

L8: Environment Model

SWS3012: Structure and Interpretation of Computer Programs

Martin Henz

July 12, 2023

Outline

- State ([3.1](#))
- Mutable Data ([3.3](#))
- Environment Model ([3.2](#))

Announcements

- New **game chapters** are out in Source Academy
 - Back to a weekly schedule

Outline

- State ([3.1](#))
- Mutable Data (3.3)
- Environment Model (3.2)

Recap: Variable Declaration Statement

```
let name = expression;
```

- Declares a **variable** *name* in the current scope and initializes its value to the value of *expression*
- From now on, *name* will evaluate to the value of *expression*
- Note that from [Source §3](#) onwards, **function parameters** are **variables**

Recap: Assignment Statement

name = expression;

- *name* is a **variable**; not evaluated
- *expression* is evaluated, then its value is **assigned** to the variable *name*
- From now on, *name* will evaluate to the value of *expression*

Assignment: Pros

- Assignment allows us to create objects with ***state***
- State allows objects to behave differently **over time**

Assignment: Cons

- Harder to reason about programs
 - Harder to debug
 - Harder to verify correctness
- **Substitution model of evaluation breaks down!**
 - Not powerful enough to explain state
 - Need a more sophisticated model — ***Environment Model***

Substitution Model Breaks Down

- Consider

```
function make_simplified_withdraw(balance) {  
  return amount => {  
    balance = balance - amount;  
    return balance;  
  }  
}
```

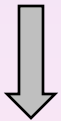
- Use **substitution model** to evaluate

```
(make_simplified_withdraw(25))(20);
```

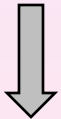
Substitution Model Breaks Down

- Use substitution model to evaluate

```
(make_simplified_withdraw(25))(20);
```



```
(amount => { balance = 25 - amount; return 25; })(20);
```



```
balance = 25 - 20; return 25; // WRONG!
```

- It returns 25, which is wrong!

Why Substitution Model Breaks Down?

- Substitution model considers a constant/variable as **just a name for a value**
 - Its value will not change
 - Therefore, one can be substituted for the other
<https://sourceacademy.nus.edu.sg/sicpjs/3.1>
- But **assignment** considers a variable as a “**container**” **holding a value**
 - The contents of the container may be **changed over time**
 - The container is maintained in a structure called an ***environment***

Outline

- State (3.1)
- Mutable Data (3.3)
- Environment Model (3.2)

Mutable Data

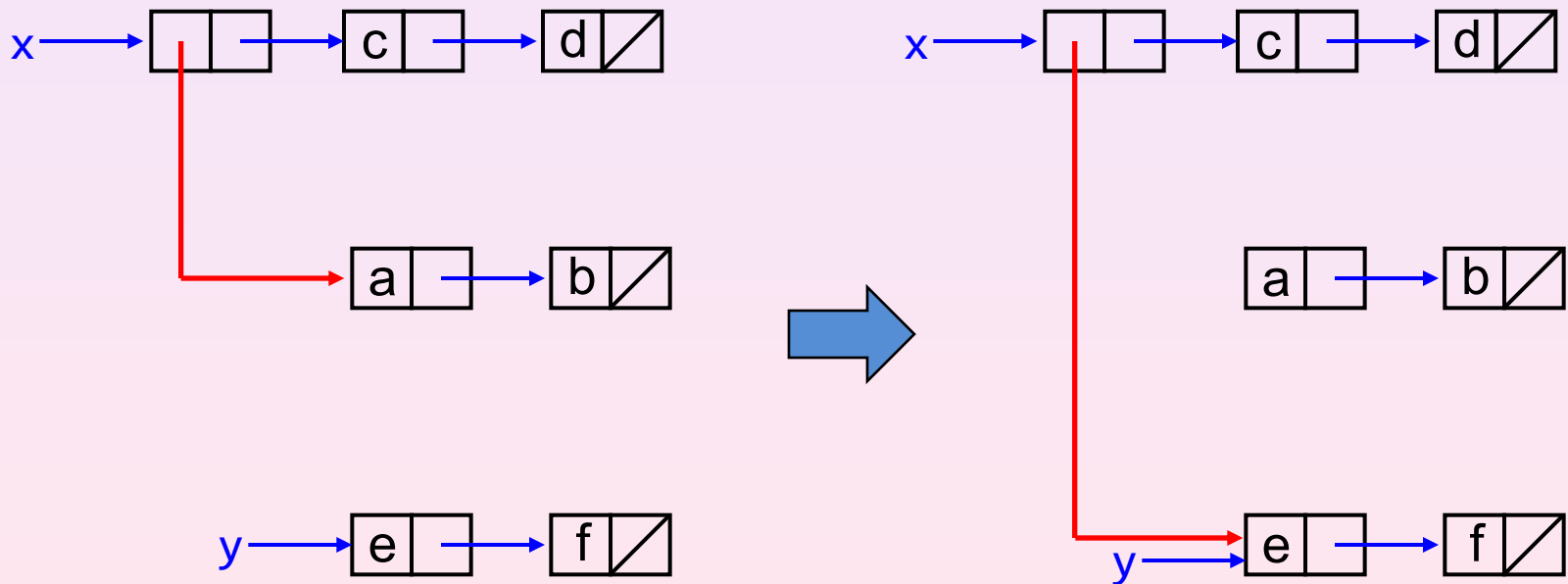
- **Assignment** gives us the ability to create *mutable* data, i.e. data that can be modified
 - E.g. bank account
- In Source §1 and §2, all our data were *immutable*. We had
 - Constructors, selectors, predicates, printers
 - But no *mutators*

Mutable Pairs

- Now we will allow *mutators* in order to create *mutable data structures*
- After creating a pair with `pair`
 - The **head** can be changed using `set_head`
 - The **tail** can be changed using `set_tail`
- Mutating mutable pairs
 - `set_head(p, x)` changes **head** of pair `p` to `x`
 - `set_tail(p, x)` changes **tail** of pair `p` to `x`

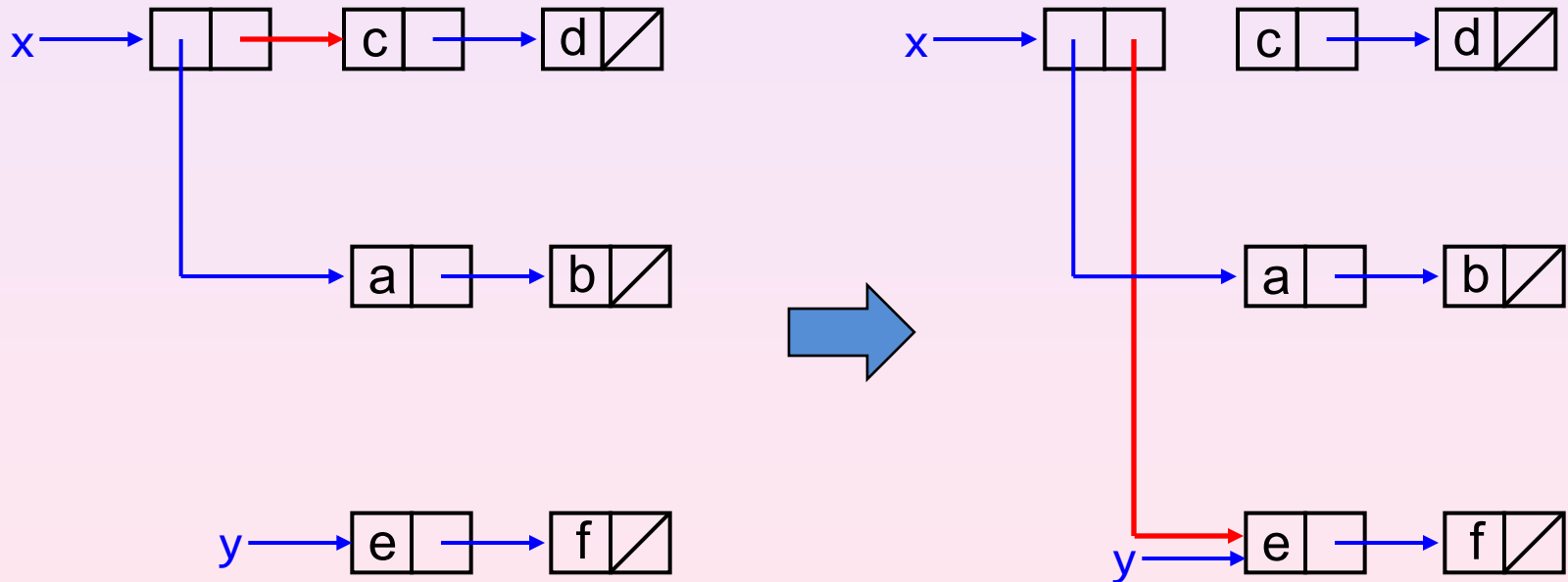
set_head Example

- Effect of `set_head(x, y)`



set_tail Example

- Effect of `set_tail(x, y)`



Be Careful with Mutators!

- **Example:**

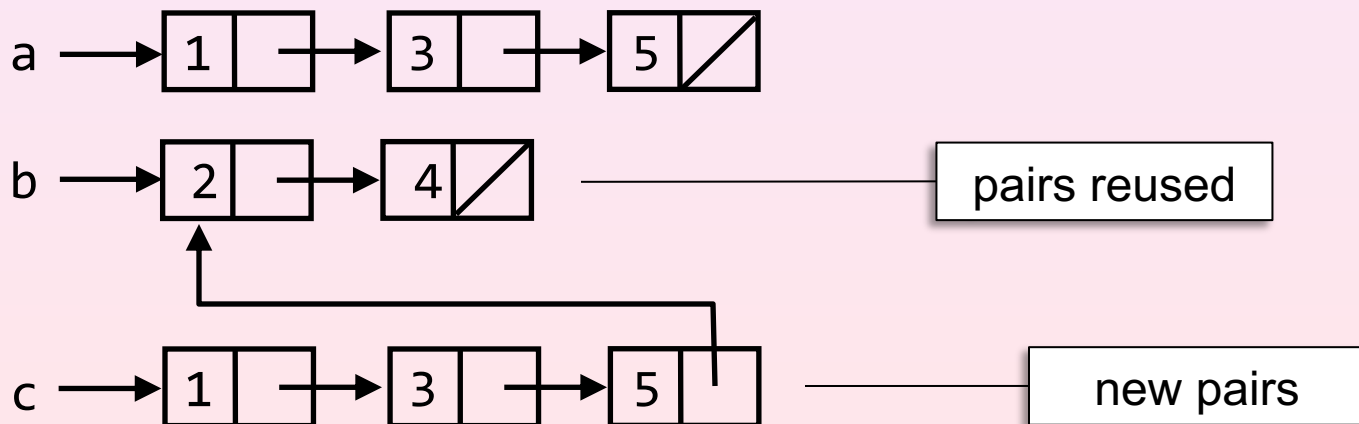
```
const a = list(1, 3, 5);  
const b = list(2, 4);  
const c = append(a, b);  
c; → [1, [3, [5, [2, [4, null]]]]]
```

```
set_head(b, 9);  
b; → [9, [4, null]]  
c; → [1, [3, [5, [9, [4, null]]]]]
```

- Mutating **b** changes **c** as well !!!
- What is happening?

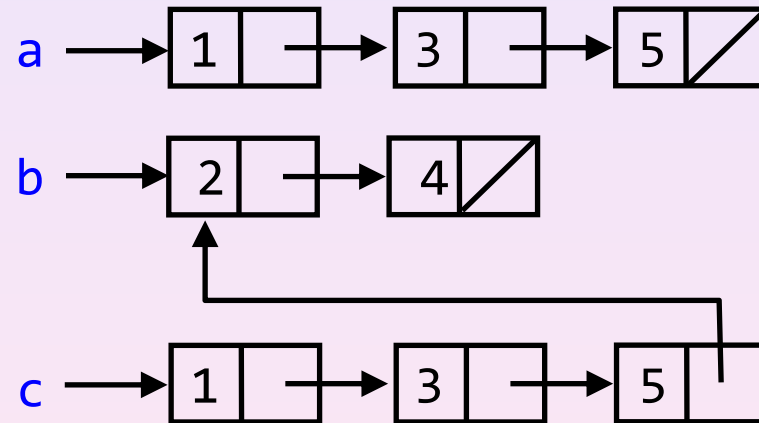
Recall the **append** function

```
function append(xs, ys) {
  return is_null(xs)
    ? ys
    : pair(head(xs), append(tail(xs), ys));
}
const a = list(1, 3, 5);
const b = list(2, 4);
const c = append(a, b);
```

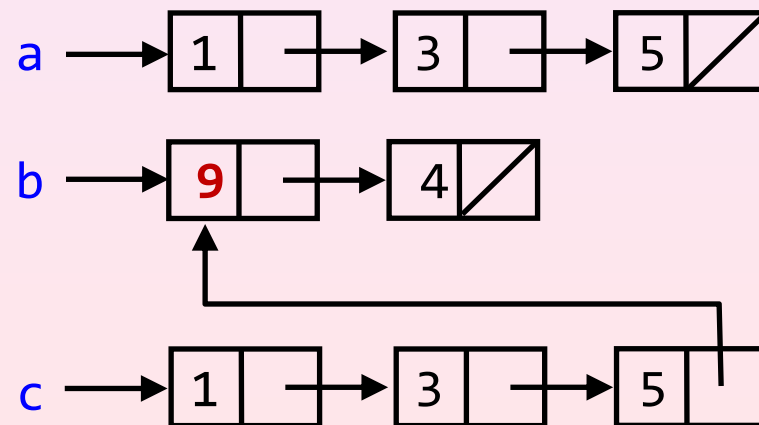


Mutation and Sharing

- Before `set_head(b, 9)`



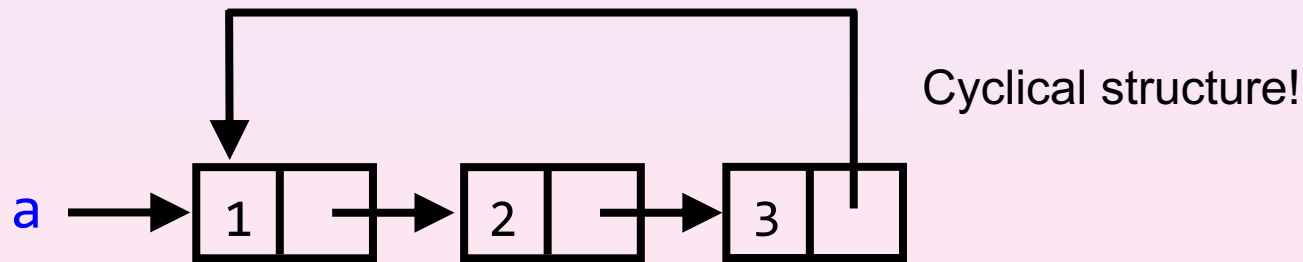
- After `set_head(b, 9)`



Be Careful with Mutators!

- Another example:

```
const a = list(1, 2, 3);  
set_tail(tail(tail(a)), a);
```



- What is `length(a)`?!

Mutable (“Destructive”) List Processing — Append

- **Wanted:**

A function to **append** two lists and return the result list

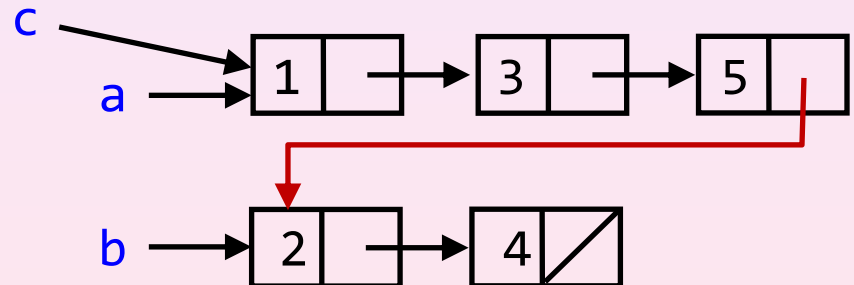
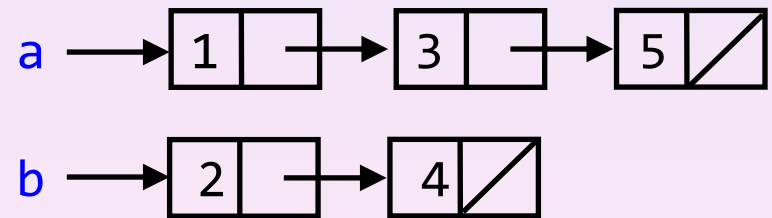
- **No new pair must be created**
- Result list is constructed from existing pairs of input lists

“Destructive” Append

- Example:

```
const a = list(1, 3, 5);
const b = list(2, 4);
```

```
const c = d_append(a, b);
```



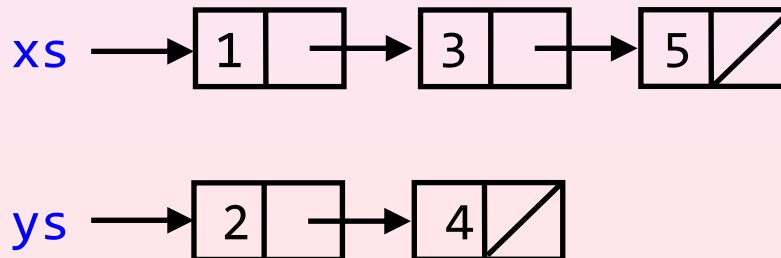
```
c; ➔ [1, [3, [5, [2, [4, null]]]]]
a; ➔ [1, [3, [5, [2, [4, null]]]]]
b; ➔ [2, [4, null]]
```

“Destructive” Append

- Implementation:

```
function d_append(xs, ys) {  
  if (is_null(xs)) {  
    return ys;  
  } else {  
    set_tail(xs, d_append(tail(xs), ys));  
    return xs;  
  }  
}
```

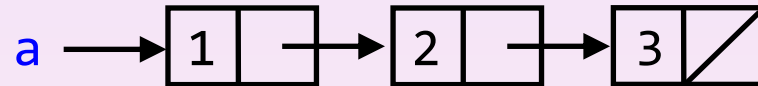
[Show in
Playground](#)



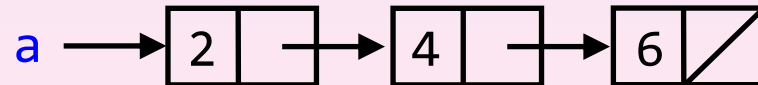
“Destructive” Map

- Example:

```
const a = list(1, 2, 3);
```



```
d_map(x => x * 2, a);
```



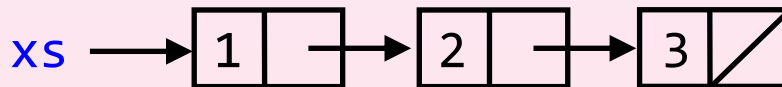
```
a; → [2, [4, [6, []]]]
```


“Destructive” Map

- Implementation:

```
function d_map(fun, xs) {  
  if (!is_null(xs)) {  
    set_head(xs, fun(head(xs)));  
    d_map(fun, tail(xs));  
  } else { }  
}
```

[Show in
Playground](#)



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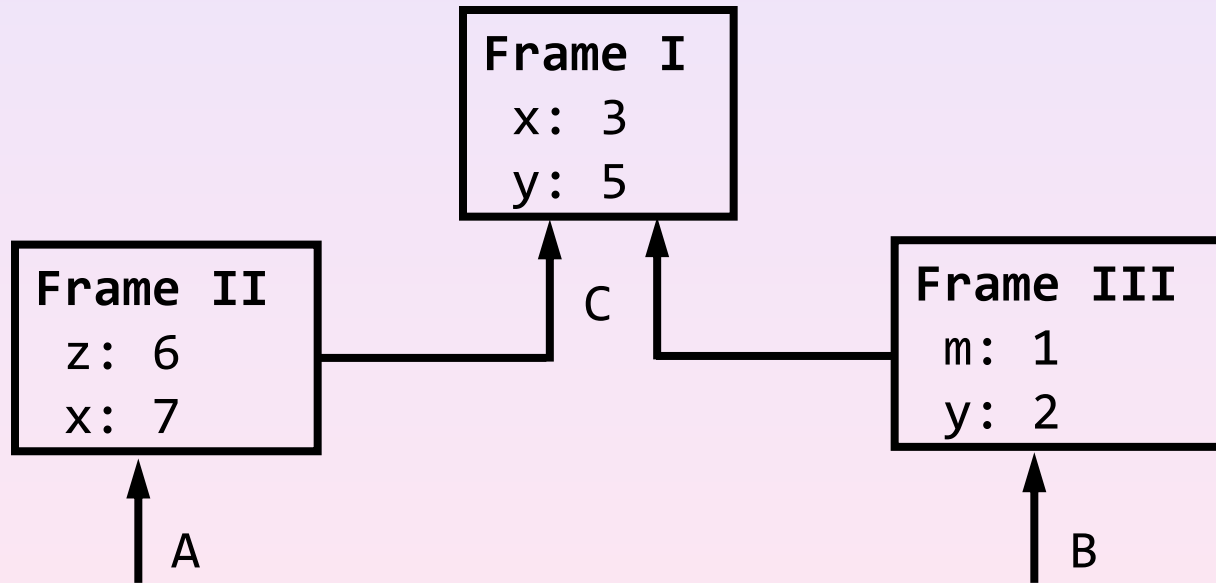
Environment

- An ***environment*** determines the context in which an expression should be evaluated
- Every **expression** is evaluated with respect to the **current** environment

Environments

- An **environment** is a **sequence of frames**
- Each frame contains **bindings** of values to names
- A frame points to its **enclosing environment**, the next one in the sequence
 - Unless the frame is the **global frame**
- **Extending** an environment means **adding** a new frame “inside” the old one

Example Environments



- A, B, C are environments
- Value of **x** in environment **A** is **7**
- Value of **x** in environment **B** is **3**
- **x** in Frame II **shadows** **x** in Frame I
- Value of **m** in environment **A**? **unbound**

Accessing an Environment

- To evaluate a **name**, look up its value in the **current frame**
- If **not found in current frame**, then look in **enclosing environment**, and so on
- If **still not found**, then the name is said to be ***unbound*** in the environment

Global Environment

- The **global environment** consists of a single **frame** with bindings of **primitive** and **pre-declared** functions and constants
 - Functions' names are bound to their implementations

global
environment



```
error:= <implementation>
display:= <implementation>
pair:= <implementation>
list:= <implementation>
math_PI:= 3.14...
...
```

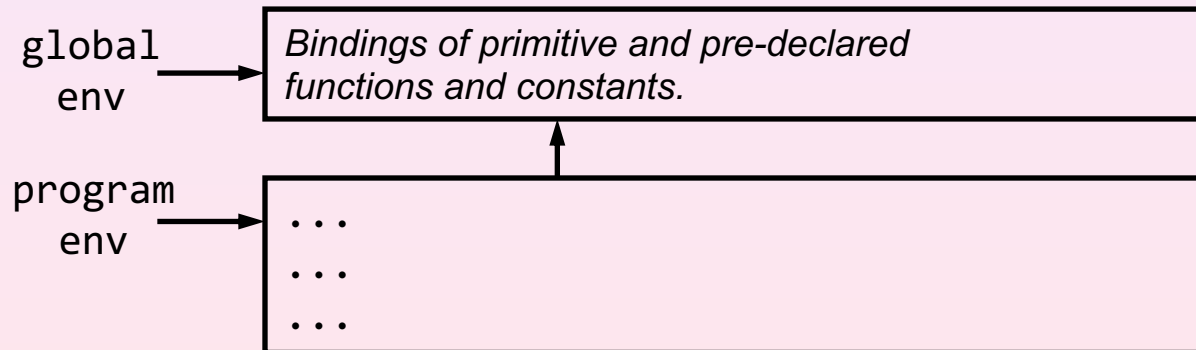
use **:=** for constants, and **:** for variables

Evaluation of Blocks

- Evaluating a **block** { *statements* } extends the environment by adding a **new frame**
 - New frame contains bindings of the **constants** and **variables** declared in the block
 - Initially, these **constants** and **variables** don't have associated values (**unassigned**)
 - When the **declarations** are evaluated, the **names** are **assigned** their values
- New **frame is not created** if the block has **no constant & variable declaration**
 - There are no empty frames

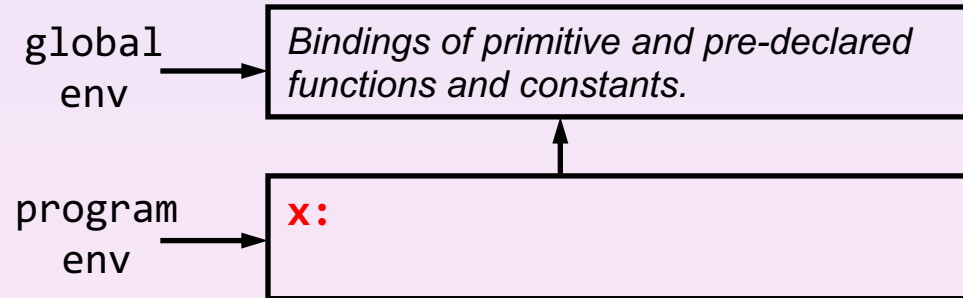
Evaluation of Blocks

- **User program** is considered in an implicit **program block**
 - It directly extends the **global environment**
 - It is called the **program environment**



Example for Evaluating Blocks

```
let x = 3;  
  
{  
  const x = 5;  
  let y = 8;  
}
```



Example for Evaluating Blocks

```
let x = 3;
```

```
{
```

```
  const x = 5;
```

```
  let y = 8;
```

```
}
```

global
env

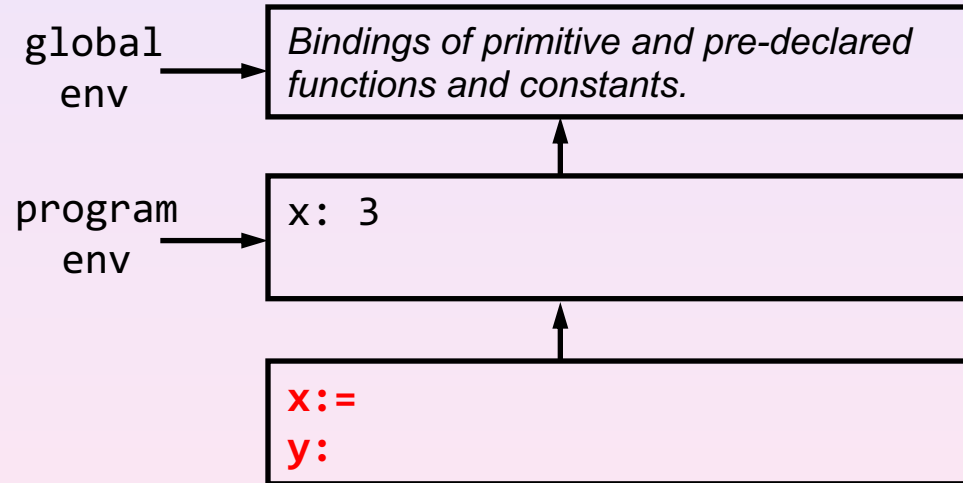
*Bindings of primitive and pre-declared
functions and constants.*

program
env

x: 3

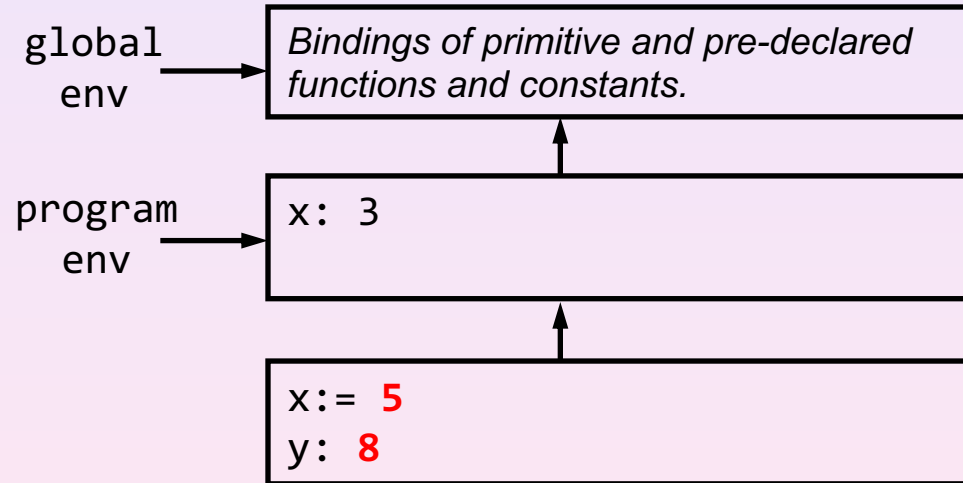
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Example for Evaluating Blocks

```
let x = 3;  
  
{  
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  let y = 8;  
}
```



Evaluation of Declarations and Assignments

- **Constant / variable** declaration statements:

const name = value;

let name = value;

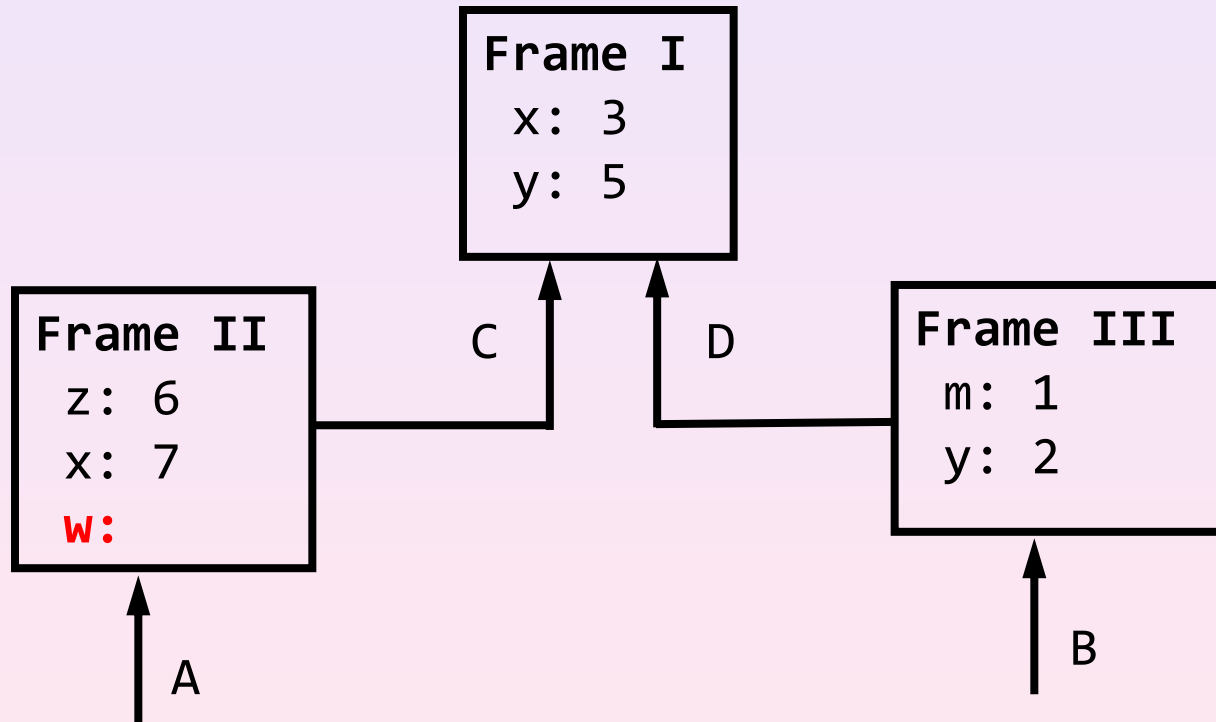
- Changes binding of *name* to *value* in the **current frame**

- **Assignment** statement:

name = new_value;

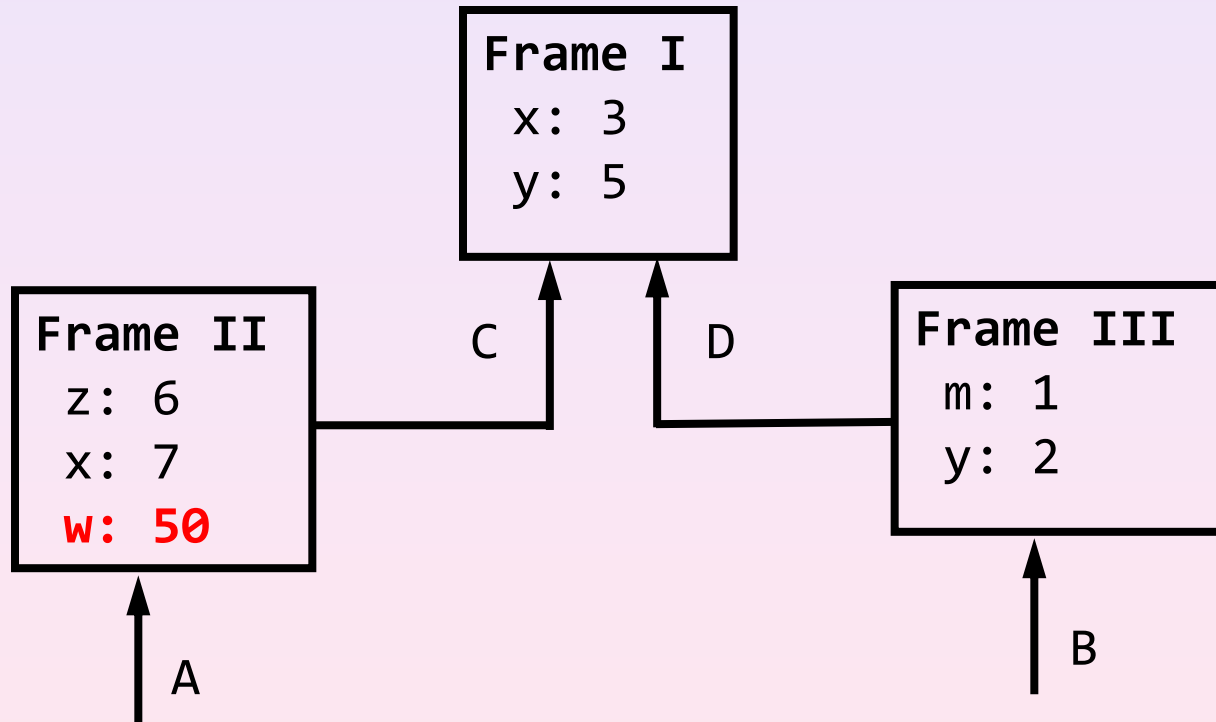
- Searches for *name* beginning from **current frame**, then changes its binding to *new_value*
 - If *name* is a **constant**, signal an error
 - If *name* is **unbound** or **unassigned**, signal an error

Example Environments



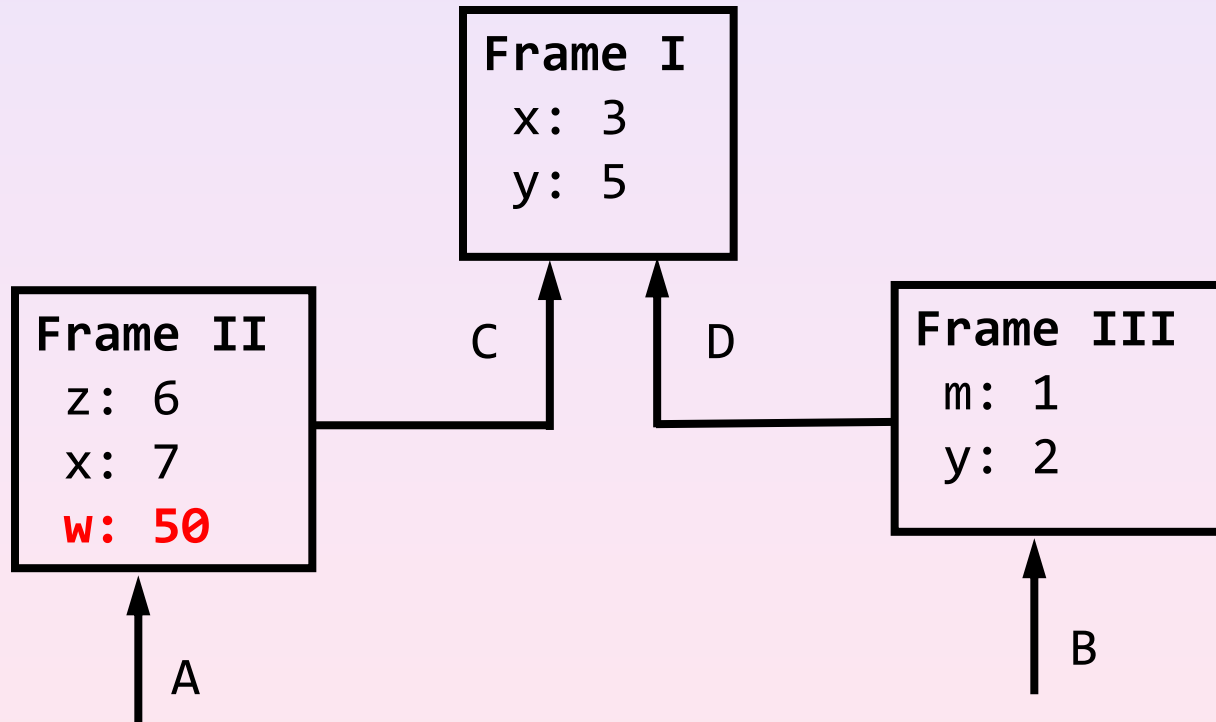
- Evaluate in **A**: `let w = 50;`

Example Environments



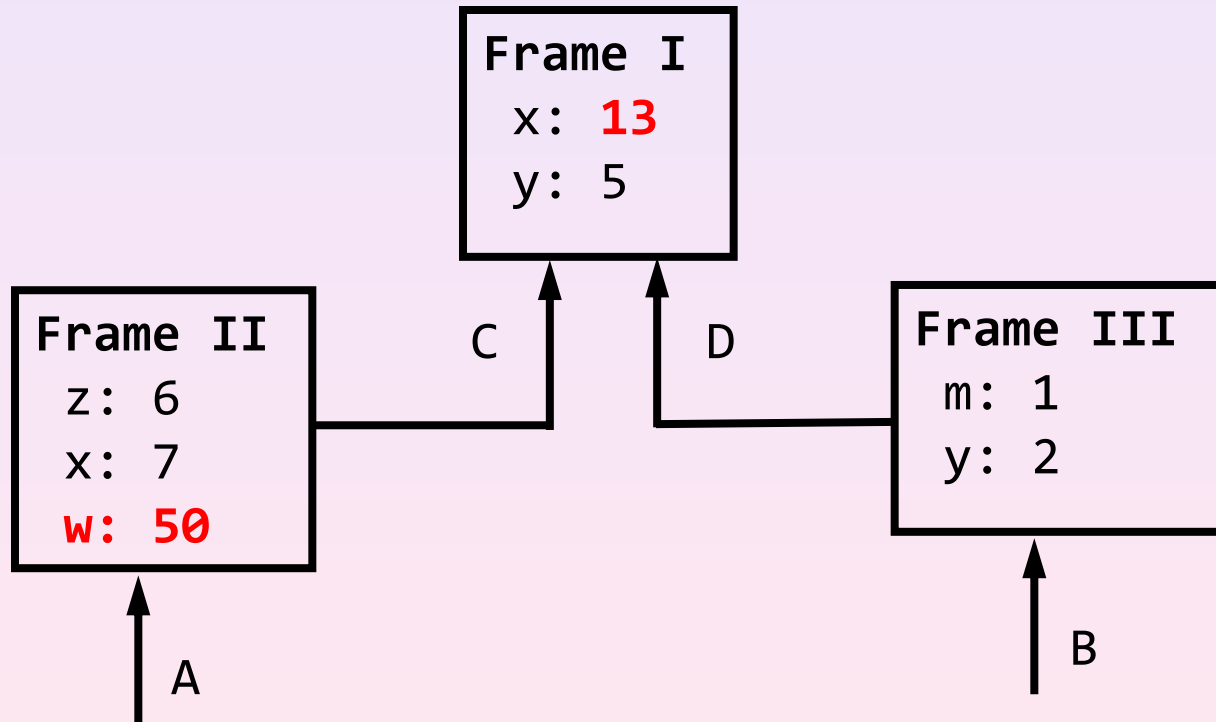
- Evaluate in **A**: `let w = 50;`
 - `w` in Frame II is assigned with the value **50**

Example Environments



- Evaluate in **B**: $x = x + 10$;

Example Environments



- Evaluate in **B**: $x = x + 10$;
 - The value of x in Frame I is updated from 3 to 13

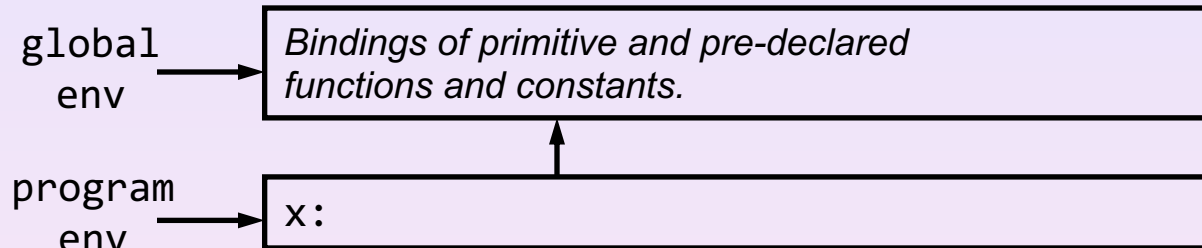
Block Example

```
let x = 3;

if (x <= 1) {
  let y = x;
  x = y + 1;
} else {
  let z = x;
  x = z * z;
}
{
  x = x + 1;
}
{
  let x = 5;
  let y = 8;
  {
    let x = 20 + y;
  }
  x = x + 1;
}
x;
```

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Playground](#)

Environment Model

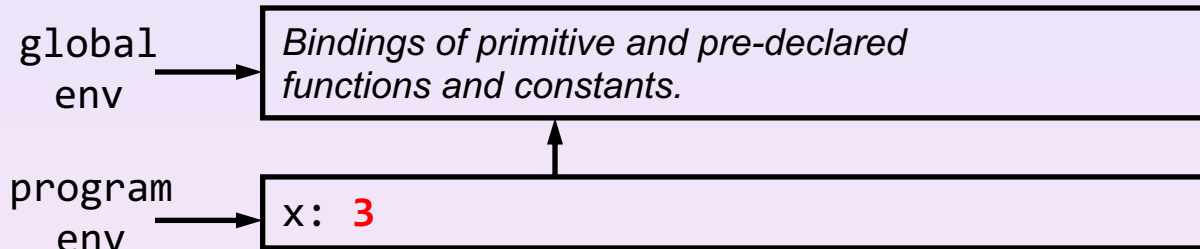


```
let x = 3;

if (x <= 1) {
    let y = x;
    x = y + 1;
} else {
    let z = x;
    x = z * z;
}

{
    x = x + 1;
}

{
    let x = 5;
    let y = 8;
    {
        let x = 20 + y;
    }
    x = x + 1;
}
x;
```



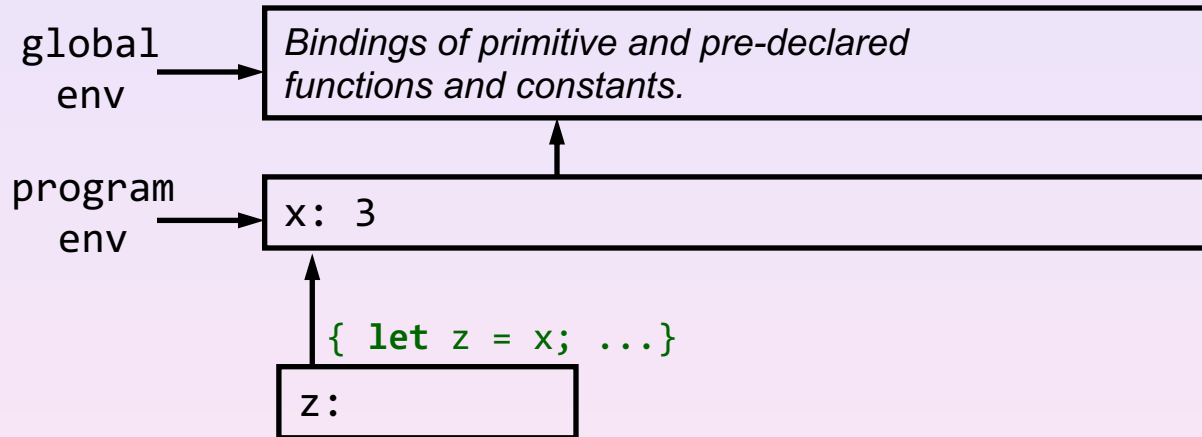
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{
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}

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    let y = 8;
    {
        let x = 20 + y;
    }
    x = x + 1;
}
x;
    
```



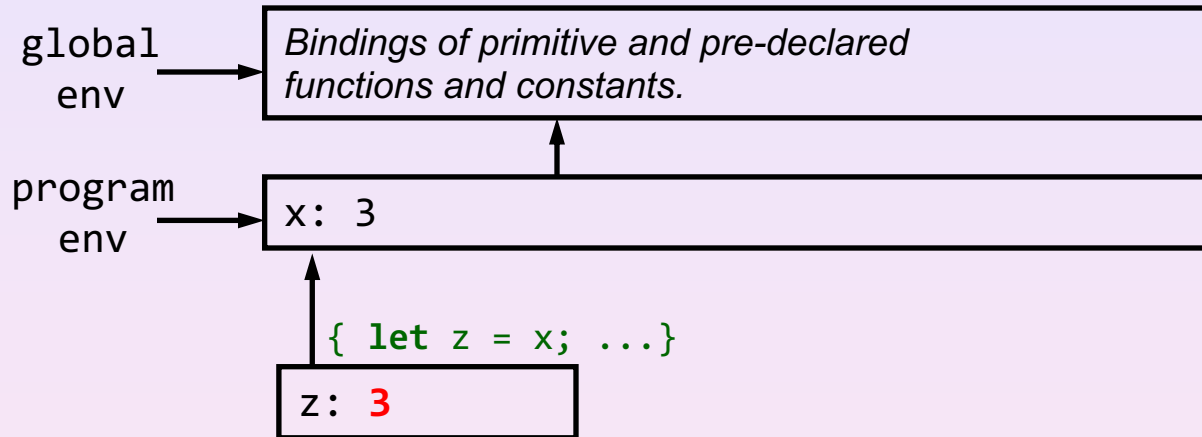
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}
x;
    
```



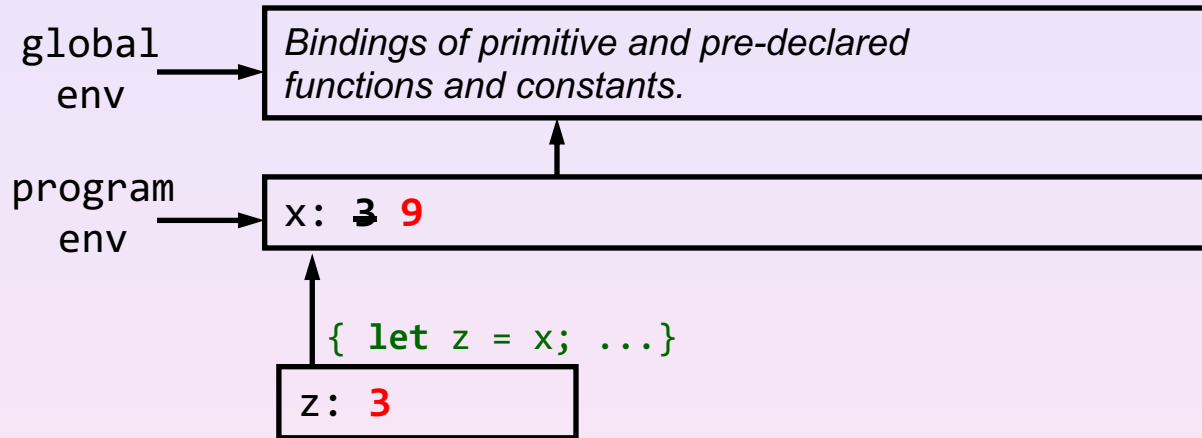
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    let y = 8;
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        let x = 20 + y;
    }
    x = x + 1;
}
x;
    
```



```

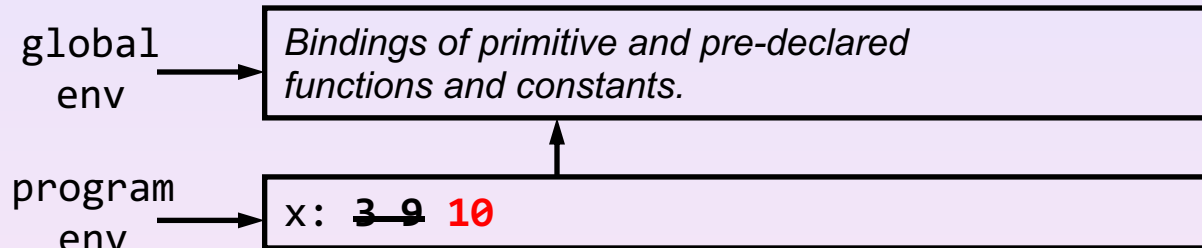
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        let x = 20 + y;
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}
x;
    
```


Environment Model



```
let x = 3;

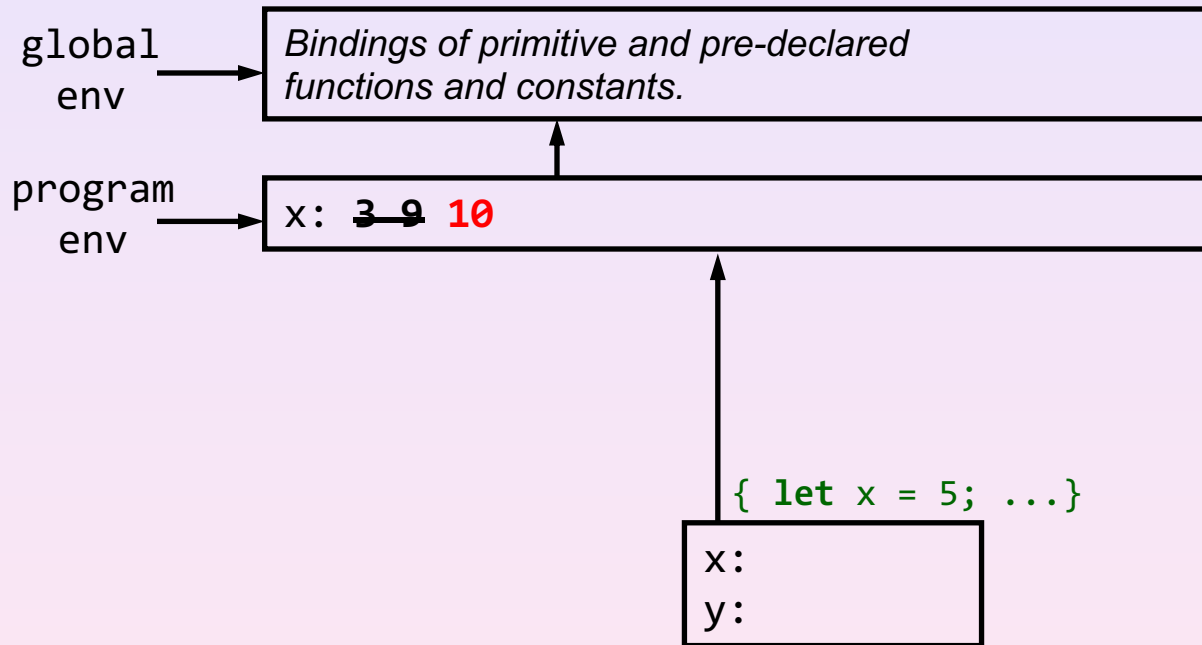
if (x <= 1) {
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}

x;
```

Environment Model



```
let x = 3;

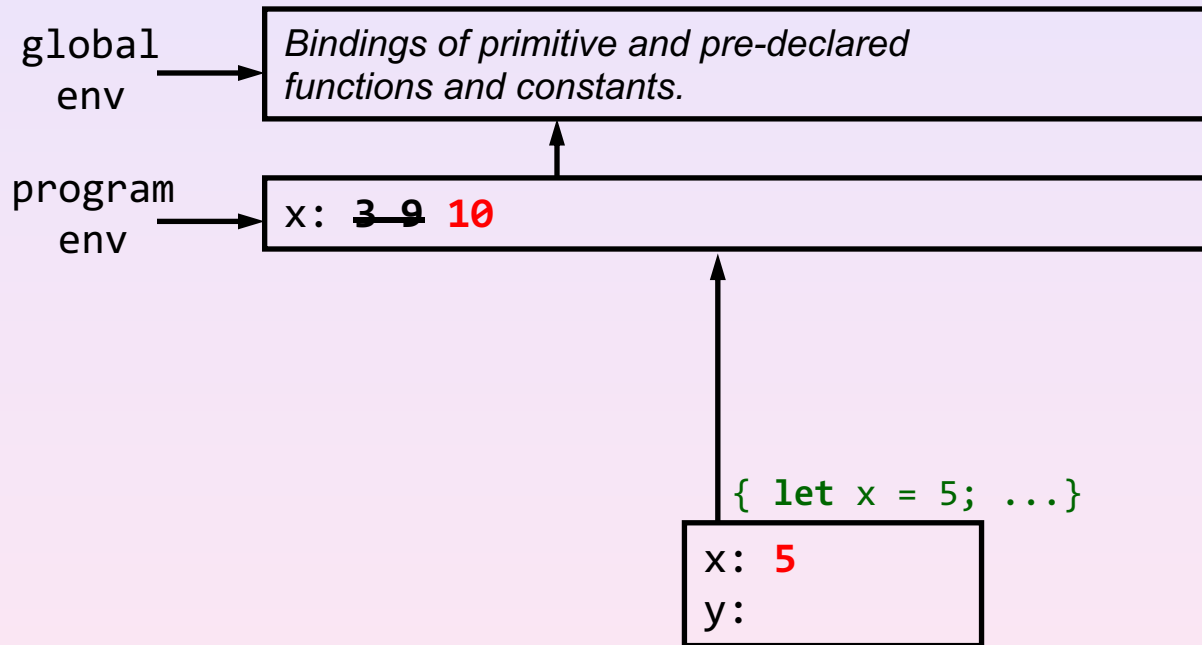
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}

x;
```

Environment Model



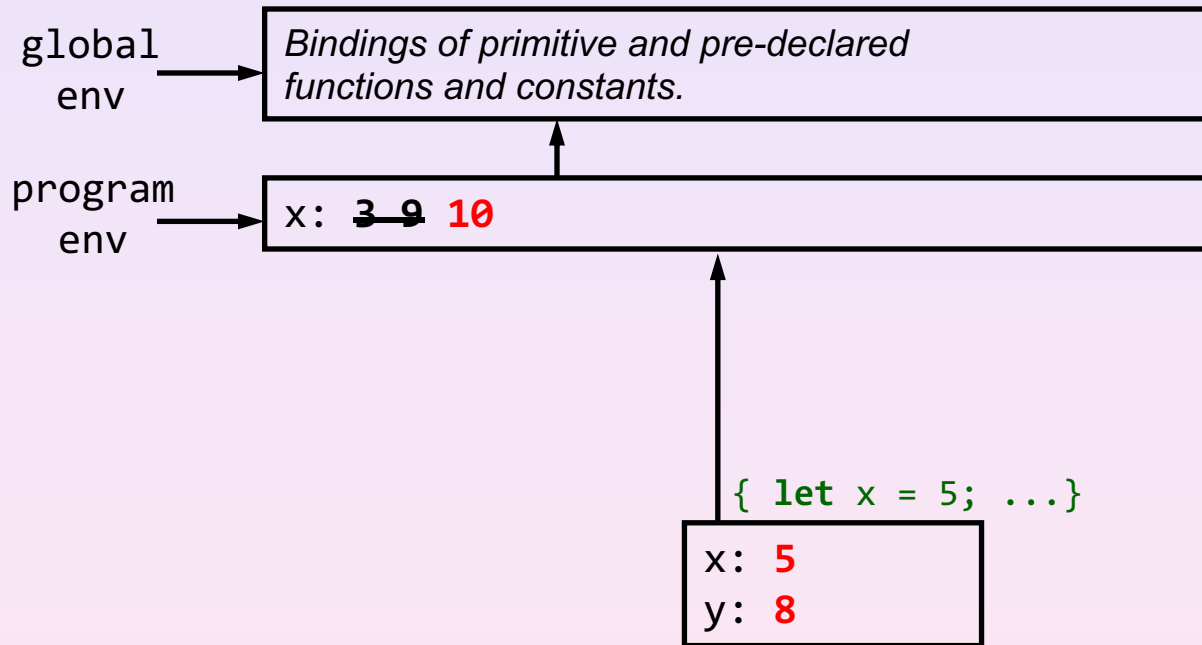
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    let x = 20 + y;
}
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}
x;
```

Environment Model



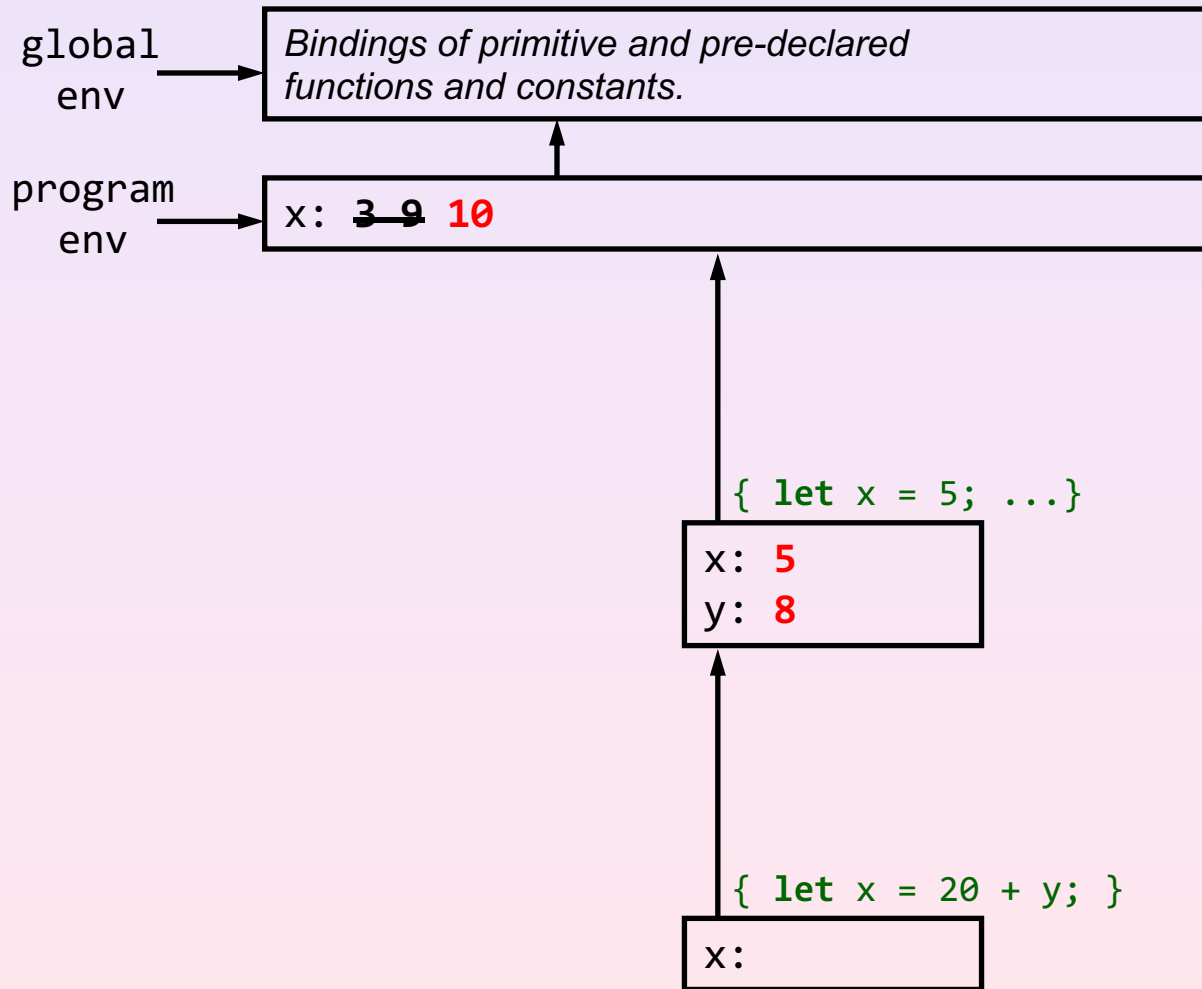
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Environment Model



```

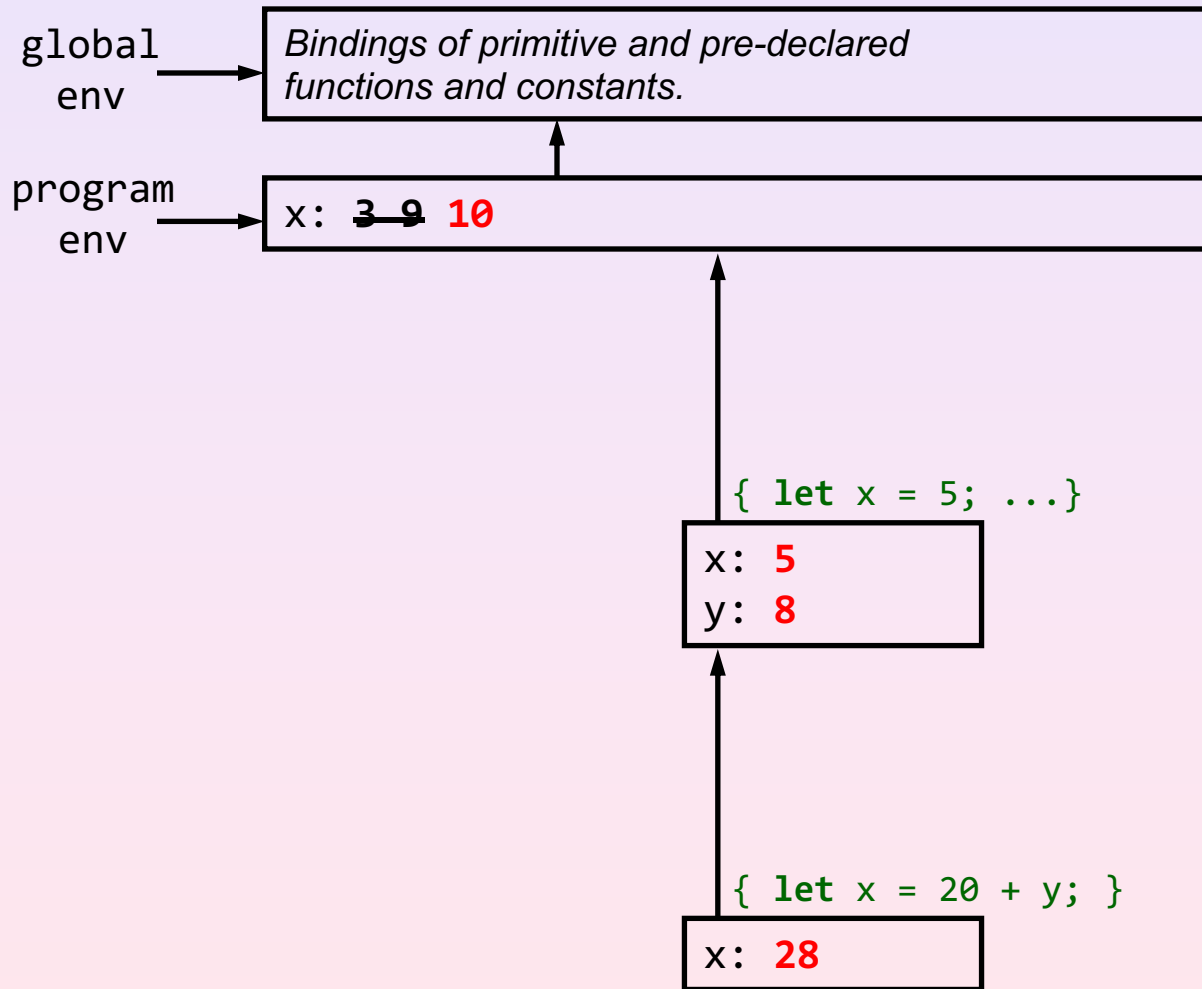
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Environment Model



```

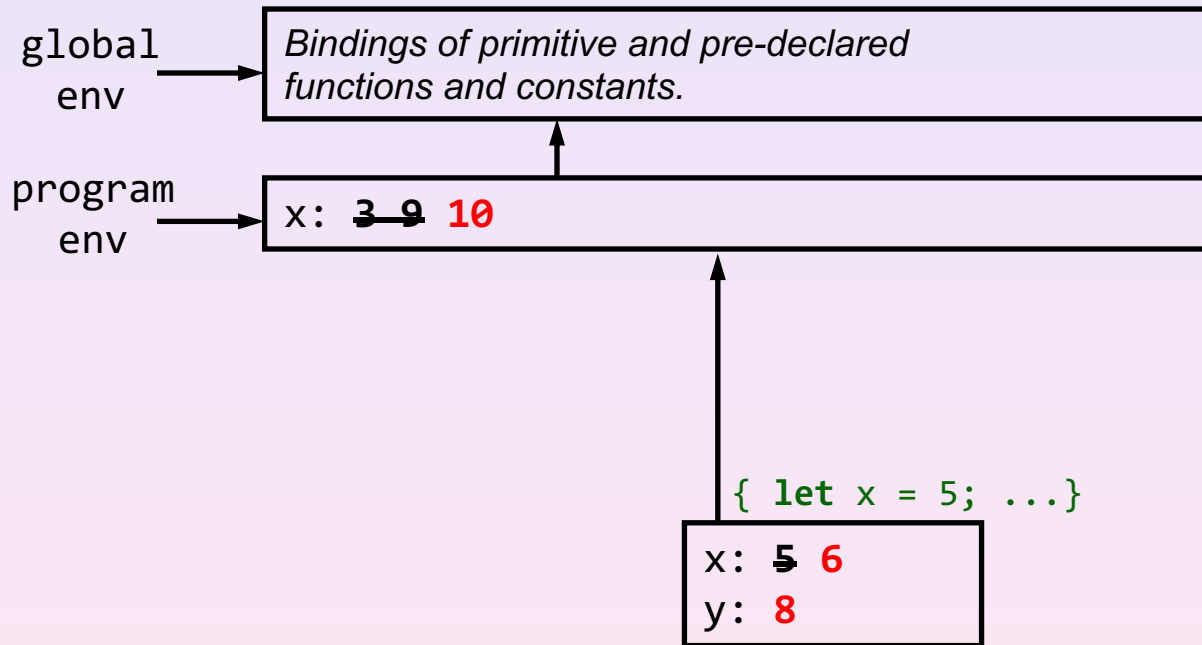
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}

x = x + 1;

let x = 5;
let y = 8;
{
    let x = 20 + y;
}
x = x + 1;
x;
    
```

Environment Model



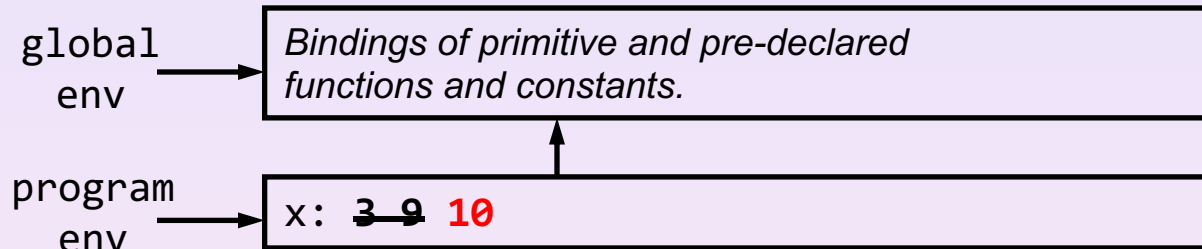
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}

{
    x = x + 1;
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    let y = 8;
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        let x = 20 + y;
    }
    x = x + 1;
}
x;
```

Environment Model



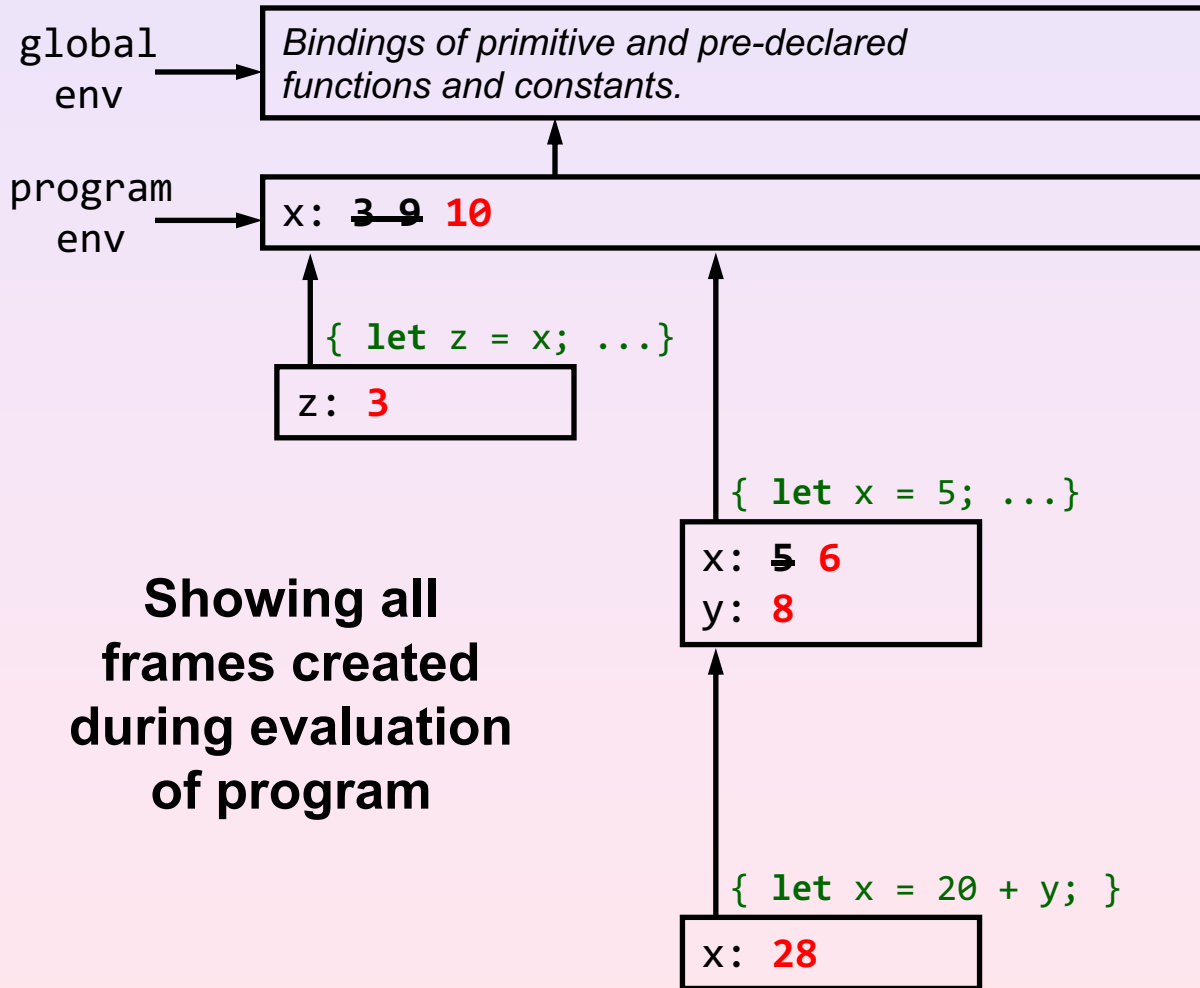
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} else {
    let z = x;
    x = z * z;
}

{
    x = x + 1;
}

{
    let x = 5;
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        let x = 20 + y;
    }
    x = x + 1;
}

x;
```

```

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}

{
    let x = 5;
    let y = 8;
    {
        let x = 20 + y;
    }
    x = x + 1;
}
x;
    
```

Evaluation of Function Applications

- Every **function application** extends the environment in which the *function was created*
- The **new frame** contains bindings of **parameter variables** to actual arguments
 - New **frame is not created** if the function has **no parameter**
- The **function body block** is evaluated in this new environment

Evaluation of Lambda Expressions

- Lambda expression:

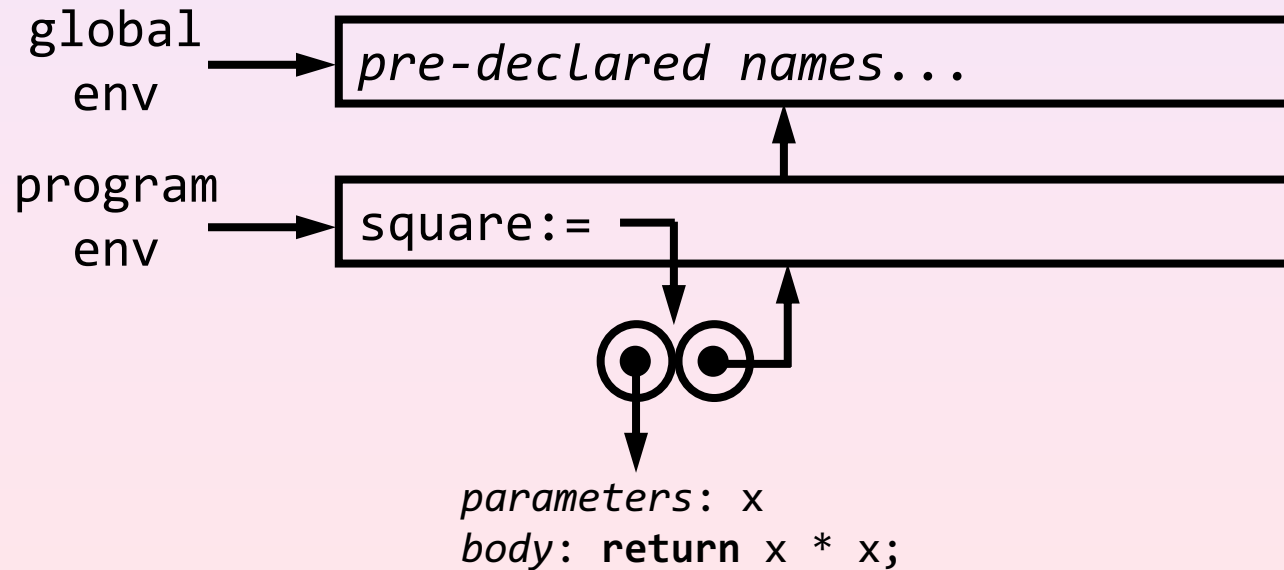
$(var1, var2, \dots) \Rightarrow \{ body \}$

- Creates a **function object**, represented as **two circles**
- One points to *body* (text)
- The other points to the environment **in which the function expression was evaluated**
- A function therefore remembers the **environment in which it was created**

Example for Evaluating Lambda Expressions

```
const square = x => x * x;
square(5);
```

- After evaluating the constant declaration in **program env**:



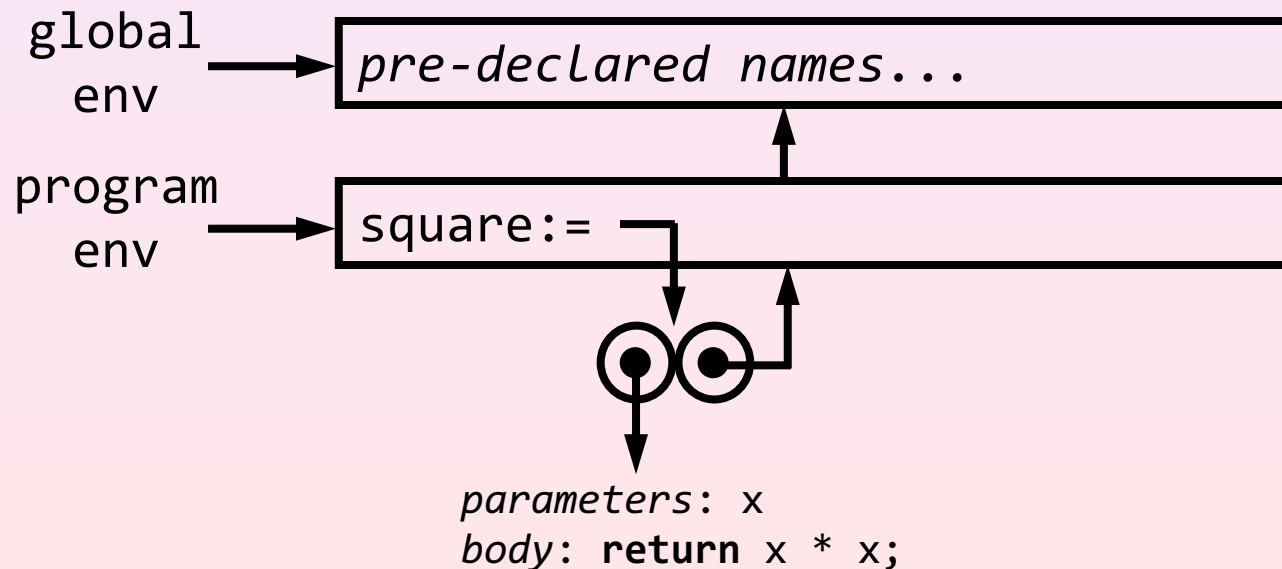
Example for Evaluating Function Declarations

```
function square(x) { return x * x; }
```

which is equivalent to*

```
const square = x => x * x;
```

- After evaluating the function declaration in **program env**:



Evaluation of Function Applications

- To evaluate a **function application**:

fun(arg1, arg2, ..., argn);

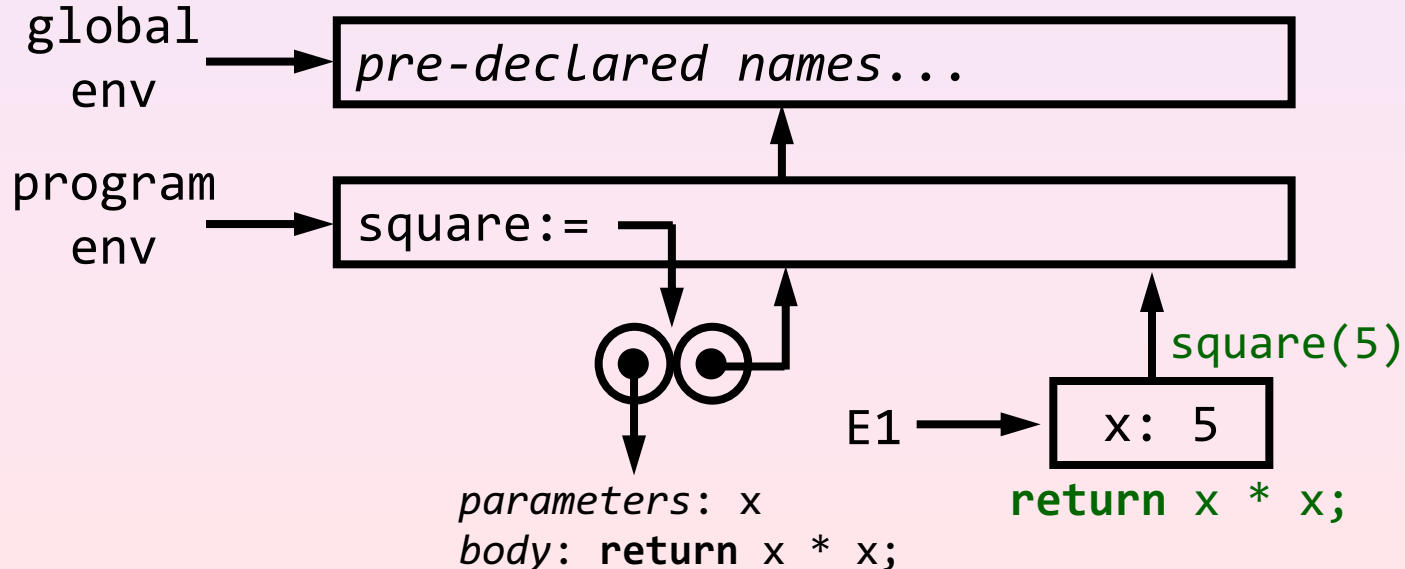
- Evaluate the **subexpressions** in the **current environment**
- Apply the value of the **function subexpression** (*fun*) to the values of the **argument subexpressions** (*arg1, arg2, ..., argn*)

Applying a Function Value to Argument Values

- Create a new frame that **points to the environment of the function**
 - **The environment of the function** is the environment in which the function was created
- In this new frame, **bind the formal parameters to the actual arguments**
- Evaluate the **body block** of the function in the new environment

Example: `square(5);`

- Evaluate subexpressions `square` and `5` in **program env**
- Create new frame that points to environment of `square`
- In this frame, bind value 5 to parameter `x`
- Evaluate body `return x * x;`



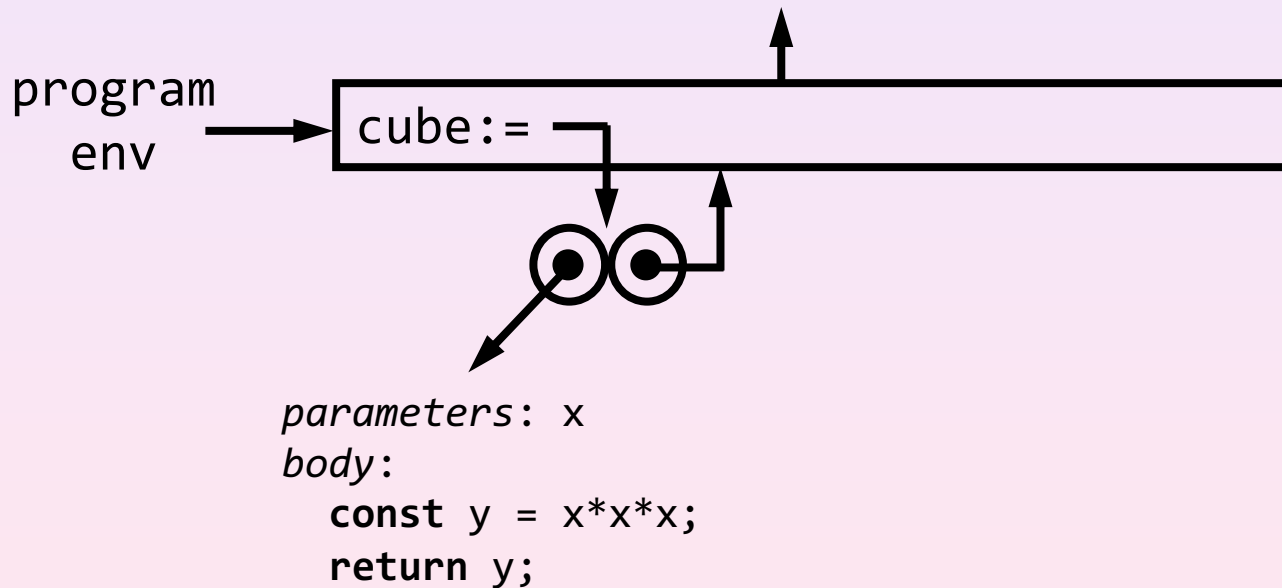
Example with Local Declaration

- The `cube` function has a **local constant** `y`

```
function cube(x) {  
    const y = x * x * x;  
    return y;  
}  
cube(3); → 27
```

cube Example

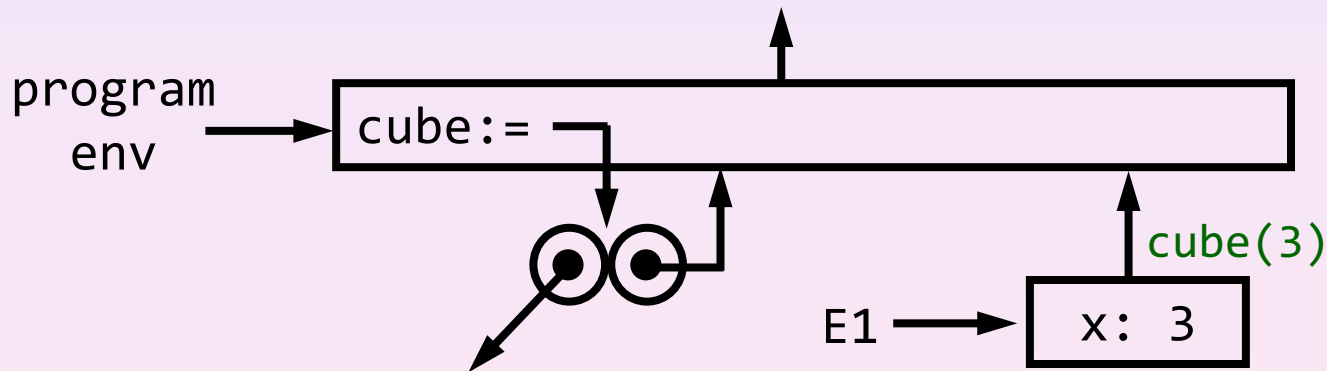
- After evaluating the declaration of function `cube`



```
function cube(x) {
  const y = x * x * x;
  return y;
}
cube(3);
```

cube Example

- Evaluate `cube(3)` in **program env**

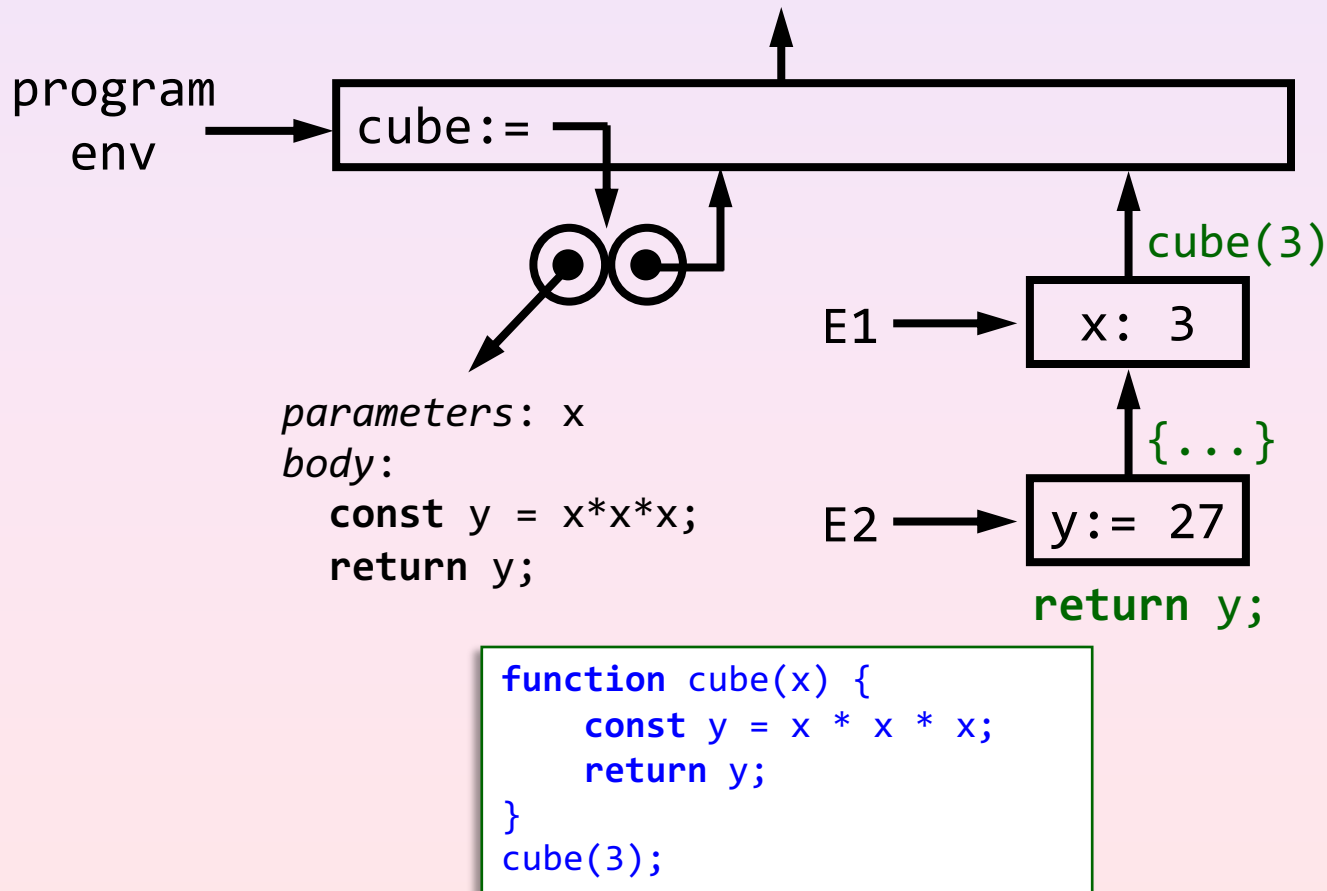


parameters: x
body:
 const y = x*x*x;
 return y;

```
function cube(x) {
  const y = x * x * x;
  return y;
}
cube(3);
```

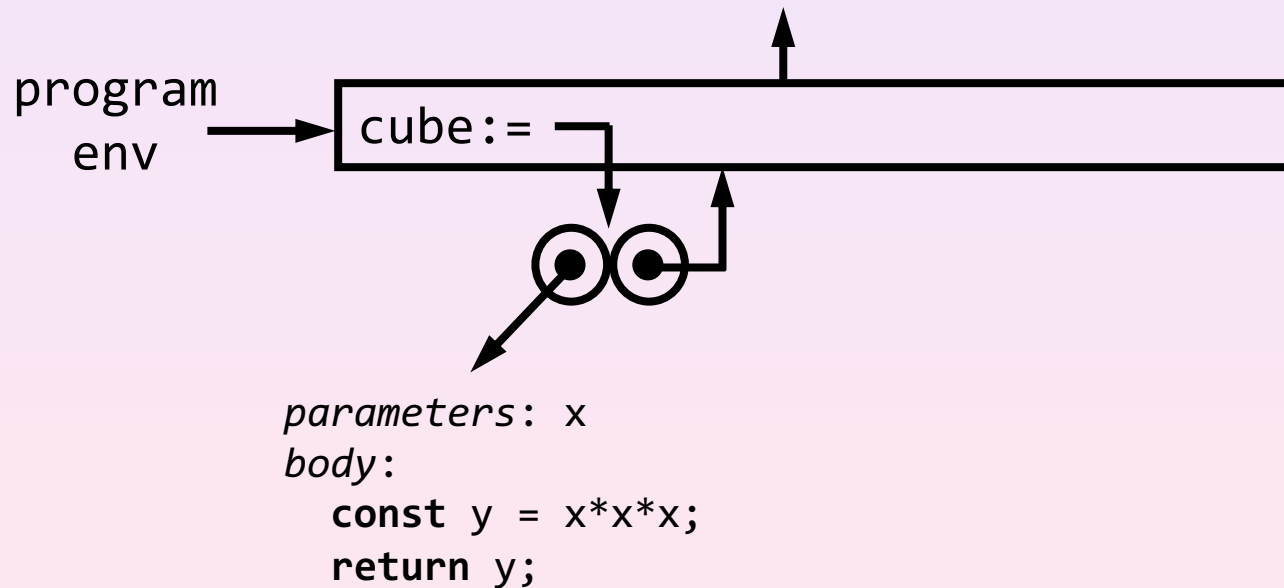
cube Example

- Evaluate function **body block**



cube Example

- Environment after evaluation of `cube(3)`



```
function cube(x) {
  const y = x * x * x;
  return y;
}
cube(3);
```

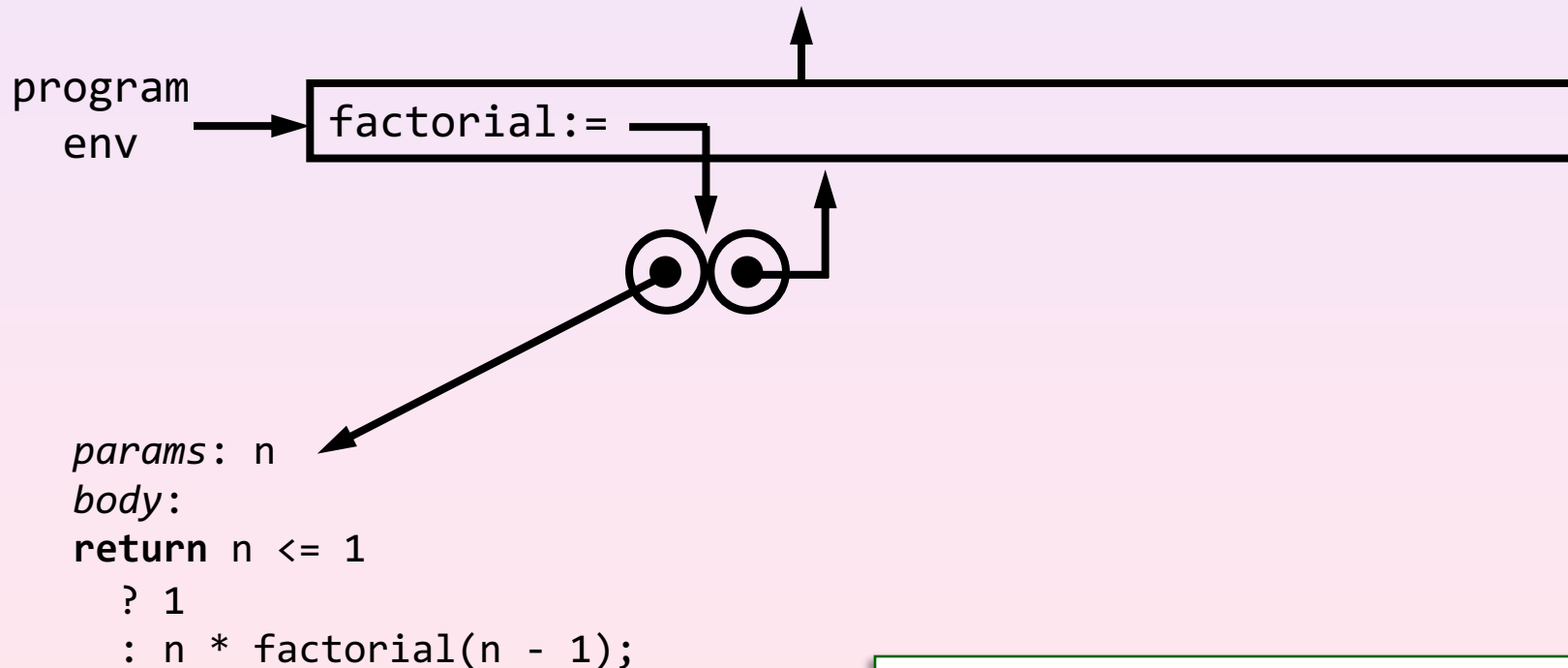
Example with Recursion

- Recursive `factorial` function

```
function factorial(n) {  
    return n <= 1  
        ? 1  
        : n * factorial(n - 1);  
}  
factorial(2); → 2
```

factorial Example

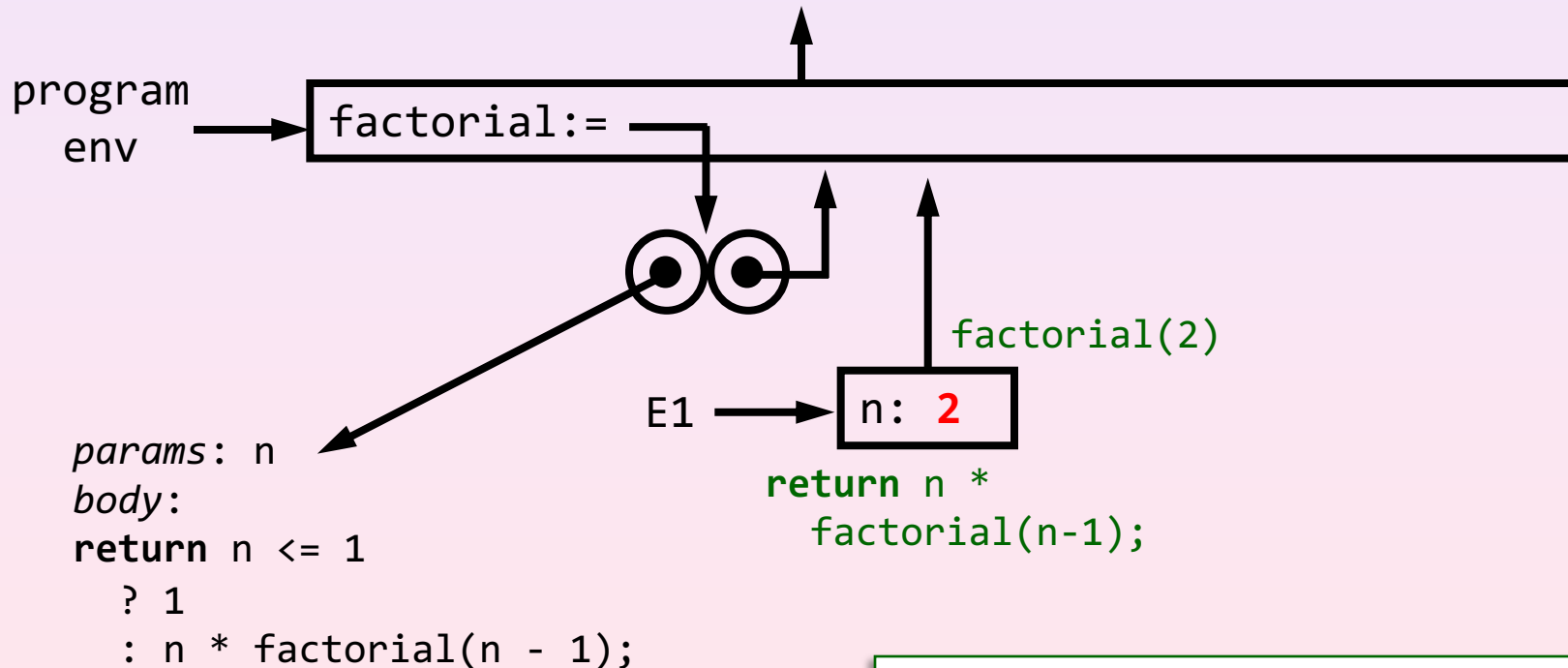
- After evaluating the declaration of function `factorial`



```
function factorial(n) {
  return n <= 1 ? 1 : n * factorial(n - 1);
}
factorial(2);
```

factorial Example

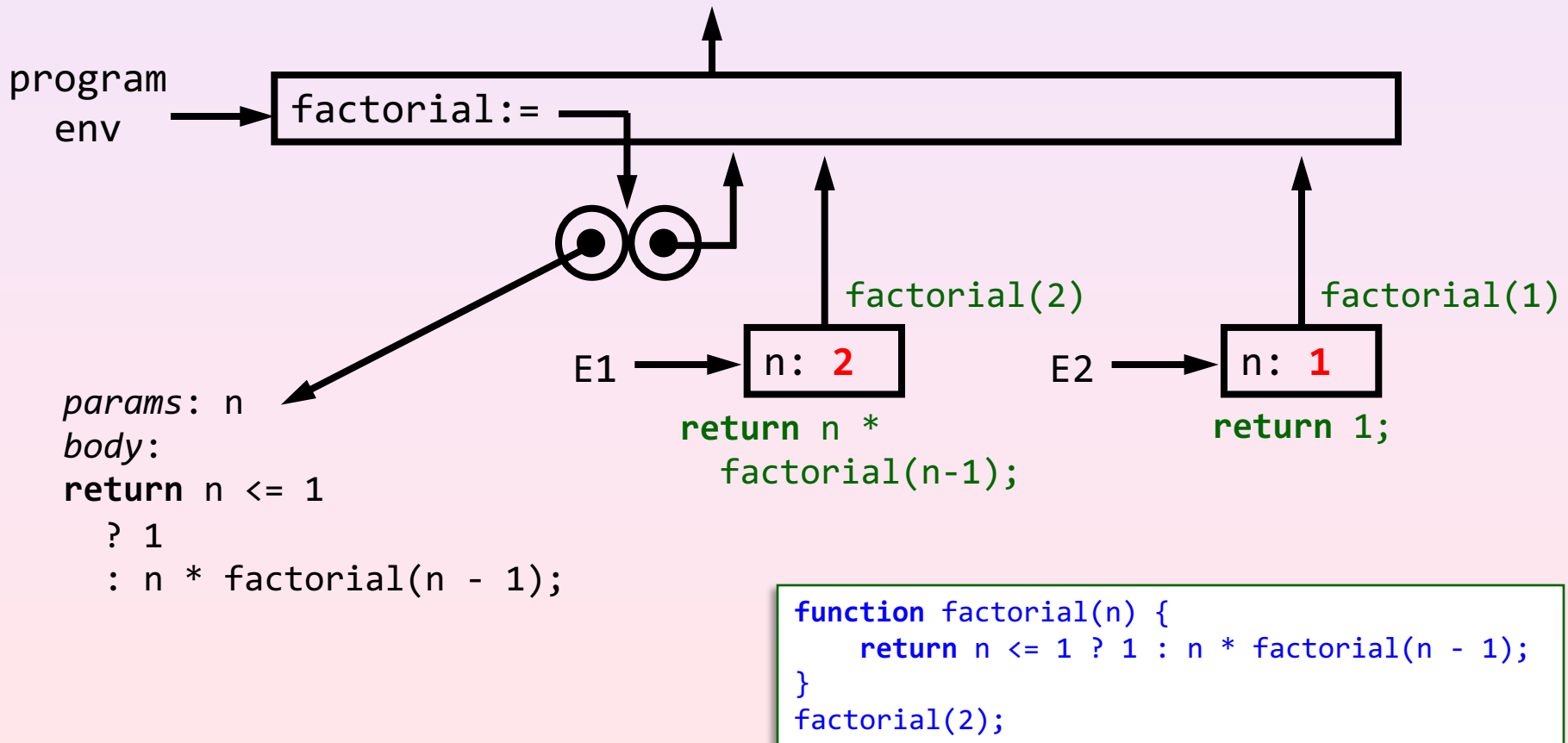
- Evaluate `factorial(2)` in **program env**



```
function factorial(n) {
    return n <= 1 ? 1 : n * factorial(n - 1);
}
factorial(2);
```

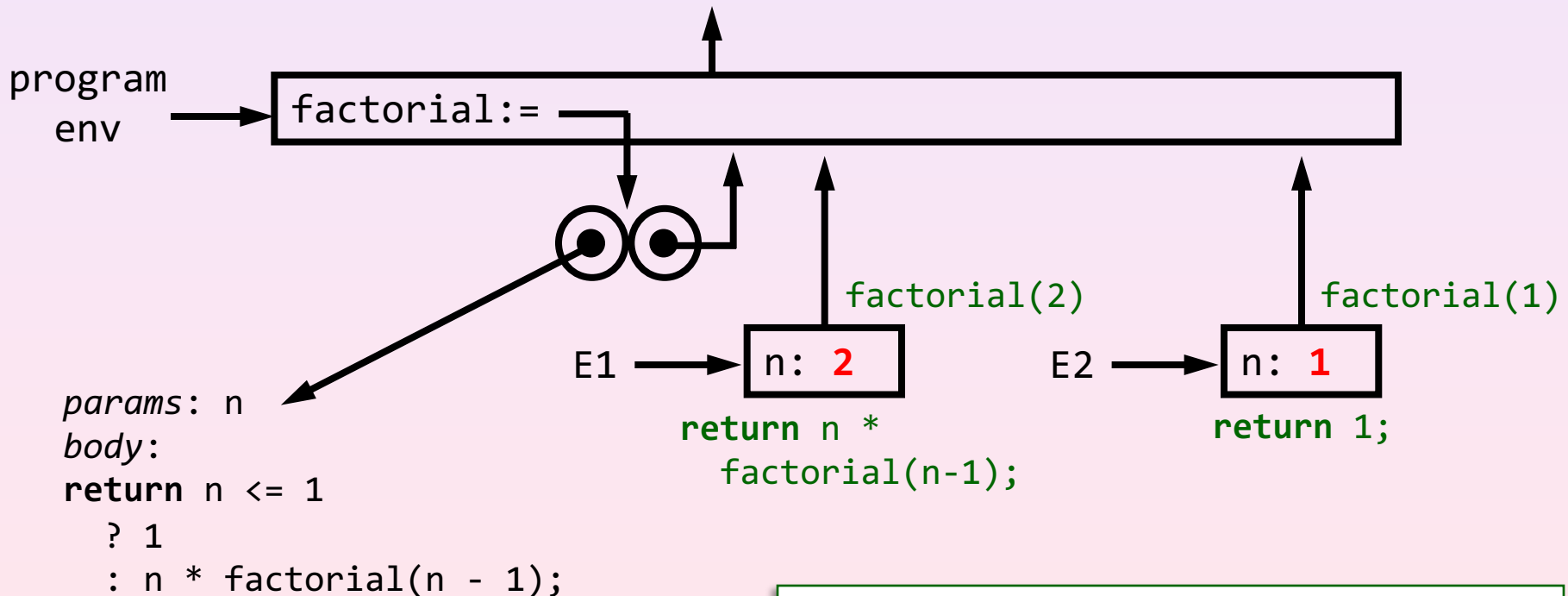

factorial Example

- Evaluate `factorial(n-1)` in **Environment E1**



factorial Example

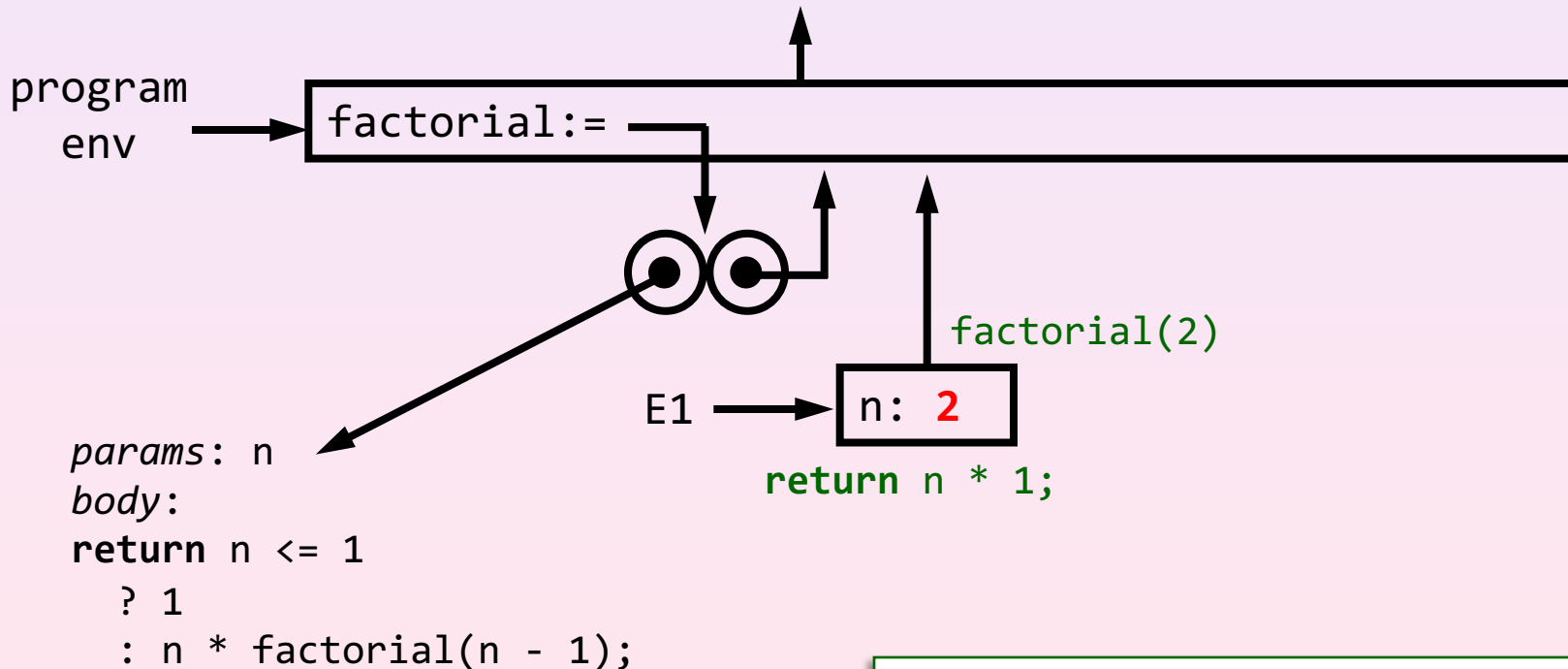
- Evaluate `return 1;` in **Environment E2**



```
function factorial(n) {
  return n <= 1 ? 1 : n * factorial(n - 1);
}
factorial(2);
```

factorial Example

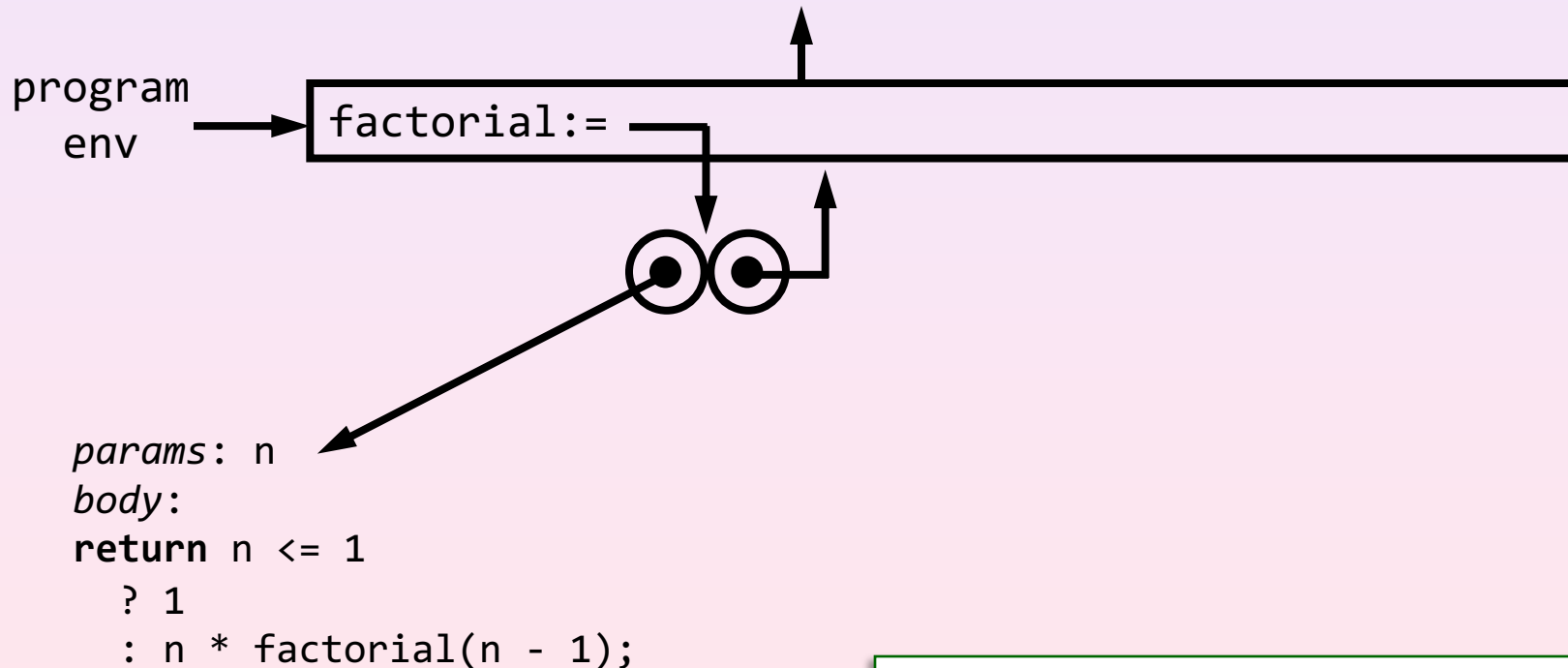
- Evaluate `return n * 1;` in **Environment E1**



```
function factorial(n) {
    return n <= 1 ? 1 : n * factorial(n - 1);
}
factorial(2);
```

factorial Example

- Environment after evaluation of `factorial(2)`



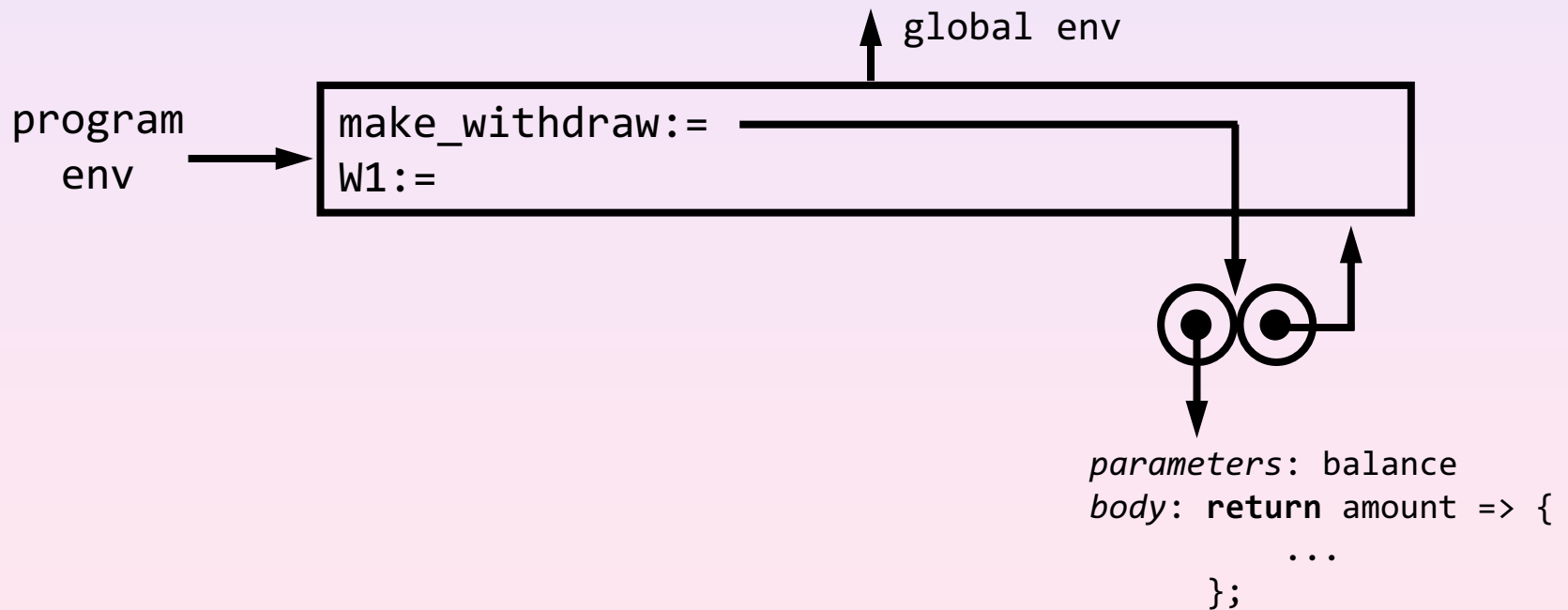
```
function factorial(n) {
    return n <= 1 ? 1 : n * factorial(n - 1);
}
factorial(2);
```

make_withdraw Example

```
function make_withdraw(balance) {  
  return amount => {  
    if (balance >= amount) {  
      balance = balance - amount;  
      return balance;  
    } else {  
      return "Insufficient funds";  
    }  
  };  
}  
  
const w1 = make_withdraw(100);  
w1(40); ➔ 60
```

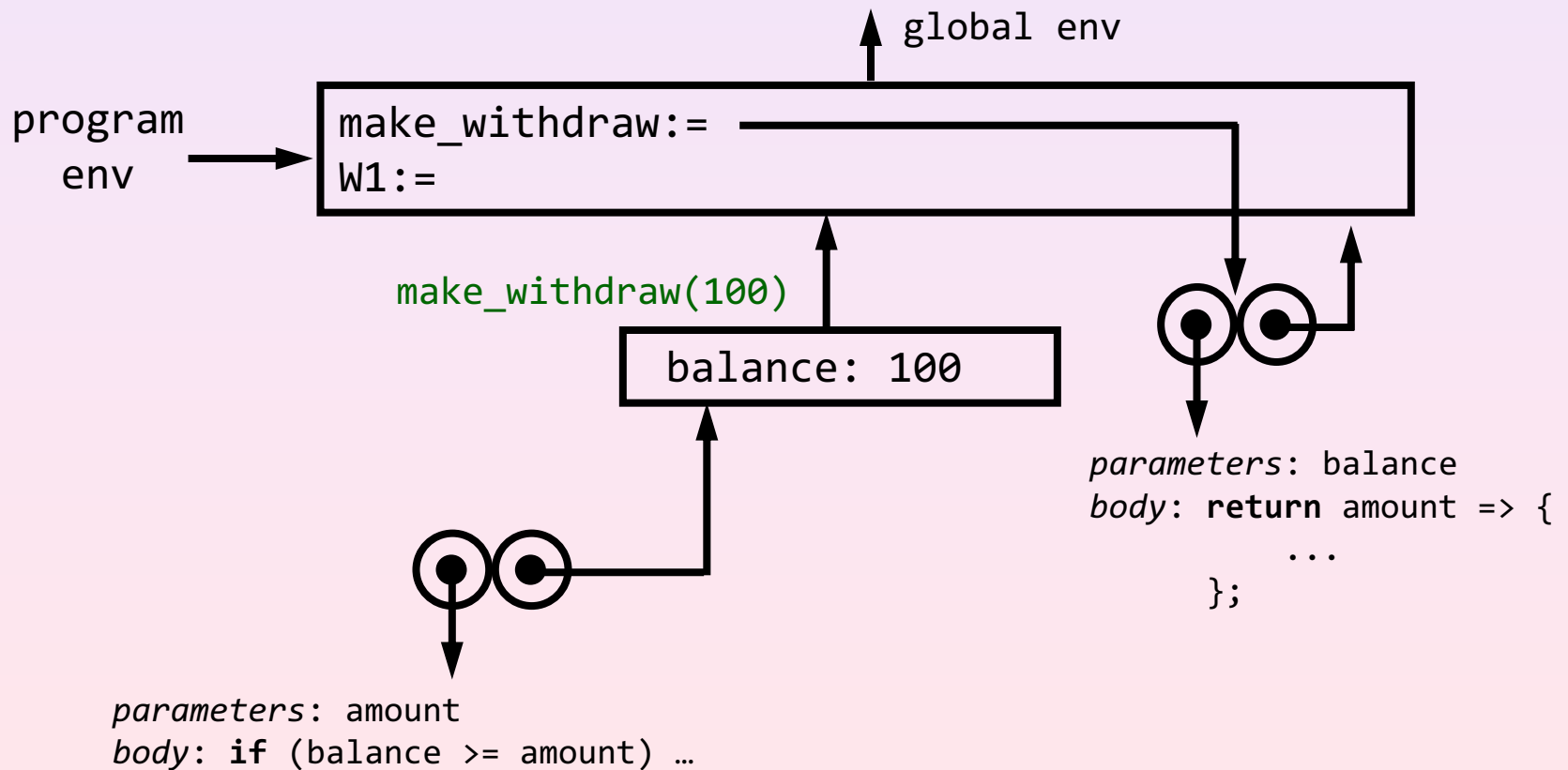
make_withdraw Example

- After declaring function `make_withdraw`



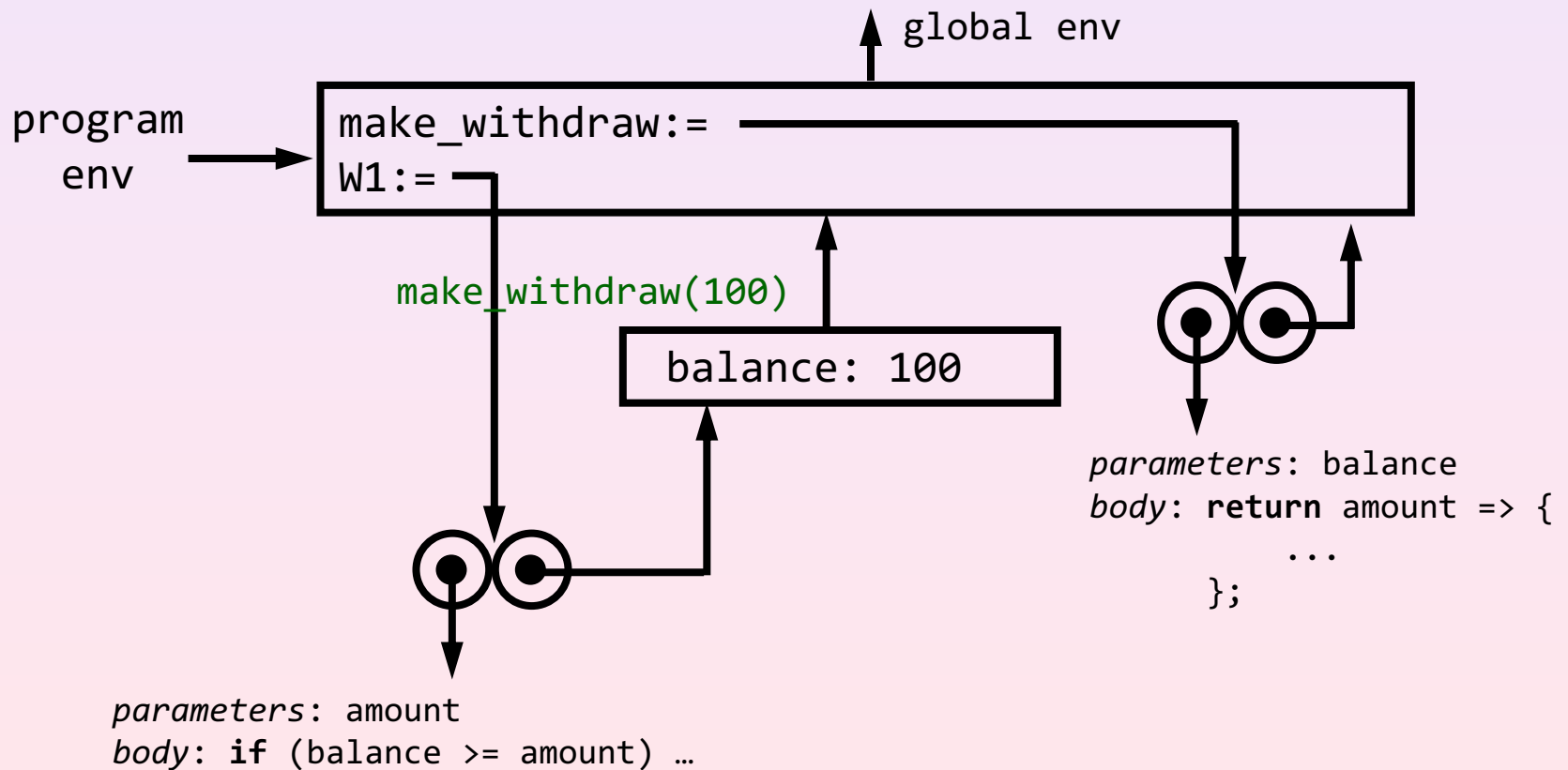
make_withdraw Example

- Evaluate `make_withdraw(100)`



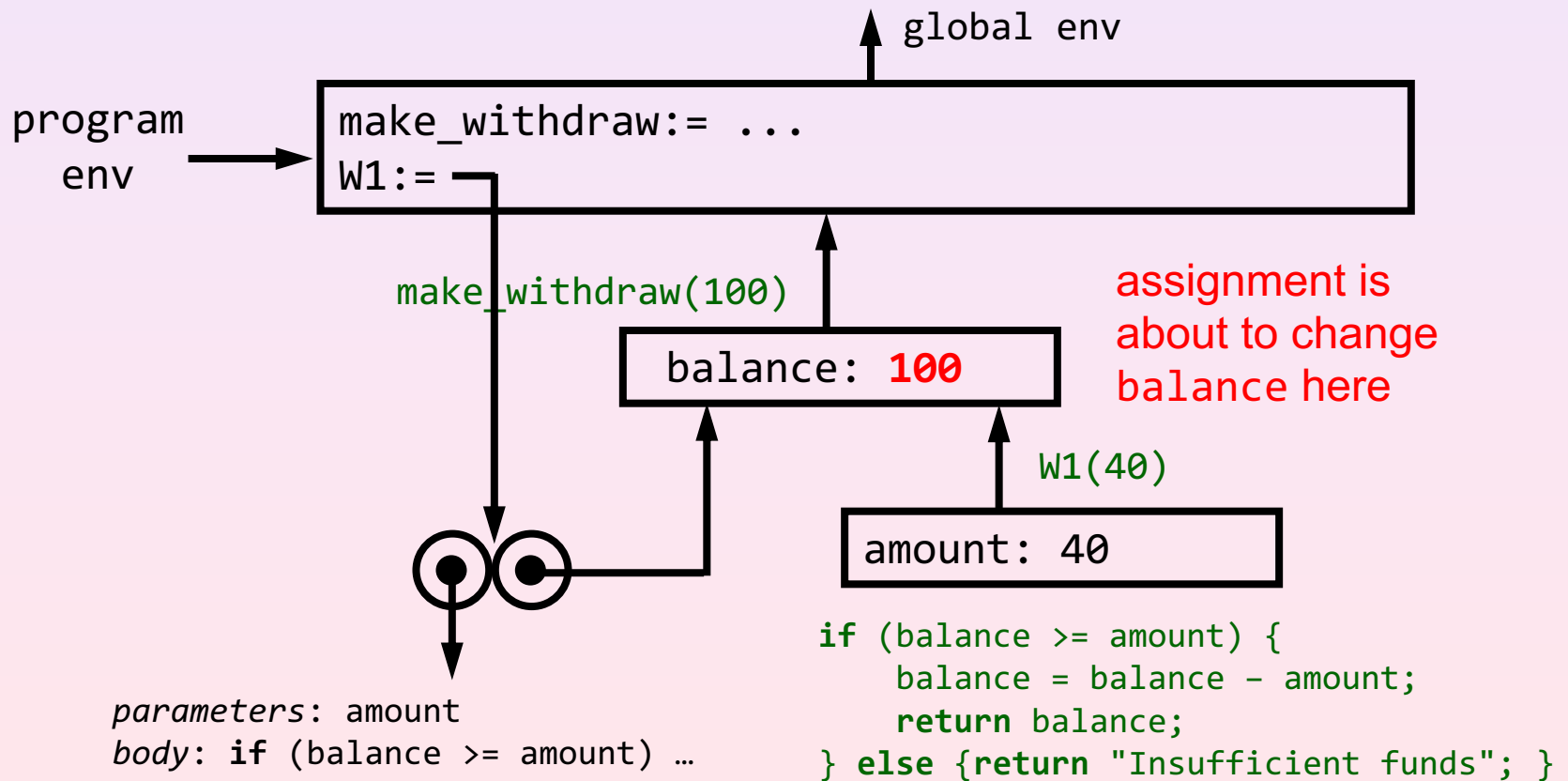
make_withdraw Example

- Evaluate `const W1 = make_withdraw(100);`



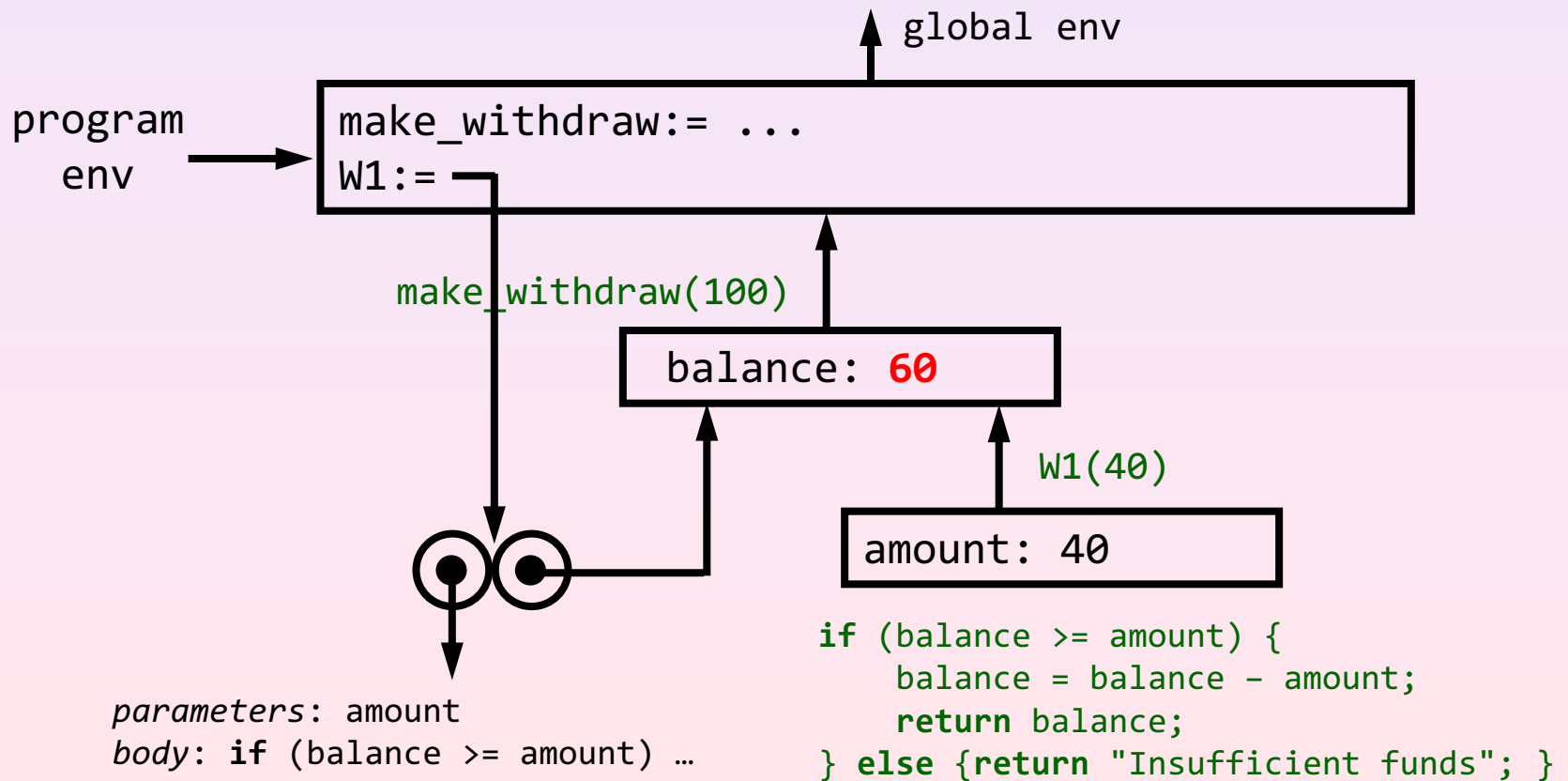
make_withdraw Example

- Evaluate `W1(40)`



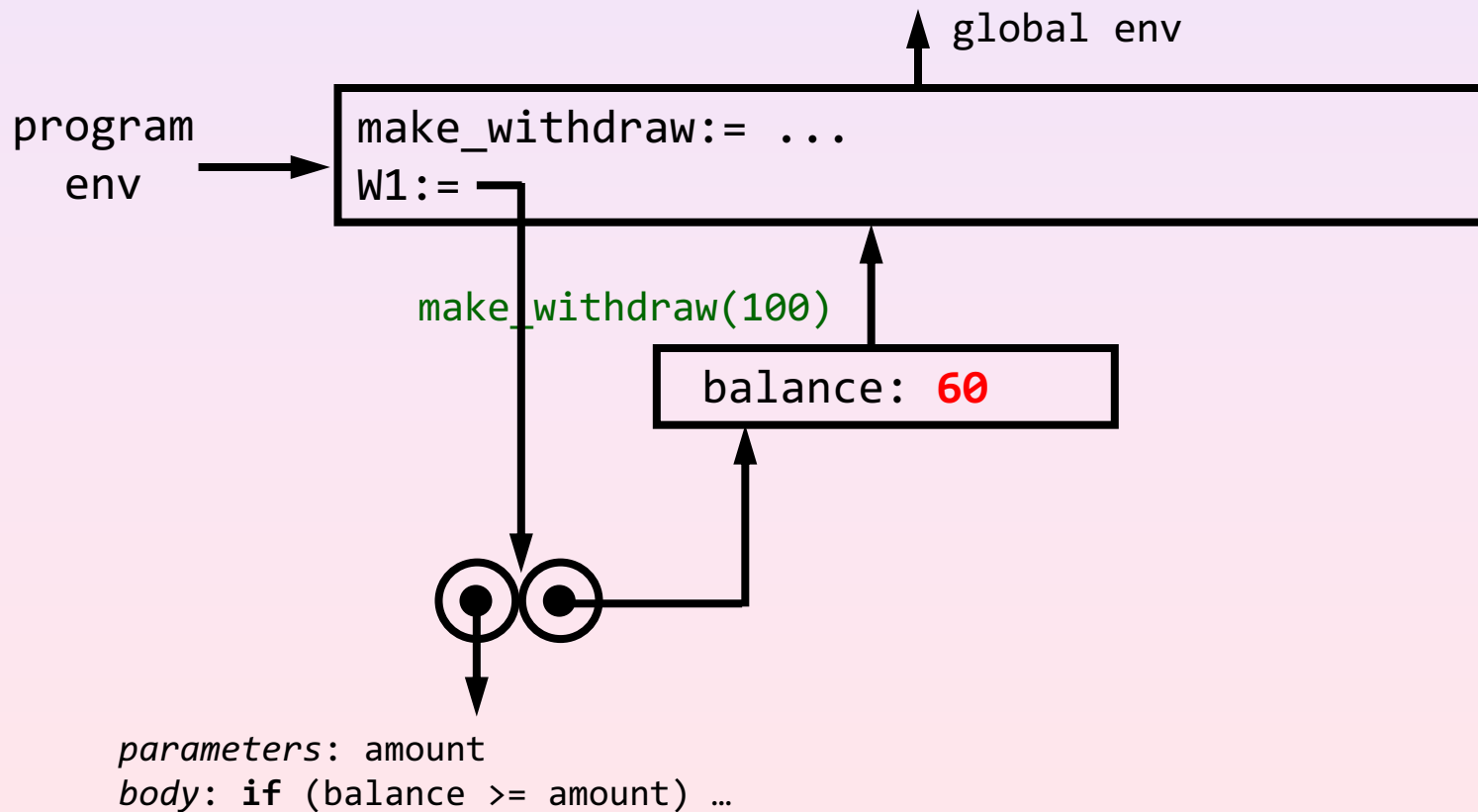
make_withdraw Example

- Evaluate `W1(40)`



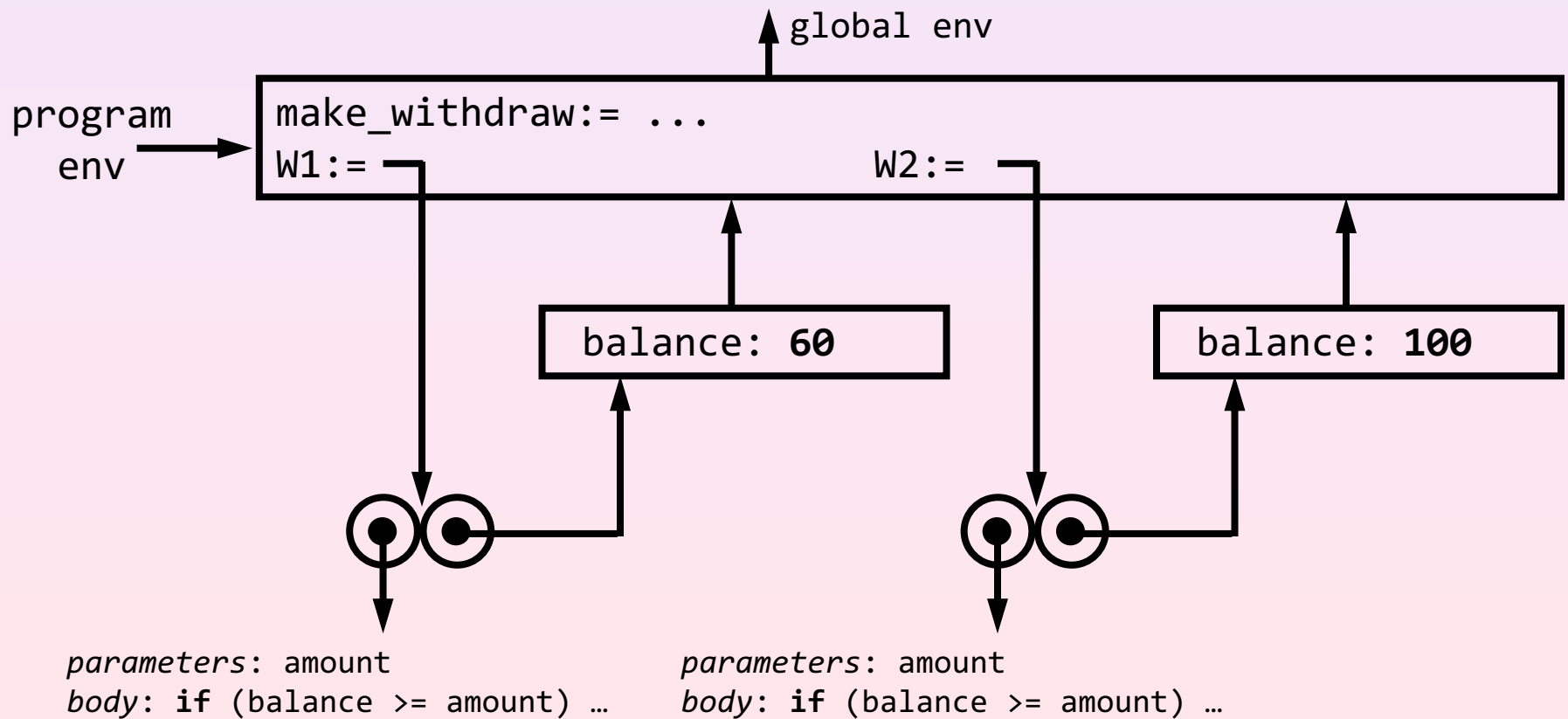
make_withdraw Example

- Environment after evaluating `W1(40);`



make_withdraw Example

- Environment after evaluating `const W2 = make_withdraw(100);`
- `W2` has its own state variable



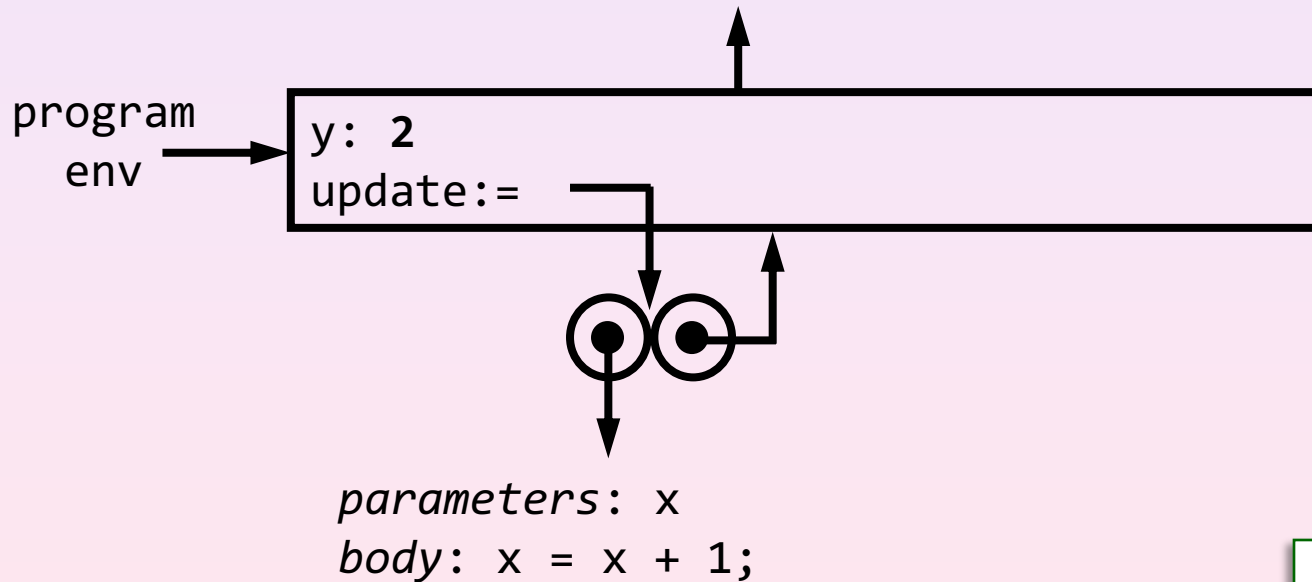
update Example

- What is the result of this program?

```
function update(x) {  
    x = x + 1;  
}  
let y = 2;  
update(y);  
y; ➔ 2
```

update Example

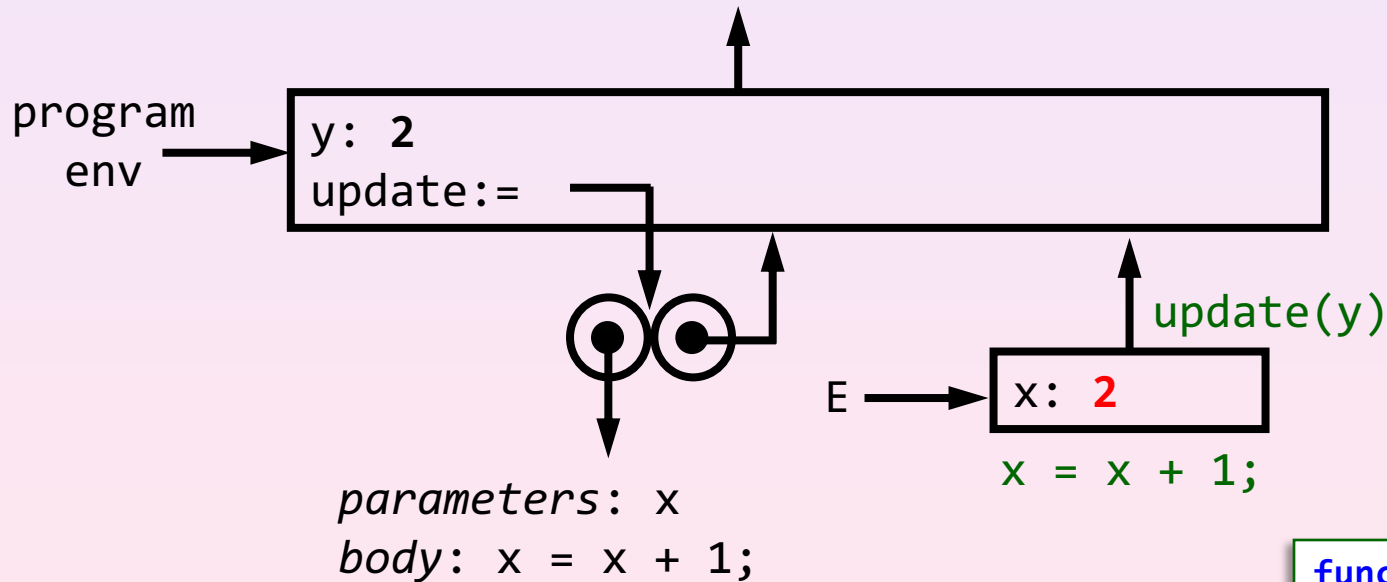
- After evaluating declarations of function `update` and variable `y`



```
function update(x) {
    x = x + 1;
}
let y = 2;
update(y);
y;
```

update Example

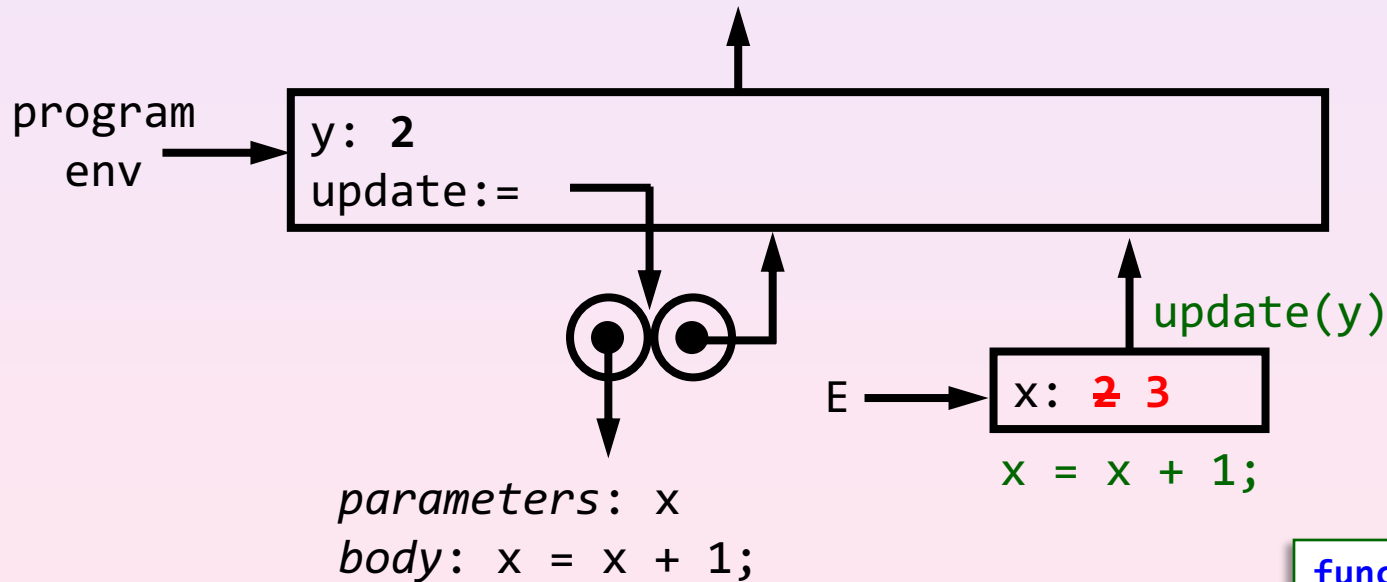
- Evaluate `update(y)` in **program env**



```
function update(x) {
    x = x + 1;
}
let y = 2;
update(y);
y;
```

update Example

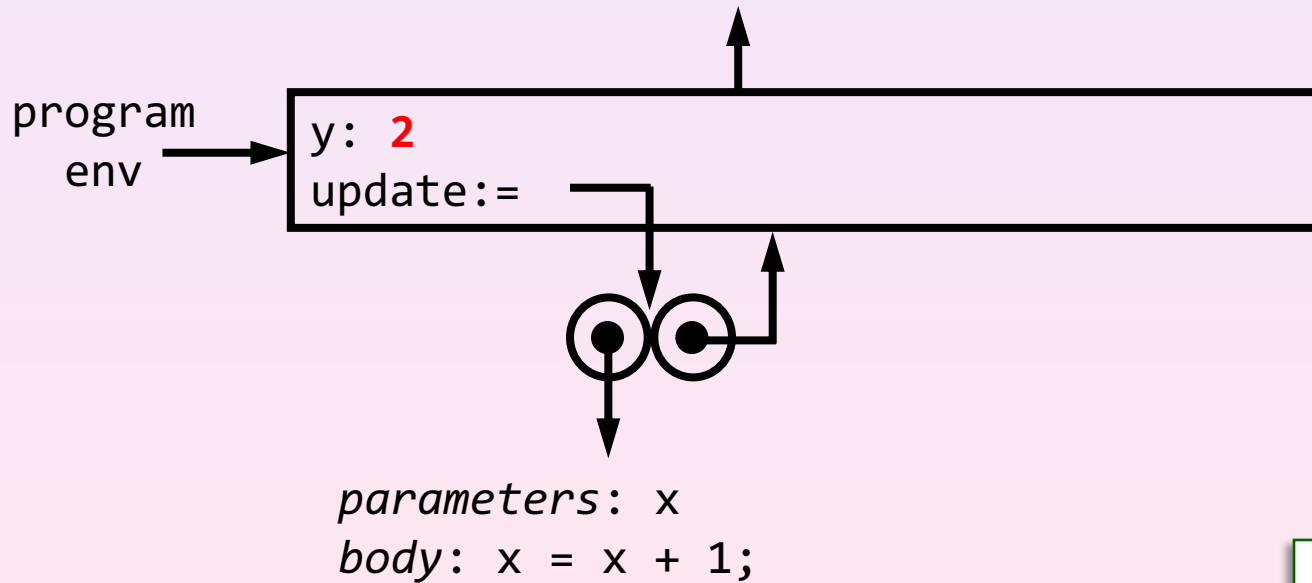
- Evaluate $x = x + 1$; in **Environment E**



```
function update(x) {
    x = x + 1;
}
let y = 2;
update(y);
y;
```


update Example

- Evaluate *y* in **program env**



```
function update(x) {
    x = x + 1;
}
let y = 2;
update(y);
y;
```

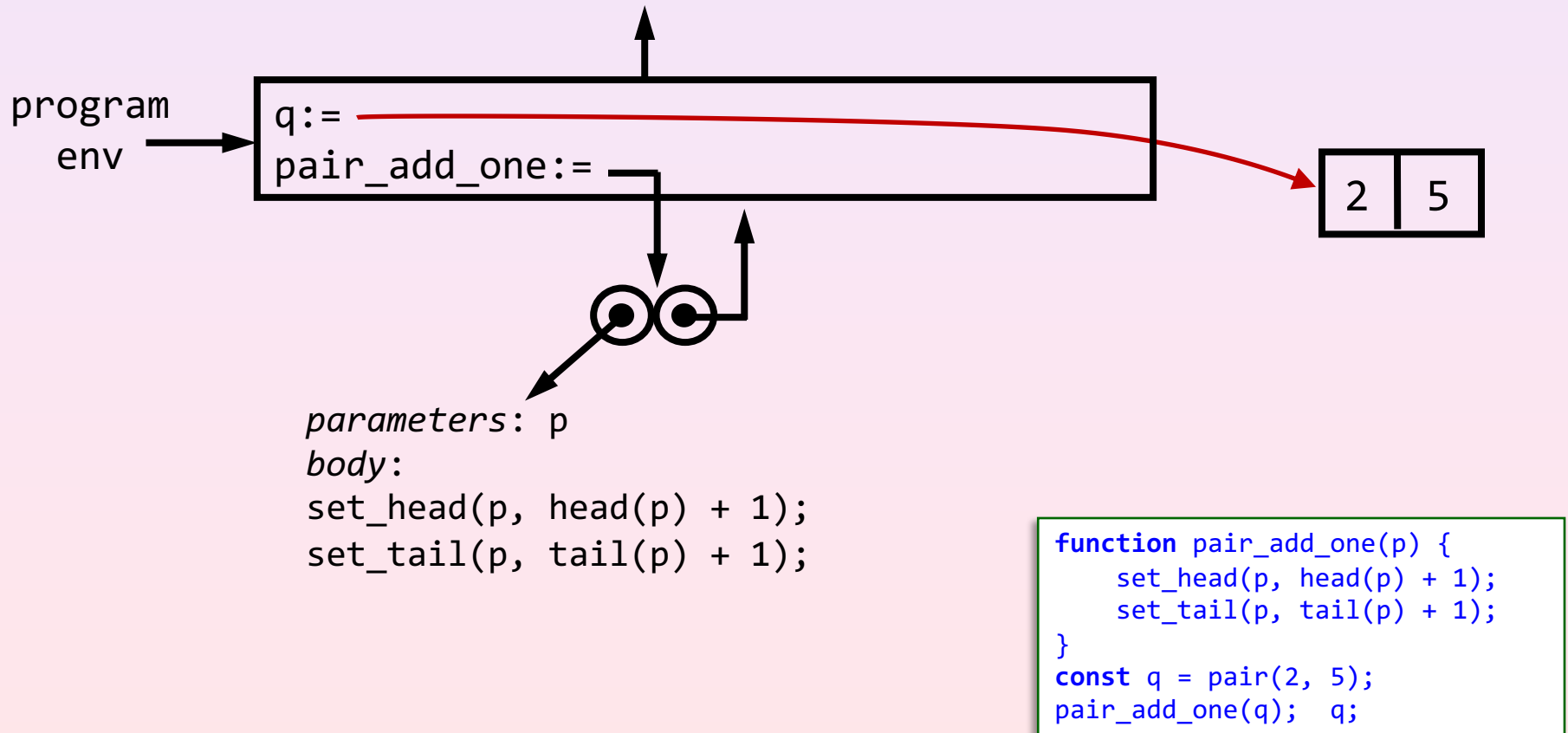
pair_add_one Example

- What is the result of this program?

```
function pair_add_one(p) {  
    set_head(p, head(p) + 1);  
    set_tail(p, tail(p) + 1);  
}  
const q = pair(2, 5);  
pair_add_one(q);  
q; → [3, 6]
```

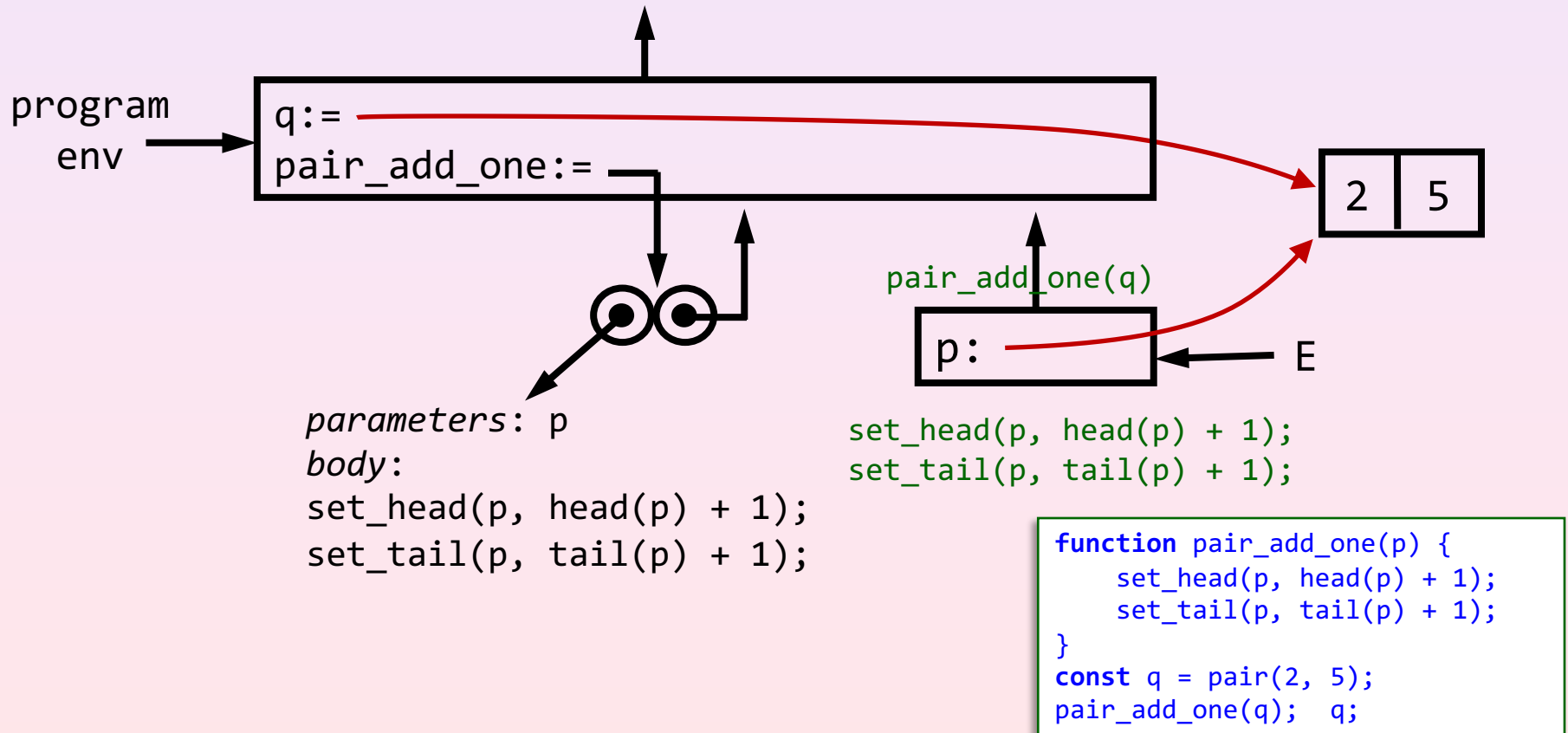
pair_add_one Example

- After evaluating declarations of function `pair_add_one` and constant `q`



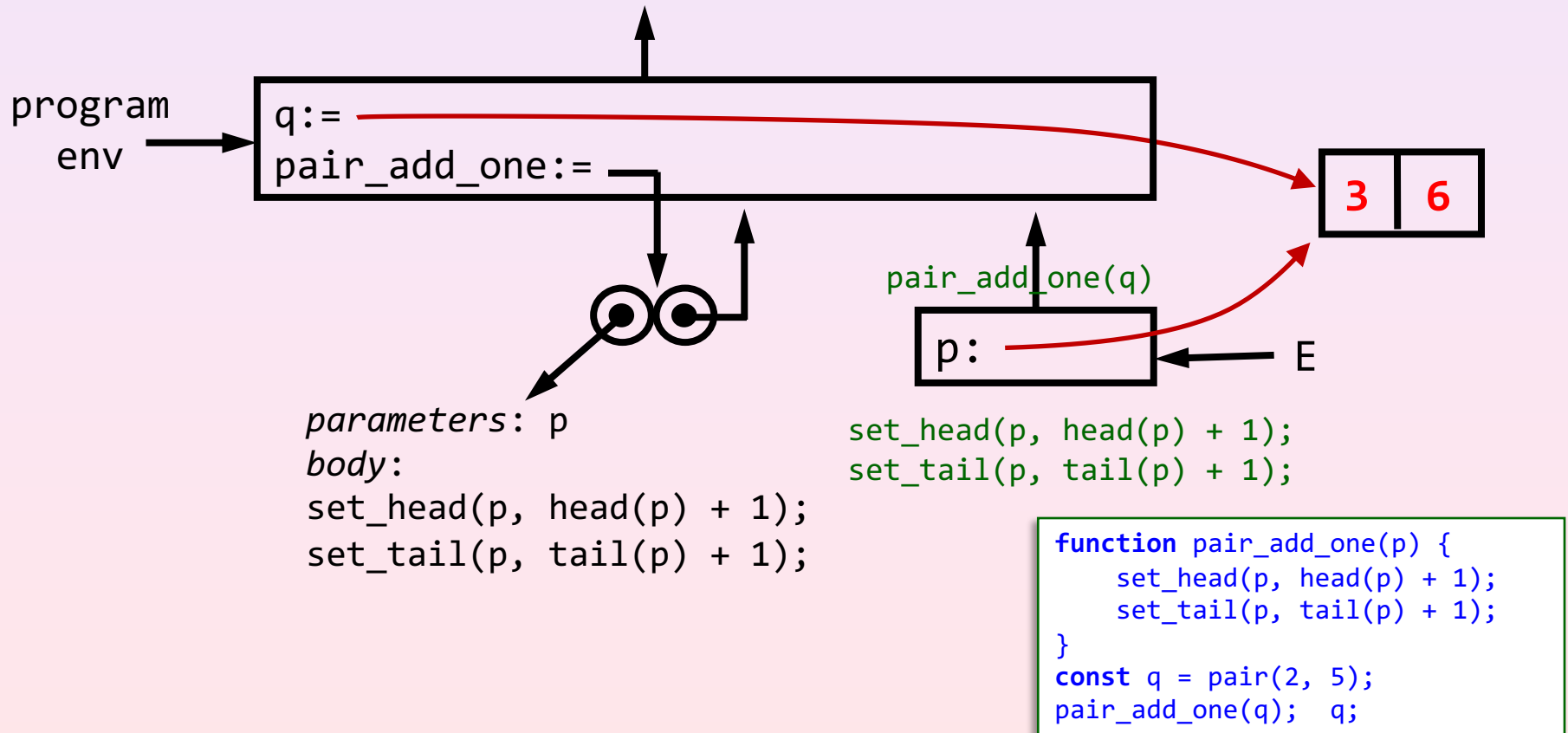
pair_add_one Example

- Evaluate `pair_add_one(q)` in **program env**



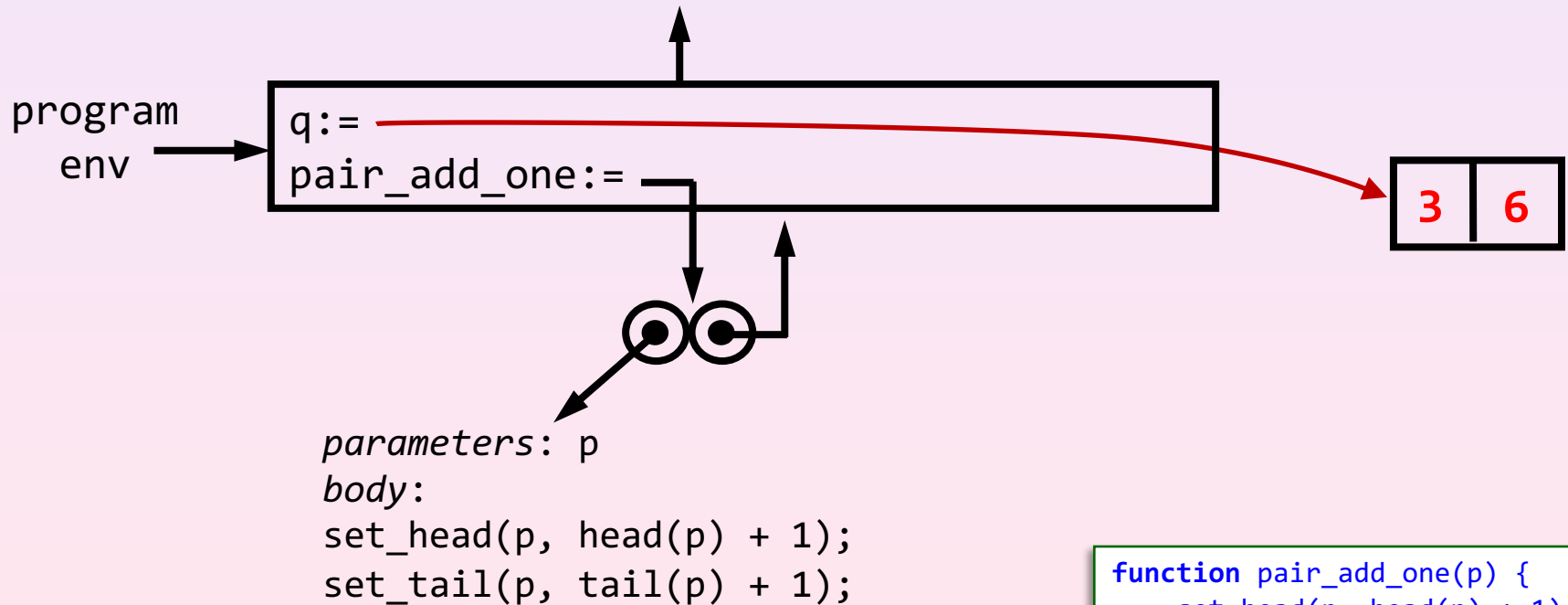
pair_add_one Example

- Evaluate `set_head(p, head(p) + 1)` and `set_tail(p, tail(p) + 1)` in **Environment E**



pair_add_one Example

- Evaluate `q` in **program env**



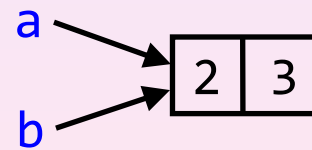
```
function pair_add_one(p) {
    set_head(p, head(p) + 1);
    set_tail(p, tail(p) + 1);
}
const q = pair(2, 5);
pair_add_one(q); q;
```

Sharing and Identity

- In `pair_add_one` example, `q` and `p` **share** the same **pair**
- Checking whether **two names** are **sharing** the same **pair**

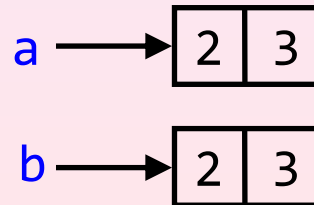
- **Example:**

```
const a = pair(2, 3);  
const b = a;  
a === b; ➔ true
```



- **Example:**

```
const a = pair(2, 3);  
const b = pair(2, 3);  
a === b; ➔ false
```

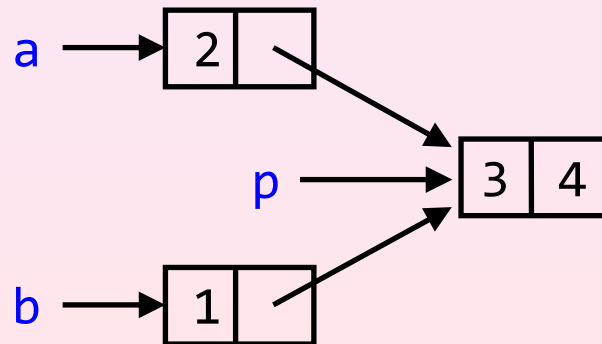


Sharing and Identity

- **Structures** sharing pairs

- **Example:**

```
const p = pair(3, 4);  
const a = pair(2, p);  
const b = pair(1, p);  
tail(a) === tail(b); ➔ true
```

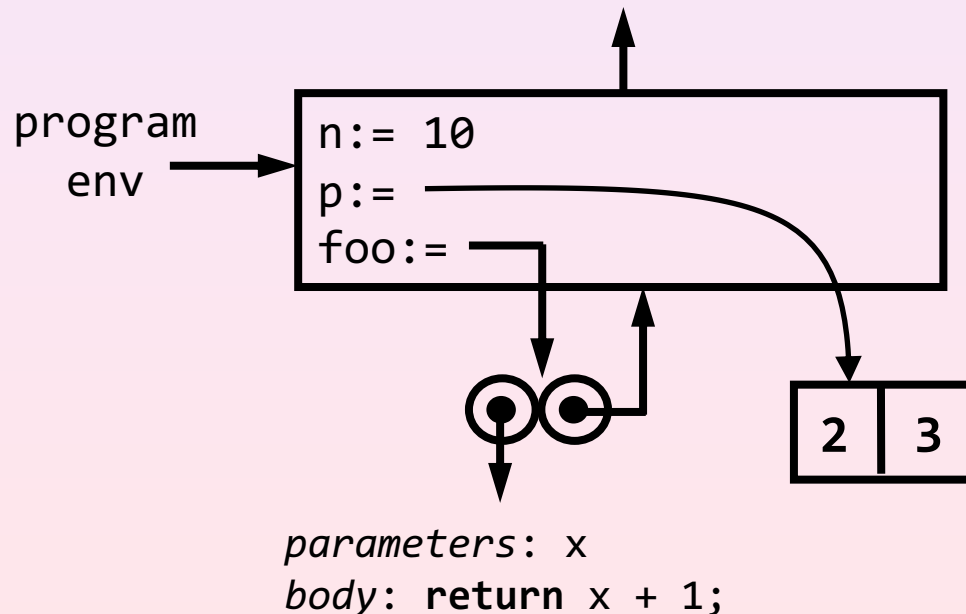


Identity in Source §3: ===

- For `a === b`, we say “is `a` *identical* to `b`”
 - `true`, `false`, `null`, `undefined` — each is identical to itself and nothing else
 - **Numbers** — two numbers are identical iff they have the same representation in the double-precision floating-point representation
 - **Strings** — two strings are identical if they have the same characters in the same order
 - **Functions** — functions are made by function expressions, and their creation bestows an identity upon them
 - **Pairs** — pairs are made by the `pair` function, and their creation bestows an identity upon them

Drawing Compound Structures

- **Primitive values** (e.g. numbers, strings, Boolean values, `null`) in bindings are drawn **inside** frames
- **Compound structures** (e.g. pairs, function objects) are drawn **outside** frames

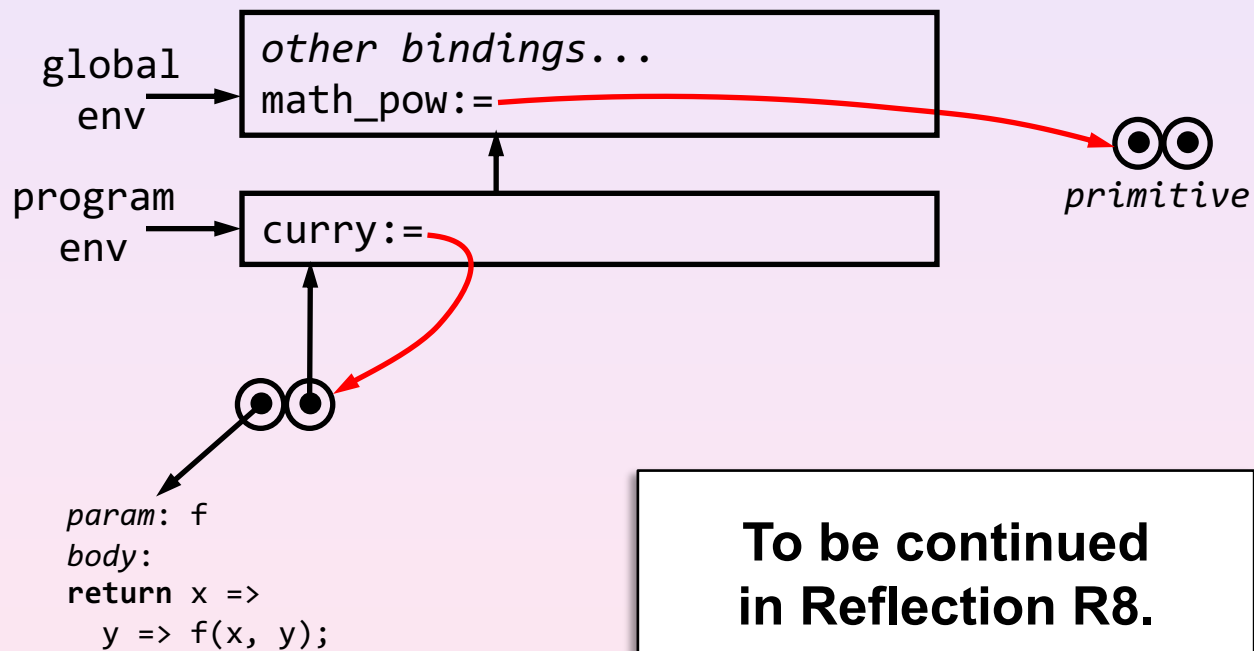


curry Example

- Consider this program

```
function curry(f) {  
    return x => y => f(x, y);  
}  
curry(math_pow)(3)(4);
```

curry Example



```
function curry(f) {
  return x => y => f(x, y);
}
curry(math_pow)(3)(4);
```

**To be continued
in Reflection R8.**

Summary

- **Assignment** allows us to create **state**
- **Substitution model** breaks down with **assignment**
- **Environment model** replaces **substitution model**
- When a **lambda expression** is evaluated, it creates a **function object** that remembers the **current environment**
- To evaluate a **function application**
 - 1) Create a **new frame** in the function's environment (right pointer of function object)
 - 2) Bind **parameters** in new frame
 - 3) Evaluate **body block** of function in the new environment