,1] = MIN(B1(1,4), B1(4,2) + B,(2,1) = MIN(6,1+3) = 4 32[4,3] = MJN (9, (4,2) + 9, (2,3) = MJN (4,1+2) = 3 (B3[1,4] = MIN [Q2(1,4), Q2(1,3) + Q2(3,4)] = MIN(0,9+5)=14 (93[2,4] = MIN [92(2,4), 92(2,3) + 92(3,4)] = MIN(x,2+5)=7 93[4,4] = MIN[B2(4,4), Q2(4,3)+9,13,4)] = MIN(x,3+5)=8 (94[2,2] = MIN [93(2,2), 93(2,4) + 93(4,2)] = MIN(10,7+1)=8 94[3,1] = MIN[83[3,1], 93(3,4)+93(4,1)]=MIN(4,5+4)=9 94[3,2] = MIN[Q3(3,2), Q3(3,4)+Q3(4,2)] = MIN(X,5+1)=6 94[3,3] = MIN[Q3(3,3), Q3(3,4)+Q3(4,3)] = MIN(X,5+3)=8