#### QUEST FOR THE MNEMON SEED

#### THE THREE R'S OF KEY MANAGEMENT: REPRODUCTION, ROTATION, AND RECOVERY

# DAD (DECENTRALIZED AUTONOMIC DATA) DERIVED-DID

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### Background

Decentralized Autonomic Data and the Three R's of Key Management RWOT Spring 2018

https://github.com/WebOfTrustInfo/rebooting-the-web-of-trust-spring2018/blob/master/final-documents/ DecentralizedAutonomicData.md

#### Motivation

Data streaming applications:

Analytics and instrumentation for distributed web, internet of things (IoT), or self-sovereign reputation systems.

Maintain a provenance chain for data undergoing various processing stages that follows diffuse trust security principles including signed at rest and in motion.

#### DAD: Decentralized Autonomic Data

Provenance for decentralized streaming data applications including transformations

Decentralized: governance of the data may not reside with a single party, trust in the data provenance is diffuse, DID based.

Autonomic: self-managing or self-regulating. Self-managing includes cryptographic techniques for maintaining data provenance that make the data self-identifying, self-certifying, and self-securing.

Autonomic implies the use of cryptographic signatures to provide a root of trust for data integrity and to maintain that trust over transformation of that data

Key management is thus first order property of DAD items.

Reproduction, Rotation, and Recovery

Pre-rotation & Hybrid recovery methods

#### DID and DDID

DID = Decentralized Identifier https://w3c-ccg.github.io/did-spec/

did:\*method\*:\*idstring\*

did:dad:Xq5YqaL6L48pf0fu7IUhL0JRaU2\_RxFP0AL43wYn148=

did:dad:Xq5YqaL6L48pf0fu7IUhL0JRaU2\_RxFP0AL43wYn148=:blue

did:dad:Xq5YqaL6L48pf0fu7IUhL0JRaU2\_RxFP0AL43wYn148=?who=me

DDO = DID Document, provides meta-data about DID

DAD Streaming = Multiplicity of data items and associated identifiers

DID/DDO pair per DAD item may not be practical

DDID = Derived DID = Unique DID format identifier derived from one root

DID/DDO that provides meta-data for a large number of DDIDs

### Example Signed DAD

```
{
    "id": "did:dad:Xq5YqaL6L48pf0fu7IUhL0JRaU2_RxFP0AL43wYn148=",
    "data":
    {
        "name": "John Smith",
        "nation": "USA"
    }
}
\r\n\r\n
u72j9aKHgz99f0K8pSkMnyqwvEr_3rpS_z2034L99sTWrMIIJGQPbVuIJ1cupo6cfIf_KCB5ecVRYoFRzAPnAQ==
```

### Minimally Sufficient Means

Streaming data applications may impose significant performance demands on the processing of the associated data

Desire efficient mechanisms for providing the autonomic properties of DADs

Find minimally sufficient means for managing cryptographic integrity

Key Reproduction, Rotation, and Recovery

Representation and Processing

### Reproduction

Simple Privacy via unique cryptonym (DDID) per pair-wise interaction context

More sophisticated methods such as zero knowledge proofs may not minimally sufficient

DDIDS derived via some type of hierarchically deterministic algorithm allow for simple method to generated large numbers public DDIDS without having to store the associated private keys only the root private key

did:dad:Xq5YqaL6L48pf0fu7IUhL0JRaU2\_RxFP0AL43wYn148=?chain=0\1\2

did:dad:Qt27fThWoNZsa88VrTkep6H-4HA8tr54sHON1vWl6FE=

### DDID Management

#### DDID Database

```
{
   "did:dad:Qt27fThWoNZsa88VrTkep6H-4HA8tr54sHON1vWl6FE=":
   "did:dad:Xq5YqaL6L48pf0fu7IUhL0JRaU2_RxFP0AL43wYn148=?chain=0\1\2",
   ...
}
```

#### DDID NameSpacing

```
did:dad:Xq5YqaL6L48pf0fu7IUhL0JRaU2_RxFP0AL43wYn148=:blue?chain=0/1
did:dad:Xq5YqaL6L48pf0fu7IUhL0JRaU2_RxFP0AL43wYn148=:red?chain=0/1
```

#### DDID Sequencing

did:dad:Qt27fThWoNZsa88VrTkep6H-4HA8tr54sHON1vWl6FE=/10057

### Change Detection

Prevent replay attacks:

sequence number in DDID

changed field with monotonically increasing sequence number or date time

```
{
    "id": "did:dad:Qt27fThWoNZsa88VrTkep6H-4HA8tr54sHON1vWl6FE=/10057",
    "changed" : "2000-01-01T00:00:00+00:00",
    "data":
    {
        "temp": 50,
        "time": "12:15:35"
    }
}
\r\n\r\n
u72j9aKHgz99f0K8pSkMnyqwvEr_3rpS_z2034L99sTWrMIIJGQPbVuIJ1cupo6cfIf_KCB5ecVRYoFRzAPnAQ==
```

#### DDID Generation

#### On the fly DDIDs:

Data source is not identified so receiver generates DDID that is later correlated to or claimed by the data source

#### Public Derivation:

Client communicates with large number of public services

DDID is derived from root private key and public service identifier

Client does not need to store DDID but can re-derive on demand

#### Rotation

Rotation = Revoke and Replace

Revoke only = Revoke and replace with null key

Replace only = security vulnerability

In general revocation is special case of rotation

### Why Rotate?

Public key vulnerable to attack due to time exposure

Continued use of private key increases vulnerability to exploit

Challenge is rotation after suspected exploit, exposure, or capture of private key.

Rotation operation that is signed by exploited private key is problematic.

Seek minimally sufficient means.

Using another key for rotation adds complexity and in a sense is just displacing the rotation problem without solving it.

#### Pre-Rotation

Likelihood of exploit is a function of exposure to probing or continued monitoring.

Narrowly restricting the the opportunity for exploit in time, place and method, such as, one time only, minimizes the vulnerability to exploit.

Exploiter has to either predict the one-time event or has to continuously universally monitor all events.

Create and declare the next "rotated" key (prerotated key) at the one-time only inception event for the original key. The private key of the pre-rotated key is then stored securely and is not exposed until a rotation event is needed.

The rotation event is now another one-time event determined by the key owner

Exploit of the original key cannot change the pre-rotated key or spoof a rotation event.

Rotation is self-contained

### Pre-Rotation Requirements

Inception event is one-time only. Contains two public keys signed by original key.

Rotation event contains three public keys. Original, Pre-rotated, next pre-rotated. Signed by original and pre-rotated.

Private: pair-wise interaction requires recipient maintain log database of inception and rotation events. Can detect any replay attacks or attempts to spoof inception event.

Public: redundant, immutable, persistent, time-stamped log database of rotation event history for a given original (inception) key on replicant servers

# Examples with DID DAD

#### Inception Event

```
"id": "did:dad:Qt27fThWoNZsa88VrTkep6H-4HA8tr54sHON1vWl6FE=",
    "changed": "2000-01-01T00:00:00+00:00",
    "ensuer": "Xq5YqaL6L48pf0fu7IUhL0JRaU2 RxFP0AL43wYn148="
\r\n\r\n
u72j9aKHgz99f0K8pSkMnyqwvEr_3rpS_z2034L99sTWrMIIJGQPbVuIJ1cupo6cfIf_KCB5ecVRYoFRzAPnAQ==
Rotation Event
    "id": "did:dad:Qt27fThWoNZsa88VrTkep6H-4HA8tr54sHON1vWl6FE=",
    "changed": "2000-01-01T00:00:00+00:00",
    "erster": "Qt27fThWoNZsa88VrTkep6H-4HA8tr54sHON1vWl6FE=",
    "signer": "Xq5YqaL6L48pf0fu7IUhL0JRaU2_RxFP0AL43wYn148=",
    "ensuer": "dZ74MLZXD-1QHoa73w9pQ9GroAvxqFi2RTZWlkC0raY="
\r\n\r\n
jc3ZXMA5GuypGWFEsxrGVOBmKDtd0J34UKZyTIYUMohoMYirR8AgH5O28PSHyUB-UlwfWaJlibIPUmZVPTG1DA==
\r\n\r\n
efIU4jplMtZzjgaWc85gLjJpmmay6QoFvApMuinHn67UkQZ2it17ZPebYFvmCEKcd0weWQONaTO-ajwQxJe2DA==
```

### Listed Pre-rotation Structure

#### Inception Event

```
"id": "did:dad:Qt27fThWoNZsa88VrTkep6H-4HA8tr54sHON1vWl6FE=",
   "changed": "2000-01-01T00:00:00+00:00",
    "signer": 0,
    "signers":
        "Qt27fThWoNZsa88VrTkep6H-4HA8tr54sHON1vWl6FE=",
        "Xq5YqaL6L48pf0fu7IUhL0JRaU2_RxFP0AL43wYn148=",
\r\n\r\n
jc3ZXMA5GuypGWFEsxrGVOBmKDtd0J34UKZyTIYUMohoMYirR8AgH5O28PSHyUB-UlwfWaJlibIPUmZVPTG1DA==
Rotation Event
    "id": "did:dad:Qt27fThWoNZsa88VrTkep6H-4HA8tr54sHON1vWl6FE=",
    "changed": "2000-01-01T00:00:00+00:00",
    "signer": 1,
    "signers":
        "Qt27fThWoNZsa88VrTkep6H-4HA8tr54sHON1vWl6FE=",
        "Xq5YqaL6L48pf0fu7IUhL0JRaU2_RxFP0AL43wYn148=",
        "dZ74MLZXD-1QHoa73w9pQ9GroAvxqFi2RTZWlkC0raY="
\r\n\
jc3ZXMA5GuypGWFEsxrGVOBmKDtd0J34UKZyTIYUMohoMYirR8AgH5O28PSHyUB-UlwfWaJlibIPUmZVPTG1DA==
\r\n\
efIU4jplMtZzjgaWc85gLjJpmmay6QoFvApMuinHn67UkQZ2it17ZPebYFvmCEKcd0weWQONaTO-ajwQxJe2DA==
```

### Listed Pre-rotation Continued

```
Rotation Event
    "id": "did:dad:Qt27fThWoNZsa88VrTkep6H-4HA8tr54sHON1vWl6FE=",
    "changed": "2000-01-01T00:00:00+00:00",
    "signer": 2,
    "signers":
        "Qt27fThWoNZsa88VrTkep6H-4HA8tr54sHON1vWl6FE=",
        "Xq5YqaL6L48pf0fu7IUhL0JRaU2 RxFP0AL43wYn148=",
        "dZ74MLZXD-1QHoa73w9pQ9GroAvxqFi2RTZWlkC0raY=",
        "3syVH2woCpOvPF0SD9Z0bu OxNe2ZgxKjTQ961L1MnA="
\r\n\r\n
AeYbsHotOpmdWAcgTo5sD8iAuSQAfnH5U6wiIGpVNJQQoYKBYrPPxAoIc1i5SHCIDS8KFFgf8iOtDq8XGizaCg==
\r\n\r\n
o9yjuKHHNJZFi0QD9K6Vpt6fP0XgXlj8z_4D-7s3CcYmuoWAh6NVtYaf_GWw_2sCrHBAA2mAEsml3thLmu50Dw==
```

# Multi-signature Pre-rotation

```
Inception Event
    "id": "did:dad:Qt27fThWoNZsa88VrTkep6H-4HA8tr54sHON1vWl6FE=",
    "changed" : "2000-01-01T00:00:00+00:00",
    "signer": [0,1],
    "signers":
        "Qt27fThWoNZsa88VrTkep6H-4HA8tr54sHON1vWl6FE=:blue",
        "Xq5YqaL6L48pf0fu7IUhL0JRaU2_RxFP0AL43wYn148=:red",
        "dZ74MLZXD-1QHoa73w9pQ9GroAvxqFi2RTZWlkC0raY=:blue",
        "3syVH2woCpOvPF0SD9Z0bu_OxNe2ZgxKjTQ961LlMnA=:red"
\r\n\r\n
AeYbsHotOpmdWAcgTo5sD8iAuSQAfnH5U6wiIGpVNJQQoYKBYrPPxAoIc1i5SHCIDS8KFFgf8iOtDq8XGizaCg==
\r\n\r\n
o9yjuKHHNJZFi0QD9K6Vpt6fP0XgXlj8z_4D-7s3CcYmuoWAh6NVtYaf_GWw_2sCrHBAA2mAEsml3thLmu50Dw==
Rotation Event
    "id": "did:dad:Qt27fThWoNZsa88VrTkep6H-4HA8tr54sHON1vWl6FE=",
    "changed": "2000-01-01T00:00:00+00:00",
    "signer": [2,3],
    "signers":
        "Qt27fThWoNZsa88VrTkep6H-4HA8tr54sHON1vWl6FE=:blue",
        "Xq5YqaL6L48pf0fu7IUhL0JRaU2_RxFP0AL43wYn148=:red",
        "dZ74MLZXD-1QHoa73w9pQ9GroAvxqFi2RTZWlkC0raY=:blue",
        "3syVH2woCpOvPF0SD9Z0bu_OxNe2ZgxKjTQ961LlMnA=:red"
        "rTkep6H-4HA8tr54sHON1vWl6FEQt27fThWoNZsa88V=:blue",
        "7IUhL0JRaU2_RxFP0AL43wYn148Xq5YqaL6L48pf0fu=:red",
\r\n\r\n
AeYbsHot0pmdWAcgTo5sD8iAuSQAfnH5U6wiIGpVNJQQoYKBYrPPxAoIc1i5SHCIDS8KFFgf8i0tDq8XGizaCg==
o9yjuKHHNJZFi0QD9K6Vpt6fP0XgXlj8z_4D-7s3CcYmuoWAh6NVtYaf_GWw_2sCrHBAA2mAEsml3thLmu50Dw==
\r\n\r\n
GpVNJQQoYKBYrPPxAoIc1i5SHCIDS8KFFgf8i0tDq8XGizaCgAeYbsHot0pmdWAcgTo5sD8iAuSQAfnH5U6wiI==
\r\n\r\n
8z_4D-7s3CcYmuoWAh6NVtYaf_GWw_2sCrHBAA2mAEsml3thLmu50Dwo9yjuKHHNJZFi0QD9K6Vpt6fP0XgXlj==
```

# COSI Collective Signatures

Typically the maximum length of a collective signature is about double the length of a non-collective signature and does not increase significantly as more signatures are added to the collective.

Collective signatures are a type of Schnorr multi-signature or Schnorr threshold signature.

Fork of LibSodium includes CoSI support for Ed255i9 signatures

Draft IETF standard for collective signatures Cost

https://tools.ietf.org/id/draft-ford-cfrg-cosi-oo.html)

https://github.com/dedis/cothority/wiki/CoSi

https://arxiv.org/pdf/1503.08768.pdf

http://dedis.cs.yale.edu/dissent/pres/160524-sp-cosi.pdf

### DDID Pre-rotation

See Paper

## Recovery

Human memory incompatible with long strings of random numbers.

Humans ill equipped for task.

# Recovery Methods

#### Physical

Shift memory burden from recalling key to recalling physical location of physical backup

Shift exploit surface from network security to physical security

Good backup method

#### Mnemonic

Shift memory burden to domain that humans are good at

#### Secret Sharing

Shift recovery burden to third parties

Requires multi-party interactive infra-structure

#### Hybrid

Minimally sufficient means

# Cryptographic Security

information-theoretic security:

cryptosystem that cannot be broken algorithimically even if the adversary has nearly unlimited computing power including quantum computing. It must be broken by brute force if at all. Brute force means that in order to guarantee success the adversary must search every combination of key or seed.

perfect security:

special case of information-theoretic security where cipher text provides no information about the key

# Practical Perfect Security

Secret Sharing or Secret Splitting:

Simple N of N

Threshold M of N

One-time Pad:

Pad = string of random characters used to encode/decode message. Pad is as long as message.

Can be performed manually

# Brute Force Cryptographic Strength

Systems with perfect security the critically parameter is the number of bits of entropy needed to resist any brute force attack

Entropy in bits per Shannon information theory

number of bits = log<sub>2</sub>(number of random possibilities)

Convention is that 128 bits is sufficient =  $2^{128}$  possibilities or random choices

1 million super computers each performing 1 quadrillion trials per second will take 8,589,934,592 years for brute force exploit

#### Mnemonics

mnemonic = device or technique to aid human memory

memory task = recall 128-bit random number as a key or seed

recovery complicated need to recover private key rotation history not just a single private key use a 128-bit random number as a seed (CSPRNG) to a cryptosystem that hides and recovers the whole rotation history, such as, one-time pad.

Cryptographic strength governed by seed

elaborative encoding: use additional sensory cues to enhance recall, such as, story, path, spatial relationships, auditory, visual, color, fantastical juxtaposition

Example: method of loci or memory palace.

#### DiceWare

randomly select words from 7776 word list

 $log_2(7776) = 12.9 bits per word$ 

128 bits = 10 words in order

Non-trivial memory load. Requires rehearsal.

# Threshold Secret Sharing

Perfect security

M of N shards to recover

Requires multi-party interaction to recover shards

Not self-contained

Good backup method

# One-time pad

Perfect security

Book cypher

Hiding in plain sight

Recovery task means no need to exchange one-time pad = self-contained

Complication with private key rotation history. Recover multiple keys

Modify to recovery of seed for CSPRNG that deterministically generates a one-time pad.

Strength bounded by seed entropy.

Rotation history much much shorter than CSPRNG period

# Simple Secret-Splitting

Create multiple random shards

XOR shards together to generate/recover secret

Order independent combination of shards. Don't have to recall order

Recall task is N shards in any order

Cryptographic strength is now the entropy in bits due to the number of combinations of N possibilities per shard taken K shards at a time.

$$C(N,K) = \begin{pmatrix} N \\ K \end{pmatrix} = \frac{N!}{K!(N-K)!}$$

$$\left(\frac{N}{K}\right)^{K} \leq \frac{N!}{K!(N-K)!} \leq \left(\frac{Ne}{K}\right)^{K}$$

# Example

Suppose each shard has 35 bits of entropy = 235 possibilities = billions

Four shards combined with secret splitting (XOR) means combinations of 235 taken 4 at a time

$$\left(\frac{N}{K}\right)^{K}, N = 2^{35}, K = 4 = 2^{2}$$

$$\left(\frac{2^{35}}{2^2}\right)^4 = \left(2^{33}\right)^4 = 2^{132} = 132 \text{ bits of entropy}$$

# Hybrid Recovery

- Mneumonic for each of multiple shards.
- In combination shards provide 128 bits of entropy for seed
- Seed used with CSPRNG to generate one-time pad
- One-time pad used to encrypt private key rotation history
- Because method employs *near* perfect security the encrypted private key rotation history may be stored in the open. Depends on value of data controlled by key
- Recovery task is to recall shards, combine shards to recover seed, use seed to reproduce onetime pad, use pad to decrypt key rotation history
- Remaining challenge is finding sources of mnemonically recoverable sources of seed material.

# Hiding seed shards in plain sight

Github.com repository commit hashes. (20 byte SHA-1 hash)

80 million repositories \* 1000 commits per repository = 80 billion choices. Four choices > 128 bits

Recall project name, repository name, commit date, which commit on a given data

Mnemonic is GitHub search tags

Flickr.com Image Database (hash of display image scraped from site)

10 billion images Four choices > 124 bits

Mnemonic is Flickr search tags and image contents

FamilySearch.org genelogical database (hash of primary record fields)

6 billion records each with 7 event types. Four choices > 128 bits

Mnemonic is name, event type, event place, event date

Google Maps world database (hash of latitude, longitude)

Inhabited land area at 4 decimal places = 10 m resolution = 40 bits of entropy. 4 choices > 128 bits

Inhabited land area at 5 decimal places = 1 m resolution = 47 bits of entropy. 3 choices > 128 bits

Mnemonic is country, region, city, neighborhood and geospatial details

# Quest for the Mnemon Seed

Virtual world role play game that generates seed through playing

Play game for rehearsal and recovery. Two modes.

Maximizes elaborative encoding mnemonic potential

Minimizes mnemonic load

Locale with 256 sites =  $2^8 = 8$  bits of entropy (map)

Each site interior has 32 spots = 25 = 5 bits of entropy

Each spot has 8 actions =  $2^3$  = 3 bits of entropy

Each spot action choice = 8 bits of entropy

Three spot-actions per site in order gives a total of 8 + 3\*8 = 32 bits of entropy

Four sites each with triple spot-action choices = 128 bits of entropy

So only have to remember how to play the easiest level in a game like Legend of Zelda

Are you smarter than a ten year old





### Conclusion & Discussion