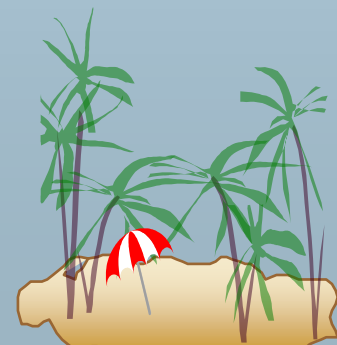


Chapter 9: Object-Relational Databases

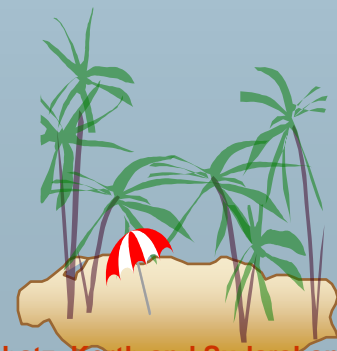
- ❑ Nested Relations
- ❑ Complex Types and Object Orientation
- ❑ Querying with Complex Types
- ❑ Creation of Complex Values and Objects
- ❑ Comparison of Object-Oriented and Object-Relational Databases





Object-Relational Data Models

- ❑ Extend the relational data model by including object orientation and constructs to deal with added data types.
- ❑ Allow attributes of tuples to have complex types, including non-atomic values such as nested relations.
- ❑ Preserve relational foundations, in particular the declarative access to data, while extending modeling power.
- ❑ Upward compatibility with existing relational languages.





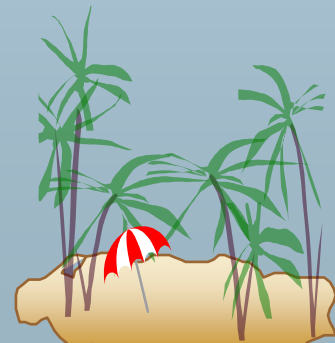
Nested Relations

□ Motivation:

- Permit non-atomic domains (atomic \equiv indivisible)
- Example of non-atomic domain: set of integers, or set of tuples
- Allows more intuitive modeling for applications with complex data

□ Intuitive definition:

- allow relations whenever we allow atomic (scalar) values
— relations within relations
- Retains mathematical foundation of relational model
- Violates first normal form.

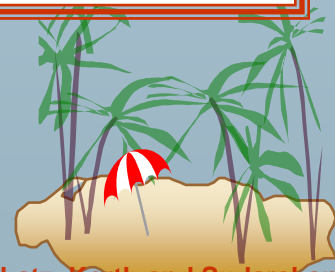




Example of a Nested Relation

- Example: library information system
- Each book has
 - title,
 - a set of authors,
 - Publisher, and
 - a set of keywords
- Non-1NF relation *books*

<i>title</i>	<i>author-set</i>	<i>publisher</i>	<i>keyword-set</i>
		(<i>name, branch</i>)	
Compilers	{Smith, Jones}	(McGraw-Hill, New York)	{parsing, analysis}
Networks	{Jones, Frick}	(Oxford, London)	{Internet, Web}



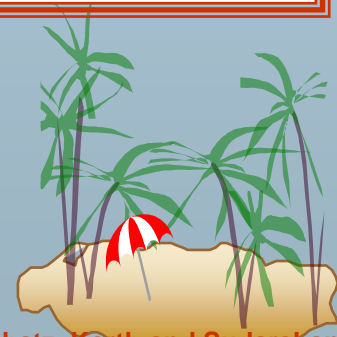


1NF Version of Nested Relation

□ 1NF version of *books*

<i>title</i>	<i>author</i>	<i>pub-name</i>	<i>pub-branch</i>	<i>keyword</i>
Compilers	Smith	McGraw-Hill	New York	parsing
Compilers	Jones	McGraw-Hill	New York	parsing
Compilers	Smith	McGraw-Hill	New York	analysis
Compilers	Jones	McGraw-Hill	New York	analysis
Networks	Jones	Oxford	London	Internet
Networks	Frick	Oxford	London	Internet
Networks	Jones	Oxford	London	Web
Networks	Frick	Oxford	London	Web

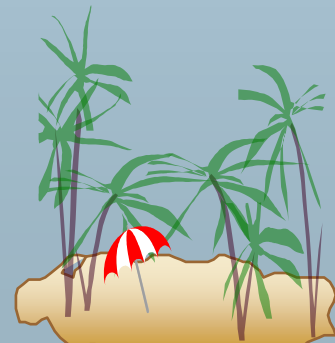
flat-books





4NF Decomposition of Nested Relation

- Remove awkwardness of *flat-books* by assuming that the following multivalued dependencies hold:
 - $title \twoheadrightarrow author$
 - $title \twoheadrightarrow keyword$
 - $title \twoheadrightarrow pub-name, pub-branch$
- Decompose *flat-doc* into 4NF using the schemas:
 - $(title, author)$
 - $(title, keyword)$
 - $(title, pub-name, pub-branch)$





4NF Decomposition of *flat-books*

<i>title</i>	<i>author</i>
Compilers	Smith
Compilers	Jones
Networks	Jones
Networks	Frick

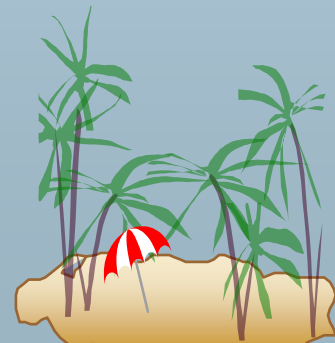
authors

<i>title</i>	<i>keyword</i>
Compilers	parsing
Compilers	analysis
Networks	Internet
Networks	Web

keywords

<i>title</i>	<i>pub-name</i>	<i>pub-branch</i>
Compilers	McGraw-Hill	New York
Networks	Oxford	London

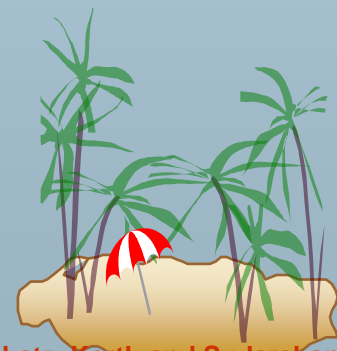
books4





Problems with 4NF Schema

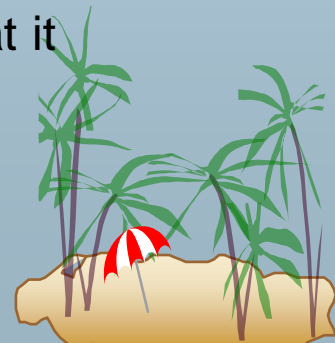
- ❑ 4NF design requires users to include joins in their queries.
- ❑ 1NF relational view *flat-books* defined by join of 4NF relations:
 - eliminates the need for users to perform joins,
 - but loses the one-to-one correspondence between tuples and documents.
 - And has a large amount of redundancy
- ❑ Nested relations representation is much more natural here.





Complex Types and SQL:1999

- Extensions to SQL to support complex types include:
 - Collection and large object types
 - ★ Nested relations are an example of collection types
 - Structured types
 - ★ Nested record structures like composite attributes
 - Inheritance
 - Object orientation
 - ★ Including object identifiers and references
- Our description is mainly based on the SQL:1999 standard
 - Not fully implemented in any database system currently
 - But some features are present in each of the major commercial database systems
 - ★ Read the manual of your database system to see what it supports
 - We present some features that are not in SQL:1999
 - ★ These are noted explicitly



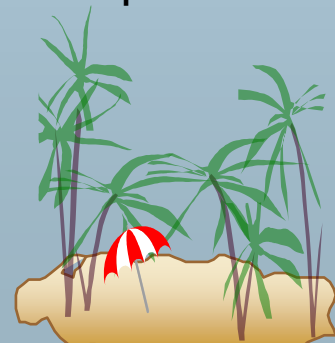


Collection Types

- Set type (not in SQL:1999)

```
create table books (  
    ....  
    keyword-set setof(varchar(20))  
    .....  
)
```

- Sets are an instance of collection types. Other instances include
 - Arrays (are supported in SQL:1999)
 - ★ E.g. *author-array* **varchar**(20) **array**[10]
 - ★ Can access elements of array in usual fashion:
 - E.g. *author-array*[1]
 - Multisets (not supported in SQL:1999)
 - ★ I.e., unordered collections, where an element may occur multiple times
 - Nested relations are sets of tuples
 - ★ SQL:1999 supports arrays of tuples





Large Object Types

□ Large object types

- **clob**: Character large objects

book-review **clob**(10KB)

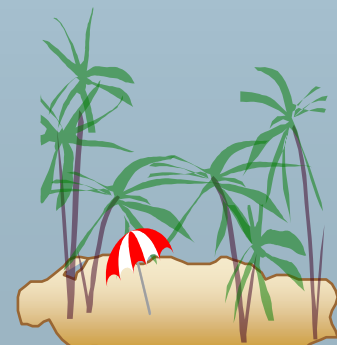
- **blob**: binary large objects

image **blob**(10MB)

movie **blob** (2GB)

□ JDBC/ODBC provide special methods to access large objects in small pieces

- Similar to accessing operating system files
- Application retrieves a **locator** for the large object and then manipulates the large object from the host language





Structured and Collection Types

- **Structured types** can be declared and used in SQL

create type *Publisher* **as**

(*name* **varchar**(20),
 branch **varchar**(20))

create type *Book* **as**

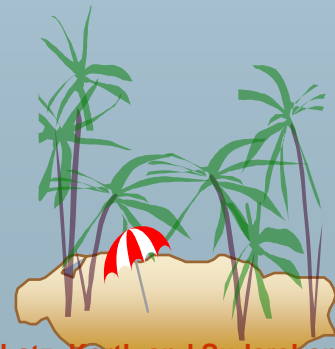
(*title* **varchar**(20),
 author-array **varchar**(20) **array** [10],
 pub-date **date**,
 publisher *Publisher*,
 keyword-set **setof**(**varchar**(20)))

- Note: **setof** declaration of keyword-set is not supported by SQL:1999
- Using an array to store authors lets us record the order of the authors

- Structured types can be used to create tables

create table *books* **of** *Book*

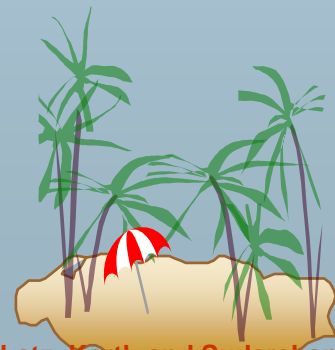
- Similar to the nested relation books, but with array of authors instead of set



Structured and Collection Types (Cont.)

- Structured types allow composite attributes of E-R diagrams to be represented directly.
- Unnamed row types can also be used in SQL:1999 to define composite attributes
 - **E.g.** we can omit the declaration of type *Publisher* and instead use the following in declaring the type *Book*

```
publisher row (name varchar(20),  
                branch varchar(20))
```
- Similarly, collection types allow multivalued attributes of E-R diagrams to be represented directly.





Structured Types (Cont.)

- We can create tables without creating an intermediate type
 - For example, the table *books* could also be defined as follows:

```
create table books  
  (title varchar(20),  
   author-array varchar(20) array[10],  
   pub-date date,  
   publisher Publisher  
   keyword-list setof(varchar(20)))
```

- Methods can be part of the type definition of a structured type:

```
create type Employee as (  
  name varchar(20),  
  salary integer)  
method giveraise (percent integer)
```

- We create the method body separately

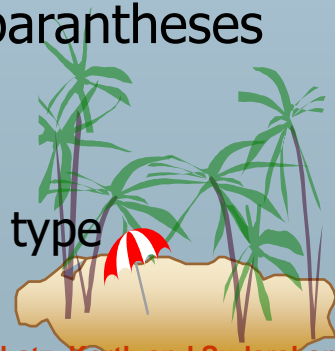
```
create method giveraise (percent integer) for Employee  
begin  
  set self.salary = self.salary + (self.salary * percent) / 100;  
end
```





Creation of Values of Complex Types

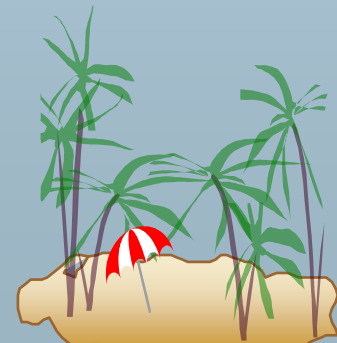
- Values of structured types are created using constructor functions
 - E.g. *Publisher*('McGraw-Hill', 'New York')
 - Note: a value is **not** an object
- SQL:1999 constructor functions
 - E.g.
create function *Publisher* (*n* **varchar**(20), *b* **varchar**(20))
returns *Publisher*
begin
 set *name*=*n*;
 set *branch*=*b*;
end
 - Every structured type has a default constructor with no arguments, others can be defined as required
- Values of **row** type can be constructed by listing values in parantheses
 - E.g. given row type **row** (*name* **varchar**(20),
 branch **varchar**(20))
 - We can assign ('McGraw-Hill', 'New York') as a value of above type





Creation of Values of Complex Types

- Array construction
array ['Silberschatz', 'Korth', 'Sudarshan']
- Set value attributes (not supported in SQL:1999)
 - **set**(v1, v2, ..., vn)
- To create a tuple of the *books* relation
('Compilers', **array**['Smith', 'Jones'],
 Publisher('McGraw-Hill', 'New York'),
 set('parsing', 'analysis'))
- To insert the preceding tuple into the relation *books*
insert into books
values
 ('Compilers', **array**['Smith', 'Jones'],
 Publisher('McGraw Hill', 'New York'),
 set('parsing', 'analysis'))





Inheritance

- Suppose that we have the following type definition for people:

```
create type Person  
  (name varchar(20),  
   address varchar(20))
```

- Using inheritance to define the student and teacher types

```
create type Student  
under Person  
  (degree varchar(20),  
   department varchar(20))
```

```
create type Teacher  
under Person  
  (salary integer,  
   department varchar(20))
```

- Subtypes can redefine methods by using **overriding method** in place of **method** in the method declaration





Multiple Inheritance

- ❑ SQL:1999 does not support multiple inheritance
- ❑ If our type system supports multiple inheritance, we can define a type for teaching assistant as follows:
create type *Teaching Assistant*
under *Student, Teacher*
- ❑ To avoid a conflict between the two occurrences of *department* we can rename them

create type *Teaching Assistant*
under
Student with (department as student-dept),
Teacher with (department as teacher-dept)

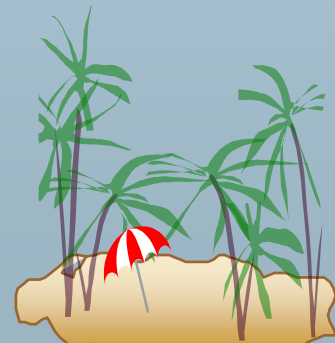




Table Inheritance

- Table inheritance allows an object to have multiple types by allowing an entity to exist in more than one table at once.
- *E.g. people table: **create table people of Person***
- We can then define the *students* and *teachers* tables as **subtables** of *people*

create table students of Student
under people
create table teachers of Teacher
under people

- Each tuple in a subtable (e.g. *students* and *teachers*) is implicitly present in its supertables (e.g. *people*)
- Multiple inheritance is possible with tables, just as it is possible with types.

create table teaching-assistants of Teaching Assistant
under students, teachers

- Multiple inheritance not supported in SQL:1999

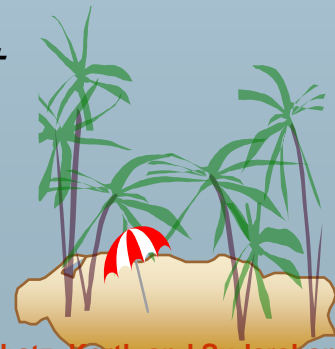




Table Inheritance: Roles

- Table inheritance is useful for modeling **roles**
- permits a value to have multiple types, without having a **most-specific type** (unlike type inheritance).
 - e.g., an object can be in the *students* and *teachers* subtables simultaneously, without having to be in a subtable *student-teachers* that is under both *students* and *teachers*
 - object can gain/lose roles: corresponds to inserting/deleting object from a subtable
- **NOTE:** SQL:1999 requires values to have a most specific type
 - so above discussion is not applicable to SQL:1999

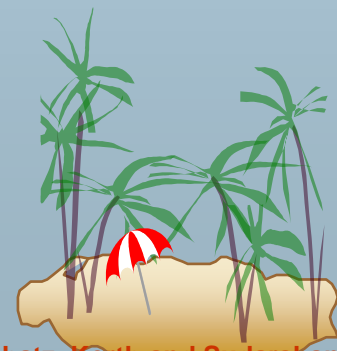




Table Inheritance: Consistency Requirements

- Consistency requirements on subtables and supertables.
 - Each tuple of the supertable (e.g. *people*) can correspond to at most one tuple in each of the subtables (e.g. *students* and *teachers*)
 - Additional constraint in SQL:1999:

All tuples corresponding to each other (that is, with the same values for inherited attributes) must be derived from one tuple (inserted into one table).

 - ★ That is, each entity must have a most specific type
 - ★ We cannot have a tuple in *people* corresponding to a tuple each in *students* and *teachers*

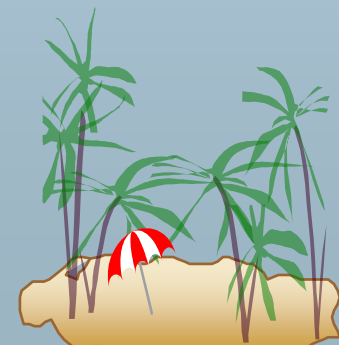
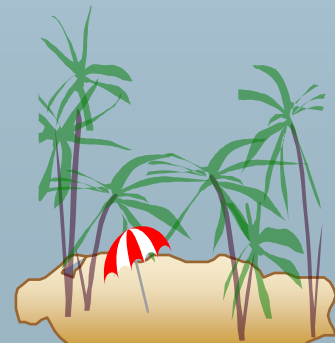




Table Inheritance: Storage Alternatives

□ Storage alternatives

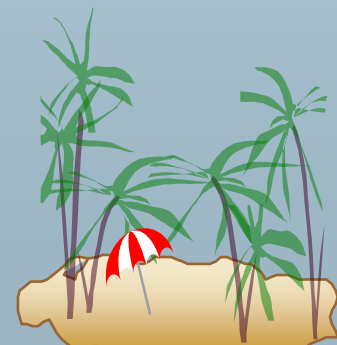
1. Store only local attributes and the primary key of the supertable in subtable
 - ★ Inherited attributes derived by means of a join with the supertable
2. Each table stores all inherited and locally defined attributes
 - ★ Supertables implicitly contain (inherited attributes of) all tuples in their subtables
 - ★ Access to all attributes of a tuple is faster: no join required
 - ★ If entities must have most specific type, tuple is stored only in one table, where it was created
 - Otherwise, there could be redundancy





Reference Types

- ❑ Object-oriented languages provide the ability to create and refer to objects.
- ❑ In SQL:1999
 - References are to tuples, and
 - References must be scoped,
 - ★ I.e., can only point to tuples in one specified table
- ❑ We will study how to define references first, and later see how to use references





Reference Declaration in SQL:1999

- E.g. define a type *Department* with a field *name* and a field *head* which is a reference to the type *Person*, with table *people* as scope

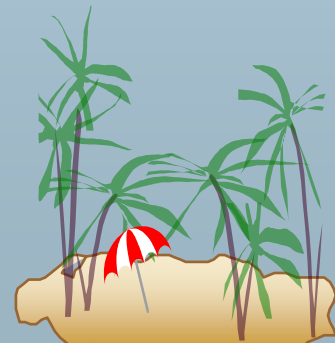
```
create type Department(  
    name varchar(20),  
    head ref(Person) scope people)
```

- We can then create a table *departments* as follows

```
create table departments of Department
```

- We can omit the declaration **scope** *people* from the type declaration and instead make an addition to the create table statement:

```
create table departments of Department  
    (head with options scope people)
```

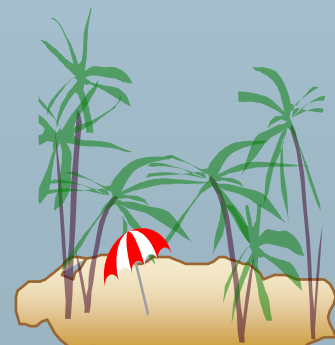




Initializing Reference Typed Values

- In Oracle, to create a tuple with a reference value, we can first create the tuple with a null reference and then set the reference separately by using the function **ref(p)** applied to a tuple variable
- E.g. to create a department with name CS and head being the person named John, we use

```
insert into departments
  values (`CS`, null)
update departments
  set head = (select ref(p)
              from people as p
              where name=`John`)
where name = `CS`
```





Initializing Reference Typed Values (Cont.)

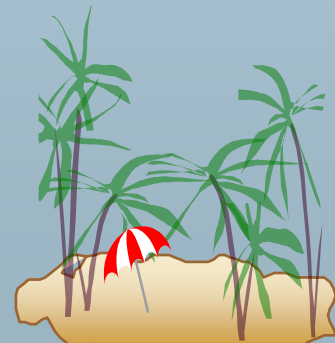
- ❑ SQL:1999 does not support the **ref()** function, and instead requires a special attribute to be declared to store the object identifier
- ❑ The self-referential attribute is declared by adding a **ref is** clause to the create table statement:

```
create table people of Person  
ref is oid system generated
```

➤ Here, *oid* is an attribute name, not a keyword.

- ❑ To get the reference to a tuple, the subquery shown earlier would use

```
select p.oid  
instead of select ref(p)
```





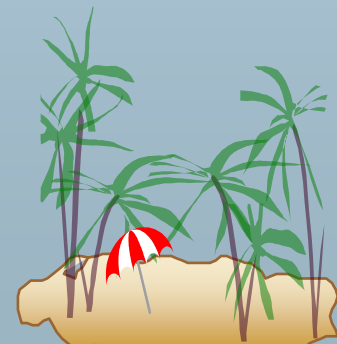
User Generated Identifiers

- SQL:1999 allows object identifiers to be user-generated
 - The type of the object-identifier must be specified as part of the type definition of the referenced table, and
 - The table definition must specify that the reference is user generated
 - E.g.

```
create type Person  
    (name varchar(20)  
    address varchar(20))  
    ref using varchar(20)  
create table people of Person  
    ref is oid user generated
```

- When creating a tuple, we must provide a unique value for the identifier (assumed to be the first attribute):

```
insert into people values  
    ('01284567', 'John', '23 Coyote Run')
```





User Generated Identifiers (Cont.)

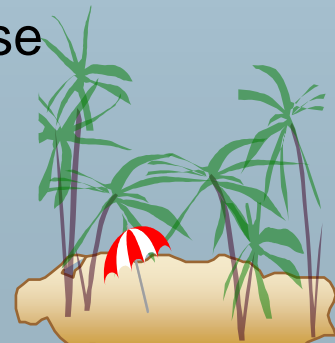
- We can then use the identifier value when inserting a tuple into *departments*
 - Avoids need for a separate query to retrieve the identifier:

E.g. **insert into** *departments*
values(`CS`, `02184567`)

- It is even possible to use an existing primary key value as the identifier, by including the **ref from** clause, and declaring the reference to be **derived**

```
create type Person  
  (name varchar(20) primary key,  
   address varchar(20))  
  ref from(name)  
create table people of Person  
  ref is oid derived
```

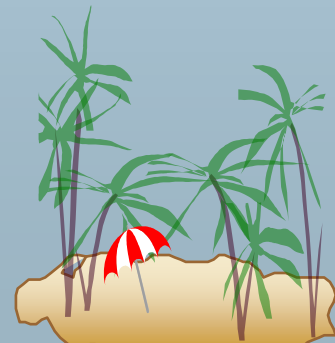
- When inserting a tuple for *departments*, we can then use
insert into *departments*
values(`CS`, `John`)





Path Expressions

- Find the names and addresses of the heads of all departments:
select *head* → *name*, *head* → *address*
from *departments*
- An expression such as “*head* → *name*” is called a **path expression**
- Path expressions help avoid explicit joins
 - If department head were not a reference, a join of *departments* with *people* would be required to get at the address
 - Makes expressing the query much easier for the user



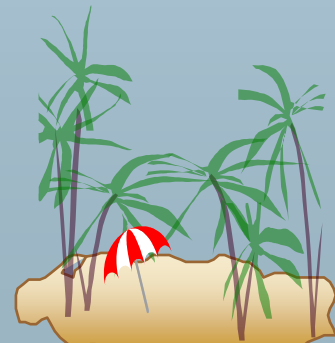


Querying with Structured Types

- Find the title and the name of the publisher of each book.

```
select title, publisher.name  
from books
```

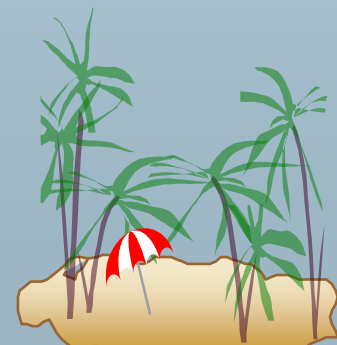
Note the use of the dot notation to access fields of the composite attribute (structured type) *publisher*





Collection-Value Attributes

- ❑ Collection-valued attributes can be treated much like relations, using the keyword **unnest**
 - The *books* relation has array-valued attribute *author-array* and set-valued attribute *keyword-set*
- ❑ To find all books that have the word “database” as one of their keywords,
select title
from books
where ‘database’ in (unnest(keyword-set))
 - Note: Above syntax is valid in SQL:1999, but the only collection type supported by SQL:1999 is the array type
- ❑ To get a relation containing pairs of the form “title, author-name” for each book and each author of the book
select B.title, A
from books as B, unnest (B.author-array) as A

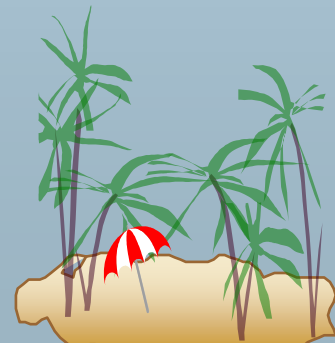




Collection Valued Attributes (Cont.)

- We can access individual elements of an array by using indices
 - E.g. If we know that a particular book has three authors, we could write:

```
select author-array[1], author-array[2], author-array[3]  
from books  
where title = 'Database System Concepts'
```



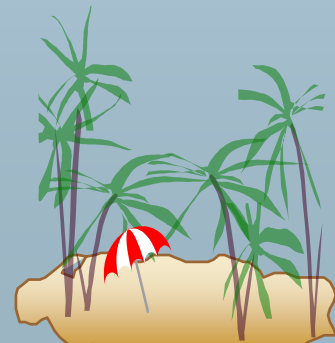


Unnesting

- The transformation of a nested relation into a form with fewer (or no) relation-valued attributes is called **unnesting**.

- E.g.

```
select title, A as author, publisher.name as pub_name,  
        publisher.branch as pub_branch, K as keyword  
  
from books as B, unnest(B.author-array) as A, unnest (B.keyword-  
list) as K
```





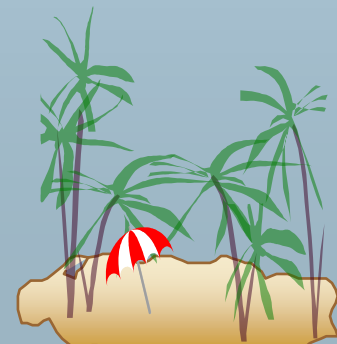
Nesting

- ❑ **Nesting** is the opposite of unnesting, creating a collection-valued attribute
- ❑ NOTE: SQL:1999 does not support nesting
- ❑ Nesting can be done in a manner similar to aggregation, but using the function `set()` in place of an aggregation operation, to create a set
- ❑ To nest the *flat-books* relation on the attribute *keyword*:

```
select title, author, Publisher(pub_name, pub_branch) as publisher,  
        set(keyword) as keyword-list  
from flat-books  
groupby title, author, publisher
```

- ❑ To nest on both authors and keywords:

```
select title, set(author) as author-list,  
        Publisher(pub_name, pub_branch) as publisher,  
        set(keyword) as keyword-list  
from flat-books  
groupby title, publisher
```



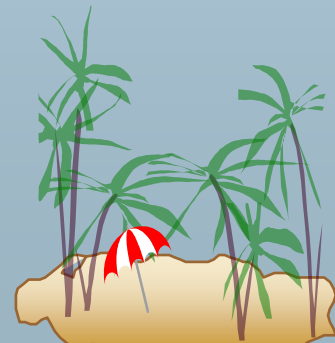


Nesting (Cont.)

- Another approach to creating nested relations is to use subqueries in the select clause.

```
select title,  
      (select author  
        from flat-books as M  
        where M.title=O.title) as author-set,  
      Publisher(pub-name, pub-branch) as publisher,  
      (select keyword  
        from flat-books as N  
        where N.title = O.title) as keyword-set  
from flat-books as O
```

- Can use **orderby** clause in nested query to get an ordered collection
 - Can thus create arrays, unlike earlier approach





Functions and Procedures

- SQL:1999 supports functions and procedures
 - Functions/procedures can be written in SQL itself, or in an external programming language
 - Functions are particularly useful with specialized data types such as images and geometric objects
 - ★ E.g. functions to check if polygons overlap, or to compare images for similarity
 - Some databases support **table-valued functions**, which can return a relation as a result
- SQL:1999 also supports a rich set of imperative constructs, including
 - Loops, if-then-else, assignment
- Many databases have proprietary procedural extensions to SQL that differ from SQL:1999





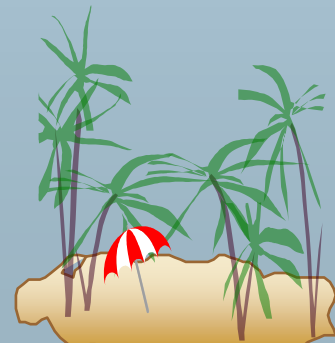
SQL Functions

- Define a function that, given a book title, returns the count of the number of authors (on the 4NF schema with relations *books4* and *authors*).

```
create function author-count(name varchar(20))  
returns integer  
begin  
    declare a-count integer;  
    select count(author) into a-count  
    from authors  
    where authors.title=name  
    return a=count;  
end
```

- Find the titles of all books that have more than one author.

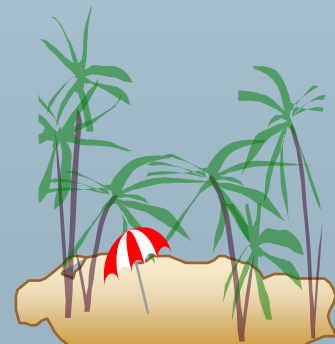
```
select name  
from books4  
where author-count(title)> 1
```





SQL Methods

- Methods can be viewed as functions associated with structured types
 - They have an implicit first parameter called **self** which is set to the structured-type value on which the method is invoked
 - The method code can refer to attributes of the structured-type value using the **self** variable
 - ★ E.g. **self.a**





SQL Functions and Procedures (cont.)

- The *author-count* function could instead be written as procedure:

```
create procedure author-count-proc (in title varchar(20),  
                                     out a-count integer)
```

```
begin
```

```
    select count(author) into a-count
```

```
    from authors
```

```
    where authors.title = title
```

```
end
```

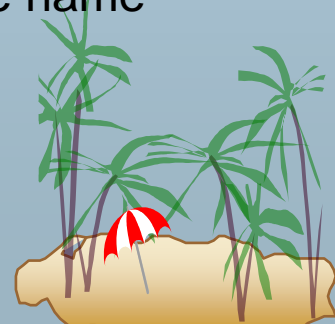
- Procedures can be invoked either from an SQL procedure or from embedded SQL, using the call statement.

- E.g. from an SQL procedure

```
declare a-count integer;
```

```
call author-count-proc(`Database systems Concepts`, a-count);
```

- SQL:1999 allows more than one function/procedure of the same name (called name **overloading**), as long as the number of arguments differ, or at least the types of the arguments differ





External Language Functions/Procedures

- ❑ SQL:1999 permits the use of functions and procedures written in other languages such as C or C++
- ❑ Declaring external language procedures and functions

```
create procedure author-count-proc(in title varchar(20),  
                                out count integer)
```

```
language C
```

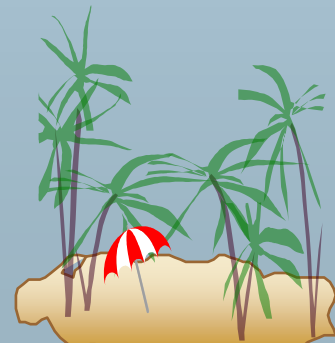
```
external name '/usr/avi/bin/author-count-proc'
```

```
create function author-count(title varchar(20))
```

```
returns integer
```

```
language C
```

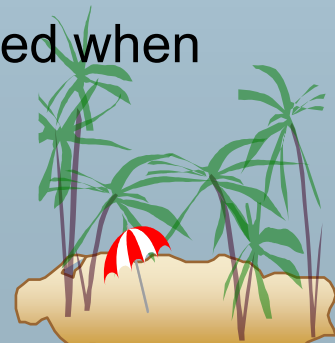
```
external name '/usr/avi/bin/author-count'
```





External Language Routines (Cont.)

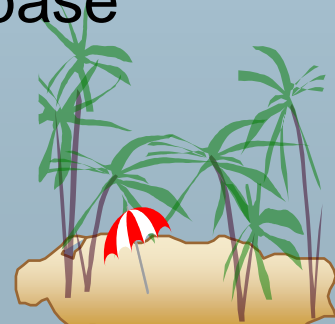
- Benefits of external language functions/procedures:
 - more efficient for many operations, and more expressive power
- Drawbacks
 - Code to implement function may need to be loaded into database system and executed in the database system's address space
 - ★ risk of accidental corruption of database structures
 - ★ security risk, allowing users access to unauthorized data
 - There are alternatives, which give good security at the cost of potentially worse performance
 - Direct execution in the database system's space is used when efficiency is more important than security





Security with External Language Routines

- To deal with security problems
 - Use **sandbox** techniques
 - ★ that is use a safe language like Java, which cannot be used to access/damage other parts of the database code
 - Or, run external language functions/procedures in a separate process, with no access to the database process' memory
 - ★ Parameters and results communicated via inter-process communication
- Both have performance overheads
- Many database systems support both above approaches as well as direct executing in database system address space





Procedural Constructs

- ❑ SQL:1999 supports a rich variety of procedural constructs
- ❑ Compound statement
 - is of the form **begin ... end**,
 - may contain multiple SQL statements between **begin** and **end**.
 - Local variables can be declared within a compound statements

- ❑ While and repeat statements

```
declare  $n$  integer default 0;
```

```
while  $n < 10$  do
```

```
    set  $n = n + 1$ 
```

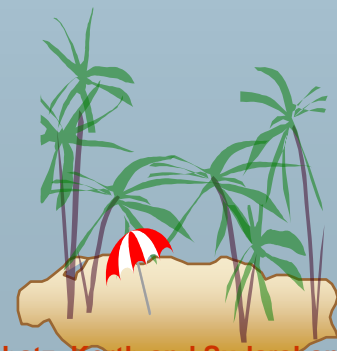
```
end while
```

```
repeat
```

```
    set  $n = n - 1$ 
```

```
until  $n = 0$ 
```

```
end repeat
```



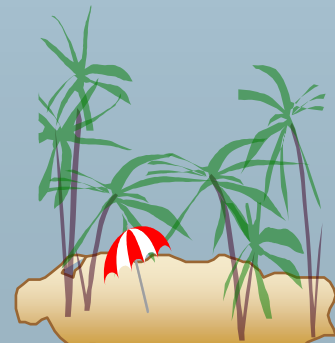


Procedural Constructs (Cont.)

□ For loop

- Permits iteration over all results of a query
- E.g. find total of all balances at the Perryridge branch

```
declare n integer default 0;  
for r as  
    select balance from account  
    where branch-name = 'Perryridge'  
do  
    set n = n + r.balance  
end for
```





Procedural Constructs (cont.)

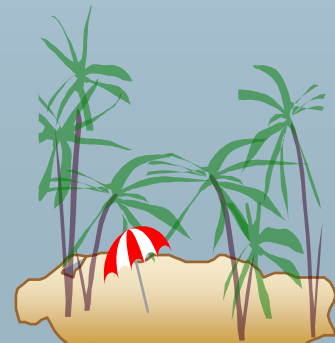
- Conditional statements (if-then-else)
E.g. To find sum of balances for each of three categories of accounts
(with balance <1000, >=1000 and <5000, >= 5000)

```
if r.balance < 1000
  then set l = l + r.balance
elseif r.balance < 5000
  then set m = m + r.balance
else set h = h + r.balance
end if
```

- SQL:1999 also supports a **case** statement similar to C case statement
- Signaling of exception conditions, and declaring handlers for exceptions

```
declare out_of_stock condition
declare exit handler for out_of_stock
begin
...
.. signal out-of-stock
end
```

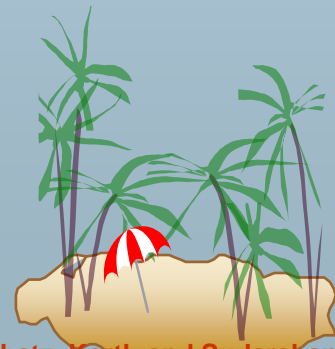
- The handler here is **exit** -- causes enclosing begin..end to be exited
- Other actions possible on exception





Comparison of O-O and O-R Databases

- Summary of strengths of various database systems:
- **Relational systems**
 - simple data types, powerful query languages, high protection.
- **Persistent-programming-language-based OODBs**
 - complex data types, integration with programming language, high performance.
- **Object-relational systems**
 - complex data types, powerful query languages, high protection.
- Note: Many real systems blur these boundaries
 - E.g. persistent programming language built as a wrapper on a relational database offers first two benefits, but may have poor performance.





Finding all employees of a manager

- Procedure to find all employees who work directly or indirectly for *mgr*
- Relation *manager(empname, mgrname)* specifies who directly works for whom
- Result is stored in *empl(name)*

```
create procedure findEmp(in mgr char(10))  
begin
```

```
    create temporary table newemp(name char(10));
```

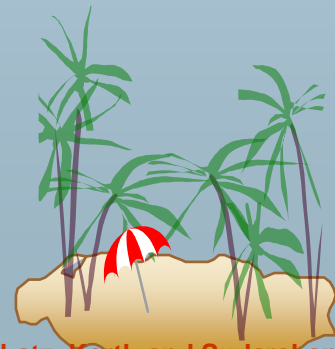
```
    create temporary table temp(name char(10));
```

```
    insert into newemp    -- store all direct employees of mgr in newemp
```

```
        select empname
```

```
        from manager
```

```
        where mgrname = mgr
```





Finding all employees of a manager(cont.)

```
repeat
  insert into empl      -- add all new employees found to empl
    select name
    from newemp;

  insert into temp      -- find all employees of people already found
    (select manager.empname
     from newemp, manager
     where newemp.empname = manager.mgrname;
    )
  except (              -- but remove those who were found earlier
    select empname
    from empl
  );

  delete from newemp;   -- replace contents of newemp by contents of temp
  insert into newemp
    select *
    from temp;
  delete from temp;

until not exists(select* from newemp) -- stop when no new employees are found
end repeat;
end
```



End of Chapter



A Partially Nested Version of the *flat-books* Relation

<i>title</i>	<i>author</i>	<i>publisher</i>	<i>keyword-set</i>
		(<i>pub-name</i> , <i>pub-branch</i>)	
Compilers	Smith	(McGraw-Hill, New York)	{parsing, analysis}
Compilers	Jones	(McGraw-Hill, New York)	{parsing, analysis}
Networks	Jones	(Oxford, London)	{Internet, Web}
Networks	Frick	(Oxford, London)	{Internet, Web}

