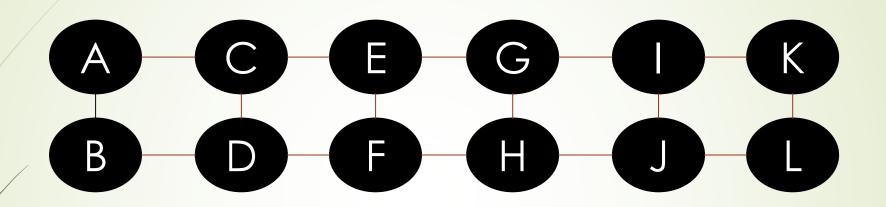
Super Nodes and Junction Trees

Presented by: Mehran Kazemi

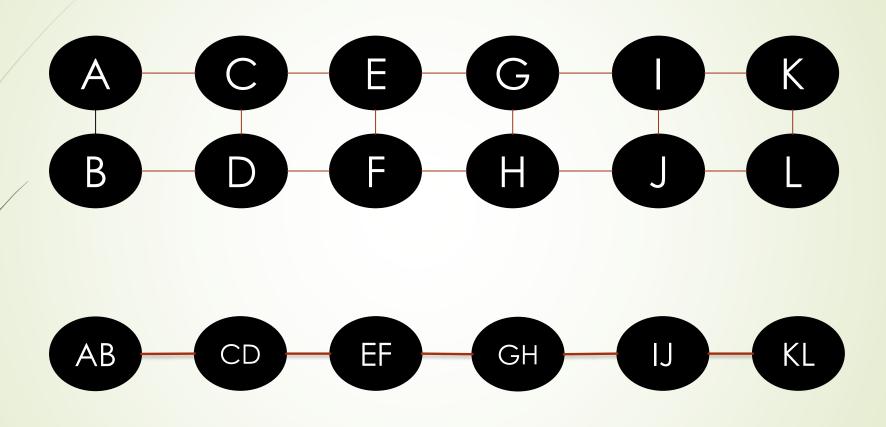
Outline

- Super Nodes
- Variable Elimination (VE)
- From VE to Junction Trees (Jtree)
- Calculating marginal probabilities in Jtrees

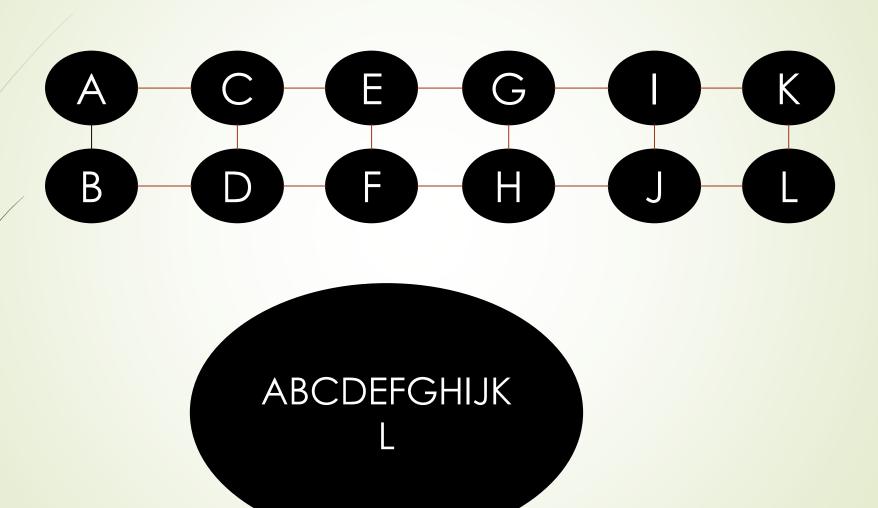
Super Nodes



Super Nodes



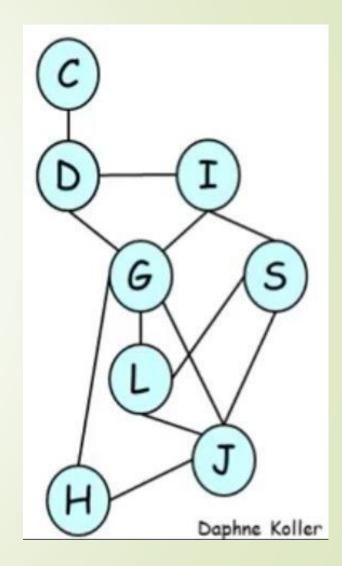
Super Nodes



Variable Elimination

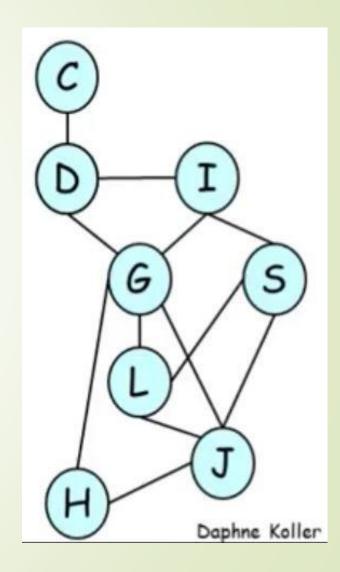
- We focus on calculating Z
- If we know how to calculate Z for a network, we can calculate all marginal probabilities.
- P(C=true | G=false)

$$= \frac{Z(Network|C=true,G=false)}{Z(Network|G=false)}$$



Variable Elimination

- $-\phi_1(C)$
- $-\phi_2(C,D)$
- $-\phi_3(D,I,G)$
- $-\phi_4(S,I)$
- $-\phi_5(H,G,J)$
- $-\phi_6(G,L)$
- $-\phi_7(S,L,J)$



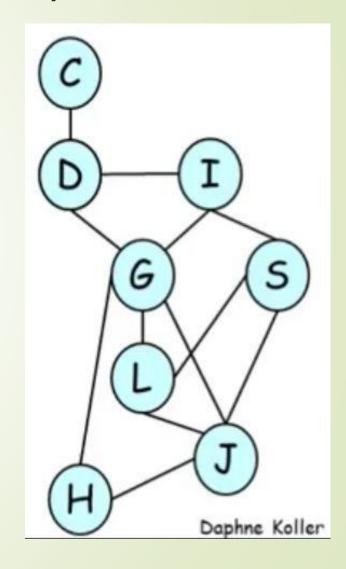
Variable Elimination (inference)

$$Z = \sum_{J} \sum_{L} \sum_{S} \sum_{G} \sum_{H} \sum_{I} \sum_{D} \sum_{C}$$

$$\phi_{1}(C)\phi_{2}(C,D)\phi_{3}(D,I,G)\phi_{4}(S,I)$$

$$\phi_{5}(H,G,J)\phi_{6}(G,L)\phi_{7}(S,L,J)$$

- Elimination Order: $\psi = \langle C, D, I, H, G, S, L, J \rangle$
- Z = $\sum_{J} \sum_{L} \sum_{S} \phi_{7}(S, L, J) \sum_{G} \phi_{6}(G, L) \sum_{H} \phi_{5}(H, G, J)$ $\sum_{I} \phi_{4}(S, I) \sum_{D} \phi_{3}(D, I, G) \sum_{C} \phi_{1}(C) \phi_{2}(C, D)$



Variable Elimination

$$Z = \sum_{J} \sum_{L} \sum_{S} \phi_{7}(S, L, J) \sum_{G} \phi_{6}(G, L) \sum_{H} \phi_{5}(H, G, J)$$

$$\sum_{I} \phi_{4}(S, I) \sum_{D} \phi_{3}(D, I, G) \sum_{C} \phi_{1}(C) \phi_{2}(C, D)$$

$$-\phi_1(C)$$

$$\phi_2(C,D)$$

$$\phi_2(C,D) \qquad \lambda_1(C,D) = \phi_1(C)\phi_2(C,D) \quad \tau_1(D) = \sum_D \lambda(C,D)$$

$$\tau_1(D) = \sum_D \lambda(C, D)$$

	С	Value
/	Τ	2
	F	1.2

С	D	Value
T	T	0.5
Τ	F	1
F	Т	1
F	F	2

С	D	Value
Τ	Т	1
Τ	F	2
F	T	1.2
F	F	2.4

$$Z = \sum_{J} \sum_{L} \sum_{S} \phi_7(S, L, J) \sum_{G} \phi_6(G, L) \sum_{H} \phi_5(H, G, J)$$

$$\sum_{I} \phi_4(S, I) \sum_{D} \phi_3(D, I, G) \tau_1(D)$$

Variable Elimination (inference)

$$Z = \sum_{J} \sum_{L} \sum_{S} \phi_7(S, L, J) \sum_{G} \phi_6(G, L) \sum_{H} \phi_5(H, G, J)$$

$$\sum_{I} \phi_4(S, I) \sum_{D} \phi_3(D, I, G) \tau_1(D)$$

$$Z = \sum_{J} \sum_{L} \sum_{S} \phi_7(S, L, J) \sum_{G} \phi_6(G, L) \sum_{H} \phi_5(H, G, J)$$

$$\sum_{I} \phi_4(S,I) \tau_2(G,I)$$

$$ightharpoonup Z = \sum_{I} \sum_{L} \sum_{S} \phi_{7}(S, L, J) \sum_{G} \phi_{6}(G, L) \tau_{3}(S, G) \tau_{4}(G, J)$$

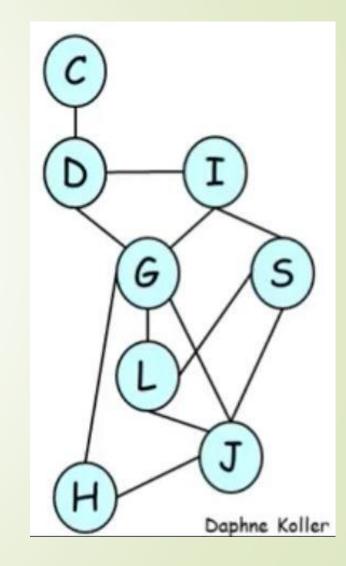
$$ightharpoonup Z = \sum_I \tau_{n-1}(I)$$

Variable Elimination (decoding)

- Elimination Order:

$$\psi = \langle C, D, I, H, G, S, L, J \rangle$$

 $ightharpoonup Argmax_I \dots Argmax_C \phi_1(C) \phi_2(C, D)$



Variable Elimination (decoding)

 $ightharpoonup Argmax_{I} ... Argmax_{C} \phi_{1}(C) \phi_{2}(C, D)$

$$\phi_1(C)$$

$$\phi_2(C,D)$$

	С	Value
/	T	2
	F	1.2

С	D	Value
T	Т	0.5
Τ	F	1
F	Τ	1
F	F	2

С	D	Value
Τ	Τ	1
T	F	2
F	T	1.2
F	F	2.4

D	Val(C)
Τ	F
F	F

 $ightharpoonup Argmax_I ... \tau_1(D)$

Time Complexity (Assuming binary variables)

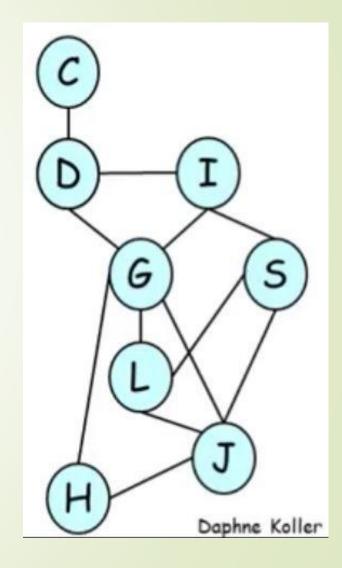
- Let $\phi_1, \phi_2, ..., \phi_m$ be potentials containing a variable V.
- Let $\tau(C_1, C_2, ..., C_w) = \sum_V \phi_1 \phi_2 ... \phi_m$
- $ightharpoonup au(C_1, C_2, ..., C_w)$ can be calculated in $O(2^w)$
- Let $w_1, w_2, ... w_n$ correspond to the number of variables in $\tau_1, \tau_2, ... \tau_n$ given a specific elimination order ψ .
- $\text{Let } \omega = \max(w_1, w_2, ..., w_n)$
- ▶ Variable elimination with elimination order ψ is $O(n2^{\omega})$
- $-\omega$ is called the width of ψ .

ω depends on the elimination order

- $Z = \sum_{C} \sum_{D} \sum_{I} \sum_{G} \sum_{S} \sum_{L} \sum_{J} \sum_{H}$ $\phi_{1}(C)\phi_{2}(C,D)\phi_{3}(D,I,G)\phi_{4}(S,I)$ $\phi_{5}(H,G,J)\phi_{6}(G,L)\phi_{7}(S,L,J)$
- Elimination Order:

$$\psi = < G, ... >$$

- $Z = \sum ... \sum_{G} \phi_3(D, I, G) \phi_5(H, G, J) \phi_6(G, L)$
- $ightharpoonup Z = \sum ... \tau(D, I, H, J, L)$

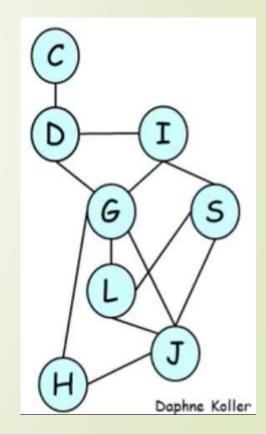


Time Complexity

- Let $\{\psi_1, \psi_2, ..., \psi_t\}$ represent all possible elimination orders, and $\{\omega(\psi_1), \omega(\psi_2), ..., \omega(\psi_t)\}$ represent the widths of these elimination orders.
- Define treewidth = $\min_{\psi \in \{\psi_1, \psi_2, ..., \psi_t\}} \omega(\psi)$
- Variable elimination is then $O(n2^{treewidth})$
- Finding a ψ with $\omega(\psi)$ =treewidth is NP-Hard.

- Variable Elimination is query sensitive: we must specify the query variable in advance. Each time we run a new query, we must re-run the entire algorithm.
- The junction tree algorithms generalizes VE to avoid this; they compile the UGM into a data structure which supports simultaneous execution of queries.

$$Z = \sum ... \sum_{D} \phi_3(D, I, G) \sum_{C} \phi_1(C) \phi_2(C, D)$$

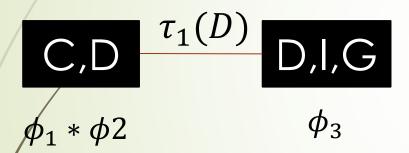




$$\phi_1 * \phi_2$$

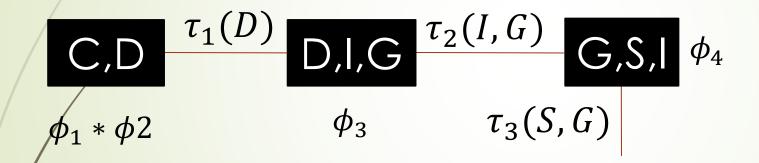
$$\tau_1(D)$$

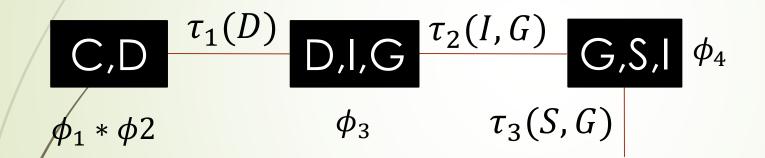
$$ightharpoonup Z = \sum ... \sum_{D} \phi_3(D, I, G) \tau_1(D)$$



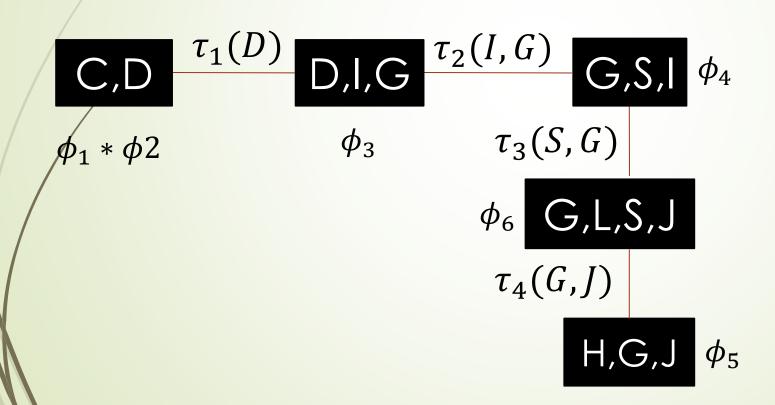
$$ightharpoonup Z = \sum ... \tau_2(I, G)$$

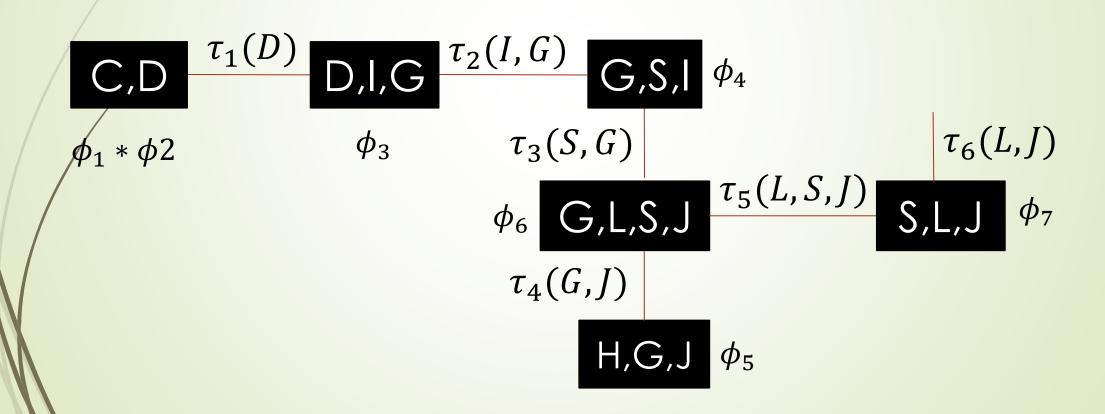
$$\tau_1(D)$$
 D,I,G $\tau_2(I,G)$ $\phi_1 * \phi_2$ ϕ_3

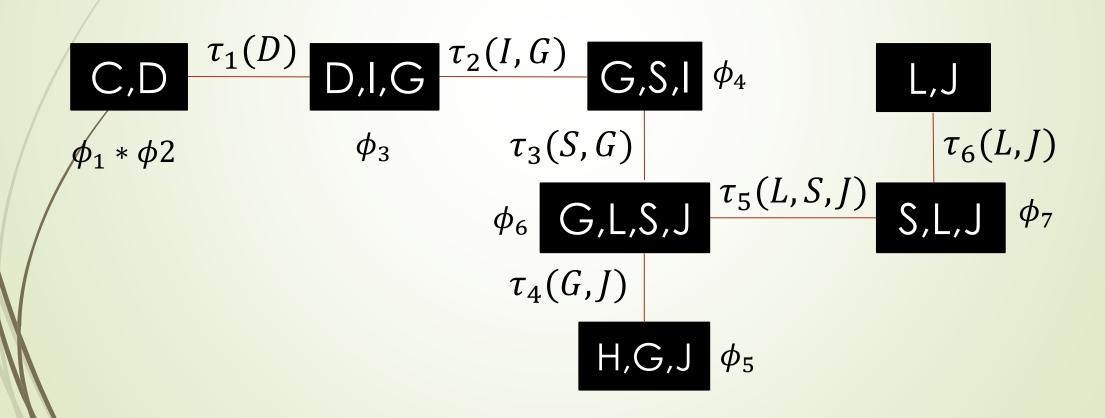


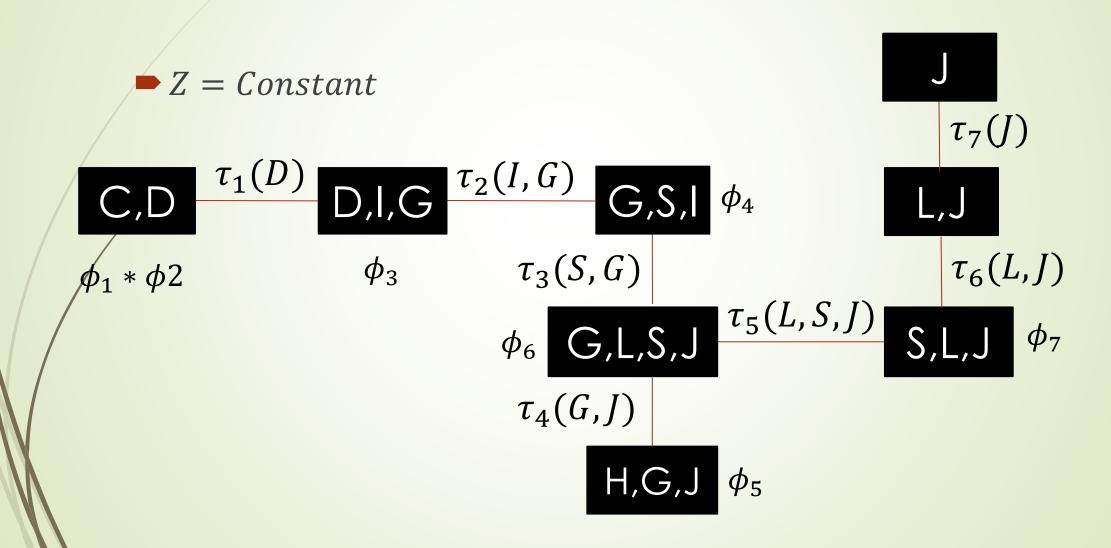


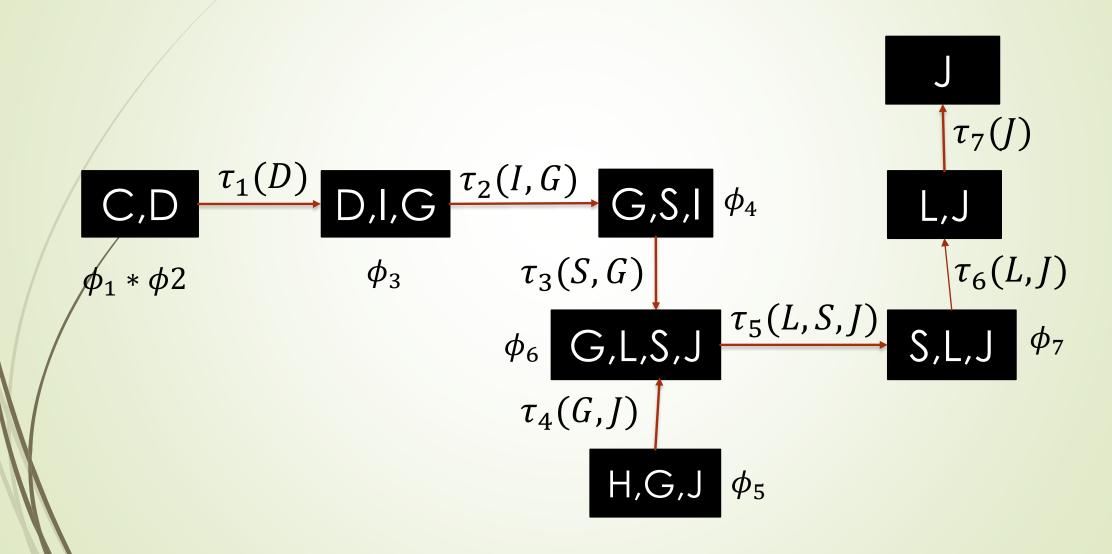
$$au_4(G,J)$$
H,G,J ϕ_1

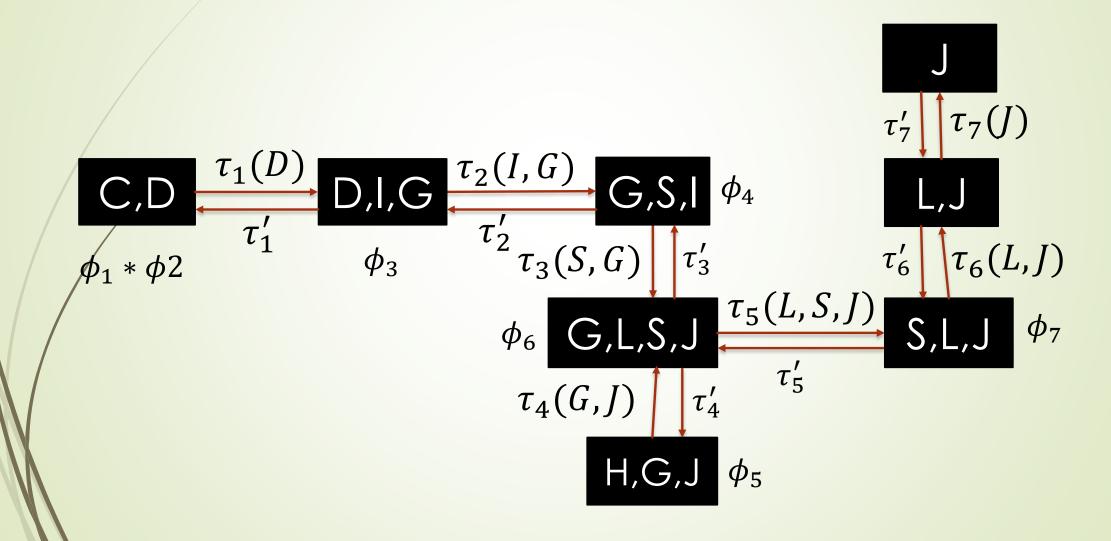




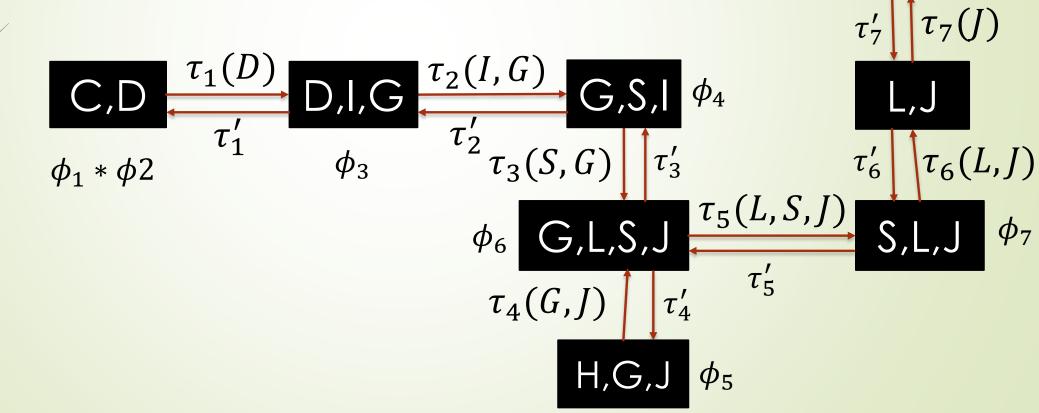




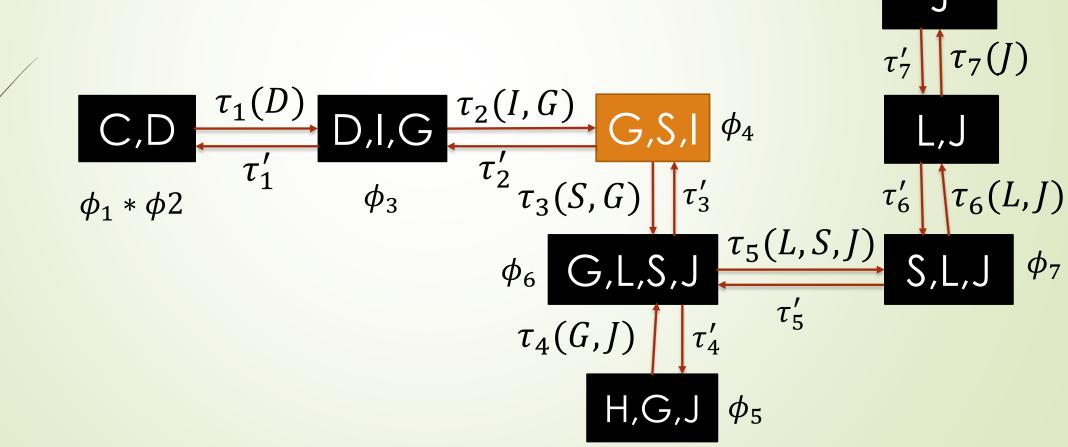




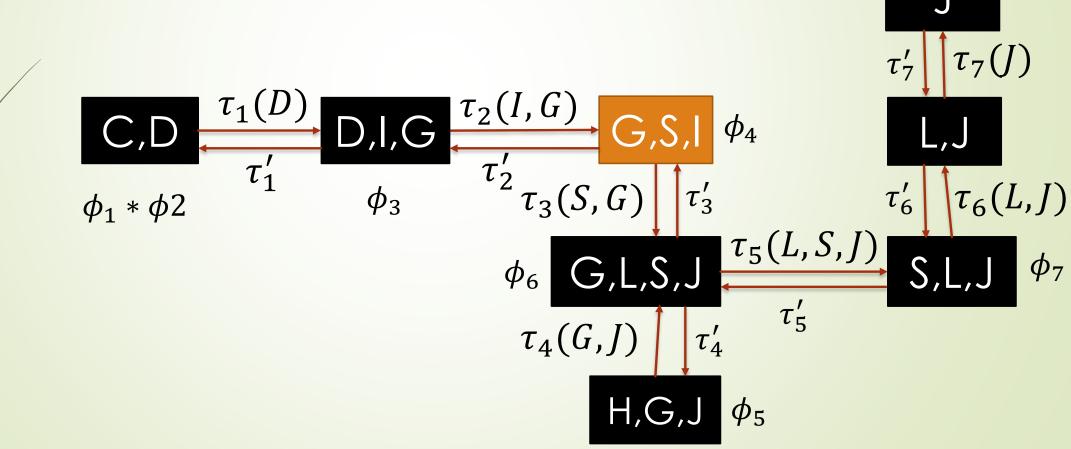
The joint probability of the variables in each cluster is proportional to the product of its potentials and its incoming messages.



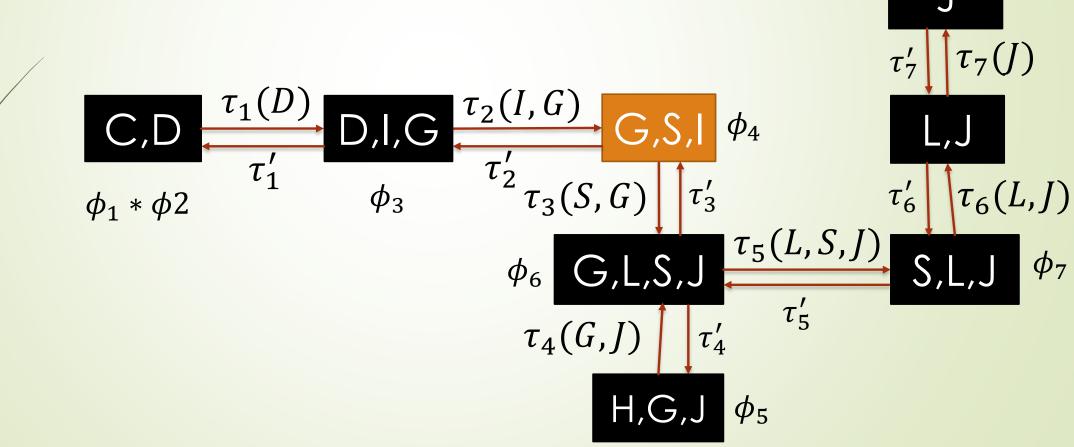
$$P(G,S,I) \propto \phi_4 \tau_2(I,G) \tau_3'(S,G)$$



$$P(G,S,I) \propto \phi_4(\sum_D \phi_3 \tau_1(D)) \tau_3'(S,G)$$



 $P(G,S,I) \propto \phi_4(\sum_D \phi_3 \sum_C \phi_1 \phi_2) \tau_3'(S,G)$



$$P(G,S,I) \propto \phi_{4}(\sum_{D}\phi_{3}\sum_{C}\phi_{1}\phi_{2})(\sum_{L,J}\phi_{6}\tau_{4}(G,J)\tau'_{5}(L,S,J))$$

$$C,D \xrightarrow{\tau_{1}(D)} D,I,G \xrightarrow{\tau_{2}(I,G)} G,S,I \phi_{4}$$

$$\phi_{1}*\phi_{2} \xrightarrow{\tau'_{1}} \phi_{3} \xrightarrow{\tau'_{2}} \tau_{3}(S,G) \xrightarrow{\tau'_{3}} \tau'_{5}(L,S,J)$$

$$\phi_{6} G,L,S,J \xrightarrow{\tau_{5}(L,S,J)} S,L,J \phi$$

$$\tau_{4}(G,J) \xrightarrow{\tau'_{4}} \tau'_{5}$$

$$P(G, S, I) \propto \phi_{4}(\Sigma_{D}\phi_{3}\Sigma_{C}\phi_{1}\phi_{2})(\Sigma_{L,J}\phi_{6}(\Sigma_{H}\phi_{5})\tau'_{5}(L, S, J))$$

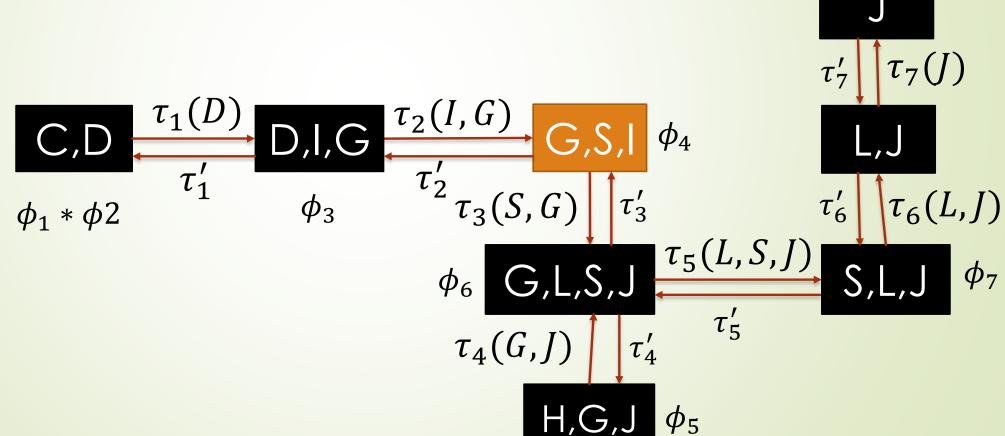
$$C,D \xrightarrow{\tau_{1}(D)} D,J,G \xrightarrow{\tau_{2}(I,G)} G,S,J \phi_{4}$$

$$\phi_{1}*\phi_{2} \xrightarrow{\tau'_{1}} \phi_{3} \xrightarrow{\tau'_{2}} \tau_{3}(S,G) \xrightarrow{\tau'_{3}} \tau_{5}(L, S, J)$$

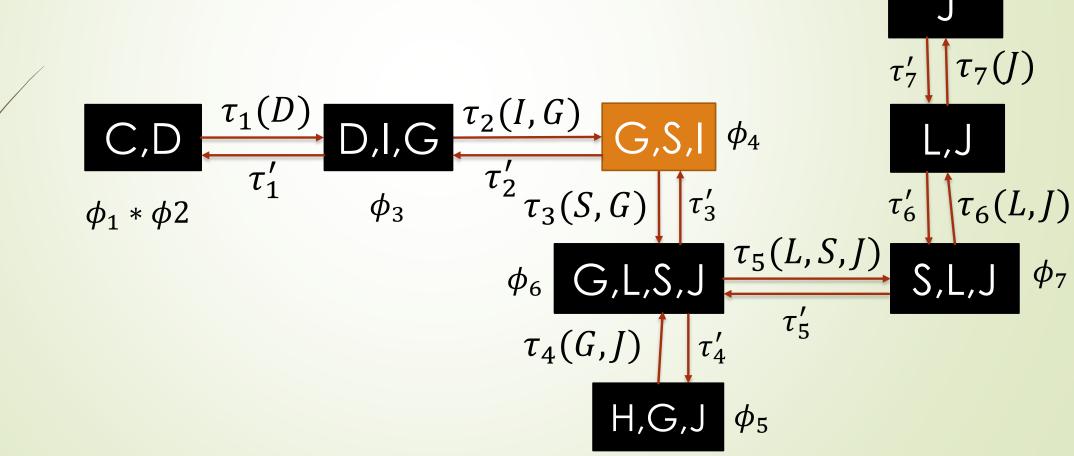
$$\phi_{6} G,L,S,J \xrightarrow{\tau_{5}(L,S,J)} S,L,J \phi$$

$$\tau_{4}(G,J) \xrightarrow{\tau'_{4}} \tau'_{5}$$

$$P(G,S,I) \propto \phi_4(\Sigma_D \phi_3 \Sigma_C \phi_1 \phi_2)(\Sigma_{L,J} \phi_6(\Sigma_H \phi_5) \phi_7 \tau_6'(L,J))$$



 $P(G,S,I) \propto \phi_4(\Sigma_D \phi_3 \Sigma_C \phi_1 \phi_2)(\Sigma_{L,J} \phi_6(\Sigma_H \phi_5) \phi_7)$



- For every elimination order ψ , we will get a different Jtree
- The time complexity of sending messages in each direction in a Jtree generated with elimination order ψ is $O(n2^{\omega(\psi)})$.
- With spending twice the time of VE, we can have the probabilities for all random variables.

Demos

Go though the GraphCuts demo

